MRA-Wa^{v^el}et subspace architecture for logic, probability, and symbolic sequence processing

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Abstract: The linear subspaces of a multiresolution analysis (MRA) and the linear subspaces of the *wavelet analysis* induced by the MRA, together with the set inclusion relation \subseteq , form a very special lattice of subspaces which herein is called a *primorial lattice*. This paper introduces an operator **R** that extracts a set of 2^{N-1} element Boolean lattices from a 2^{N} element Boolean lattice. Used recursively, a sequence of Boolean lattices with decreasing order is generated—a structure that is similar to an MRA. A second operator, which is a special case of a "difference operator", is introduced that operates on consecutive Boolean lattices L_2^n and L_2^{n-1} to produce a sequence of *orthocomplemented lattices.* These two sequences, together with the subset ordering relation \subseteq , form a *primorial lattice* \mathbb{P} . A *logic* or *probability* constructed on a Boolean lattice L_{2}^{N} likewise induces a primorial lattice \mathbb{P} . Such a logic or probability can then be rendered at N different "resolutions" by selecting any one of the N Boolean lattices in \mathbb{P} and at N different "frequencies" by selecting any of the N different orthocomplemented lattices in \mathbb{P} . Furthermore, \mathbb{P} can be used for symbolic sequence analysis by projecting sequences of symbols onto the sublattices in \mathbb{P} using one of three lattice projectors introduced. \mathbb{P} can be used for symbolic sequence processing by judicious rejection and selection of projected sequences. Examples of symbolic sequences include sequences of logic values, sequences of probabilistic events, and genomic sequences (as used in "genomic signal processing").

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1 Background: lattices

1.1 Order

1.1.1 Order relations

Definition 1.1 ¹	Let	t X	be a set.Let 2^{XX}	be the set of all relations on	X.	A relation \leq is an
order relation in	12^{XX}	if			_	

1. $x \leq x$	$\forall x {\in} X$	(reflexive)	and	preorder
2. $x \le y$ and $y \le z \implies x \le z$	$\forall x, y, z \in X$	(transitive)	and	
3. $x \le y$ and $y \le x \implies x = y$	$\forall x, y \in X$	(anti-symmetric)	-	

An **ordered set** is the pair (X, \leq) . The set *X* is called the **base set** of (X, \leq) . If $x \leq y$ or $y \leq x$, then elements *x* and *y* are said to be **comparable**, denoted $x \sim y$. Otherwise they are **incomparable**, denoted x||y. The relation \leq is the relation $\leq \setminus =$ ("less than but not equal to"), where \setminus is the *set difference* operator, and = is the equality relation.

Definition 1.2 ² Let (X, \leq) be an *ordered set* (Definition 1.1 page 3). Let 2^{XX} be the set of all relations on *X*. The relations \geq , <, > $\in 2^{XX}$ are defined as follows:

² 📃 [139], page 2

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 $^{1 \}cong [113]$, page 470, $\cong [12]$, page 1, [] [105], page 156, $\langle I, II, (1) \rangle$, [] [38], page 373, $\langle I-III \rangle$. An *order relation* is also called a **partial order relation**. An *ordered set* is also called a **partially ordered set** or **poset**.

Example 1.3

_	order relation		dual order relation
\leq	(integer less than or equal to)	≥	(integer greater than or equal to)
\subseteq	(subset)	⊇	(super set)
	(divides)		(divided by)
\Rightarrow	(implies)	\Rightarrow	(implied by)

Definition 1.4 ³ A relation \leq is a **linear order relation** on *X* if

1. \leq is an *order relation* (Definition 1.1 page 3) and

2. $x \leq y$ or $y \leq x \quad \forall x, y \in X$ (comparable).

A **linearly ordered set** is the pair (X, \leq) .

A linearly ordered set is also called a **totally ordered set**, a **fully ordered set**, and a **chain**.

1.1.2 Representation

Definition 1.5 ⁴ *y* **covers** *x* in the ordered set (X, \leq) if 1. $x \leq y$ (*y* is greater than *x*) and 2. $(x \leq z \leq y) \implies (z = x \text{ or } z = y)$ (there is no element between *x* and *y*). The case in which *y* covers *x* is denoted $x \prec y$.

An ordered set can be represented in any of three ways:

- Hasse diagram
 a set of ordered pairs of order relations
 (Definition 1.6 page 4)
 (Definition 1.1 page 3)
- a set of ordered pairs of *cover relations* (Definition 1.5 page 4)

Definition 1.6 Let (X, \leq) be an ordered pair. A diagram is a **Hasse diagram** of (X, \leq) if it satisfies the following criteria:

- A Each element in *X* is represented by a dot or small circle.
- For each *x*, *y* ∈ *X*, if *x* ≺ *y*, then *y* appears at a higher position than *x* and a line connects *x* and *y*.

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³ 🖱 [113], page 470, 📃 [133], page 410

⁴ [14], page 445

Example 1.7 Here are three ways of representing the ordered set $(2^{\{x,y\}}, \subseteq)$;

(1) **Hasse diagrams**: If two elements are comparable, then the lesser of the two is drawn lower on the page than the other with a line connecting them.



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- (2) Sets of ordered pairs specifying *order relations* (Definition 1.1 page 3): $\subseteq = \left\{ \begin{array}{c} (\emptyset, \emptyset), \quad (\{x\}, \{x\}), \quad (\{y\}, \{y\}), \quad (\{x, y\}, \{x, y\}), \\ (\emptyset, \{x\}), \quad (\emptyset, \{y\}), \quad (\emptyset, \{x, y\}), \quad (\{x\}, \{x, y\}), \quad (\{y\}, \{x, y\}) \end{array} \right\}$
- (3) Sets of ordered pairs specifying *covering relations*: $\prec = \{ (\emptyset, \{x\}), (\emptyset, \{y\}), (\{x\}, \{x, y\}), (\{y\}, \{x, y\}) \}$

1.1.3 Decomposition

Definition 1.8 ⁵ The tupple (Y, \bigotimes) is a **subposet** of the ordered set (X, \leq) if 1. $Y \subseteq X$ (*Y* is a subset of *X*) and

1. $Y \subseteq X$ (*Y* is a subset of *X*) 2. $\bigotimes = (\le \cap Y^2)$ (\bigotimes is the relation \le restricted to $Y \times Y$)

Example 1.9

Subposets of include include

Example 1.10 Let

$$(X, \leq) \triangleq \left(\{0, a, b, c, p, 1\}, \begin{cases} (0, 0), (a, a), (b, b), (c, c), (p, p), (1, 1), \\ (0, a), (0, b), (0, c), (0, p), (0, 1), \\ (a, b), (a, c), (a, 1), (p, 1), \end{cases} \right)$$

$$(Y, \otimes) \triangleq \left(\{0, a, c, p, 1\}, \begin{cases} (0, 0), (a, a), (c, c), (p, p), (1, 1), \\ (0, a), (0, c), (0, p), (0, 1), \\ (a, c), (a, 1), (p, 1), (c, 1), (p, 1) \end{cases} \right)$$

Then (Y, \bigotimes) is a subposet of (X, \leq) because $Y \subseteq X$ and $\bigotimes = (\leq \cap Y^2)$.

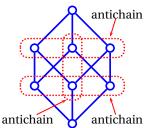
5 🖱 [72], page 2

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A *chain* is an ordered set in which every pair of elements is *comparable* (Definition 1.4 page 4). An antichain is just the opposite-it is an ordered set in which no pair of elements is comparable (next definition).

Definition 1.11 ⁶ The subposet (A, \bigotimes) in the ordered set (X, \leq) is an **antichain** if all elements in A are *incomparable* (Definition 1.1 page 3), such that

x || y $\forall x, y \in A$



Definition 1.12 ⁷ The length $\ell(L)$ of a *chain* (Definition 1.4 page 4) L with N elements is N - 1. The length of an ordered set (Definition 1.1 page 3) is the length of the longest chain in the ordered set. The width of an ordered set is the number of elements in the largest antichain in the ordered set.

Theorem 1.13 (Dilworth's theorem) ⁸ Let (X, \leq) be an ordered set.

 $\begin{cases} \text{WIDTH } N \text{ of } (X, \leq) \\ \text{is FINITE} \end{cases} \end{cases} \implies \begin{cases} 1. \text{ there exists a partition of } (X, \leq) \text{ into } N \text{ chains } and \\ 2. \text{ there does not exist any partition} \\ \text{of } (X, \leq) \text{ into less than } N \text{ chains} \end{cases}$

Definition 1.14 ⁹ Let *X* and *Y* be disjoint sets. Let $P \triangleq (X, \bigotimes)$ and $Q \triangleq (Y, \lessdot)$ be ordered sets on *X* and *Y*. The **direct sum** of *P* and *Q* is defined as

 $\boldsymbol{P} + \boldsymbol{Q} \triangleq (X \cup Y, \leq)$ where $x \leq y$ if 1. $x, y \in X$ and $x \otimes y$ or 2. $x, y \in Y$ and $x \lessdot y$

The direct sum operation is also called the **disjoint union**. The notation nP is defined as

 $n\boldsymbol{P} \triangleq \boldsymbol{P} + \boldsymbol{P} + \cdots + \boldsymbol{P}$. n-1 "+" operations

Definition 1.15 ¹⁰ Let X and Y be disjoint sets. Let $P \triangleq (X, \bigotimes)$ and $Q \triangleq (Y, \blacktriangleleft)$ be ordered sets on *X* and *Y*. The **direct product** of *P* and *Q* is defined as

 $\boldsymbol{P} \times \boldsymbol{Q} \triangleq (X \times Y, \leq)$ where $(x_1, y_1) \leq (x_2, y_2)$ if $x_1 \otimes x_2$ and $y_1 \otimes y_2$.

