

On Two-generated Non-commutative Algebras Subject to the Affine Relation^{*}

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Abstract. We consider algebras over a field \mathbb{K} , generated by two variables x and y subject to the single relation $yx = qxy + \alpha x + \beta y + \gamma$ for $q \in \mathbb{K}^*$ and $\alpha, \beta, \gamma \in \mathbb{K}$. We prove, that among such algebras there are precisely five isomorphism classes. The representatives of these classes, which are ubiquitous operator algebras, are called model algebras. We derive explicit multiplication formulas for $y^m \cdot x^n$ in terms of standard monomials $x^i y^j$ for many algebras of the considered type. Such formulas are used in e. g. establishing formulas of binomial type and in an implementation of non-commutative multiplication in a computer algebra system. By using the formulas we also study centers and ring-theoretic properties of the non-commutative model algebras.

In this paper we study non-commutative algebras in two generators obeying single affine relation. Many operator algebras, coming from different areas of natural sciences, are built from algebras in two generators, see Sect. 2.1 for examples. One of generators, say x , often corresponds to the operator of the multiplication with the function x . Another operator, say y , corresponds to a linear operator, acting on functions in the variable x .

In the main Theorem we identify precisely five types of non-isomorphic algebras, which we call model algebras, among them. Despite the fact that many such algebras have been studied in the literature (see e. g. [3,7,2,1], many aspects and properties are too scattered in the existing literature. Another point of this note is to search systematically for closed form of multiplication formulas on monomials. Such closed forms are needed, among other, in computer algebra, where many sophisticated algorithms heavily rely on basic multiplication among monomials. It is not enough to have such formulas just for model algebras, since isomorphisms do not preserve monomials but turn them into polynomials. It turned out, that there are still several cases, where we were not able to derive

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closed formulas in terms of standard monomials. With our approach, however, one is still able to derive formulas of certain type for them.

1 Preliminaries

Let \mathbb{K} be a field. Moreover, let A be an associative \mathbb{K} -algebra and $q \in \mathbb{K}^*$. We use the following notations: $[a, b]_q := ab - q \cdot ba$ is a q -commutator of $a, b \in A$. The commutator or the Lie bracket is $[a, b] := [a, b]_1 = ab - ba$. We also write $[n] = [n]_q = \frac{q^n - 1}{q - 1}$ for the q -number, $(a; q)_n := \prod_{k=0}^{n-1} (1 - aq^k)$ for the q -Pochhammer symbol, $[n]_q^{\underline{k}} = \frac{(q^{n-k+1}; q)_k}{(1-q)^k}$ for the q -falling factorial and $\begin{bmatrix} n \\ k \end{bmatrix} = \begin{bmatrix} n \\ k \end{bmatrix}_q = \frac{[n]!}{[n-k]![k]!}$

for the q -binomial coefficient. Note, that $\begin{bmatrix} n \\ k \end{bmatrix} = 0$ for $k > n$.

Lemma 1. $\forall a, b, c \in A$ and $\lambda, \mu \in \mathbb{K}$ the following identities hold.

$$\begin{aligned} & - [a, b]_q = -q(ba - \frac{1}{q}ab) = -q[b, a]_{q^{-1}}, \quad [a, a]_q = (1-q)a^2 \\ & - [a + \lambda, b]_q = [a, b]_q + \lambda(1-q)b, \quad [a, b + \mu]_q = [a, b]_q + \mu(1-q)a \\ & - [ab, c]_q = a[b, c]_q + q \cdot [a, c]b = a[b, c] + [a, c]_q b \end{aligned}$$

In particular, $[a, b] = -[b, a]$ and $[ab, c] = a[b, c] + [a, c]b$.

We study two-generated non-commutative \mathbb{K} -algebras with affine relations

$$A(q, \alpha, \beta, \gamma) := \mathbb{K}\langle x, y \mid yx = q \cdot xy + \alpha x + \beta y + \gamma \rangle$$

for $q \in \mathbb{K}^*$ and $\alpha, \beta, \gamma \in \mathbb{K}$. The scalar q plays an important role and we distinguish two major cases. If $q = 1$, an algebra is of *Lie type*, that is it is isomorphic to a factor-algebra of the universal enveloping algebra of a finite-dimensional Lie algebra. If $q \neq 1$, in the research of *quantum algebras* one distinguishes two situations (which lead to different behaviour of algebras): either q is transcendental over some subfield $k \subset \mathbb{K}$ or q is a root of unity in \mathbb{K} . Without assumptions on q we will write $\mathbb{K}(q)$ in general (thus encompassing the case $q \in \mathbb{K}^*$ as well), while in the case $q = 1$ just \mathbb{K} will be used. For $(a, b) \in \mathbb{N}_0^2$ we call an element $x^a y^b$ a (standard) **monomial**. If an algebra A possesses a \mathbb{K} -basis, consisting of monomials, the latter basis is often called a **Poincaré-Birkhoff-Witt** basis. The following Lemma is well-known.

Lemma 2. $A(q, \alpha, \beta, \gamma)$ has $\{x^a y^b \mid (a, b) \in \mathbb{N}_0^2\}$ as a $\mathbb{K}(q)$ -basis.

Indeed, this Lemma is a consequence of the more general statement, which can be easily proved by using Gröbner bases in the free associative algebra $\mathbb{K}\langle x, y \rangle$. The latter algebra has a \mathbb{K} -basis, consisting of *words*, that is of elements from the free monoid $F = \langle x, y \rangle$. The empty word from F is written as 1 in $\mathbb{K}\langle x, y \rangle$. The free monoid F can be totally well-ordered with an ordering, which is compatible with the bilateral multiplication on F . We call an ordering, having such properties, a *monomial* ordering on $\mathbb{K}\langle x, y \rangle$. See e. g. [8,4,9] for the Gröbner bases theory for $\mathbb{K}\langle x, y \rangle$.

