

# Connectivity for bridge-alterable graph classes

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## Abstract

A collection  $\mathcal{A}$  of graphs is called bridge-alterable if, for each graph  $G$  with a bridge  $e$ ,  $G$  is in  $\mathcal{A}$  if and only if  $G - e$  is. For example the class  $\mathcal{F}$  of forests is bridge-alterable. For a random forest  $F_n$  sampled uniformly from the set  $\mathcal{F}_n$  of forests on vertex set  $\{1, \dots, n\}$ , a classical result of Rényi (1959) shows that the probability that  $F_n$  is connected is  $e^{-\frac{1}{2}+o(1)}$ .

Recently Addario-Berry, McDiarmid and Reed (2012) and Kang and Panagiotou (2013) independently proved that, given a bridge-alterable class  $\mathcal{A}$ , for a random graph  $R_n$  sampled uniformly from the graphs in  $\mathcal{A}$  on  $\{1, \dots, n\}$ , the probability that  $R_n$  is connected is at least  $e^{-\frac{1}{2}+o(1)}$ . Here we give a stronger non-asymptotic form of this result, with a more straightforward proof. We see that the probability that  $R_n$  is connected is at least the minimum over  $n/3 < t \leq n$  of the probability that  $F_t$  is connected.

**Keywords:** random graph, connectivity, bridge-addable, bridge-alterable

## 1 Introduction

A collection  $\mathcal{A}$  of graphs is *bridge-addable* if for each graph  $G$  in  $\mathcal{A}$  and pair of vertices  $u$  and  $v$  in different components, the graph  $G + uv$  obtained by adding the edge (bridge)  $uv$  is also in  $\mathcal{A}$ ; that is, if  $\mathcal{A}$  is closed under adding bridges. This property was introduced in [8] (under the name ‘weakly addable’). If also  $\mathcal{A}$  is closed under deleting bridges we call  $\mathcal{A}$  *bridge-alterable*. Thus  $\mathcal{A}$  is bridge-alterable exactly when, for each graph  $G$  with a bridge  $e$ ,  $G$  is in  $\mathcal{A}$  if and only if  $G - e$  is in  $\mathcal{A}$ . The class  $\mathcal{F}$  of forests is bridge-alterable, as for example is the class of series-parallel graphs, the class of planar graphs, and indeed the class of graphs embeddable in any

given surface. All natural examples of bridge-addable classes seem to satisfy the stronger condition of being bridge-alterable.

Given a class  $\mathcal{A}$  of graphs we let  $\mathcal{A}_n$  denote the set of graphs in  $\mathcal{A}$  on vertex set  $[n] := \{1, \dots, n\}$ . Also, we use the notation  $R_n \in_u \mathcal{A}$  to mean that  $R_n$  is a random graph sampled uniformly from  $\mathcal{A}_n$  (where we assume implicitly that  $\mathcal{A}_n \neq \emptyset$ ). For a random forest  $F_n \in_u \mathcal{F}$ , a classical result of Rényi (1959) shows that, as  $n \rightarrow \infty$

$$\mathbb{P}(F_n \text{ is connected}) = e^{-\frac{1}{2}+o(1)}. \quad (1)$$

It was conjectured [9] in 2006 that, when  $\mathcal{A}$  is bridge-addable, for  $R_n \in_u \mathcal{A}$

$$\mathbb{P}(R_n \text{ is connected}) \geq e^{-\frac{1}{2}+o(1)}? \quad (2)$$

and this conjecture was strengthened (see Conjecture 1.2 of [2], Conjecture 5.1 of [1], or Conjecture 6.2 of [7]) to the non-asymptotic form

$$\mathbb{P}(R_n \text{ is connected}) \geq \mathbb{P}(F_n \text{ is connected})? \quad (3)$$

Early progress was made on conjecture (2) by Balister, Bollobás and Gerke [2, 3], and very recently Norin [11] has made significant further progress, showing that  $\mathbb{P}(R_n \text{ is connected}) \geq e^{-\frac{2}{3}+o(1)}$ ; but the full conjecture remains open.

Recently Addario-Berry, McDiarmid and Reed (2012) and Kang and Panagiotou (2013) independently proved the special case of conjecture (2) when  $\mathcal{A}$  is bridge-alterable.