¹⁰ 🔁 [155], pages 100–101, 🖻 [154], page 43

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⁶ 🖱 [72], page 2 ⁷ 🔁 [72], page 2, 🛸 [18], page 5 ⁸ [47], page 161, ⁶ [48], ^[] [56], page 4

⁹ 🖱 [155], page 100

The direct product operation is also called the **cartesian product**. The order relation \leq is called a **coordinate wise** order relation. The notation P^n is defined as

 $\boldsymbol{P}^{n} \triangleq \underbrace{\boldsymbol{P} \times \boldsymbol{P} \times \cdots \times \boldsymbol{P}}_{n-1 \text{ "} \times \text{" operations}}.$

Definition 1.16 ¹¹ Let *X* and *Y* be disjoint sets. Let $P \triangleq (X, \bigotimes)$ and $Q \triangleq (Y, \blacktriangleleft)$ be ordered sets on *X* and *Y*. The **ordinal sum** of *P* and *Q* is defined as

 $\boldsymbol{P} \oplus \boldsymbol{Q} \triangleq (X \cup Y, \leq)$ where $x \leq y$ if 1. $x, y \in X$ and $x \otimes y$ or 2. $x, y \in Y$ and $x \lessdot y$

> 3. $x \in X$ and $y \in Y$.

Definition 1.17 ¹² Let X and Y be disjoint sets. Let $P \triangleq (X, \bigotimes)$ and $Q \triangleq (Y, \blacktriangleleft)$ be ordered sets on X and Y. The ordinal product of P and Q is defined as

 $\boldsymbol{P} \otimes \boldsymbol{Q} \triangleq (X \times Y, \leq)$

where $(x_1, y_1) \le (x_2, y_2)$ if $\begin{cases} 1. & x_1 \ne x_2 \text{ and } x_1 \otimes x_2 \text{ or } \\ 2. & x_1 = x_2 \text{ and } y_1 \lessdot y_2 \end{cases}$ The order relation \le is called a **lexicographical** order relation, **dictionary** order relation,

or

or **alphabetic** order relation.

Definition 1.18 ¹³ Let $P \triangleq (X, \leq)$ be an ordered set. Let \geq be the dual order relation of \leq . The **dual** of **P** is defined as $P^* \triangleq (X, \geq)$

Definition 1.19 ¹⁴ Let X and Y be disjoint sets. Let $P \triangleq (X, \otimes)$ and $Q \triangleq (Y, \ll)$ be ordered sets on X and Y. $Q^{P} \triangleq (\{f \in Y^{X} | f \text{ is order preserving}\}, \leq)$ where $f \leq g$ if $f(x) \leq g(x)$ $\forall x \in X$. The order relation \leq is called a **pointwise order** relation.

Theorem 1.20 (cardinal arithmetic) ¹⁵ Let $P \triangleq (X, \leq)$ be an ordered set.

1.	P+Q	=	Q + P	(COMMUTATIVE)
2.	P imes Q	=	$Q \times P$	(COMMUTATIVE)
З.	(P + Q) + R	=	P + (Q + R)	(ASSOCIATIVE)
4.	$(\boldsymbol{P} \times \boldsymbol{Q}) \times \boldsymbol{R}$	=	$\boldsymbol{P} \times (\boldsymbol{Q} \times \boldsymbol{R})$	(ASSOCIATIVE)
5.	$P \times (Q + R)$	=	$(\boldsymbol{P} \times \boldsymbol{Q}) + (\boldsymbol{P} \times \boldsymbol{R})$	(DISTRIBUTIVE)
6.	R^{P+Q}	=	$R^P imes R^Q$	
7.	$(\boldsymbol{P}^{\boldsymbol{Q}})^{\boldsymbol{R}}$	=	₽ ^{Q×R}	

¹¹ 🖻 [155], page 100

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¹² 🖻 [155], page 101, 🖻 [154], page 44, 🖻 [81], page 58, 🖱 [82], page 54

¹³ 🖱 [155], page 101

¹⁴ 🔁 [155], page 101

¹⁵ 🖱 [155], page 102

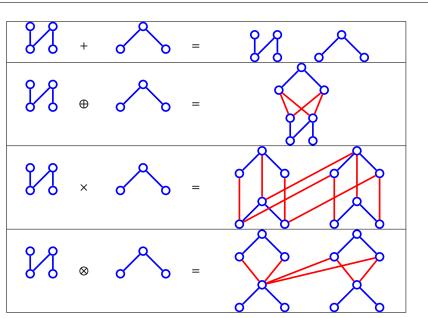


Figure 1: Operations on ordered sets (Example 1.23 page 8)

Definition 1.21 The ordered set L_1 is defined as $(\{x\}, \leq)$, for some value x. It is illustrated by the Hasse diagram to the right.

Definition 1.22 The ordered set L_2 is defined as $L_2 \triangleq L_1^2$. It is illustrated by the Hasse diagram to the right.

1.1.4 Decomposition examples

Example 1.23 Figure 1 (page 8) illustrates the four ordered set operations $+, \times, \oplus$, and \otimes .

Example 1.24 ¹⁶The ordered set nL_1 is the *anti-chain* with *n* elements. The ordered set $4L_1$ is illustrated to the right. $\circ \circ \circ \circ$

Example 1.25 The ordered set L_1^n is the *chain* with *n* elements. The ordered set L_1^4 is illustrated to the right.

Examples of the *Boolean lattices* (Definition 1.69 page 18) L_2^1 , L_2^2 , L_2^3 , L_2^4 and L_2^5 are illustrated in Example 1.74 (page 21).

¹⁶ 🔁 [155], page 100

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1 BACKGROUND: LATTICES			Daniel J. Greenhoe				page 9	
longest antichain p q r q r q r	partition 1: { partition 2: { partition 3: { partition 4: { partition 5: { partition 6: { }	$\{0, a\}, \\ \{0, a, p\}, \\ \{0, a, p, 1\}, \\ \{0, b, p, 1\}, \\ \{0, c, r, 1\}, \\ \{0, c, q, 1\}, \end{cases}$	${b}, {a}, {a}, {a, p},$	$\{c,q\},\ \{c,q\},\ \{c,q\},\ \{b\},$				

examples of partitions of chains

Figure 2: Lattice of width 4 and examples of minimal order partitions of chains (see Example 1.26 page 9)

Example 1.26 ¹⁷ The longest *antichain* (Definition 1.11 page 6) in the lattice illustrated in Figure 2 (page 9) has 4 elements giving this ordered set a *width* (Definition 1.12 page 6) of 4. The longest chain also has 4 elements, giving the ordered set a *length* (Definition 1.12 page 6) of 3. By *Dilworth's theorem* (Theorem 1.13 page 6), the smallest partition consists of four *chains* (Definition 1.4 page 4). Examples of such minimal order partitions those listed in Figure 2.

Definition 1.27 Let (X, \leq) be an ordered set and 2^X the power set of *X*. For any set $A \in 2^X$, *c* is an **upper bound** of *A* in (X, \leq) if

1. $x \leq c \quad \forall x \in A$.

An element *b* is the **least upper bound**, or **LUB**, of *A* in (X, \leq) if

2. *b* and *c* are *upper bounds* of $A \implies b \le c$.

The least upper bound of the set *A* is denoted $\bigvee A$. It is also called the **supremum** of *A*, which is denoted sup *A*. The **join** $x \lor y$ of *x* and *y* is defined as $x \lor y \triangleq \bigvee \{x, y\}$.

Definition 1.28 Let (X, \leq) be an ordered set and 2^X the power set of *X*. For any set $A \in 2^X$, *p* is a **lower bound** of *A* in (X, \leq) if

1.
$$p \le x \quad \forall x \in A$$
.

An element *a* is the **greatest lower bound**, or **GLB**, of *A* in (X, \leq) if

2. *a* and *p* are *lower bounds* of $A \implies p \le a$.

The greatest lower bound of the set *A* is denoted $\bigwedge A$. It is also called the **infimum** of *A*, which is denoted inf *A*. The **meet** $x \land y$ of *x* and *y* is defined as $x \land y \triangleq \bigwedge \{x, y\}$.

Proposition 1.29 Let $(X, \lor, \land; \le)$ be an ORDERED SET (Definition 1.1 page 3).

$$x \leq y \iff \left\{\begin{array}{rrr} 1. & x \wedge y = x & and \\ 2. & x \vee y = y \end{array}\right\} \quad \forall x, y \in X$$

¹⁷ [56], page 4

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Proposition 1.30 Let 2^X be the POWER SET of a set X.

$$A \subseteq B \implies \left\{ \begin{array}{rrr} 1. \quad \bigvee A & \leq \quad \bigvee B \quad and \\ 2. \quad \bigwedge A & \leq \quad \bigwedge B \end{array} \right\} \qquad \forall A, B \in 2^{\lambda}$$

1.2 Lattices

1.2.1 Definition

The structure available in an *ordered set* (Definition 1.1 page 3) tends to be insufficient to ensure "well-behaved" mathematical systems. This situation is greatly remedied if every pair of elements in the ordered set has both a *least upper bound* and a *greatest lower bound* (Definition 1.28 page 9) in the set; in this case, that ordered set is a *lattice* (next definition). Gian-Carlo Rota (1932–1999) has illustrated the advantage of lattices over simple ordered sets by pointing out that the *ordered set of partitions of an integer* "is fraught with pathological properties", while the *lattice* of partitions of a set "remains to this day rich in pleasant surprises".¹⁸

Definition 1.31 ¹⁹ An algebraic structure $L \triangleq (X, \lor, \land; \leq)$ is a **lattice** if

1.	(X, \leq) is an ordered set	$((X, \leq)$ is a partially or totally ordered set)	and
2.	$\exists x \lor y \in X \forall x, y \in X$	(every pair of elements in <i>X</i> has a <i>least upper bound</i> in <i>X</i>)	and

3. $\exists x \land y \in X \quad \forall x, y \in X$ (every pair of elements in *X* has a *greatest lower bound* in *X*). The algebraic structure $L^* \triangleq (X, \oslash, \odot; \ge)$ is the **dual** lattice of *L*, where \odot and \odot are determined by \ge . The *lattice L* is *linear* if (X, \le) is a *chain* (Definition 1.4 page 4).