Lemma 3. *Let $\mathbb{K}\langle x, y \rangle$ be the free associative algebra and \prec be a monomial ordering. Consider a polynomial $p = c \cdot yx + q(x, y)$, $c \in \mathbb{K} \setminus \{0\}$ such that every word of $q(x, y)$ is smaller than yx with respect to \prec . Then $\{p\}$ is a two-sided Gröbner basis of the two-sided ideal $\langle p \rangle$.*

Proof. Since there are no words $u, v \in F$ of degree less than 2, such that $u \cdot yx = yx \cdot v$ holds, there are no generalized s -polynomials in the free algebra [4]. Hence the set $\{p\}$ is already a two-sided Gröbner basis.

Writing $A(q, \alpha, \beta, \gamma) \cong \mathbb{K}\langle x, y \rangle / \langle -yx + q \cdot xy + \alpha x + \beta y + \gamma \rangle$, we see that the previous Lemma proves Lemma 2. Consider a monomial ordering, which satisfies (a) $w \prec u$ implies $\deg w \leq \deg u$ and (b) $xy \prec yx$, for instance, degree right lexicographical ordering. Then by the previous Lemma $\{-yx + q \cdot xy + \alpha x + \beta y + \gamma\}$ is a two-sided Gröbner basis, hence, the basis of the factor algebra $A(q, \alpha, \beta, \gamma)$ is spanned by all words which do not contain yx as a subword. And such words are precisely the standard monomials.

A product of two monomials is, in general, not a monomial, but a polynomial, that is a sum of monomials.

2 Main Theorem and Applications

Theorem 1. *$A(q, \alpha, \beta, \gamma)$ is isomorphic to one of the five **model algebras**:*

1. *the commutative algebra $\mathbb{K}[x, y]$,*
2. *the first Weyl algebra $A_1 = \mathbb{K}\langle x, d \mid dx = xd + 1 \rangle$ (the algebra of linear differential operators with coefficients from $\mathbb{K}[x]$),*
3. *the shift algebra $S_1 = \mathbb{K}\langle x, s \mid sx = xs + s \rangle$ (the universal enveloping algebra of the non-abelian solvable two-dimensional Lie algebra; the algebra of linear shift operators with coefficients from $\mathbb{K}[x]$),*
4. *the q -commutative algebra $\mathbb{K}_q[x, y] := \mathbb{K}(q)\langle x, y \mid yx = q \cdot xy \rangle$ (Manin's quantum plane; the algebra of linear q -shift operators with coeff's from $\mathbb{K}(q)[x]$)*
5. *the first q -Weyl algebra $A_1^{(q)} = \mathbb{K}(q)\langle x, \partial \mid \partial x = q \cdot x\partial + 1 \rangle$ (the algebra of linear q -differential operators with coefficients from $\mathbb{K}(q)[x]$).*

Moreover, the model algebras are pairwise non-isomorphic (see Prop. 3).

In Tables 1 and 2 we write isomorphisms to model algebras and write formulas for the multiplication in every concrete class of algebras. In some cases we also write down the recurrence formulas for the coefficients in the expansion of $y^m x^n$ in terms of standard monomials $x^a y^b$. For some algebras we put simpler formulas for $y^m x$ and yx^n as well as a part of our proof.

By writing **not known yet** in the table we mean, that up to now, no explicit formula in terms of standard monomials is known to us. However, by applying an isomorphism (for instance, the one we give explicitly in the table) to the explicit multiplication formula of the corresponding model algebra (Algebra Class in the table), we obtain a non-expanded formula for any algebra in the table.