**Theorem 1.1.** [1, 5] *Let  $\mathcal{A}$  be a bridge-alterable class of graphs, and let  $R_n \in_u \mathcal{A}$ . Then*

$$\mathbb{P}(R_n \text{ is connected}) \geq e^{-\frac{1}{2}+o(1)}.$$

Here we give a short and reasonably straightforward proof of the following non-asymptotic form of this result, which together with (1) gives Theorem 1.1. This is a first step towards conjecture (3) for a bridge-alterable class.

**Theorem 1.2.** *Let  $\mathcal{A}$  be a bridge-alterable class of graphs, let  $n$  be a positive integer, let  $R_n \in_u \mathcal{A}$ , and let  $F_t \in_u \mathcal{F}$  for  $t = 1, 2, \dots$ . Then*

$$\mathbb{P}(R_n \text{ is connected}) \geq \min_{n/3 < t \leq n} \mathbb{P}(F_t \text{ is connected}). \quad (4)$$

The value  $n/3$  can be increased towards  $n/2$ , see the final section of the paper.

## 2 Proof of Theorem 1.2

We use two lemmas in the proof.

**Lemma 2.1.** *Let  $\mathcal{A}$  be a bridge-alterable class of graphs, let  $n$  be a positive integer, let  $R_n \in_u \mathcal{A}$ , and let  $F_t \in_u \mathcal{F}$  for  $t = 1, 2, \dots$ . Then*

$$\mathbb{P}(R_n \text{ is connected}) \geq \min_{t=1, \dots, n} \max\{e^{-\frac{t}{n}}, \mathbb{P}(F_t \text{ is connected})\}. \quad (5)$$

**Lemma 2.2.** *For each  $n = 2, 3, \dots$*

$$\mathbb{P}(F_n \text{ is connected}) < e^{-\frac{1}{3}}.$$

To deduce Theorem 1.2 from these lemmas, observe that by Lemma 3.1, for each  $2 \leq t \leq n/3$

$$e^{-\frac{t}{n}} \geq e^{-\frac{1}{3}} \geq \mathbb{P}(F_n \text{ is connected}),$$

and so the right side in (5) is at least the right side in (4).

**Proof of Lemma 2.1** Our proof initially follows the lines of the proofs in [1] and [5] of Theorem 1.1, in that we aim to lower bound the probability of connectedness for the random graph  $F^n$  introduced below. Consider a fixed  $n \geq 2$ .

Given a graph  $G$ , let  $b(G)$  be the graph obtained by removing all bridges from  $G$ . We say  $G$  and  $G'$  are *equivalent* if  $b(G) = b(G')$ . This is an equivalence relation on graphs, and if a graph  $G$  is in  $\mathcal{A}_n$  then so is the whole equivalence class  $[G]$ . Thus  $\mathcal{A}_n$  is a union of disjoint equivalence classes. To prove the theorem we consider an arbitrary (fixed) equivalence class.

Fix a bridgeless graph  $G$  on vertex set  $[n]$  and let  $\mathcal{B} = [G]$ . Let  $G$  have  $t$  components, with  $n_1, \dots, n_t$  vertices, where  $n = \sum_{i=1}^t n_i$ . We use  $\mathbf{n} = (n_1, \dots, n_t)$  to define probabilities. First, given a forest  $F \in \mathcal{F}_t$ , let

$$\text{mass}(F) = \prod_{i=1}^t n_i^{d_F(i)},$$

where  $d_F(i)$  denotes the degree of vertex  $i$  in  $F$ . For  $\mathcal{F}' \subseteq \mathcal{F}_t$  let  $\text{mass}(\mathcal{F}') = \sum_{F \in \mathcal{F}'} \text{mass}(F)$ . Now let

$$\mathbb{P}(F^n = F) = \frac{\text{mass}(F)}{\text{mass}(\mathcal{F}_t)} \quad \text{for each } F \in \mathcal{F}_t.$$

By Lemma 2.3 of [1], for a uniformly random element  $R^{\mathcal{B}}$  of  $\mathcal{B}$ ,

$$\mathbb{P}(R^{\mathcal{B}} \text{ is connected}) = \mathbb{P}(F^{\mathbf{n}} \text{ is connected}).$$

Hence to prove the theorem it suffices to consider  $F^{\mathbf{n}}$ , and show that

$$\mathbb{P}(F^{\mathbf{n}} \text{ is connected}) \geq \max\{e^{-\frac{t}{n}}, \mathbb{P}(F_t \text{ is connected})\}. \quad (6)$$

To see this, observe that then the probability that  $R_n$  is connected is an average of values each at least the right side of (6) for some  $t$ , and so it is at least the right side in (5).

