Prime Power Divisors of the Number of $n \times n$ Alternating Sign Matrices

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Abstract

We let A(n) equal the number of $n \times n$ alternating sign matrices. From the work of a variety of sources, we know that

$$A(n) = \prod_{\ell=0}^{n-1} \frac{(3\ell+1)!}{(n+\ell)!}.$$

We find an efficient method of determining $ord_p(A(n))$, the highest power of p which divides A(n), for a given prime p and positive integer n, which allows us to efficiently compute the prime factorization of A(n). We then use our method to show that for any nonnegative integer k, and for any prime p > 3, there are infinitely many positive integers n such that $ord_p(A(n)) = k$. We show a similar but weaker theorem for the prime p = 3, and note that the opposite is true for p = 2.

1 Background

An $n \times n$ alternating sign matrix is an $n \times n$ matrix consisting entirely of 1s, 0s, and -1s with the property that the sum of the entries of each row and each column must be 1 and the signs of the nonzero entries in each row and column must alternate. For example, the matrix

$$\left(\begin{array}{rrrr} 0 & 0 & 1 & 0 \\ 1 & 0 & 0 & 0 \\ 0 & 1 & -1 & 1 \\ 0 & 0 & 1 & 0 \end{array}\right)$$

is one of the 42 4×4 alternating sign matrices.

Let A(n) denote the number of $n \times n$ alternating sign matrices. Because each permutation matrix is also an alternating sign matrix, it is clear that $A(n) \ge n!$.

Recently, Zeilberger [9] proved that

$$A(n) = \prod_{\ell=0}^{n-1} \frac{(3\ell+1)!}{(n+\ell)!}.$$
(1)

The reader interested in the development of (1) is encouraged to see the survey articles of Robbins [8] and Bressoud and Propp [4], as well as the award–winning text of Bressoud [2].

In [5], we gave a characterization of the values of n for which A(n) is odd. In this paper, we give a method for determining $ord_p(A(n))$, the highest power of a given prime p that divides A(n) for given n. As a consequence we obtain a fast method for calculating A(n) for large n. (In general, A(n)is difficult to compute using (1) if n is large.) Our main theorem however, is Theorem 1.1 which comes as a consequence of our method. Theorem 1.1 is an interesting contrast to the situation when p = 2 where the opposite is true (see [6, Corollary 3.8]). We also give a characterization of when a given prime divides A(n) based on its base p representation.

The basic machinery is given in Section 2 and was largely developed in [5], (where the authors were primarily interested in the parity of A(n)), but is greatly simplified and generalized here. As this paper was in preparation for publication, we were contacted by Cartwright and Kupka and made aware of their paper [3] which essentially duplicates our method for determining $ord_p(A(n))$. However, we include that result here anyway as our proof emphasizes the periodicity of the function $c_{p,k}(n)$ which is used in the proof of Theorem 1.1 and is somewhat shorter and easier (though Cartwright and Kupka's perspective is broader than ours).

Theorem 1.1. If p is a prime greater than 3, then for each nonnegative integer k there exist infinitely many positive integers n for which $ord_p(A(n)) = k$.

2 Calculation of A(n)

We begin by stating the following lemma which is critical to our discussion. For a proof of this lemma, see [7, Theorem 2.29].

Lemma 2.1. For any prime p and any positive integer N,

$$ord_p(N!) = \sum_{k \ge 1} \left\lfloor \frac{N}{p^k} \right\rfloor.$$

Applying Lemma 2.1 to (1) then gives us

$$ord_p(A(n)) = \sum_{\ell=0}^{n-1} \sum_{k \ge 1} \left\lfloor \frac{3\ell+1}{p^k} \right\rfloor - \sum_{\ell=0}^{n-1} \sum_{k \ge 1} \left\lfloor \frac{n+\ell}{p^k} \right\rfloor = \sum_{k \ge 1} c_{p,k}(n), \quad (2)$$

where

where

$$c_{p,k}(n) := \sum_{\ell=0}^{n-1} \left(\left\lfloor \frac{3\ell+1}{p^k} \right\rfloor + \left\lfloor \frac{\ell}{p^k} \right\rfloor \right) - \sum_{\ell=0}^{2n-1} \left\lfloor \frac{\ell}{p^k} \right\rfloor.$$
(3)

$$\left(\text{Note that } \sum_{\ell=0}^{n-1} \left\lfloor \frac{n+\ell}{p^k} \right\rfloor = \sum_{\ell=0}^{2n-1} \left\lfloor \frac{\ell}{p^k} \right\rfloor - \sum_{\ell=0}^{n-1} \left\lfloor \frac{\ell}{p^k} \right\rfloor. \right)$$