Table 1. Multiplication Formulas for Algebras of Lie Type

Algebra Type	Relation	Algebra Class
(1, 0, 0, 0)	$yx = xy$	$YX = XY$
Isomorphism: $X \rightarrow x, Y \rightarrow y$ $y^m x^n = x^n y^m$		
(1, α , 0, 0)	$yx = xy + \alpha x$	$YX = XY + Y$
Isomorphism: $X \rightarrow -\alpha^{-1}y, Y \rightarrow x$ $y^m x^n = x^n (y + n\alpha)^m = \sum_{k=0}^m \binom{m}{k} (n\alpha)^{m-k} x^n y^k,$ Coeff. recurrence: $C_k = \frac{(k+1)n\alpha}{m-k} C_{k+1}$		
(1, 0, β , 0)	$yx = xy + \beta y$	$YX = XY + Y$
Isomorphism: $X \rightarrow \beta^{-1}x, Y \rightarrow y$ $y^m x^n = (x + m\beta)^n y^m = \sum_{k=0}^n \binom{n}{k} (m\beta)^{n-k} x^k y^m,$ Coeff. recurrence: $C_k = \frac{(k+1)m\beta}{n-k} C_{k+1}$		
(1, α , β , 0)	$yx = xy + \alpha x + \beta y$	$YX = XY + Y$
Isomorphism: $X \rightarrow -\alpha^{-1}y, Y \rightarrow \alpha x + \beta y$ $yx^n = \frac{1}{\beta} ((x + \beta)^n (\alpha x + \beta y) - \alpha x^{n+1}), y^m x = \frac{1}{\alpha} ((\alpha x + \beta y)(y + \alpha)^m - \beta y^{m+1}),$ $y^m x^n =$ not known yet		
(1, 0, 0, γ)	$yx = xy + \gamma$	$YX = XY + 1$
Isomorphism: $X \rightarrow x, Y \rightarrow \gamma^{-1}y$ $yx^n = x^{n-1} (xy + n\gamma), y^m x = (xy + m\gamma) y^{m-1},$ $y^m x^n = \sum_{k=0}^n \binom{m}{k} n^k \gamma^k x^{n-k} y^{m-k} = \sum_{k=0}^{\min\{m,n\}} \frac{m!n! \gamma^k x^{n-k} y^{m-k}}{k!(m-k)!(n-k)!},$ Coeff. recurrence: $C_k = \frac{(m-k+1)(n-k+1)\gamma}{k} C_{k-1}$		
(1, α , 0, γ)	$yx = xy + \alpha x + \gamma$	$YX = XY + Y$
Isomorphism: $X \rightarrow -\alpha^{-1}y, Y \rightarrow \alpha x + \gamma$ $yx^n = x^n y + n x^{n-1} (\alpha x + \gamma), y^m x = \frac{1}{\alpha} ((\alpha x + \gamma)(y + \alpha)^m - \gamma y^m),$ $y^m x^n = \frac{1}{\alpha^n} \sum_{i=0}^n \binom{n}{i} (-\gamma)^{n-i} (\alpha x + \gamma)^i (y + i\alpha)^m$		
(1, 0, β , γ)	$yx = xy + \beta y + \gamma$	$YX = XY + Y$
Isomorphism: $X \rightarrow \beta^{-1}x, Y \rightarrow \beta y + \gamma$ $y^m x = xy^m + m y^{m-1} (\beta y + \gamma), yx^n = \frac{1}{\beta} ((x + \beta)^n (\beta y + \gamma) - \gamma x^n),$ $y^m x^n = \frac{1}{\beta^m} \sum_{i=0}^m \binom{m}{i} (-\gamma)^{m-i} (x + i\beta)^n (\beta y + \gamma)^i$		
(1, α , β , γ)	$yx = xy + \alpha x + \beta y + \gamma$	$YX = XY + Y$
Isomorphism: $X \rightarrow -\alpha^{-1}y, Y \rightarrow \alpha x + \beta y + \gamma$ $y^m x^n =$ not known yet		

Table 2. Multiplication Formulas for Quantum Algebras

Algebra Type	Commutation	Algebra Class
$(q, 0, 0, 0)$	$yx = qxy$	$YX = qXY$
Isomorphism: $X \rightarrow x$, $Y \rightarrow y$ $y^m x^n = q^{mn} x^n y^m$		
$(q, \alpha, 0, 0)$	$yx = qxy + \alpha x$	$YX = qXY$
Isomorphism: $X \rightarrow x$, $Y \rightarrow y - \alpha(1 - q)^{-1}$ $y^m x^n = x^n (q^n y + [n]\alpha)^m$		
$(q, 0, \beta, 0)$	$yx = qxy + \beta y$	$YX = qXY$
Isomorphism: $X \rightarrow x - \beta(1 - q)^{-1}$, $Y \rightarrow y$ $y^m x^n = (q^m x + [m]\beta)^n y^m$		
$(q, \alpha, \beta, 0)$	$yx = qxy + \alpha x + \beta y$	$YX = qXY$
Isomorphism: $X \rightarrow x - \beta(1 - q)^{-1}$, $Y \rightarrow y - \alpha(1 - q)^{-1}$ $y^m x = x(qy + \alpha)^m + \beta \sum_{k=1}^m y^k \alpha^{m-k} \sum_{i=0}^{k-1} \binom{m-k+i}{i} q^i$, $y^m x^n =$ not known yet		
$(q, 0, 0, \gamma)$	$yx = qxy + \gamma$	$YX = qXY + 1$
Isomorphism: $X \rightarrow x$, $Y \rightarrow \gamma^{-1} y$ $y^m x^n = \sum_{k=0}^n \begin{bmatrix} m \\ k \end{bmatrix} [n]^k q^{(n-k)(m-k)} \gamma^k x^{n-k} y^{m-k}$.		
$(q, \alpha, 0, \gamma)$	$yx = qxy + \alpha x + \gamma$	$YX = qXY + 1$
Isomorphism: $X \rightarrow \gamma^{-1} x$, $Y \rightarrow y - \alpha(1 - q)^{-1}$ $y^m x^n = \sum_{k=0}^n \sum_{j=0}^{m-k} \begin{bmatrix} n \\ k \end{bmatrix} \gamma^k \left(\frac{\alpha}{1-q} \right)^{m-j-k} c_{j,k,m,n} x^{n-k} y^j$, where $c_{j,k,m,n} = \sum_{i=0}^{m-j-k} (-1)^i \binom{m}{i+j+k} \binom{i+j}{j} [i+j+k]^k q^{(i+j)(n-k)}$		
$(q, 0, \beta, \gamma)$	$yx = qxy + \beta y + \gamma$	$YX = qXY + 1$
Isomorphism: $X \rightarrow x - \beta(1 - q)^{-1}$, $Y \rightarrow \gamma^{-1} y$ $y^m x^n = \sum_{k=0}^n \sum_{j=0}^{n-k} \begin{bmatrix} m \\ k \end{bmatrix} \gamma^k \left(\frac{\beta}{1-q} \right)^{n-j-k} c_{j,k,m,n} x^j y^{m-k}$, where $c_{j,k,m,n} = \sum_{i=0}^{n-j-k} (-1)^i \binom{n}{i+j+k} \binom{i+j}{j} [i+j+k]^k q^{(i+j)(m-k)}$		
$(q, \alpha, \beta, \gamma)$	$yx = qxy + \alpha x + \beta y + \gamma$	$YX = qXY + 1$
Isomorphism: $X \rightarrow x - \beta(1 - q)^{-1}$, $Y \rightarrow ((1 - q)y - \alpha)(\gamma(1 - q) + \alpha\beta)^{-1}$ $y^m x^n =$ not known yet		