Theorem 1.32 ²⁰ $(X, \lor, \land; \leq)$ *is a* LATTICE \Leftarrow

$\int x \lor x$	=	x	$x \wedge x$	=	x	$\forall x \in X$	(IDEMPOTENT)	and	
$x \lor y$		2	$x \wedge y$		2		(COMMUTATIVE)		
$(x \lor y) \lor z$	=	$x \lor (y \lor z)$	$(x \land y) \land z$	=	$x \wedge (y \wedge z)$	$\forall x, y, z \in X$	(ASSOCIATIVE)	and	
$\int x \vee (x \wedge y)$	=	x	$x \wedge (x \vee y)$	=	x	$\forall x, y \in X$	(ABSORPTIVE).	J	

Lemma 1.33 ²¹ Let $L \triangleq (X, \lor, \land; \leq)$ be LATTICE (Definition 1.31 page 10).

 $x \le y \qquad \Longleftrightarrow \qquad x = x \land y \qquad \forall x, y \in L$

Proof:

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¹⁸ [148], page 1440, ((illustration)), [147], page 498, (partitions of a set)

¹⁹ 🖱 [113], page 473, 🖱 [17], page 16, 📃 [133], 📃 [14], page 442, 🖱 [116], page 1

²⁰ ► [113], pages 473–475, 〈LEMMA 1, THEOREM 4〉, ► [23], pages 4–7, ► [16], pages 795–796,

 $[\]boxed{133}$, page 409, $\langle (\alpha) \rangle$, $\boxed{14}$, page 442, $\boxed{138}$, pages 371–372, $\langle (1)-(4) \rangle$ 21 $\boxed{88}$

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- (1) Proof for \implies case: by left hypothesis and definition of \land (Definition 1.28 page 9).
- (2) Proof for \leftarrow case: by right hypothesis and definition of \land (Definition 1.28 page 9).

Proposition 1.34 (Monotony laws) ²² Let $(X, \lor, \land; \le)$ be a lattice. $\begin{cases} a \le b \text{ and} \\ x \le y \end{cases} \implies \begin{cases} a \land x \le b \land y \text{ and} \\ a \lor x \le b \lor y \end{cases}$

Theorem 1.35 (Minimax inequality) ²³ Let $(X, \lor, \land; \leq)$ be a lattice.

$$\bigvee_{i=1}^{m} \bigwedge_{j=1}^{n} x_{ij} \leq \bigwedge_{j=1}^{n} \bigvee_{i=1}^{m} x_{ij} \qquad \forall x_{ij} \in X$$

maxmini: largest of the smallest

minimax: smallest of the largest

Special cases of the minimax inequality include three distributive inequalities (next theorem). If for some lattice any one of these inequalities is an equality, then all three are equalities (Theorem 1.54 page 15); and in this case, the lattice is a called a *distributive* lattice (Definition 1.53 page 15).