The proof of (6) breaks into two parts, and the first is standard. By Lemma 3.2 of [1], for  $i = 1, \dots, t-1$

$$\mathbb{P}(\kappa(F^{\mathbf{n}}) = i+1) \leq \frac{1}{i} \frac{t}{n} \mathbb{P}(\kappa(F^{\mathbf{n}}) = i),$$

and thus

$$\mathbb{P}(\kappa(F^{\mathbf{n}}) = i+1) \leq \frac{1}{i!} \left(\frac{t}{n}\right)^i \mathbb{P}(\kappa(F^{\mathbf{n}}) = 1).$$

Hence

$$1 = \sum_{i=0}^{t-1} \mathbb{P}(\kappa(F^{\mathbf{n}}) = i+1) \leq \sum_{i=0}^{t-1} \frac{1}{i!} \left(\frac{t}{n}\right)^i \mathbb{P}(\kappa(F^{\mathbf{n}}) = 1) < e^{\frac{t}{n}} \cdot \mathbb{P}(\kappa(F^{\mathbf{n}}) = 1)$$

and so  $\mathbb{P}(F^{\mathbf{n}} \text{ is connected}) > e^{-\frac{t}{n}}$  (as noted at the end of Section 3 of [1]).

It remains to show that

$$\mathbb{P}(F^{\mathbf{n}} \text{ is connected}) \geq \mathbb{P}(F_t \text{ is connected}). \quad (7)$$

We may assume that  $t \geq 2$ . Let  $\mathcal{T}$  be the class of trees. Then

$$\text{mass}(\mathcal{T}_t) = \prod_{i=1}^t n_i \cdot n^{t-2}. \quad (8)$$

This result is proved for example in [1] (see the proof of Lemma 4.2) and in [5], though in fact it has long been known, see Theorem 6.1 of Moon [10] (1970), and see also Problems 5.3 and 5.4 of Lovász [6]. We let  $N = \prod_{i=1}^t n_i$  and rewrite (8) as

$$\text{mass}(\mathcal{T}_t) = N \binom{n}{t}^{t-2} \cdot |\mathcal{T}_t|. \quad (9)$$

For the case  $t = 2$ ,  $\text{mass}(\mathcal{T}_2) = n_1 n_2$  and  $\text{mass}(\mathcal{F}_2) = \text{mass}(\mathcal{T}_2) + 1$ , so

$$\mathbb{P}(F^{\mathbf{n}} \text{ is connected}) = \frac{n_1 n_2}{n_1 n_2 + 1} \geq \frac{1}{2} = \mathbb{P}(F_2 \text{ is connected}).$$

Thus we may assume from now on that  $t \geq 3$ .

For each integer  $k$  with  $1 \leq k \leq t$  let  $\mathcal{F}_t^k$  be the set of forests in  $\mathcal{F}_t$  with  $k$  components. We shall show that for each such  $k$

$$\text{mass}(\mathcal{F}_t^k) \leq N \left(\frac{n}{t}\right)^{t-2} \cdot |\mathcal{F}_t^k|. \quad (10)$$

Summing over  $k$  will then give

$$\text{mass}(\mathcal{F}_t) \leq N \left(\frac{n}{t}\right)^{t-2} \cdot |\mathcal{F}_t|,$$

and so, using also (9)

$$\mathbb{P}(F^{\mathbf{n}} \text{ is connected}) = \frac{\text{mass}(\mathcal{T}_t)}{\text{mass}(\mathcal{F}_t)} \geq \frac{|\mathcal{T}_t|}{|\mathcal{F}_t|} = \mathbb{P}(F_t \text{ is connected}).$$

This will complete the proof of (7) and thus of the lemma. Hence it remains now to prove (10).