Proof. While for some of the above cases the explicit formulas for $y^m x^n$ are rather simple (and therefore easily found), others are quite complicated and required some work. A good strategy for finding a general formula for $y^m x^n$ is to study the special cases yx^n and $y^m x$ first. Once this is done, further multiplications by y (and x , respectively) lead to the general formula. However, for the most general commutation rules (e.g., $yx = xy + \alpha x + \beta y + \gamma$), this strategy fails.

All the formulas for $y^m x$, yx^n , and $y^m x^n$ have in common that they are easily proved by induction. As an example, consider the algebra $(1, 0, \beta, \gamma)$. We have stated above that $y^m x = xy^m + my^{m-1}(\beta y + \gamma)$. For $m = 1$ this reduces just to the given commutation relation $yx = xy + \beta y + \gamma$. Now consider $y^{m+1}x = y \cdot (y^m x)$ which by induction hypothesis is $yx y^m + my^m(\beta y + \gamma) = (xy + \beta y + \gamma)y^m + my^m(\beta y + \gamma)$ which after collecting powers gives the desired formula $xy^{m+1} + (m+1)y^m(\beta y + \gamma)$. Similarly, the general formula

$$y^m x^n = \frac{1}{\beta^m} \sum_{i=0}^m \binom{m}{i} (-\gamma)^{m-i} (x + i\beta)^n (\beta y + \gamma)^i$$

can be shown (now we use induction on n). A straightforward calculation shows that this formula for $n = 1$ reduces to the one given above for $y^m x$. Thus it has to be investigated what happens after multiplying another x from the right:

$$\begin{aligned} (\beta y + \gamma)^i x &= \sum_{j=0}^i \binom{i}{j} (\beta y)^j \gamma^{i-j} x = \sum_{j=0}^i \binom{i}{j} \beta^j (xy^j + jy^{j-1}(\beta y + \gamma)) \gamma^{i-j} \\ &= x \sum_{j=0}^i \binom{i}{j} \beta^j y^j \gamma^{i-j} + \beta \left(\sum_{j=0}^i \binom{i-1}{j-1} i (\beta y)^{j-1} \gamma^{i-j} \right) (\beta y + \gamma) \\ &= x(\beta y + \gamma)^i + \beta i (\beta y + \gamma)^{i-1} (\beta y + \gamma) \\ &= (x + \beta i) (\beta y + \gamma)^i \end{aligned}$$

We have additionally checked the validity of the formulas above with our respective implementations in computer algebra systems SINGULAR:PLURAL [6] and MATHEMATICA [5].

2.1 Operator Algebras and Model Algebras

Fix a constant $c \in \mathbb{K}^*$. Then the c -shift operator acts as $s_c(f(x)) = f(x - c)$. The corresponding c -shift algebra is $\mathbb{K}\langle x, s_c \mid s_c \cdot x = x \cdot s_c - cs_c \rangle$. For $c = 1$ one recovers discrete shift operator. If $c < 0$ (resp. $c > 0$), s_c is called an *advance operator* (resp. a *time-delay operator*) in both discrete and continuous settings. The corresponding algebras are of the type $(1, 0, \beta = -c, 0)$ and thus they are isomorphic to $\mathbb{K}\langle X, Y \mid YX = XY + Y \rangle$, the model shift algebra.

Let $c = (c_1, c_2)$ for $c_i \in \mathbb{K}^*$. The c -difference operator acts as $\Delta_c(f(x)) = \frac{f(x+c_1) - f(x)}{c_2}$. The corresponding c -difference algebra is

$$\mathbb{K}\langle x, \Delta_c \mid \Delta_c \cdot x = x \cdot \Delta_c + c_1 \Delta + \frac{c_1}{c_2} \rangle.$$

For $c = (1, 1)$ one recovers discrete difference operator; for $c = (\Delta x, \Delta x)$ the first-order divided difference operator. The corresponding algebras are of the type $(1, 0, \beta = c_1, \gamma = c_1 c_2^{-1})$ and hence they are isomorphic to $\mathbb{K}\langle X, Y \mid YX = XY + Y \rangle$, the model shift algebra.

Following Chyzak and Salvy [2], the q -dilation and q -shift operators give rise to the same operator algebra, the q -commutative model algebra $\mathbb{K}_q[x, y]$. Both continuous and discrete q -difference operators [2] give rise to the algebra $\mathbb{K}(q)\langle x, y \mid yx = qxy + (q - 1)x \rangle$ of the type $(q, \alpha = q - 1, 0, 0)$. Hence it is isomorphic to the q -commutative model algebra $\mathbb{K}_q[x, y]$.