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Theorem 1.36 (distributive inequalities) <sup>24</sup> (X, \lor, \land; \leq) is a lattice \implies
```

 $\begin{array}{lll} x \wedge (y \lor z) &\geq & (x \wedge y) \lor (x \wedge z) \\ x \lor (y \wedge z) &\leq & (x \lor y) \land (x \lor z) \end{array} \qquad \qquad \forall x, y, z \in X \quad (\text{JOIN SUPER-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y, z \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y \in X \quad (\text{MEET SUB-DISTRIBUTIVE}) \quad and \\ \forall x, y \in X \quad (x \to X$ $(x \land y) \lor (x \land z) \lor (y \land z) \leq (x \lor y) \land (x \lor z) \land (y \lor z) \quad \forall x, y, z \in X \quad (\text{median inequality}).$

Besides the distributive property, another consequence of the minimax inequality is the modularity inequality (next theorem). A lattice in which this inequality becomes equality is said to be *modular* (Definition 1.47 page 14).

Theorem 1.37 (Modular inequality) ²⁵ Let $(X, \lor, \land; \leq)$ be a LATTICE (Definition 1.31 page 10). $x \leq y$ $x \lor (y \land z) \le y \land (x \lor z)$ \implies

Theorem 1.32 (page 10) gives 4 necessary and sufficient pairs of properties for a structure $(X, \lor, \land; \leq)$ to be a *lattice*. However, these 4 pairs are actually *overly* sufficient (they are not *independent*), as demonstrated next.

²³ 🖱 [17], pages 19–20

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²² 🔁 [68], page 39, 📃 [50], pages 97–99, 🐿 [78], (§4.2)

²⁴ 🔊 [36], page 85, 🔊 [72], page 38, 📃 [14], page 444, 📃 [105], page 157, 🛸 [125], page 13, (terminology)

²⁵ 🔊 [17], page 19, 🛸 [23], page 11, 📃 [38], page 374

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Theorem 1.38 ²⁶

($X, \vee, \wedge; \leq$	is a	lattice	\Leftrightarrow						
	$x \lor y$	=	$y \lor x$	$x \wedge y$	=	$y \wedge x$	$\forall x, y \in X$	(COMMUTATIVE)	and	
<	$(x \lor y) \lor z$	=	$x \lor (y \lor z)$	$(x \land y) \land z$	=	$x \wedge (y \wedge z)$	$\forall x, y, z \in X$	(ASSOCIATIVE)	and	}
	$x \lor (x \land y)$	=	x	$x \land (x \lor y)$	=	x	$\forall x, y \in X$	(ABSORPTIVE)	J	

1.2.2 Bounded lattices

Let $L \triangleq (X, \lor, \land; \leq)$ be a lattice. By the definition of a *lattice* (Definition 1.31 page 10), the *upper* bound $(x \lor y)$ and *lower bound* $(x \land y)$ of any two elements in X is also in X. But what about the upper and lower bounds of the entire set X ($\bigvee X$ and $\bigwedge X$) (Definition 1.27 page 9, Definition 1.28 page 9)? If both of these are in X, then the lattice L is said to be *bounded* (next definition). All *finite* lattices are bounded (next proposition). However, not all lattices are bounded—for example, the lattice (\mathbb{Z}, \leq) (the lattice of integers with the standard integer ordering relation) is *unbounded*.

Definition 1.39 Let $L \triangleq (X, \lor, \land; \le)$ be a lattice. Let $\bigvee X$ be the least upper bound of (X, \le) and let $\bigwedge X$ be the greatest lower bound of (X, \le) .

L is upper boundedif $(\bigvee X) \in X$.L is lower boundedif $(\bigwedge X) \in X$.L is boundedif L is both upper and lower bounded.

A *bounded* lattice is optionally denoted $(X, \lor, \land, 0, 1; \leq)$, where $0 \triangleq \bigwedge X$ and $1 \triangleq \bigvee X$.

Proposition 1.40 Let $L \triangleq (X, \lor, \land; \le)$ be a lattice. { *L* is FINITE} \implies { *L* is BOUNDED}

Proposition 1.41 ²⁷ Let $L \triangleq (X, \lor, \land; \leq)$ be a lattice with $\bigvee X \triangleq 1$ and $\bigwedge X \triangleq 0$.

 $\{L \text{ is BOUNDED}\} \implies \begin{cases} x \lor 1 = 1 \quad \forall x \in X \quad (upper bounded) \quad and \\ x \land 0 = 0 \quad \forall x \in X \quad (lower bounded) \quad and \\ x \lor 0 = x \quad \forall x \in X \quad (join-identity) \quad and \\ x \land 1 = x \quad \forall x \in X \quad (meet-identity) \end{cases} \}$

Definition 1.42 ²⁸ Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be a *bounded lattice* (Definition 1.39 page 12). The **height** h(x) of a point $x \in L$ is the *least upper bound* of the *lengths* (Definition 1.12 page 6) of all the *chains* that have 0 and in which x is the *least upper bound*. The **height** h(L) of the lattice L is defined as

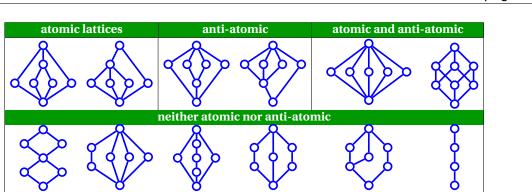
 $\mathsf{h}(\boldsymbol{L}) \triangleq \mathsf{h}(1) \ .$

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²⁶ \square [136], pages 7–8, \square [12], page 5, \blacksquare [120], page 24, \blacksquare [77], (Theorem 1.22), \square [78], (§4.4) ²⁷ \blacksquare [77], (§1.2.2), \square [78], (§4.5)

²⁸ 🖱 [18], page 5

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Figure 3: Selected *atomic, anti-atomic,* and neither atomic nor anti-atomic lattices (see Example 1.45 page 13)

Example 1.43 The *height* of the lattice illustrated in Figure 2 (page 9) is 3 because

$$h(L) \triangleq h(1)$$

$$\triangleq \bigvee \{\ell(C) | C \text{ is a chain in } L \text{ containing both } 0 \text{ and } 1 \}$$

$$= \bigvee \{\ell(\{0, a, p, 1\}, \leq), \ell(\{0, b, p, 1\}, \leq), \ell(\{0, c, p, 1\}, \leq), \ell(\{0, c, q, 1\}, \leq), \ell(\{0, c, r, 1\}, \leq), \}$$

$$= \bigvee \{4 - 1, 4 - 1, 4 - 1, 4 - 1\}$$

$$= \bigvee \{3, 3, 3, 3, 3\}$$

$$= 3$$

1.2.3 Atomic lattices

Definition 1.44 ²⁹ Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be a *bounded lattice* (Definition 1.39 page 12). *x* is an **atom** of *L* if *x* covers (Definition 1.5 page 4) 0.

- *x* is an **anti-atom** of *L* if *x* is *covered by* 1.
- *L* is **atomic** if every $x \in X \setminus 0$ can be represented as joins of atoms of *L*.
- *L* is **anti-atomic** if every $x \in X \setminus 1$ can be represented as meets of anti-atoms of *L*.

Example 1.45 Figure 3 (page 13) illustrates some examples of lattices that are *atomic*, *anti-atomic*, both, and neither.

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²⁹ [[108], page 178, [[16], page 800, (see footnote ‡)

1.2.4 Modular Lattices

Definition 1.46 ³⁰ Let $(X, \lor, \land; \le)$ be a lattice. Let 2^{XX} be the set of all *relation*s in X^2 . The **modularity** relation $\circledast \in 2^{XX}$ and the **dual modularity** relation $\circledast^* \in 2^{XX}$ are defined as

 $\stackrel{\text{def}}{\Longleftrightarrow} \quad \left\{ (x,y) \in X^2 \mid a \le y \quad \Longrightarrow \quad y \land (x \lor a) = (y \land x) \lor a \quad \forall a \in X \right\}$ $x \otimes v$ $x \circledast^* y \quad \stackrel{\mathsf{def}}{\Longleftrightarrow} \quad \left\{ (x, y) \in X^2 \ | a \geq y \quad \Longrightarrow \quad y \lor (x \land a) = (y \lor x) \land a \quad \forall a \in X \right\}.$

A pair $(x, y) \in \mathbb{B}$ is alternatively denoted as $(x, y) \otimes$, and is called a **modular** pair. A pair $(x, y) \in \mathbb{O}^*$ is alternatively denoted as $(x, y) \mathbb{O}^*$, and is called a **dual modular** pair. A pair (x, y) that is *not* a modular pair $((x, y) \notin \mathbb{Q})$ is denoted $x \oplus y$. A pair (x, y) that is *not* a dual modular pair is denoted $x \mathfrak{D}^* y$.

Modular lattices are a generalization of *distributive lattices* (Definition 1.53 page 15) in that all distributive lattices are modular, but not all modular lattices are distributive (Example 1.61 page 16, Example 1.62 page 17).

Definition 1.47 ³¹ A lattice $(X, \lor, \land; \leq)$ is modular if $x \otimes y$ $\forall x, y \in X.$

Theorem 1.48 ³² Let $L \triangleq (X, \lor, \land; \leq)$ be a lattice.

L is modular $\iff \{x \le y \implies x \lor (z \land y) = (x \lor z) \land y\} \forall x, y, z \in X$ $x \lor [(x \lor y) \land z] = (x \lor y) \land (x \lor z)$ $\forall x, y, z \in X$ $\iff x \wedge [(x \wedge y) \lor z] = (x \wedge y) \lor (x \wedge z)$ $\forall x, y, z \in X$

Definition 1.49 (N5 lattice/pentagon) ³³ The N5 lattice is the ordered set ($\{0, a, b, p, 1\}, \leq$) with cover relation

 $\prec = \{(0, a), (a, b), (b, 1), (p, 1), (0, p)\}.$