We introduce the increment

$$c_{p,k}(n+1) - c_{p,k}(n) = \left\lfloor \frac{3n+1}{p^k} \right\rfloor + \left\lfloor \frac{n}{p^k} \right\rfloor - \left\lfloor \frac{2n+1}{p^k} \right\rfloor - \left\lfloor \frac{2n}{p^k} \right\rfloor$$

of $c_{p,k}(n)$ and note that the increment is periodic of period p^k as Proposition 2.2 shows.

Proposition 2.2. Let k, n be positive integers and p be prime. Then

$$c_{p,k}(p^{k}+n+1) - c_{p,k}(p^{k}+n) = c_{p,k}(n+1) - c_{p,k}(n).$$

Proof.

$$\begin{aligned} c_{p,k}(p^{k}+n+1) - c_{p,k}(p^{k}+n) \\ &= \left\lfloor \frac{3(p^{k}+n)+1}{p^{k}} \right\rfloor + \left\lfloor \frac{p^{k}+n}{p^{k}} \right\rfloor - \left\lfloor \frac{2p^{k}+2n+1}{p^{k}} \right\rfloor - \left\lfloor \frac{2p^{k}+2n}{p^{k}} \right\rfloor \\ &= \left\lfloor 3 + \frac{3n+1}{p^{k}} \right\rfloor + \left\lfloor 1 + \frac{n}{p^{k}} \right\rfloor - \left\lfloor 2 + \frac{2n+1}{p^{k}} \right\rfloor - \left\lfloor 2 + \frac{2n}{p^{k}} \right\rfloor \\ &= \left\lfloor \frac{3n+1}{p^{k}} \right\rfloor + \left\lfloor \frac{n}{p^{k}} \right\rfloor - \left\lfloor \frac{2n+1}{p^{k}} \right\rfloor - \left\lfloor \frac{2n}{p^{k}} \right\rfloor + 3 + 1 - 2 - 2 \\ &= c_{p,k}(n+1) - c_{p,k}(n). \end{aligned}$$

If p > 2 and $0 \le n < p^k$, we can easily see that

$$c_{p,k}(n+1) - c_{p,k}(n) = \begin{cases} 0 & \text{if } 0 \le n < \left\|\frac{p^k}{3}\right\|, n = \left\lfloor\frac{p^k}{2}\right\rfloor, \text{ or } \left\|\frac{2p^k}{3}\right\| \le n < p^k \\ 1 & \text{if } \left\|\frac{p^k}{3}\right\| \le n < \frac{p^k}{2} \\ -1 & \text{if } \frac{p^k}{2} < n < \left\|\frac{2p^k}{3}\right\| \end{cases}$$
(4)

where ||t|| is the integer nearest to t. Hence, since the number of +1 increments is equal to the number of -1 increments, $c_{p,k}(p^k) = 0$ which means that $c_{p,k}(n)$ is itself periodic of period p^k . If p = 2, then equation (4) is nearly the same except that the increment at $\frac{p^k}{2}$ is -1 rather than 0.

Theorem 2.3. If $0 \le n < p^k$. Then

$$c_{p,k}(n) = max\left\{0, \frac{p^k}{2} - \left\|\frac{p^k}{3}\right\| - \left|n - \frac{p^k}{2}\right|\right\}.$$

Consequently, $c_{p,k}(n) = c_{p,k}(p^k - n)$.