Let $c = (c_1, c_2)$ for $c_i \in \mathbb{K}(q)^*$ with $q^{c_i} \neq 0$. The c - q -differential operator acts as $\Delta_c^{(q)}(f(x)) = \frac{f(q^{c_1}x) - f(x)}{(q^{c_2} - 1)x}$. The corresponding c - q -differential algebra is

$$\mathbb{K}(q)\langle x, \Delta_c^{(q)} \mid \Delta_c \cdot x = q^{c_1}x \cdot \Delta_c + (q^{c_1} - 1) \cdot (q^{c_2} - 1)^{-1} \rangle.$$

For $c = (1, 1)$ one recovers the q -differential operator $D_q(f(x)) = \frac{f(qx) - f(x)}{qx - x}$. Otherwise, we use Table 2 and by sending $x \rightarrow X, \Delta_c^{(q)} \rightarrow Y := (q^{c_2} - 1)(q^{c_1} - 1)^{-1} \Delta_c^{(q)}$ we obtain the isomorphic algebra $\mathbb{K}(q)\langle X, Y \mid YX = q^{c_1}XY + 1 \rangle$. Let $\tilde{q} = q^{c_1}$, then the subalgebra $\mathbb{K}(\tilde{q})\langle X, Y \mid YX = \tilde{q}XY + 1 \rangle$ of the previous algebra is the first \tilde{q} -Weyl model algebra.

Consider the differentiation $y = \frac{d}{dt}$ and the operator $x(f(t)) := e^{\lambda t} \cdot f(t)$ for $\lambda \in \mathbb{K}^*$. Then the algebra, generated by x, y has the relation $yx = xy + \lambda x$ and it is isomorphic to the model shift algebra.

Of course, there are operators obeying relations, which are not affine. Consider the integration operator $I(f(x)) := \int_0^x f(t)dt$. Its relation with x reads as $Ix = xI - I^2$. Similarly, let $x = t^{-1}$ and $y = \frac{d}{dt}$. Then the relation is $yx = xy - x^2$. Both algebras can be realized as G -algebras. It is interesting to study model algebras for non-affine relations.

Remark 1. Note, that isomorphy of q -shift and q -commutative algebras does not have an analogue in the classical situation, since for $q = 1$ the model shift algebra is not isomorphic to the model commutative algebra. Thus the following question arises: is there a quantum algebra (clearly, with non-affine relation), which becomes shift model algebra in the limit $q \rightarrow 1$?

2.2 Binomial Theorems

Notation: in a noncommutative algebra A , for two elements $a, b \in A \setminus \mathbb{K}$, we define

$$[a + b]^n := \sum_{i=0}^n \binom{n}{i} a^i b^{n-i}. \text{ Respectively, we define } [a + b]_q^n := \sum_{i=0}^n \begin{bmatrix} n \\ i \end{bmatrix}_q a^i b^{n-i}.$$

Then, if x, y commute, one expresses the binomial theorem as $(x + y)^n = [x + y]^n$. Respectively, if $yx = qxy$, we obtain $(x + y)^n = [x + y]_q^n$.

Using the formulas obtained above, we can provide formulas of binomial type, which are important in applications. Among the variety of possible presentations in such formulas we aim at those, which express $(x + y)^n$ in terms of standard monomials $x^i y^j$.

In the free associative algebra $\mathbb{K}\langle a, b \rangle$, we can write $(a + b)^n = \sum_{w \in \langle a, b \rangle_n} w$, that is w run through all words of length n in the free monoid $\langle a, b \rangle$. One defines a *misordering index* [1] of w to be the number of operations, each of them exchanges two neighbour non-equal letters, needed to move all a 's to the left (thus finishing when a standard monomial has been achieved), starting from the last letter in w . For example, the misordering index of a standard monomial is 0, while the misordering index of $bbbab$ is 3, since the sequence of exchange operations is $bbbab, bbabb, babbb, abbbb$. We say also that $bbbab$ converges to $abbbb$ here. It is known, that in any algebra $A(q, \alpha, \beta, \gamma)$ the leading monomial of a polynomial $y^m \cdot x^n$ is $x^n y^m$. Hence, the coefficients of a standard monomial $x^a y^b$ of degree $a + b$ will appear from the multiplication, applied on every word, which converges to $x^a y^b$. And closed formulas for multiplication allow to perform this task symbolically.

Lemma 4. *Let $A = A_1$ be the first Weyl algebra, where $\partial x = x\partial + 1$ holds. Then the following binomial theorem takes place:*

$$(x + d)^n - [x + d]^n = \sum_{k=0}^{n-2} \sum_{j=0}^{n-k-2} \binom{n}{j} \binom{n-j}{k} g(n-j-k) x^k d^j$$

where $g(n) := (n-1)!!$, if n is even and 0 otherwise. Alternatively we can write

$$\begin{aligned} (x + d)^n - [x + d]^n &= \sum_{0 \leq k \leq n-2} \sum_{\substack{0 \leq j \leq n-k-2 \\ n-j-k \text{ even}}} \binom{n}{j} \binom{n-j}{k} (n-j-k-1)!! x^k d^j \\ &= \sum_{0 \leq k \leq n-2} \sum_{\substack{0 \leq j \leq n-k-2 \\ n-j-k \text{ even}}} \frac{n!}{j! k! \left(\frac{n-j-k}{2}\right)!} \left(\frac{1}{2}\right)^{\frac{n-j-k}{2}} x^k d^j. \end{aligned}$$

Lemma 5. *For the shift algebra S_1 , where $xs = sx + s$ holds, we obtain the following binomial theorem:*

$$(x + s)^n = [x + s]^n + \sum_{k=0}^{n-1} \sum_{j=0}^{n-k-1} \binom{n}{k} S(n-k, j) x^k s^j$$

where $S(n, k)$ denote the Stirling numbers of the second kind.

We omit the technical proofs for both Lemmas. They can be done by induction, using the multiplication formulas.