The N5 lattice is also called the **pentagon**. The N5 lattice is illustrated by the Hasse diagram to the right.

Theorem 1.50 ³⁴ Let *L* be a LATTICE (Definition 1.31 page 10).

L is MODULAR (Definition 1.47 page 14) \iff *L* does NOT contain the N5 LATTICE (Definition 1.49) page 14).

Theorem 1.51 ³⁵ Let $\mathbf{A} \triangleq (X, \lor, \land; \leq)$ be an algebraic structure. $\begin{cases} (x \land y) \lor (x \land z) = [(z \land x) \lor y] \land x \quad \forall x, y, z \in X \quad and \\ [x \lor (y \lor z)] \land z = z \quad \forall x, y, z \in X \end{cases} \iff \begin{cases} \mathbf{A} \text{ is } a \\ \mathbf{modular lattice} \end{cases}$ ³⁰ S [157], page 11, S [116], page 1, (Definition (1.1)), S [117], page 248 ³¹ S [18], page 82, S [116], page 3, (Definition (1.7)) ³² 🖱 [136], page 39, 🖱 [133], page 413, ((2)), 🖱 [78], (Theorem 5.1) ³³ [12], pages 12–13, [[38], pages 391–392, ((44) and (45)) ³⁴ S [23], page 11, S [71], page 70, [[38], (cf Stern 1999 page 10), S [78], (Theorem 5.1)

³⁵ 🔁 [136], pages 42–43, 📃 [145]

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Examples of *modular lattices* are provided in Example 1.61 (page 16) and Example 1.62 (page 17).

1.2.5 Distributive Lattices

Definition 1.52 ³⁶ Let $(X, \lor, \land; \le)$ be a *lattice* (Definition 1.31 page 10). Let 2^{XXX} be the set of all *relations* in X^3 . The **distributivity** relation $@ \in 2^{XXX}$ and the **dual distributivity** relation $@^* \in 2^{XXX}$ are defined as

 $\widehat{\mathbb{D}} \stackrel{\text{def}}{=} \left\{ (x, y, z) \in X^3 | x \land (y \lor z) = (x \land y) \lor (x \land z) \right\}$ (each (x, y, z) is *disjunctive distributive*) and $\widehat{\mathbb{D}}^* \stackrel{\text{def}}{=} \left\{ (x, y, z) \in X^3 | x \lor (y \land z) = (x \lor y) \land (x \lor z) \right\}$ (each (x, y, z) is *conjunctive distributive*). A triple $(x, y, z) \in \widehat{\mathbb{D}}$ is alternatively denoted as (x, y, z), and is a **distributive** triple. A triple $(x, y, z) \in \widehat{\mathbb{D}}^*$ is alternatively denoted as (x, y, z), and is a **dual distributive** triple.

Definition 1.53 ³⁷ A lattice $(X, \lor, \land; \leq)$ is **distributive** if $(x, y, z) \in \bigcirc \forall x, y, z \in X$

Not all lattices are *distributive*. But if a lattice *L* does happen to be distributive (Definition 1.53 page 15)—that is all triples in *L* satisfy the *distributive* property (Definition 1.53 page 15)—then all triples in *L* also satisfy the *dual distributive* property, as well as another property called the *median property*. The converses also hold (next theorem).

Theorem 1.54 ³⁸ Let $L \triangleq (X, \lor, \land; \leq)$ be a LATTICE (Definition 1.31 page 10).

L is di	L is DISTRIBUTIVE (Definition 1.53 page 15)										
\Leftrightarrow	$x \land (y \lor z) = (x \land y) \lor (x \land z)$	$\forall x, y, z \in X$	(DISJUNCTIVE DISTRIBUTIVE)								
\Leftrightarrow	$x \lor (y \land z) = (x \lor y) \land (x \lor z)$	$\forall x, y, z \in X$	(CONJUNCTIVE DISTRIBUTIVE)								
\Leftrightarrow	$(x \lor y) \land (x \lor z) \land (y \lor z) = (x \land y) \lor (x \land z) \lor (y \land z)$	$\forall x, y, z \in X$	(MEDIAN PROPERTY)								

Definition 1.55 (M3 lattice/diamond) ³⁹ The **M3 lattice** is the ordered set $(\{0, p, q, r, 1\}, \leq)$ with covering relation

 $\leq = \{(p, 1), (q, 1), (r, 1), (0, p), (0, q), (0, r)\}.$ The M3 lattice is also called the **diamond**, and is illustrated by the Hasse diagram to the right.



 $^{^{36} \}cong [116]$, page 15, (Definition 4.1), [] [62], page 67, [] [130], page 32, (Definition 5.1), [] [37], page 314, (*disjunctive distributive* and *conjunctive distributive* functions)

³⁹ [12], pages 12–13, (105], page 157, $\langle p_1 \equiv x, p_2 \equiv y, p_3 \equiv z, g \equiv 1, 0 \equiv 0 \rangle$

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³⁷ **►** [23], page 10, **►** [17], page 133, **■** [133], page 414, *(arithmetic axiom)*, **■** [14], page 453, **►** [9], page 48, (Definition II.5.1)

³⁸ [[49], page 237, < [23], page 10, [[133], page 416, ⟨(7),(8), Theorem 3⟩, [[134], ⟨cf Gratzer 2003 page 159⟩, [153], page 286, ⟨cf Birkhoff(1948)p.133⟩, [[105], ⟨cf Birkhoff(1948)p.133⟩, [[78], ⟨Theorem 6.1⟩

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 $\left\{ \begin{array}{c} Lemma 1.56 & {}^{40} \\ \left\{ \begin{array}{c} L \text{ is an} \\ M3 \text{ lattice} \end{array} \right\} \implies \left\{ \begin{array}{c} 1. \quad L \text{ is NOT distributive} & \text{(Definition 1.53 page 15)} & and \\ 2. \quad L \text{ Is modular} & \text{(Definition 1.47 page 14)} \end{array} \right\}$

Theorem 1.57 (Birkhoff distributivity criterion) ⁴¹ Let $L \triangleq (X, \lor, \land; \le)$ be a LATTICE. *L* is DISTRIBUTIVE $\iff \begin{cases} L \text{ does not contain N5 as a sublattice} & & \text{and} \\ L \text{ does not contain M3 as a sublattice} & & & & \end{cases}$

Distributive lattices are a special case of modular lattices. That is, all distributive lattices are modular, but not all modular lattices are distributive (next theorem). An example is the *M3 lattice*—it is modular, but yet it is not *distributive*.

Theorem 1.58 ⁴² Let $(X, \lor, \land; \le)$ be a lattice. $\{(X, \lor, \land; \le) \text{ is DISTRIBUTIVE}\} \xrightarrow{\rightleftharpoons} \{(X, \lor, \land; \le) \text{ is MODULAR}\}$

Theorem 1.59 ⁴³ Let $L \triangleq (X, \lor, \land; \leq)$ be a LATTICE (Definition 1.31 page 10).

 $\left\{ \begin{array}{ll} 1. \quad L \text{ is DISTRIBUTIVE} \quad and \\ 2. \quad x \lor a = x \lor b \qquad and \\ 3. \quad x \land a = x \land b \end{array} \right\} \implies \{a = b\} \quad \forall x, a, b \in X$

Proposition 1.60 ⁴⁴ Let X_n be a finite set with order $n = |X_n|$. Let l_n be the number of unlabeled lattices on X_n , m_n the number of unlabeled modular lattices on X_n , and d_n the number of unlabeled distributive lattices on X_n .

n		0	1	2	3	4	5	6	7	8	9	10	11	12	13	14
l_n		1	1	1	1	2	5	15	53	222	1078	5994	37622	262776	2018305	16873364
m	n	1	1	1	1	2	4	8	16	34	72	157	343	766	1718	3899
d_n	- 11	1	1	1	1	2	3	5	8	15	26	47	82	151	269	494

Example 1.61 ⁴⁵ There are a total of 5 unlabeled lattices on a five element set. Of these, 3 are *distributive* (Proposition 1.60 page 16, and thus also *modular*), one is *modular* but *non*-

⁴⁵ [[54], pages 4–5, ► [78], (Example 6.2)

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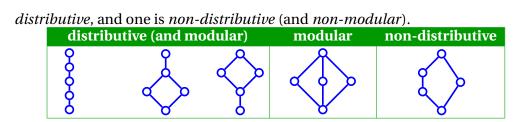
⁴⁰ 🖱 [17], page 6, 🖱 [23], page 11, 🖱 [105], page 157, (cf Salii1988 p. 37)

⁴¹ 🖱 [23], page 12, 🖱 [17], page 134, 📃 [19] 🖱 [78], (Theorem 6.2)

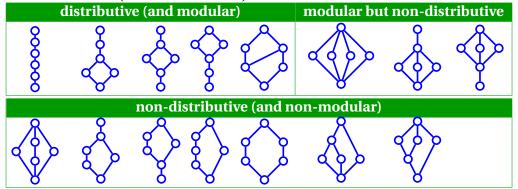
⁴² • [17], page 134, • [23], page 11 [[77], (Theorem 1.37), • [78], (§6.2.3)

^{43 🔁 [113],} pages 484–485

⁴⁴ 및 [2] (http://oeis.org/A006966), 및 [2] (http://oeis.org/A006982), 및 [2] (http://oeis.org/A006981), [84], (*l_n*), [54], page 17, (*d_n*), [160]



Example 1.62 ⁴⁶ There are a total of 15 unlabeled lattices on a six element set. Of these, 5 are *distributive* (Proposition 1.60 page 16, and *modular*), 3 are *modular* but *non-distributive*, and 7 are *non-distributive* (and *non-modular*).



1.2.6 Complemented lattices

Definition 1.63 ⁴⁷ Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be a *bounded lattice* (Definition 1.39 page 12). An element $x' \in X$ is a **complement** of an element x in L if

1. $x \wedge x' = 0$ (non-contradiction) and

2.
$$x \lor x' = 1$$
 (excluded middle)

An element x' in L is the *unique complement* of x in L if x' is a *complement* of x and y' is a *complement* of $x \implies x' = y'$. L is **complemented** if every element in X has a complement in X. L is **uniquely complemented** if every element in X has a unique complement in X. A complemented lattice that is *not* uniquely complemented is **multiply complemented**.

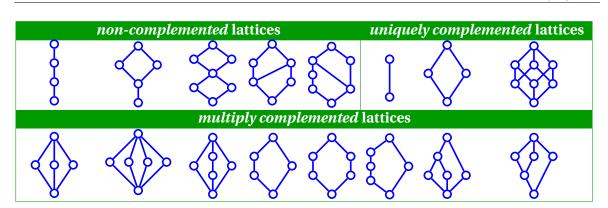
Example 1.64 Here are some examples:

```
<sup>46</sup> ● [78], ⟨Example 5.6⟩
<sup>47</sup> ● [157], page 9, ● [17], page 23
```

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Example 1.65 Of the 53 unlabeled lattices on a 7 element set, 0 are *uniquely complemented*, 17 are *multiply complemented*, and 36 are *non-complemented*.

Theorem 1.66 (next) is a landmark theorem in mathematics.

Theorem 1.66 ⁴⁸ For every lattice *L*, there exists a lattice *U* such that

- 1. $L \subseteq U$ (*L* is a sublattice of U) and
- 2. **U** is uniquely complemented.

Corollary 1.67 ⁴⁹ Let $L \triangleq (X, \lor, \land; \leq)$ be a lattice.

 $\left\{\begin{array}{cc}
1. \quad \boldsymbol{L} \text{ is DISTRIBUTIVE} & and \\
2. \quad \boldsymbol{L} \text{ is COMPLEMENTED}
\end{array}\right\}$

```
{ L is UNIQUELY COMPLEMENTED }
```

Theorem 1.68 (Huntington properties) ⁵⁰ Let *L* be a lattice.

 $\left\{ \begin{array}{c} L \ is \\ \text{UNIQUELY} \\ \text{COMPLEMENTED} \end{array} \right\} and \left\{ \begin{array}{c} L \ is \ \text{MODULAR} & or \\ L \ is \ \text{ATOMIC} & or \\ L \ is \ \text{OTHOCOMPLEMENTED} & or \\ L \ has \ \text{FINITE WIDTH} & or \\ L \ is \ \text{DECOMPLEMENTED} \end{array} \right\} \implies \left\{ \begin{array}{c} L \ is \\ \text{DISTRIBUTIVE} \end{array} \right\}$

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1.2.7 Boolean lattices

Definition 1.69 ⁵¹ A *lattice* (Definition 1.31 page 10) *L* is **Boolean** if

- 1. *L* is *bounded* (Definition 1.39 page 12) and
- 2. *L* is *distributive* (Definition 1.53 page 15) and
- 3. *L* is *complemented* (Definition 1.63 page 17).

⁵¹ 🖱 [113], page 488, 🖱 [97]

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⁴⁸ 📃 [46], page 123, 🍽 [151], page 51, 🍽 [72], page 378, (Corollary 3.8)

^{49 🔁 [113],} page 488, 🛸 [151], page 30, (Theorem 10)

⁵⁰ 🖻 [146], page 103, 🐃 [3], page 79, 🐃 [151], page 40, 📃 [46], page 123, 🐃 [73], page 698

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DAGRGROUND, LATTICES	

In this case, *L* is a **Boolean algebra** or a **Boolean lattice**.

In this paper, a *Boolean lattice* with 2^N elements is sometimes denoted L_2^N .

The next theorem presents the classic properties of any Boolean algebra. The first 4 pairs of properties are true for any lattice (Theorem 1.32 page 10). The *bounded*, *distributive*, and *complemented* properties are true by definition of a *Boolean lattice* (Definition 1.69 page 18).

Theorem 1.70 (classic 10 Boolean properties) 52 Let $A \triangleq (X, \lor, \land, 0, 1; \le)$ be an algebraic structure. In the event that A is a BOUNDED LATTICE (Definition 1.39 page 12), let x' represent a COMPLEMENT (Definition 1.63 page 17) of an element x in A.

0											
$x \lor x$	=	x	$x \wedge x$	=	x	(IDEMPOTENT)	and				
$x \lor y$	=	$y \lor x$	$x \wedge y$	=	$y \wedge x$	(COMMUTATIVE)	and				
$x \lor (y \lor z)$	=	$(x \lor y) \lor z$	$x \wedge (y \wedge z)$	=	$(x \land y) \land z$	(ASSOCIATIVE)	and				
$x \lor (x \land y)$	=	x	$x \wedge (x \lor y)$	=	x	(ABSORPTIVE)	and				
$x \lor 1$	=	1	$x \wedge 0$	=	0	(BOUNDED)	and				
$x \lor 0$	=	x	$x \wedge 1$	=	x	(IDENTITY)	and				
$x \lor (y \land z)$	=	$(x \lor y) \land (x \lor z)$	$x \land (y \lor z)$	=	$(x \land y) \lor (x \land z)$	(DISTRIBUTIVE)	and				
$x \lor x'$	=	1	$x \wedge x'$	=	0	(COMPLEMENTED)	and				
$(x \lor y)'$	=	$x' \wedge y'$	$(x \wedge y)'$	=	$x' \lor y'$	(de Morgan)	and				
		(x')	y' = x			(INVOLUTORY)					
disju	disjunctive properties				conjunctive properties						

A is a **Boolean algebra** $\iff \forall x, y, z \in X$

Proposition 1.71 (Huntington's fourth set) ⁵³ Let $A \triangleq (X, \lor, \land; \le)$ be an ALGEBRAIC STRUCTURE. A is a Boolean algebra \iff

ſ	1.	$x \lor x$	=	x	$\forall x \in X$	(IDEMPOTENT)	and]
J	2.	$x \lor y$	=	$y \lor x$	$\forall x, y \in X$	(COMMUTATIVE)	and
		$(x \lor y) \lor z$		$x \lor (y \lor z)$	$\forall x, y, z \in X$	(ASSOCIATIVE)	and (
l	4.	$(x' \lor y')' \lor (x' \lor y)'$	=	x	$\forall x, y \in X.$	(Huntington's axiom)	J

1.3 Orthocomplemented Lattices

Orthocomplemented lattices (Definition 1.72 page 20) are a kind of generalization of *Boolean algebras*. The relationship between lattices of several types, including orthocomplemented and Boolean lattices, is stated in Theorem 1.86 (page 26) and illustrated in Figure 4 (page 20).

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⁵² [[89], pages 292–293, ("1st set"), [[90], page 280, ("4th set"), [∞] [113], page 488, [∞] [68], page 10, [∞] [124], pages 20–21, [∞] [153], [∞] [167], pages 35–37

⁵³ [] [90], page 280, \langle "4th set" \rangle

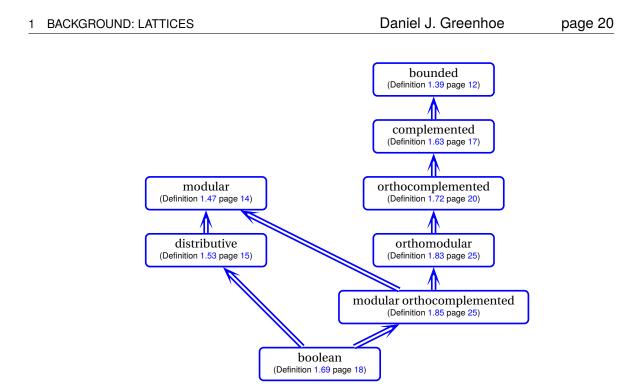


Figure 4: relationships between selected lattice types (see Theorem 1.86 page 26)

1.3.1 Definition

Definition 1.72 ⁵⁴ Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be a *bounded lattice* (Definition 1.39 page 12). An element $x^{\perp} \in X$ is an **orthocomplement** of an element $x \in X$ if

1.	$x^{\perp\perp}$	=	x	$\forall x \in X$	(involutory)	and
2.	$x \wedge x^{\perp}$	=	0	$\forall x \in X$	(non-contradiction)	and
3.	$x \leq y$	\Rightarrow	$y^{\perp} \leq x^{\perp}$	$\forall x, y \in X$	(antitone).	

The lattice *L* is **orthocomplemented** (*L* is an **orthocomplemented lattice**) if every element *x* in *X* has an *orthocomplement*. The elements $\{x, y\}$ are **orthocomplemented pairs** in *L* if $y = x^{\perp}$.

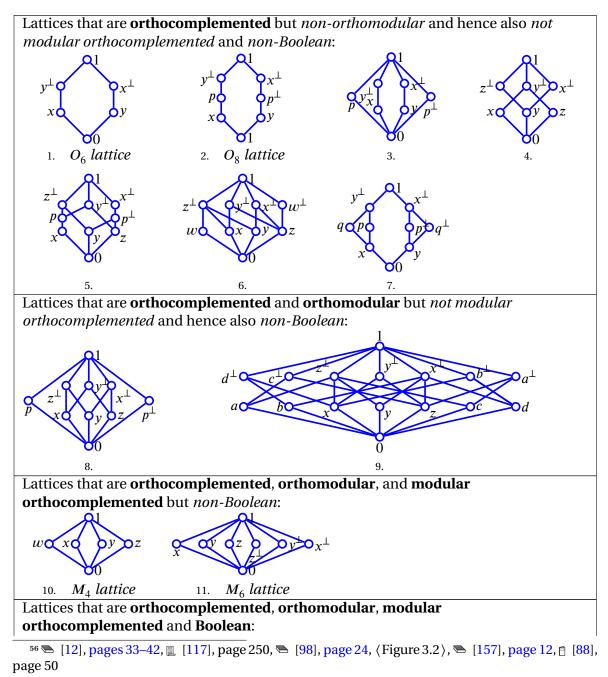
Definition 1.73 ⁵⁵ The **O**₆ **lattice** is the ordered set $(\{0, p, q, p^{\perp}, q^{\perp}, 1\}, \leq)$ with cover relation $q^{\perp} \circ q^{\perp} = \{(0, p), (0, q), (p, q^{\perp}), (q, p^{\perp}), (p^{\perp}, 1), (q^{\perp}, 1)\}.$ The O_6 lattice is illustrated by the Hasse diagram to the right.

⁵⁴ ∞ [157], page 11, ∞ [12], page 28, ∞ [98], page 16, ∞ [79], page 76, ∞ [112], page 3, <u>■</u> [20], page 830, ⟨L71–L73⟩

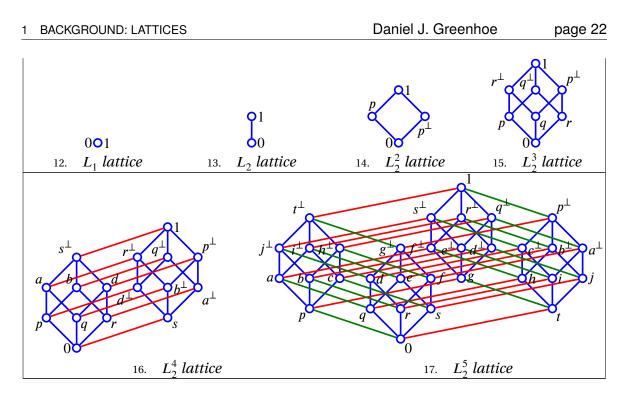
⁵⁵ [98], page 22, \bigcirc [88], page 50, \bigtriangledown [12], page 33, \backsim [157], page 12. The O_6 lattice is also called the **Benzene ring** or the **hexagon**.

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Example 1.74 ⁵⁶ There are a total of 10 **orthocomplemented lattices** with 8 elements or less. These 10, along with 3 other orthocomplemented lattices with 10 elements, are illustrated next:

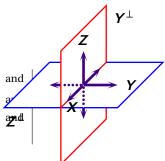


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Example 1.75 The structure $(2^{\mathbb{R}^N}, +, \cap, \emptyset, H; \subseteq)$ is an **orthocomplemented lattice** where

- $\mathcal{P} \mathbb{R}^N$ is an **Euclidean space** with dimension N
- $2^{\mathbb{R}^{N}}$ is the set of all subspaces of \mathbb{R}^{N}
- $\mathcal{P} = \mathbf{V} + \mathbf{W}$ is the *Minkowski sum* of subspaces \mathbf{V} and \mathbf{W}
- $\mathcal{D} \cap W$ is the *intersection* of subspaces V and W.



Example 1.76 The structure $(2^H, \oplus, \cap, \emptyset, H; \subseteq)$ is an **orthocomplemented lattice** where H is a **Hilbert space**, 2^H is the set of all closed subspaces of H, X + Y is the *Minkowski* sum of subspaces X and Y, $X \oplus Y \triangleq (X + Y)^-$ is the *closure* of X + Y, and $X \cap Y$ is the *intersection* of subspaces X and Y.

1.3.2 Properties

Theorem 1.77 57 Let x^{\perp} be the ORTHOCOMPLEMENT (Definition 1.72 page 20) of an element x in a BOUNDED LATTICE $L \triangleq (X, \lor, \land, 0, 1; \leq)$.

⁵⁷ **(12)**, pages 30–31, **(20)**, page 830, (L74), **(29)**, page 37, (3B.13. Theorem)

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BACKGROUND: LATTIC	ES				Daniel	J. Greenhoe	page 23
$\left. \begin{array}{c} L \ is \\ ortho- \\ complemented \end{array} \right\} \Longrightarrow$	$ \left\{\begin{array}{c} (1).\\ (2).\\ (3).\\ (4).\\ (5).\\ \end{array}\right. $	0^{\perp} 1^{\perp} $(x \lor y)^{\perp}$ $(x \land y)^{\perp}$ $x \lor x^{\perp}$	= = =	1 0 $x^{\perp} \wedge y^{\perp}$ $x^{\perp} \vee y^{\perp}$ 1	$\forall x, y \in X$ $\forall x, y \in X$ $\forall x \in X$	(BOUNDARY CONDITION) (BOUNDARY CONDITION) (DISJUNCTIVE DE MORGAN) (CONJUNCTIVE DE MORGAN, (EXCLUDED MIDDLE).	and and and) and

SPROOF: Let $x^{\perp} \triangleq \neg x$, where \neg is an *ortho negation* function (Definition 2.14 page 29). Then this theorem follows directly from Theorem 2.21 (page 30). Ð