Proof. The increment of $c_{p,k}(n)$ in (4) tells us that if p > 2, the maximum value of $c_{p,k}(n)$ is $\frac{p^k}{2} - \left\|\frac{p^k}{3}\right\| - \frac{1}{2}$ (when $n = \frac{p^k}{2} \pm \frac{1}{2}$) and that it decrements by 1 moving away from these values in either direction. (If p = 2, the maximum value is $\frac{p^k}{2} - \left\|\frac{p^k}{3}\right\|$ at $n = \frac{p^k}{2}$.) Note that $\left\|\frac{2p^k}{3}\right\| - \frac{p^k}{2} = \frac{p^k}{2} - \left\|\frac{p^k}{3}\right\|$ so we have complete symmetry about $\frac{p^k}{2}$. If n is within $\left\|\frac{p^k}{3}\right\|$ of $\frac{p^k}{2}$ then we see that $c_{p,k}(n)$ is the difference of this maximum value

(plus $\frac{1}{2}$ if $p \neq 2$), and the distance between n and $\frac{p^k}{2}$. If not, then we have passed into an interval where both the value of $c_{p,k}(n)$ and the increment is 0. Thus, we have our result.

3 Infiniteness Results

In this section, we will prove Theorem 1.1. We first mention two propositions whose proofs we omit since they are straightforward computations.

Proposition 3.1. If p is a prime and $p \equiv 1 \pmod{6}$ then

$$\left\| \frac{p^k}{3} \right\| = \left\| \frac{p^{k-i}}{3} \right\| + \left\| \frac{p^i}{3} \right\| \cdot p^{k-i} \text{ and}$$
$$\left\| \frac{2p^k}{3} \right\| = \left\| \frac{2p^{k-i}}{3} \right\| + \left(\left\| \frac{2p^i}{3} \right\| - 1 \right) \cdot p^{k-i}$$

Proposition 3.2. If p is a prime and $p \equiv 5 \pmod{6}$, then

$$\begin{split} \left\| \frac{p^{k}}{3} \right\| &= \left\| \frac{p^{k-2i}}{3} \right\| + \left\| \frac{p^{2i}}{3} \right\| \cdot p^{k-2i}, \\ \left\| \frac{2p^{k}}{3} \right\| &= \left\| \frac{2p^{k-2i}}{3} \right\| + \left(\left\| \frac{2p^{2i}}{3} \right\| - 1 \right) p^{k-2i}, \\ \left\| \frac{p^{k}}{3} \right\| &= \left\| \frac{2p^{k-2i-1}}{3} \right\| + \left(\left\| \frac{p^{2i+1}}{3} \right\| - 1 \right) p^{k-2i-1}, \text{ and} \\ \left\| \frac{2p^{k}}{3} \right\| &= \left\| \frac{p^{k-2i-1}}{3} \right\| + \left\| \frac{2p^{2i+1}}{3} \right\| \cdot p^{k-2i-1}. \end{split}$$

We are now ready to prove Theorem 1.1.

Proof of Theorem 1.1. If k = 0, then we know that $c_{p,i}(p^j) = 0$ for all positive j, so we have established that $ord_p(A(n)) = 0$ for infinitely many positive integers n. Assume then that k > 0.

positive integers n. Assume then that k > 0. If $p \equiv 1 \pmod{6}$, consider $n = \left\| \frac{p^k}{3} \right\| + 1$. Then $n = \left(\left\| \frac{p^{k-i}}{3} \right\| + 1 \right) + \left\| \frac{p^i}{3} \right\| \cdot p^{k-i}$ for $0 \le i \le k-1$, so by the periodicity of $c_{p,i}$ for each i, we know that $c_{p,k-i}(n) = c_{p,k-i} \left(\left\| \frac{p^{k-i}}{3} \right\| + 1 \right) = 1$ for $0 \le i \le k-1$. Moreover, since $n < \left\| \frac{p^j}{3} \right\|$ for all j > k, $c_{p,j}(n) = 0$ for all j > k. Hence, $ord_p(A(n)) = k$.