3 Application in Computer Algebra Implementation

As described in [6], a general multiplication in a non-commutative G -algebra boils down to the multiplication of $y^m \cdot x^n$ for a couple of variables x, y such that $x^a y^b$ is a standard word. In general, the polynomial $y^m \cdot x^n$ involves other variables as well, but the case, when x, y generate a subalgebra of the type

$A(q, \alpha, \beta, \gamma)$, appears often enough. Suppose from now on we are in such situation.

In [6] it has been proposed to address each pair of non-commuting and non- q -commuting variables separately. To each such pair a matrix M is assigned, such that $M_{ij} = y^i \cdot x^j$ is a polynomial, written in terms of standard monomials. There is a general multiplication algorithm, which uses matrix entries of lower degree in order to compute the higher degrees on demand.

There are several different strategies on the usage of the formulas for enhancing the polynomial multiplication. Of course, this problem barely has an analogue in the commutative case. Initialization of non-commutative relation between y and x saves the relation $yx = q \cdot xy + \alpha x + \beta y + \gamma$ as a part of data structure on the algebra, where the computations take place.

1. Faster computation, considerable memory usage: As proposed in [6], the results of all required multiplications $y^i \cdot x^j$ and the intermediate multiplications in lower degree will be saved. Due to the same principles applied for the search of previously computed elements of lower degree, the multiplication matrix will be filled with many elements. On the other hand, the intermediate elements will be reused intensively and this leads to fast arithmetics in the algebra.

2. Saving memory, slower computation: All required multiplications $y^i \cdot x^j$ will be done according to the formulas, the results will not be saved for the future use. Thus this way uses the least amount of memory, but can take much longer, especially if many multiplications are requested repeatedly.

3. Mixing 1 and 2 and using formulas: Computing, by utilizing formulas, the requested elements and storing them into the multiplication matrix eliminates the need to compute and store intermediate elements from the approach 1. Storing the demanded elements increases the chances for the future reuse of matrix entries. Still, there are more possibilities to develop strategies by mixing both approaches and working with multiplication matrices dynamically, like keeping (e. g. by periodic cleaning) higher degree part of the matrix sparse while being as dense as possible in the lower degree part. But the question, how to determine the value, which distinguishes high degree from low degree, is open. At last, but not at least, we have experimented with counting the requests to each matrix entry, thus having a metric for the usability of every single entry. This is useful while following the strategy, which uses periodic cleaning of multiplication matrices.

Experiments. Let us do experiments with the most general case: $A = A(q, \alpha, \beta, \gamma)$, where q, α, β, γ are transcendental over the base field. As described above, in general the product $y^a \cdot x^b$ is computed (inductively) either as $y \cdot (y^{a-1} \cdot x^b)$ or as $(y^a \cdot x^{b-1}) \cdot x$. Let us consider the products $y^i \cdot x$ and $y \cdot x^i$. The determination of the computational method for these products can be made during run-time by analyzing given q, α, β, γ . Both $y^i \cdot x$ and $y \cdot x^i$ are of the same length (with $2(i+1)$ terms), with the same leading term and of the same internal byte-size. Counting the byte-size of both expressions for $i = 1..10, 15, 20$ we obtain 8, 21, 40, 65, 96, 133, 176, 225, 280, 341, 736, 1281. Indeed this sequence

coincides with octagonal numbers⁴ shifted by 1. Hence the byte-size $s(i)$ of $y^i \cdot x$ is $3i^2 + 5i + 2$.

Further on we look for some computation-specific patterns, for instance, during the computation of a left Gröbner basis. We use the implementation of *slim Gröbner basis* algorithm in SINGULAR for highly resource-demanding and tasks like the computation of Bernstein-Sato polynomials with two different algorithms. With the latter algorithms one computes in the tensor product of model algebras (Weyl and shift algebras), what suggests using formulas. By caching below we mean the use of the multiplication table for saving once computed elements.

We experiment with the following strategies: 1. Using formulas and caching 2. Using formulas without caching the results 3. Caching the results, obtained without formulas. The timings and maximal memory usage are collected in the Table 3. Since we are interested in the efficiency of the caching, we count the requests to compute every needed elementary product $y^a \cdot x^b$ as above in a separate computation. In the process of computation of Bernstein-Sato polynomial of various Reiffen curves $f(x, y) \in \mathbb{K}[x, y]$, there is a complicated computation in the tensor product of two Weyl algebras and $\mathbb{K}[s]$. We count requests to compute $y^m x^n = x^n y^m$ in Table 4 and the number of requests to compute $\partial_x^m x^n$ in the first Weyl algebra, where $[\partial_x, x] = 1$ holds, in Table 5. All the data is available online from <http://www.mathematik.uni-kl.de/~motsak/ncSATests>. Timings and memory are given in seconds resp. in Kb. The tests were run on a PC running 64 bit Arch Linux 2.6.38, having 16 GB RAM and Intel Core i7 CPU 860 at 2.80GHz (4 Cores/8 Threads).

We gain some speedup by using both caching and formulas. It can be enhanced by optimizing the way of caching, especially for algebras with many variables.

