

On the number of co-prime-free sets.

by

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Abstract: For a variety of arithmetic properties P (such as the one in the title) we investigate the number of subsets of the positive integers $\leq x$, that have that property. In so doing we answer some questions posed by Cameron and Erdős.

1. Introduction.

In [CE] Cameron and Erdős investigated subsets of the positive integers $\leq x$ with certain given properties P ; in particular, how large such sets can be, and how many there are. The properties P that they were interested in are monotone decreasing, that is, if S has property P , and T is a subset of S , then T has property P . Thus if S is a maximal set of positive integers $\leq x$ with property P then one knows that there are $\geq 2^{|S|}$ such sets. In this paper we improve various estimates in [CE] for the number of sets satisfying certain properties P :

In Theorem 3.5 of [CE], Cameron and Erdős showed that the number of sets of positive integers $\leq x$, in which any two elements have a common factor, lies between $2^{\lfloor x/2 \rfloor}$ and $x2^{\lfloor x/2 \rfloor}$. Here we improve this to

Theorem 1. *The number of sets of integers $\leq x$, with any two elements having a common factor, is*

$$(1.1) \quad 2^{\lfloor x/2 \rfloor} + 2^{\lfloor x/2 \rfloor - N} + O\left(2^{\lfloor x/2 \rfloor - N} \exp\left(-C \frac{x}{\log^2 x \log \log x}\right)\right),$$

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for some absolute constant $C > 0$, where N , which will be defined in the proof, satisfies

$$(1.2) \quad N = (e^{-\gamma} + o(1)) \frac{x}{\log \log x},$$

and γ is the Euler-Mascheroni constant.

In Theorem 3.3 of [CE], Cameron and Erdős showed that the number of sets of positive integers $\leq x$, in which any two elements are coprime, lies between $2^{\pi(x)} e^{(1/2+o(1))\sqrt{x}}$ and $2^{\pi(x)} e^{(2+o(1))\sqrt{x}}$. Here we improve this to

Theorem 2. *The number of sets of integers $\leq x$, with every pair of elements coprime, is*

$$(1.3) \quad 2^{\pi(x)} e^{\sqrt{x}\{1+O(\log \log x/\log x)\}}.$$

In Section 4.3 of [CE], Cameron and Erdős conjectured that there are $c(s)^{x+o(x)}$ sets of integers $\leq x$, with sum of reciprocals bounded by s , for some positive constant $c(s)$. We prove a quantitative form of this conjecture here:

Theorem 3. *The number, $\nu(x)$, of sets $\{a_1, a_2, \dots, a_t\}$ of positive integers $\leq x$, with $\sum_i 1/a_i \leq s$, is*

$$(1.4) \quad c(s)^x e^{O(x^{3/4})}, \quad \text{where } c(s) = \left(1 + e^{-f(s)}\right),$$

and $f(s)$ is defined by

$$(1.5) \quad s = \int_{f(s)}^{\infty} \frac{du}{u(1+e^u)}.$$

Cameron and Erdős observed that $c(s) \geq 2^{1-e^{-s}}$, for all s . We can provide some rather more accurate estimates based on an analysis of (1.4) and (1.5): As $s \rightarrow 0$, we have

$$(1.6) \quad c(s) = 1 + s(\log(1/s) + \log \log(1/s) + O(1));$$

as $s \rightarrow \infty$, we have

$$(1.7) \quad c(s) = 2 - Ce^{-2s} + O(e^{-4s}),$$

for some constant $C \approx .8819384944\dots$

2. Sets where any two elements have a common factor.

Proof of Theorem 1: Clearly there are $2^{\lfloor x/2 \rfloor}$ such sets that contain only even numbers. We now count such sets that contain at least one odd number:

Let k be the largest integer for which $q = p_1 p_2 \dots p_k \leq x$ (where p_j is the j th smallest odd prime); by the Prime Number Theorem, $k \sim \log x / \log \log x$. Define

$$\mathcal{R} = \{n \leq x : n \text{ is divisible by } 2p_j \text{ for some } j \leq k\}.$$

Clearly any set of the form $S \cup \{q\}$, where S is a subset of \mathcal{R} , has the property that any two elements have a common factor. The number of such subsets S is $2^{|\mathcal{R}|}$, and $|\mathcal{R}| = \lfloor x/2 \rfloor - N$ where N is the number of integers $2m \leq x$ that are coprime to q . From the combinatorial sieve we obtain $N \sim \frac{\phi(q)}{q} \frac{x}{2}$, and then Mertens' Theorem implies (1.2).