Fix an integer  $k$  with  $1 \leq k \leq t$ . Given a partition  $\mathbf{U} = (U_1, \dots, U_k)$  of  $[t]$  into  $k$  unordered sets, let  $J = J(\mathbf{U}) = \{i : |U_i| \geq 2\}$ , and let  $\mathcal{F}(\mathbf{U})$  be the set of forests in  $\mathcal{F}_t^k$  such that the  $U_i$  are the vertex sets of the  $k$  component trees. For non-empty sets  $U \subseteq [t]$ , let  $p(U) = \prod_{i \in U} n_i$  and  $s(U) = \sum_{i \in U} n_i$ . Observe that the mass of a forest is the product of the masses of its component trees, and a singleton component just gives a factor 1. Now fix a partition  $\mathbf{U} = (U_1, \dots, U_k)$  as above.

If  $J = \emptyset$  then  $\text{mass}(\mathcal{F}(\mathbf{U})) = 1 = |\mathcal{F}(\mathbf{U})|$ . Now suppose that  $J \neq \emptyset$ . Then by (9)

$$\begin{aligned} \text{mass}(\mathcal{F}(\mathbf{U})) &= \prod_{i \in J} p(U_i) \left( \frac{s(U_i)}{|U_i|} \right)^{|U_i|-2} |U_i|^{|U_i|-2} \\ &\leq N \cdot \prod_{i \in J} \left( \frac{s(U_i)}{|U_i|} \right)^{|U_i|-2} \cdot \prod_{i \in J} |U_i|^{|U_i|-2} \\ &= N \cdot \prod_{i \in J} \left( \frac{s(U_i)}{|U_i|} \right)^{|U_i|-2} \cdot |\mathcal{F}(\mathbf{U})|. \end{aligned}$$

To handle the middle factor here, we can use Jensen's inequality, since  $\log(x)$  is concave: we have

$$\begin{aligned}
& \log \prod_{i \in J} \left( \frac{s(U_i)}{|U_i|} \right)^{|U_i|-2} \\
&= (t-2) \sum_{i \in J} \frac{|U_i|-2}{t-2} \log \frac{s(U_i)}{|U_i|} \\
&\leq (t-2) \sum_{i \in J} \frac{|U_i|}{t} \log \frac{s(U_i)}{|U_i|} \quad \text{since } |U_i| \leq t \\
&\leq (t-2) \sum_{i=1}^k \frac{|U_i|}{t} \log \frac{s(U_i)}{|U_i|} \\
&\leq (t-2) \log \left( \sum_{i=1}^k \frac{|U_i|}{t} \frac{s(U_i)}{|U_i|} \right) \quad \text{since } \log \text{ is concave} \\
&= (t-2) \log \frac{n}{t}.
\end{aligned}$$

Hence in each case

$$\text{mass}(\mathcal{F}(\mathbf{U})) \leq N \left( \frac{n}{t} \right)^{t-2} |\mathcal{F}(\mathbf{U})|.$$

So, summing over partitions  $\mathbf{U} = (U_1, \dots, U_k)$  of  $[t]$ ,

$$\begin{aligned}
\text{mass}(\mathcal{F}_t^k) &= \sum_{\mathbf{U}=(U_1, \dots, U_k)} \text{mass}(\mathcal{F}(\mathbf{U})) \\
&\leq \sum_{\mathbf{U}=(U_1, \dots, U_k)} N \left( \frac{n}{t} \right)^{t-2} |\mathcal{F}(\mathbf{U})| \\
&= N \left( \frac{n}{t} \right)^{t-2} |\mathcal{F}_t^k|.
\end{aligned}$$

This completes the proof of (10), and thus the proof of Lemma 2.1.  $\square$

To prove Lemma 2.2 we will use the standard inequality

$$\left(1 - \frac{j}{n}\right)^{n-j} \geq e^{-j} \quad \text{for } 1 \leq j < n. \quad (11)$$

[To see this, fix  $j$  and let  $g(x) = (x-j) \log(1 - \frac{j}{x})$  for  $x > j$ . Then

$$g'(x) = (x-j) \left( \frac{1}{x-j} - \frac{1}{x} \right) + \log(1 - \frac{j}{x}) = \frac{j}{x} + \log(1 - \frac{j}{x}) < 0,$$

and so  $g(n)$  is decreasing for  $n > j$ . But  $g(n) \rightarrow e^{-j}$  as  $n \rightarrow \infty$ , so  $g(n) > e^{-j}$  for each  $n > j$ .]