Table 3. Time and Memory Comparison for Different Multiplication Strategies

Name	Cache + Formulas		Formulas only		Cache only	
	time	memory	time	memory	time	memory
reiffen11-mod	8.04	9.912	8.14	11.999	7.97	9.912
reiffen45-3-ann	8525.40	518.262	8517.73	523.502	8547.37	518.261
reiffen45-6-ann	158.25	46.509	160.20	46.506	160.54	46.509
reiffen57-mod	36.16	11.422	39.28	11.415	36.41	11.422
reiffen59-mod	80.03	15.759	87.90	15.746	81.03	15.759
reiffen67-mod	216.81	45.369	261.30	45.361	222.15	45.369
reiffen68-mod	117.58	142.254	141.97	138.494	118.31	142.254
reiffen76-mod	389.54	99.026	630.48	99.026	389.42	99.026
reiffen86-mod	298.83	69.610	453.93	69.610	297.81	69.610

⁴ <http://oeis.org/A000567>

Table 4. Number of Requests for $y^m x^n = x^n y^m$

m	n	1	2	3	4	5	6	7	8	9	10
1	23711	18629	17628	14796	8368	2899	2444	1315	296	186	32
2	8264	4952	4806	4947	2952	728	715	549	47	26	0
3	4900	3002	3233	3202	1577	286	277	237	18	9	0
4	2084	1230	1268	1189	585	104	82	60	6	3	0
5	215	155	118	131	127	59	48	37	2	1	0
6	62	45	30	30	26	6	3	0	0	0	0
7	19	14	9	9	8	2	1	0	0	0	0
8	3	2	1	1	1	0	0	0	0	0	0

Table 5. Number of Requests for $\partial_x^m x^n = x^n \partial_x^m + \dots$

m	n	1	2	3	4	5	6	7	8	9	10	11
1	27345	22324	21914	20484	14636	5702	4076	3104	1515	1005	563	164
2	12627	10267	9799	9219	6910	2592	1888	1523	718	455	246	72
3	4271	3319	2904	2895	2544	942	763	691	300	181	90	26
4	1149	872	604	659	780	277	273	275	105	58	24	6
5	247	203	50	79	224	54	65	83	26	11	0	0

4 Centers and Ring-Theoretic Properties of Model Algebras

By using the formulas, we compute explicitly the centers of non-commutative model algebras, depending on the ground field \mathbb{K} . Recall, that for some $f \in A$ one defines the centralizer subalgebra $C(f) = \{a \in A \mid fa = af\} \supseteq \mathbb{K}[f]$.

Proposition 1. *For the algebras of Lie type one has*

- If $\text{char } \mathbb{K} = 0$, $Z(A_1) = \mathbb{K}$ and $Z(S_1) = \mathbb{K}$.
- If $\text{char } \mathbb{K} = p$, $Z(A_1) = \mathbb{K}[x^p, \partial^p]$ and $Z(S_1) = \mathbb{K}[x^p - x, s^p]$.

For the quantum algebras one has

- If q is not a root of unity, $Z(\mathbb{K}_q[x, y]) = \mathbb{K}(q)$ and $Z(A_1^{(q)}) = \mathbb{K}(q)$.
- If q is a primitive root of unity of order p over \mathbb{K} , $Z(\mathbb{K}_q[x, y]) = \mathbb{K}(q)[x^p, y^p]$ and $Z(A_1^{(q)}) = \mathbb{K}(q)[x^p, \partial^p]$.

Proof. Since all the proofs are similar, let us consider the q -Weyl algebra $A_1^{(q)}$. We compute the center as the intersection of two centralizers $Z(A) = C(x) \cap C(\partial)$. Since A is a \mathbb{Z} -graded algebra (e. g. with $\deg x = -1$, $\deg \partial = 1$), $C(x)$, $C(\partial)$ and $Z(A)$ are \mathbb{Z} -graded subalgebras. The 0-th graded part of A is $\mathbb{K}(q)[x\partial]$. For $k \in \mathbb{Z}_+$, the k -th graded part of A is $A_k = \mathbb{K}(q)[x\partial]\partial^k$ and $A_{-k} = \mathbb{K}(q)[x\partial]x^k$.

By Thm. 1, we see that $\partial^m x = q^m x \partial^m + [m]_q \partial^{m-1} \in A_{m-1}$ is homogeneous of degree $m-1$. Thus for $f = \sum_{\alpha} c_{\alpha}(x) \partial^{\alpha}$ one has $0 = fx - xf = \sum_{\alpha} c_{\alpha}(x)(\partial^{\alpha} x - x \partial^{\alpha})$. Note, that $\partial^{\alpha} x - x \partial^{\alpha} = (q^{\alpha} - 1)x \partial^{\alpha} + [\alpha]_q \partial^{\alpha-1}$ is graded. So, for all α $(q^{\alpha} - 1)x \partial^{\alpha} + [\alpha]_q \partial^{\alpha-1} = 0$, that is $q^{\alpha} = 1$ for all α . Hence $q^p = 1$ implies $C(x) = \mathbb{K}(q)[x, \partial^p]$, $C(\partial) = \mathbb{K}(q)[x^p, \partial]$ and thus $Z(A) = \mathbb{K}(q)[x^p, \partial^p]$.

It is known, that over any field $A(q, \alpha, \beta, \gamma)$ is a G -algebra (or a PBW algebra) [6,1]. Thus it is a Noetherian domain of Gel'fand-Kirillov dimension 2, which is Cohen-Macaulay and Auslander-regular [1]. However, the global homological dimension is between 1 and 2.

Proposition 2. *gl. dim $A(q, \alpha, \beta, \gamma) = 1$ if and only if $\text{char } \mathbb{K} = 0$ and $A(q, \alpha, \beta, \gamma)$ is isomorphic to the Weyl algebra.*

Proof. Let $A = A(q, \alpha, \beta, \gamma)$. Because of Cohen-Macaulay property, $\text{gl. dim } A = 2$ if and only if there exist a module M of finite dimension over $\mathbb{K}(q)$. We look for $M = A/L$ for an ideal $L \subset A$. Over $\mathbb{K}[x, y]$ and $\mathbb{K}_q[x, y]$ all 1-dimensional modules are described by ideals $\langle x - a, y - b \rangle$ for $a, b \in \mathbb{K}(q)$. In the shift algebra there are ideals $\langle x - a, s \rangle$ for $a \in \mathbb{K}$ while in the q -Weyl algebra these ideals are $\langle x - a, y - ((1 - q)a)^{-1} \rangle$ for $a \in \mathbb{K}(q)^*$.