Our proof of the rest of Theorem 1 is based on the ideas of Pomerance given in Cameron and Erdős [CE]:

We start by ordering the odd numbers $\leq x$ as $q = m_1, m_2, \dots$ so that

$$\frac{\phi(m_1)}{m_1} \leq \frac{\phi(m_2)}{m_2} \leq \frac{\phi(m_3)}{m_3} \leq \dots$$

Define $f_i(x)$ to be the number of sets of integers $\leq x$, that contain m_i but not m_{i+1}, m_{i+2}, \dots , and for which every pair of integers in the set have a common factor.

If $\frac{\phi(m_i)}{m_i} \geq 2/3$ then

$$\#\{n \leq x : (n, m_i) > 1\} \leq \sum_{p|m_i} \frac{x}{p} \leq x \sum_{p|m_i} \log \left(\frac{p}{p-1} \right) = x \log \left(\frac{m_i}{\phi(m_i)} \right) \leq x \log(3/2),$$

so that $f_i(x) \leq 2^{x \log(3/2)}$, which is part of the error term in (1.1).

Define $A_m(x)$ to be the number of even integers $\leq x$, that are coprime to m , and $c_i = A_{m_i}(x)$. Clearly $f_i(x) \leq 2^{i-1+[x/2]-c_i}$, and so, in order to complete the proof of Theorem 1, we must show that

$$(2.1) \quad c_i(x) - i - N \gg \frac{x}{\log^2 x \log \log x},$$

for all $i \geq 2$ for which $\frac{\phi(m_i)}{m_i} \leq 2/3$.

Suppose that $m = pr \leq m' = p'r \leq x$ are odd, squarefree numbers, with $p < x^{1/3}$. Noting that $|A_r(x)| = |A_m(x)| + |A_r(x/p)|$, we see that

$$(2.2) \quad |A_{m'}(x)| - |A_m(x)| = |A_r(x/p) \setminus A_r(x/p')| \gg \left(\frac{1}{p} - \frac{1}{p'}\right) \frac{x}{\log \log x} \gg \frac{x}{p^2 \log \log x}.$$

Thus, given any squarefree integer $r \leq x$, $r \neq q$, we form a sequence $r = r_0, r_1, \dots, r_j = q$, as follows: If $r_i = p_1 p_2 \dots p_h$ for some $h < k$, then then let $r_{i+1} = r_i p_{h+1}$. Otherwise we construct r_{i+1} by dividing r_i by its largest prime factor, and multiplying it by the smallest odd prime that does not yet divide it.

Now as $|A_q(x)| = N$, and as the prime factors of q are all $\ll \log x$, we find that, if $i \geq 2$ then $c_i - N \gg x/\log^2 x \log \log x$ by (2.2); which implies (2.1) for $i \ll x/\log^2 x \log \log x$.

So we are left with those i for which $\frac{\phi(m_i)}{m_i} \leq 2/3$ and $i \gg x/\log^2 x \log \log x$. We first deal with those $i \leq 100x/\log \log x$:

If \mathcal{A} is any set of $< i$ odd integers $\leq x$ then

$$\frac{\phi(m_i)}{m_i} \geq \min_{\substack{n \leq x, n \text{ odd} \\ n \notin \mathcal{A}}} \frac{\phi(n)}{n}.$$

So we choose \mathcal{A} to be the set of those $[x/\log^3 x]$ odd integers $\leq x$, with the most distinct prime factors.

Hardy and Ramanujan [HR] showed that there exists an absolute constant c such that the number of integers $\leq x$ with exactly k distinct prime factors is

$$\ll \frac{x}{\log x} \frac{(\log \log x + c)^{k-1}}{(k-1)!};$$

therefore the number of integers $\leq x$ with at least $10 \log \log x$ distinct prime factors is $\ll x / \log^{10} x$. Thus, if $n \notin \mathcal{A}$ then n has $\leq 10 \log \log x$ distinct prime factors, and so

$$\frac{\phi(n)}{n} \geq \prod_{p \leq (\log \log x)^2} \left(1 - \frac{1}{p}\right) \gg \frac{1}{\log \log \log x}$$

by Mertens' Theorem. Therefore, by the combinatorial sieve,

$$c_i \gg \frac{\phi(m_i)}{m_i} x \gg \frac{x}{\log \log \log x},$$

which implies (2.1) in this range of i .

Finally we come to those i in the range $100x / \log \log x \leq i \leq x/2$, with $\frac{\phi(m_i)}{m_i} \leq 2/3$:

An immediate consequence of Proposition 4 of [PS] is that

$$i < \frac{3}{20} x \frac{\phi(m_i)}{m_i} / \left(1 - \frac{\phi(m_i)}{m_i}\right) \leq \frac{9}{20} x \frac{\phi(m_i)}{m_i}.$$

On the other hand, by the combinatorial sieve,

$$c_i \geq \left\{ \frac{1}{2} + o(1) \right\} x \frac{\phi(m_i)}{m_i} > \left\{ \frac{10}{9} + o(1) \right\} i;$$

and thus

$$c_i - i \geq \frac{i}{10} \geq \frac{10x}{\log \log x},$$

and so (2.1) is satisfied.