Corollary 1.78 Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be a LATTICE (Definition 1.31 page 10). $\begin{cases}
L \text{ is orthocomplemented} \\
(Definition 1.72 page 20)
\end{cases} \implies \begin{cases}
L \text{ is complemented} \\
(Definition 1.63 page 17)
\end{cases}$

PROOF: This follows directly from the definition of *orthocomplemented lattices* (Definition 1.72 page 20) and complemented lattices (Definition 1.63 page 17). ø

Example 1.79



The O_6 lattice (Definition 1.73 page 20) illustrated to the left is both orthocomplemented (Definition 1.72 page 20) and multiply complemented (Definition 1.63 page 17). The lattice illustrated to the right is **multiply complemented**, but is non-orthocomplemented.

PROOF:

- (1) Proof that O_6 lattice is multiply complemented: b and q are both complements of p.
- (2) Proof that the right side lattice is multiply complemented: *a*, *p*, and *q* are all *complements* of r.

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Restrictions resulting in Boolean algebras 1.3.3

Proposition 1.8058 Let $L = (X, \lor, \land, 0, 1; \leq)$ be a BOUNDED LATTICE (Definition 1.39 page 12). $\left\{ \begin{array}{cc} 1. & L \text{ is orthocomplemented} \\ 2. & L \text{ is distributive} \end{array} \right.$ (Definition 1.53 page 15) $A = \left\{ \begin{array}{c} L \text{ is Boolean} \\ (Definition 1.69 page 18) \end{array} \right\}$ ⁵⁸ 🖱 [98], page 22

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page 24

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PROOF:

$$\left\{ \begin{array}{l} L \text{ is ortho complemented} \\ L \text{ is distributive} \end{array} \right\} \implies \left\{ \begin{array}{l} L \text{ is complemented} \\ L \text{ is distributive} \end{array} \right\} \qquad \text{by Corollary 1.78} \\ \implies \left\{ \begin{array}{l} L \text{ is Boolean} \end{array} \right\} \qquad \text{by Definition 1.69} \end{array}$$

The *center* of an *orthocomplemented lattice* is defined later, but here is a characterization involving it now anyways.

Proposition 1.81 Let $L = (X, \lor, \land, 0, 1; \le)$ be a LATTICE (Definition 1.31 page 10).

[®]Proof:

- (1) Proof that (1,2) \implies Boolean: *L* is Boolean because it satisfies Huntington's Fourth Set (Proposition 1.71 page 19), as demonstrated by the following ...
 - (a) Proof that $x \lor x = x$ (*idempotent*): *L* is a *lattice* (by definition of *L*), and all lattices are *idempotent* (Definition 1.31 page 10).
 - (b) Proof that $x \lor y = y \lor x$ (*commutative*): *L* is a *lattice* (by definition of *L*), and all lattices are *commutative* (Definition 1.31 page 10).
 - (c) Proof that $(x \lor y) \lor z = x \lor (y \lor z)$ (*associative*): *L* is a *lattice* (by definition of *L*), and all lattices are *associative* (Definition 1.31 page 10).
 - (d) Proof that $(x^{\perp} \lor y^{\perp})^{\perp} \lor (x^{\perp} \lor y)^{\perp} = x$ (*Huntington's axiom*):

 $(x^{\perp} \lor y^{\perp})^{\perp} \lor (x^{\perp} \lor y)^{\perp}$ = $(x^{\perp} \perp \land y^{\perp} \perp) \lor (x^{\perp} \perp \land y^{\perp})$ by *de Morgan* property (Theorem 1.77 page 22) = $(x \land y) \lor (x \land y^{\perp})$ by *involution* property (Definition 1.72 page 20) = x by def. of *center* (Definition 3.15 page 37)

- (2) Proof that (1) \leftarrow Boolean:
 - (a) Proof that $x \lor x^{\perp} = 1$: by definition of *Boolean algebras* (Definition 1.69 page 18).
 - (b) Proof that $x \wedge x^{\perp} = 0$: by definition of *Boolean algebras* (Definition 1.69 page 18).
 - (c) Proof that $x^{\perp \perp} = x$: by *involutory* property of *Boolean algebra* (Theorem 1.70 page 19).
 - (d) Proof that $x \le y \implies y^{\perp} \le x^{\perp}$:

 $\begin{array}{lll} y^{\perp} \leq x^{\perp} & \Longleftrightarrow \ y^{\perp} & = y^{\perp} \wedge x^{\perp} & \text{by Lemma 1.33 page 10} \\ & \Leftrightarrow \ y^{\perp \perp} & = (y^{\perp} \wedge x^{\perp})^{\perp} & \\ & \Leftrightarrow \ y^{\perp \perp} & = y^{\perp \perp} \lor x^{\perp \perp} & \text{by de Morgan property (Theorem 1.70 page 19)$} \\ & \Leftrightarrow \ y & = y \lor x & \text{by $involutory$ property (Theorem 1.70 page 19)$} \\ & \Leftrightarrow \ y & = y & \text{by $x \leq y$ hypothesis$} \end{array}$

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(3) Proof that (2) \leftarrow Boolean: for all $x, y \in L$

$$(x \land y) \lor (x \land y^{\perp}) = [(x \land y) \lor x] \land [(x \land y) \lor y^{\perp}]$$

= $x \land [(x \land y) \lor y^{\perp}]$
= $x \land [(x \lor y^{\perp}) \land (y \lor y^{\perp})]$
= $x \land (x \lor y^{\perp}) \land 1$
= x
 $\implies x © y \quad \forall x, y \in L$
 $\implies x \text{ is in the center of } L$

by *distributive* property (Theorem 1.70 page 19) by *absorptive* property (Theorem 1.70 page 19) by *distributive* property (Theorem 1.70 page 19) by *complement* property (Theorem 1.70 page 19) by *absorptive* property (Theorem 1.70 page 19) by Definition 3.9 page 36 by Definition 3.15 page 37



Example 1.82 The O_6 *lattice* (Definition 1.73 page 20) illustrated to the left is **orthocomplemented** (Definition 1.72 page 20) but **non-join-distributive** (Definition 1.53 page 15), and hence *non-Boolean*. The lattice illustrated to the right is **orthocomplemented** *and* **distributive** and hence also **Boolean** (Proposition 1.80 page 23).

1.3.4 Orthomodular lattices

Definition 1.83 ⁵⁹ Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be a *bounded lattice* (Definition 1.39 page 12). *L* is **orthomodular** if

1. *L* is orthocomplemented and 2. $x \le y \implies x \lor (x^{\perp} \land y) = y \forall x, y \in X$ (orthomodular identity)

Theorem 1.84 ⁶⁰ Let $L = (X, \lor, \land, 0, 1; \le)$ be an algebraic structure.

$$\left\{\underbrace{\begin{array}{c} \mathbf{L} \text{ is an orthomodular lattice} & and \\ (x \wedge y^{\perp})^{\perp} = y \lor (x^{\perp} \wedge y^{\perp}) \\ \overbrace{\text{ElKAN'S LAW}}^{\text{ELKAN'S LAW}} & \forall x, y \in X \end{array}\right\} \implies \left\{\begin{array}{c} \mathbf{L} \text{ is a} \\ \textbf{Boolean algebra} \\ (\text{Definition 1.69 page 18}) \end{array}\right\}$$

Definition 1.85 Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be a *bounded lattice* (Definition 1.39 page 12). *L* is a **modular orthocomplemeted lattice** if

- 1. L is orthocomplemented (Definition 1.72 page 20) and
- 2. *L* is modular (Definition 1.47 page 14)

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⁵⁹ **●** [98], page 22, **●** [110], page 90, **■** [91] ⁶⁰ **■** [144], page 72

Theorem 1.86 ⁶¹ Let *L* be a lattice.

{ <i>L is</i> Boolean}	\implies	{L is modular orthocomplemented	(Definition 1.85 page 25)}
	\implies	{ L <i>is</i> orthomodular	(Definition 1.83 page 25)}
	\Rightarrow	{ <i>L is</i> orthocomplemented	(Definition 1.72 page 20)}

{*L is* orthocomplemented

Background: functions on lattices 2

2.1 Valuations

Definition 2.1 ⁶² Let $L \triangleq (X, \lor, \land; \leq)$ be a *lattice* (Definition 1.31 page 10). A function $v \in \mathbb{R}^X$ is a valuation on \overline{L} if

 $v(x \lor y) + v(x \land y) = v(x) + v(y) \quad \forall x, y \in X$

Proposition 2.2 Let $v \in \mathbb{R}^X$ be a FUNCTION on a LATTICE $L \triangleq (X, \lor, \land; \leq)$ (Definition 1.31) page 10).

 $\{ L \text{ is LINEAR (Definition 1.31 page 10)} \} \implies \{ v \text{ is a VALUATION (Definition 2.1 page 26)} \}$

SPROOF: Let $x, y \in X$ such that $x \le y$ or $y \le x$.

$$v(x \lor y) + v(x \land y) = v(x) + v(y)$$
 because *L* is *linear*

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Example 2.3 ⁶³ Consider the *real valued lattice* $L \triangleq (\mathbb{R}, \max, \min; \leq)$. The *absolute value* function $|\cdot|$ is a *valuation* on *L*.

PROOF: *L* is *linear* (Definition 1.31 page 10), so v is a *valuation* by Proposition 2.2 (page 26). ø

Definition 2.4 ⁶⁴ Let *X* be a set and \mathbb{R}^{\vdash} the set of non-negative real numbers. A function $d \in \mathbb{R}^{\vdash X \times X}$ is a **metric** on *X* if

1.	d(x, y)	\geq	0	$\forall x, y \in X$	(non-negative)	and			
2.	d(x, y)	=	$0 \iff x = y$	$\forall x, y \in X$	(nondegenerate)	and			
3.	d(x, y)	=	d(y, x)	$\forall x, y \in X$	(symmetric)	and			