Consider the case when A is the Weyl algebra. If $\text{char } \mathbb{K} = p > 0$, from Prop. 1 follows, that $I_p = \langle x^p, \partial^p \rangle$ is a proper two-sided ideal and A/I_p is finite dimensional, thus $\text{gl. dim } A = 2$. Now let $\text{char } \mathbb{K} = 0$. A module of Gel'fand-Kirillov dimension 0 is finite-dimensional over \mathbb{K} and hence can be considered as a representation of an algebra. Assume there exists some $m \in \mathbb{N}$ and two matrices $X, D \in \text{Mat}(m, \mathbb{K})$ such that $x \mapsto X, \partial \mapsto D$ is a representation, in other words a homomorphism of left A_1 -modules. Thus $DX - XD = \text{Id}_m$ holds and also $0 = \text{tr}(DX - XD) = m$, what is a contradiction, showing that there are no finite dimensional modules and the minimal Gel'fand-Kirillov dimension of a module over A_1 is thus 1 (which encompasses for instance holonomic left modules $A_1/\langle x \rangle$ and $A_1/\langle \partial \rangle$).

Lemma 6. *For any field \mathbb{K} , there are no nonzero \mathbb{K} -algebra homomorphisms from $A_1(\mathbb{K})$ to $\mathbb{K}_q[x, y]$ or to $\mathbb{K}[x, y]$.*

Proof. Assume there is a homomorphism of \mathbb{K} -algebras $\phi : A_1(\mathbb{K}) \rightarrow \mathbb{K}_q[x, y]$. Thus there exists $X = \phi(x), D = \phi(\partial) \in \mathbb{K}_q[x, y]$, such that $DX - XD = 1$. Write $D = \sum_{\alpha} c_{\alpha} x^{\alpha_1} y^{\alpha_2}$ for $c_{\alpha} \in \mathbb{K}$ and $\mathbb{N}_0^2 \ni \alpha = (\alpha_1, \alpha_2)$. Analogously $X = \sum_{\beta} d_{\beta} x^{\beta_1} y^{\beta_2}$. Then in $\mathbb{K}_q[x, y]$ one has $DX - XD = \sum_{\alpha, \beta} c_{\alpha} d_{\beta} (q^{\beta_1 \alpha_2} - q^{\beta_2 \alpha_1}) x^{\alpha_1 + \beta_1} y^{\alpha_2 + \beta_2}$ and the coefficient by $1 = x^0 y^0$ vanishes. In the limit $q \rightarrow 1$, that is in $\mathbb{K}[x, y]$ we obtain $DX - XD = 0$. Hence the only homomorphism from $A_1(\mathbb{K})$ to $\mathbb{K}[x, y]$ or to $\mathbb{K}_q[x, y]$ is 0.

Proposition 3. *Five model algebras are pairwise non-isomorphic over any field.*

Proof. Let $\text{char } \mathbb{K} = 0$. From Prop. 1 we see that $A_1(\mathbb{K}), S_1(\mathbb{K}), A_1^{(q)}(\mathbb{K}) \not\cong \mathbb{K}[x, y]$. By Prop. 2 and Lemma 6 we conclude $S_1(\mathbb{K}), \mathbb{K}_q[x, y], A_1^{(q)}(\mathbb{K}) \not\cong A_1(\mathbb{K})$.

For any field \mathbb{K} , $A_1^{(q)}(\mathbb{K}) \not\cong \mathbb{K}_q[x, y]$: let U, V be affine subspaces of \mathbb{K}^2 of all 1-dimensional (thus irreducible) representations of both algebras in \mathbb{K} . Then U, V are zero sets of corresponding ideals $I = \langle (1 - q)ab + 1 \rangle$ and $J = \langle (1 - q)cd \rangle = \langle c \rangle \cap \langle d \rangle$, what implies $\mathbb{K}[U] \not\cong \mathbb{K}[V]$. Since the variety $W \subset \mathbb{K}^2$ of 1-dimensional representations of S_1 is $W = \{(a, b) \mid ba = ab + b\} = \{(a, 0) \mid a \in \mathbb{K}\}$ cannot be in bijection with either U or V , S_1 is not isomorphic to $A_1^{(q)}(\mathbb{K})$ or $\mathbb{K}_q[x, y]$. Also $\mathbb{K}[x, y]$ with \mathbb{K}^2 as the variety of 1-dimensional representations is not isomorphic to other model algebras for any \mathbb{K} . Now, let $\text{char } \mathbb{K} = p$. Then $A_1(\mathbb{K})$ has finite dimensional representations since $m = \text{tr}(1_{m \times m}) = \text{tr}(DX - XD) = 0$ for a $m \times m$ representation X, D of $A_1(\mathbb{K})$. Hence $p \mid m$ and the smallest irreducible representation is in dimension p . Thus $A_1(\mathbb{K})$ cannot be isomorphic to other model algebras. The remaining cases can be proved analogously.

Future work includes the study of Ore localizations of model algebras, for which no analog of the "five models" theorem is not known yet. Groups of linear endomorphisms and sets of linear antiendomorphisms of model algebras are of interest as well.

By performing Ore localization on model algebras one gets interesting and ubiquitous algebras, for which the "five models" theorem is not known yet.

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