3. Sets where any two elements are coprime.

Proof of Theorem 2: Let $t = \pi(\sqrt{x})$. For the upper bound, note that the number of composite elements of each set is $\leq t$, as these elements must all have distinct prime factors $\leq x$. All other elements are prime, and so the number of such sets is

$$\leq 2^{\pi(x)} \sum_{i=0}^t \binom{[x]}{i} \ll 2^{\pi(x)} \frac{x^t}{t!} \ll 2^{\pi(x)} \left(\frac{ex}{t}\right)^t,$$

which gives (1.3) by the Prime Number Theorem. (The proof here is the same as in Theorem 3.3 of [CE], except that they made a computational error in the final step.)

To obtain the lower bound, we shall construct (1.3) such sets. Let

$$k = t \left(1 - \frac{\log \log x}{\log x} \right)$$

and let

$$\sqrt{x} < q_1 < q_2 < \dots < q_k$$

be the k smallest primes larger than \sqrt{x} . Note that, using the Prime Number Theorem in the form $\pi(x) = \frac{x}{\log x}(1 + O(1/\log x))$ we have

$$(3.1) \quad q_j = x^{1/2}(1 + j/t + O(1/\log x))$$

and so $q_k < 2\sqrt{x}$.

We construct our sets as follows:

Each prime in the interval $(2\sqrt{x}, x]$ is in our set or not as desired, giving $2^{\pi(x) - \pi(2\sqrt{x})}$ different options.

We may put any number of the form $p_k q_k$ in the set, where p_k is a prime less than x/q_k (giving $\pi(x/q_k)$ choices).

Then any $p_{k-1} q_{k-1}$ where p_{k-1} is a prime $\leq x/q_{k-1}$ (giving $\pi(x/q_{k-1}) - 1$ choices).

We continue in this fashion, taking, in general, any $p_{k-j} q_{k-j}$, where p_{k-j} is a prime $\leq x/q_{k-j}$, not already used as some p_{k-i} , (giving us $\pi(x/q_{k-j}) - j$ choices), for $j = 0, 1, \dots, k-1$.

Thus the number of different sets constructed is

$$2^{\pi(x) + O(\sqrt{x}/\log x)} \prod_{i=0}^k \left\{ \pi \left(\frac{x}{q_i} \right) - (k-i) \right\}.$$

Now, by (3.1),

$$\begin{aligned} \pi \left(\frac{x}{q_i} \right) - (k-i) &= \pi \left(\frac{x^{1/2}}{1 + i/t} \left(1 + O \left(\frac{1}{\log x} \right) \right) \right) - (k-i) \\ &= t \left\{ \frac{1}{1 + i/t} + O \left(\frac{1}{\log x} \right) - 1 + \frac{\log \log x}{\log x} + \frac{i}{t} \right\}, \end{aligned}$$

using the Prime Number Theorem again,

$$= t \left(\frac{\log \log x}{\log x} + \frac{(i/t)^2}{1+i/t} + O\left(\frac{1}{\log x}\right) \right) \geq \frac{t}{\log x}.$$

Therefore, the number of sets is at least

$$2^{\pi(x)+O(\sqrt{x}/\log x)} \left(\frac{t}{\log x} \right)^k,$$

which gives (1.3).

Remark: With some care it is possible to replace the $\log \log x$ in (1.3) with $(\log \log \log x)^2$, but we are currently unable to do better.

4. Sets whose sum of reciprocals is bounded.

Proof of Theorem 3: Let $y = \lfloor x^{1/4} \rfloor$, $z = \lfloor x^{1/2} \rfloor$, and $x_j = \lfloor jx/z \rfloor$ for $y \leq j \leq z$. A given set of positive integers $\leq x$, whose sum of reciprocals is bounded by s , has, say, b_j integers in the interval $(x_j, x_{j+1}]$ for $y \leq j < z$ (with $0 \leq b_j \leq x/z + 1$), and so satisfies

$$(4.1) \quad \sum_{j=y}^{z-1} \frac{b_j}{x_{j+1}} \leq s.$$

So, if the b_j are fixed with these values then the number of sets of integers $\leq x$, with precisely b_j integers from the interval $(x_j, x_{j+1}]$, is

$$(4.2) \quad \leq 2^{xy/z} \prod_{y \leq j < z} \binom{\lfloor x/z \rfloor + 1}{b_j}.$$

So define $\alpha_j := b_j / (\lfloor x/z \rfloor + 1)$ for each j . Clearly $0 \leq \alpha_j \leq 1$ for each j , and, by (4.1), they must satisfy

$$(4.3) \quad \sum_{j=y}^{z-1} \frac{\alpha_j}{j} \leq s \left(1 + O\left(\frac{1}{y}\right) \right).$$