If $p \equiv 5 \pmod{6}$, consider $n = \left\| \frac{p^{2k}}{3} \right\| + 1$. Then $n = \left\| \frac{p^{2k-2i}}{3} \right\| + 1 + \left\| \frac{p^{2i}}{3} \right\| \cdot p^{2k-2i}$ for $0 \le i \le k-1$, so by the periodicity of $c_{p,i}$ for each i, we know that $c_{p,2k-2i}(n) = c_{p,2k-2i} \left(\left\| \frac{p^{2k-2i}}{3} \right\| + 1 \right) = 1$ for $0 \le i \le k-1$. Also, $n = \left\| \frac{p^{2k}}{3} \right\| + 1 = \left\| \frac{2p^{2k-2i-1}}{3} \right\| + 1 + \left(\left\| \frac{2p^{2i+1}}{3} \right\| - 1 \right) p^{2k-2i-1}$, so $c_{p,2k-2i-1}(n) = c_{p,2k-2i-1} \left(\left\| \frac{2p^{2k-2i-1}}{3} \right\| + 1 \right) = 0$ for $0 \le i \le k-1$. Moreover, since $n < \left\| \frac{p^{j}}{3} \right\|$ for all j > k, $c_{p,j}(n) = 0$ for all j > k. Hence, $ord_p(A(n)) = k$.

In both cases, for each positive integer k we have established the existence of at least one positive integer n such that $ord_p(A(n)) = k$. To get infinitely many, consider $n+p^j$ where j > k. Then since $n \mod p^i = n < \left\|\frac{p^i}{3}\right\|$ for $k+1 \le i \le j$ and $n+p^j < \left\|\frac{p^i}{3}\right\|$ for i > j, $c_{p,i}(n+p^j) = c_{p,i}(n)$ for all positive i. Hence, we have $ord_p(A(n+p^j)) = k$ for j > k.

We can prove a result similar to Theorem 1.1 for p = 3, although it is somewhat weaker. By considering integers of the form $3^j + 3^k$, where j > k, we can show that there are infinitely many positive integers n such that $3^{3^k}|A(n)$. Moreover, by considering integers of the form $3^j + 3^{k-1}$, where j > k, we can show that there are infinitely many positive integers n such that $3^{3^k} \not|A(n)$, but $3^{3^{k-1}}|A(n)$. In fact, for a given k, if one can show that there is some integer n such that $ord_3(A(n)) = k$, then there are infinitely many integers m such that $ord_3(A(m)) = k$, namely, the integers $m = 2 \cdot 3^j + n$ where $j > \log_3(n) + 1$. The situation is very different when p = 2. In fact, for any positive integer k, there are only finitely many positive integers n for which $ord_2(A(n)) = k$. For a proof of this claim, see [6, Corollary 3.8].

We close this section by noting that the intuition we used to prove Theorem 1.1 and Propositions 3.1 and 3.2 came from viewing Theorem 2.3 in base p. First, $n \mod p^k$ is simply the last k digits of the base prepresentation of n. Moreover, in base p > 2, $\left\|\frac{p^k}{3}\right\|$ and $\left\|\frac{2p^k}{3}\right\|$ are kdigit numbers of a very special form. For example, for p = 13 (which is congruent to 1 mod 6), $\left\|\frac{p^k}{3}\right\| = 444 \cdots 44_{13}$ while $\left\|\frac{2p^k}{3}\right\| = 888 \cdots 89_{13}$, and for p = 11 (which is congruent to 5 mod 6), we have $\left\|\frac{p^k}{3}\right\| = 3737 \cdots 37_{11}$ or $3737 \cdots 374_{11}$ and $\left\|\frac{2p^k}{3}\right\| = 7373 \cdots 74_{11}$ or $7373 \cdots 37_{11}$ depending on

whether k is even or odd. In general, the digits corresponding to primes $p \equiv 1 \pmod{6}$ are $\left\|\frac{p}{3}\right\|$, $\left\|\frac{2p}{3}\right\|$ and $\left\|\frac{2p}{3}\right\| - 1$, while the digits corresponding to primes $p \equiv 5 \pmod{6}$ are $\left\|\frac{p}{3}\right\|$, $\left\|\frac{2p}{3}\right\|$ and $\left\|\frac{p}{3}\right\| - 1$. If p = 3, $\left\|\frac{p^k}{3}\right\|$ is a power of three while $\left\|\frac{2p^k}{3}\right\|$ is twice a power of three (unless k = 1). One can use such information to write down a characterization of when A(n) is relatively prime to a given prime p. We do so below. See [3] for a different characterization.

Let the base p representation of a given positive integer n be represented by $n = p^m d_m + p^{m-1} d_{m-1} + \dots + p d_1 + d_0$ where $0 \le d_i < p$ for $i = 0, \dots, m$.