**Proof of Lemma 2.2** For a graph  $G$  let  $\text{frag}(G)$  be the number of vertices in  $G$  less the number of vertices in a largest component; and for integers  $n$  and  $j$  with  $1 \leq j < n$  let  $f(n, j)$  be the number of forests  $F$  on  $[n]$  with  $\text{frag}(F) = j$ . By (11), for  $1 \leq j < n/2$

$$\begin{aligned} f(n, j) &= \binom{n}{j} |\mathcal{F}_j| (n-j)^{n-j-2} \\ &= n^{n-2} \cdot \frac{|\mathcal{F}_j|}{j!} \cdot \frac{\binom{n}{j}}{n^j} \left(1 - \frac{j}{n}\right)^{n-j-2} \\ &\geq n^{n-2} \cdot \frac{|\mathcal{F}_j|}{j! e^j} \cdot \frac{\binom{n}{j}}{n^j} \left(1 - \frac{j}{n}\right)^{-2}. \end{aligned}$$

Now consider just  $j \leq 2$  and let  $n \geq 5$ . Then  $\frac{\binom{n}{j}}{n^j} \left(1 - \frac{j}{n}\right)^{-2} \geq 1$ , so

$$\frac{|\mathcal{F}_n|}{n^{n-2}} > \sum_{j=0}^2 \frac{|\mathcal{F}_j|}{j! e^j} = 1 + \frac{1}{e} + \frac{2}{2! e^2} \approx 1.5032 \approx e^{0.4076}.$$

It is easy to check that this holds also for  $n = 2, 3$  and  $4$ ; so

$$\mathbb{P}(F_n \text{ is connected}) < e^{-2/5} < e^{-1/3} \quad \text{for each } n \geq 2$$

as required. □

### 3 Concluding Remarks

We can improve on Lemma 2.2 by pushing the proof further.

**Lemma 3.1.** *If we set  $\alpha = 0.48$  then for each  $n = 2, 3, \dots$*

$$\mathbb{P}(F_n \text{ is connected}) < e^{-\alpha}.$$

*Proof.* It is easy to check that  $\frac{\binom{n}{j}}{n^j} \left(1 - \frac{j}{n}\right)^{-2} \geq 1$  for each  $j \leq 6$  and  $n > 12$ . Hence, from the proof of Lemma 2.2, for  $n > 12$

$$\frac{|\mathcal{F}_n|}{n^{n-2}} > \sum_{j=0}^6 \frac{|\mathcal{F}_j|}{j! e^j} \approx 1.6167 \approx e^{0.4804} > e^{0.48},$$

and this holds also for  $2 \leq n \leq 12$ . To check this we may for example use [13] for  $|\mathcal{F}_j|$  for  $j \leq 12$ . □

Lemma 3.1 allows us to strengthen Theorem 1.2 as follows: with the same premises, if we set  $\alpha = 0.48$  then

$$\mathbb{P}(R_n \text{ is connected}) \geq \min_{\alpha n < t \leq n} \mathbb{P}(F_t \text{ is connected}). \quad (12)$$

It is well known (see for example Flajolet and Sedgewick [4] Section II.5.3) that  $\sum_{j \geq 1} \frac{|\mathcal{T}_j|}{j! e^j} = \frac{1}{2}$  and so by the exponential formula  $\sum_{j \geq 0} \frac{|\mathcal{F}_j|}{j! e^j} = e^{\frac{1}{2}}$ . We could expect with more work to increase the value  $\alpha = 0.48$  in (12) to nearer  $\frac{1}{2}$  – but can we can go further?

- (a) Perhaps  $\mathbb{P}(F_n \text{ is connected}) \leq e^{-\frac{1}{2}}$  for each  $n \geq 2$ ? In that case we could replace the value  $\alpha = 0.48$  in (12) by the more satisfactory value  $\frac{1}{2}$ .
- (b) Indeed, perhaps  $\mathbb{P}(F_n \text{ is connected})$  is increasing for  $n \geq 5$ ? In that case we could improve the bounds in Theorem 1.2 and in (12) to

$$\mathbb{P}(R_n \text{ is connected}) \geq \mathbb{P}(F_{\lceil n/2 \rceil} \text{ is connected}), \quad (13)$$

which is getting closer to Conjecture (3).

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