Moreover, by Stirling's formula,

$$\binom{[x/z] + 1}{b_j} \ll \exp\left(O(\log x) - \frac{x}{z}(\alpha_j \log \alpha_j + (1 - \alpha_j) \log(1 - \alpha_j))\right).$$

Noting that there are no more than $(x/z + 1)^z$ choices for the b_j , we thus see that

$$\nu(x) \ll \exp\left(O(x^{3/4}) - \frac{x}{z} \min_{\substack{0 \leq \alpha_j \leq 1, y \leq j < z \\ (4.3) \text{ holds}}} \sum_{y \leq j < z} (\alpha_j \log \alpha_j + (1 - \alpha_j) \log(1 - \alpha_j))\right).$$

By the method of Lagrange multipliers we find that the minimum occurs when each $\alpha_j = 1/(1 + e^{A/j})$ for some constant $A > 0$.

Now

$$\begin{aligned} \sum_{j=y}^{z-1} \frac{\alpha_j}{j} &= \int_y^z \frac{dt}{t(1 + e^{A/t})} + O\left(\frac{1}{ye^{A/y}} + \frac{1}{ze^{A/z}}\right) \\ &= \int_{A/z}^\infty \frac{du}{u(1 + e^u)} + O\left(\frac{1}{ye^{A/y}} + \frac{1}{(A/y)e^{A/y}} + \frac{1}{ze^{A/z}}\right). \end{aligned}$$

To obtain equality in (4.3), we need to select $A = x^{1/2} f(s) + O(x^{1/4})$. Thus

$$\begin{aligned} - \sum_{y \leq j < z} \alpha_j \log \alpha_j + (1 - \alpha_j) \log(1 - \alpha_j) &= \sum_{y \leq j < z} \log\left(1 + e^{-A/j}\right) + \frac{A}{j} \frac{1}{(1 + e^{A/j})} \\ &= \int_y^z \left(\left(1 + e^{-A/u}\right) + \frac{A}{u(1 + e^{A/u})} \right) du + O(1). \end{aligned}$$

By the substitution $t = A/u$, and the fact that

$$\frac{d}{dt} \frac{\log(1 + e^{-t})}{t} = \frac{1}{t^2} \left(\log(1 + e^{-t}) + \frac{t}{1 + e^t} \right),$$

this last line is

$$z \log(1 + e^{-A/z}) + O(1),$$

and so (1.4) is an upper bound for $\nu(x)$.

To obtain a lower bound for $\nu(x)$ note that if we select integers b_j such that

$$(4.4) \quad \sum_{j=y}^{z-1} \frac{b_j}{x_j} \leq s,$$

then the sum of the reciprocals of any set consisting of b_j integers from the interval (x_j, x_{j+1}) , for each $y \leq j < z$, is $\leq s$. Clearly the number of such sets is

$$(4.5) \quad \geq \prod_{y \leq j < z} \binom{[x/z]}{b_j}.$$

Now the idea is to select each $b_j = x/z (1 + e^{A/j}) + O(1)$, for some constant $A > 0$, so that (4.4) is satisfied; this can be done by the choice $A = x^{1/2} f(s) + O(x^{1/4})$ (the proof being almost identical to that for the upper bound). Now, when we estimate (4.5) with this value for A , we proceed as in the upper bound, and show that (4.5) is at least (1.4).

Remarks: Cameron and Erdoš observed that any set of integers taken from $[x/e^s, x]$ has sum of reciprocals $\leq s$, and thus $c(s) \geq 2^{1-e^{-s}}$. We will derive (1.6) and (1.7):

If x is ‘large’ then

$$\int_x^\infty \frac{du}{u(1+e^u)} = \frac{1}{xe^x} \left\{ 1 + O\left(\frac{1}{x}\right) \right\},$$

which implies (1.6). On the other hand, (1.5) gives that, for $f(s) < 1$ (that is, s ‘large’),

$$(4.6) \quad s = C_1 + \frac{1}{2} \log(1/f(s)) - \int_0^{f(s)} \left(\frac{1}{1+e^u} - \frac{1}{2} \right) \frac{du}{u},$$

where

$$C_1 = \int_1^\infty \frac{du}{u(1+e^u)} + \int_0^1 \left(\frac{1}{1+e^u} - \frac{1}{2} \right) \frac{du}{u}.$$

Therefore $f(s) \asymp e^{-2s}$, and so the last term in (4.6) is $\ll f(s) \ll e^{-2s}$. Thus $f(s) = Ce^{-2s} + O(e^{-4s})$, where $C = e^{2C_1}$ (which can be computed explicitly), which implies (1.7).

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