Definition 3.3. Let n be a positive integer, and p any prime greater than 3. Let m be 1 or -1 as $p \equiv 1$ or $-1 \pmod{6}$. Consider a sequence of consecutive digits $d_i, d_{i-1}, \ldots, d_{i-j}$ in the base p representation of n. We call such a sequence a determining sequence if

$$\left\|\frac{p^{j+1}}{3}\right\| < p^j d_i + p^{j-1} d_{i-1} + \dots + d_{i-j} < \left\|\frac{2p^{j+1}}{3}\right\| - \frac{1+m^{j+1}}{2}.$$

Proposition 3.4. Using the notation above, the following characterizes for which values of n, (A(n), p) > 1.

- If p = 2, see [5] and [3].
- If p = 3, (A(n), 3) = 1 if and only if whenever $d_i = 1$, $d_j = 0$ for j = 0, ..., i 1.
- If $p \equiv 1 \pmod{6}$, then (A(n), p) > 1 if and only if either there is a determining sequence for n with $j \in \{0, 1\}$ or if $d_0 = \left\|\frac{2p}{3}\right\| 1$.
- If $p \equiv 5 \pmod{6}$, then (A(n), p) > 1 if and only if either there is a determining sequence for n with $j \in \{0, 1, 2\}$ or if $d_1 = \left\|\frac{2p}{3}\right\|$ and $d_0 = \left\|\frac{p}{3}\right\| - 1$.

4 Running Time Analysis

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We provide here an analysis of the speed of the method that we and Cartwright and Kupka independently developed for calculating $c_{p,k}(n)$. As mentioned in Section 2, it is clear that

$$prd_p(A(n)) = \sum_{k \ge 1} c_{p,k}(n).$$

Since A(n) is smooth for each n and only contains primes p < 3n, we can now quickly determine the prime factorization of A(n) for large values of n. As an example, we calculated A(50000). Note that $A(50000) \approx \left(\frac{27}{16}\right)^{50000^2/2}$, so that A(50000) has a base 10 representation with approximately 284,000,000 digits. The prime factorization of A(50000) is

 $2^{11102} \cdot 3^{2848} \cdot 5^{2083} \cdot \ldots \cdot 149969^{10} \cdot 149971^{10} \cdot 149993^2.$

This value would certainly be intractable if calculated using (1). The Maple code used to calculate this result can be obtained from the authors.

Calculation of A(n) via the method we describe here is much faster than calculation using (1). Indeed, since no obvious cancellation scheme appears available in (1), calculation of A(n) via (1) will involve $O(n^2)$ operations. This is because the number of operations needed to compute n! is n, (or n-1 if one ignores the multiplication by 1). But this implies that the number of multiplications needed to compute A(n) via (1) is at least

$$0 + 3 + 6 + 9 + \ldots + 3n - 3 + (n - 1) + n + (n + 1) + \ldots + (2n - 1)$$

which is $O(n^2)$ since $1 + 2 + 3 + \ldots + n = n(n + 1)/2$. The additional multiplications needed to combine all of the separate factorials into one final number will not contribute to a larger order of magnitude, so that $O(n^2)$ is indeed the number of operations involved.

In contrast, determination of the prime factorization of A(n) via the method in Section 2 involves only O(n) computations. The calculation involves finding $ord_p(A(n))$ for each $p \leq 3n$ (since any prime greater than 3n will yield $c_{p,k}(n) = 0$ for all k), and thus we are calculating a double sum. The inner sum, which, in practice, runs from k = 1 to $\lceil \log_2(n) \rceil + 1$, requires $O(\log(n))$ computations. The outer sum, which involves all primes p less than 3n, requires $O\left(\frac{n}{\log(n)}\right)$ computations since, by the Prime Number Theorem [1], if $\pi(x)$ equals the number of primes less than or equal to x, then

$$\lim_{x \to \infty} \frac{\pi(x)}{\frac{x}{\ln(x)}} = 1$$

Thus, our algorithm requires O(n) computations since $\log(n) \cdot \frac{n}{n} = n.$

$$\log(n) \cdot \frac{n}{\log(n)} = n.$$

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