4.	d(x, y)	\leq	d(x,z) + d(z,y)	$\forall x, y, z \in X$	(subadditive/triangle inequality).65				
	the many is the main (V, I). A methic is also called a distance for atten								

A metric space is the pair (X, d). A *metric* is also called a **distance function**.

[43], page 105, ((8.1.1)), [41], page 143, (§10.3), [42], page 193, (§10.3)

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⁶¹ ► [98], page 32, (20.), 🖞 [94], page 57

⁶² S [93], page 127, S [18], page 230, (Definition X.1(V1)), S [22], page 58, (Exercise 4.25),

⁶³ ► [101], page 119, (§5.7)

⁶⁴ 🖻 [45], page 28, 🖻 [31], page 21, 🖻 [82], page 109, 🖱 [64], 🖱 [63], page 30

page 27

Definition 2.566Let (X, d) be a metric space (Definition 2.4 page 26).An open ball centered at x with radius ris the set $B(x, r) \triangleq \{y \in X | d(x, y) \leq r\}.$ A closed ball centered at x with radius ris the set $\overline{B}(x, r) \triangleq \{y \in X | d(x, y) \leq r\}.$ A unit ball centered at xis the set $\overline{B}(x, r) \triangleq \{y \in X | d(x, y) \leq r\}.$ A closed unit ball centered at xis the set $\overline{B}(x, 1).$ A closed unit ball centered at xis the set $\overline{B}(x, 1).$

Theorem 2.6 ⁶⁷ Let $v \in \mathbb{R}^X$ be a function on a LATTICE $L \triangleq (X, \lor, \land; \leq)$ (Definition 1.31 page 10). 1. $v(x \lor y) + v(x \land y) = v(x) + v(y) \quad \forall x, y \in X$ (VALUATION) and 2. $x \le y \implies v(x) \le v(y) \quad \forall x, y \in X$ (ISOTONE) $\end{cases} \implies \begin{cases} d(x, y) \triangleq \\ v(x \lor y) - v(x \land y) \\ is \ a \ METRIC \ on \ L \end{cases}$

Definition 2.7 ⁶⁸ Let v be a *valuation* (Definition 2.1 page 26) on a *lattice* $L \triangleq (X, \lor, \land; \le)$ (Definition 1.31 page 10). Let d(x, y) be the *metric* defined in Theorem 2.6 (page 27). The pair (L, d) is called a *metric lattice*.

For *finite modular* lattices, the *height* function h(x) (Definition 1.42 page 12) can serve as the isotone valuation that induces a metric (next proposition).

Proposition 2.8 ⁶⁹ Let h(x) be the HEIGHT (Definition 1.42 page 12) of a point x in a BOUNDED LATTICE (Definition 1.39 page 12) $L \triangleq (X, \lor, \land, 0, 1; \le)$.

 $\{ 1. L is MODULAR and 2. L is FINITE \}$

\Rightarrow	<i>∫</i> 1.	$h(x \lor y) + h(x \land y) = h(x) + h(y)$	$\forall x, y \in X$	(VALUATION)	and 🔪
\rightarrow	2.	$x \lneq y \implies h(x) \lneq h(y)$	$\forall x, y \in X$	(POSITIVE)	ſ
	<i>∫</i> 1.	$h(x \lor y) + h(x \land y) = h(x) + h(y)$	$\forall x, y \in X$	(VALUATION)	and)
\rightarrow	2.	$h(x \lor y) + h(x \land y) = h(x) + h(y)$ $x \lneq y \implies h(x) \lneq h(y)$ $h(x \lor y) + h(x \land y) = h(x) + h(y)$ $x \leq y \implies h(x) \leq h(y)$	$\forall x, y \in X$	(ISOTONE)	Ĵ

Theorem 2.9 ⁷⁰ Let v be a VALUATION (Definition 2.1 page 26) on a LATTICE $L \triangleq (X, \lor, \land; \leq)$ (Definition 1.31 page 10). Let d(x, y) be the METRIC defined in Theorem 2.6 (page 27).

 $\left\{\begin{array}{c} (\boldsymbol{L},\mathsf{d}) \text{ is a METRIC LATTICE} \\ (Definition 2.7 page 27) \end{array}\right\} \implies \left\{\begin{array}{c} \boldsymbol{L} \text{ is MODULAR} \\ (Definition 1.47 page 14) \end{array}\right\}$

⁶⁸ • [43], page 105, • [18], page 231, (§X.2)

⁶⁹ 🖱 [18], page 230

⁷⁰ (18), page 232, (Theorem X.2), (43), pages 105–106, (22), page 58, (Exercise 4.25)

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⁶⁵ (55), (Book I Proposition 20)

⁶⁶ 🖱 [5], page 35

⁶⁷ **(**43], page 105, ((8.1.2)), **(**18], pages 230–231

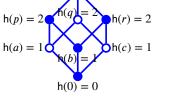
Example 2.10

The function h on the *Boolean* (and thus also *modular*) lattice L_2^3 illustrated to the right is a *valuation* (Definition 2.1 page 26) that is *positive* (and thus also *isotone*, Proposition 2.8 page 27). Therefore

$$\begin{split} \mathsf{d}(x,y) &\triangleq \mathsf{h}(x \lor y) - \mathsf{h}(x \land y) & \forall x,y \in X \\ \text{is a } metric (\text{Definition 2.7 page 27}) \text{ on } \boldsymbol{L}_2^3. \text{ For example,} \\ \mathsf{d}(b,q) &\triangleq \mathsf{h}(b \lor q) - \mathsf{h}(b \land q) = \mathsf{h}(1) - \mathsf{h}(0) = 3 - 0 = 3 \\ \text{The } closed unit ball \text{ centered at } b \text{ (Definition 2.5 page 27) and illustrated} \end{split}$$

with solid dots to the right is

 $\mathsf{B}(b,1) \triangleq \{x \in X \mid \mathsf{d}(b,x) \le 1\} = \{b, p, r, 0\}$



h(1) = 3

Example 2.11

The *height* function h (Definition 1.42 page 12) on the *orthocomplemented* but *non-modular* lattice O_6 illustrated to the right is *not* a *valuation* because for example

 $h(a \lor c) + h(a \land c) = h(1) + h(0) = 3 + 0 = 3 \neq 2 = 1 + 1 = h(a) + h(b)$. Moreover, we might expect the "distance" from *a* to *c* to be 2. However, if we attempt to use h(x) to define a metric on O_6 , then we get

 $d(a, c) \triangleq h(a \lor c) - h(a \land c) = h(1) - h(0) = 3 - 0 = 3 \neq 2.$

h(1) = 3 h(p) = 2 h(a) = 1h(0) = 0

2.2 Negation

2.2.1 Definitions

Definition 2.12 ⁷¹ Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be a *bounded lattice* (Definition 1.39 page 12). A *function* $\neg \in X^X$ is a **subminimal negation** on L if ⁷² $x \le y \implies \neg y \le \neg x \quad \forall x, y \in X \quad (antitone).$

Definition 2.13 ⁷³ Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be a *bounded lattice* (Definition 1.39 page 12).

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⁷¹ S [51], pages 4–6, S [52], pages 24–26, (2 The Kite of Negations)

⁷² In the context of natural language, D. Devidi has argued that, *subminimal negation* (Definition 2.12 page 28) is "difficult to take seriously as" a negation. For further details see [40], page 511, [39], page 568, $[[77], (\S2.1.1), \textcircled{51}, (\S11.1)$

⁷³ (51], pages 4–6, (52], pages 24–26, (2 The Kite of Negations), (161], page 4, (1.6] Intuitionism. (b), (162), page 11, (Definition 16), (70), page 21, (Definition 3.3), (132), page 50, (Definition 2.26), (131), pages 98–99, (5.4 Negations), (10), pages 155–156, $((N1) \neg 0 = 1 \text{ and } \neg 1 = 0$, $(N3) \neg \neg x = x$

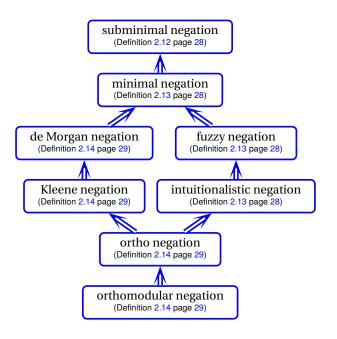
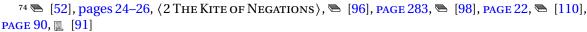


Figure 5: lattice of negations

A function $\neg \in X^X$ is a **negation**, or **minimal negation**, on *L* if 1. $x \leq y \implies \neg y \leq \neg x \quad \forall x, y \in X$ (antitone) and 2. x \leq $\forall x \in X$ (weak double negation). $\neg \neg x$ A *minimal negation* ¬ is an **intuitionistic negation** on *L* if 3. $x \wedge \neg x$ = 0 $\forall x, y \in X$ (non-contradiction). A *minimal negation* ¬ is a **fuzzy negation** on *L* if 4. ¬1 = 0 (boundary condition). **Definition 2.14** ⁷⁴ Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be a *bounded lattice* (Definition 1.39 page 12). A *minimal negation* ¬ is a **de Morgan negation** on *L* if = 5. *x* $\neg \neg x$ $\forall x \in X$ (involutory). A *de Morgan negation* ¬ is a **Kleene negation** on *L* if $y \lor \neg y$ 6. $x \wedge \neg x < d$ $\forall x, y \in X$ (*Kleene condition*). A *de Morgan negation* ¬ is an **ortho negation** on *L* if 7. $x \wedge \neg x =$ 0 $\forall x, y \in X$ (non-contradiction). A *de Morgan negation* ¬ is an **orthomodular negation** on *L* if 8. $x \wedge \neg x =$ 0 $\forall x, y \in X$ (non-contradiction) and 9. $x \leq y \implies x \lor (x^{\perp} \land y) = y \quad \forall x, y \in X \quad (orthomodular).$



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Remark 2.15 ⁷⁵ The *Kleene condition* is a weakened form of the *non-contradiction* and *excluded middle* properties in the sense $x \land \neg x = 0 \le 1 = y \lor \neg y$.

Definition 2.16 Let $L \triangleq (X, \lor, \land, \neg, 0, 1; \le)$ be a *bounded lattice* (Definition 1.39 page 12) with a function $\neg \in X^X$. If \neg is a *negation* (Definition 2.13 page 28), then L is a **lattice with negation**.

2.2.2 Properties of negations

Theorem 2.17 ⁷⁶ Let
$$\neg \in X^X$$
 be a function on a BOUNDED LATTICE $L \triangleq (X, \lor, \land, 0, 1; \le)$.

$$\begin{cases} \neg is a \\ FUZZY \text{ NEGATION} \end{cases} \implies \{ \neg 0 = 1 \text{ (BOUNDARY CONDITION)} \}$$

Theorem 2.18 ⁷⁷ Let $\neg \in X^X$ be a function on a BOUNDED LATTICE $L \triangleq (X, \lor, \land, 0, 1; \le)$.

$(\neg is an)$		(a)	-1	=	0	(BOUNDARY CONDITION)	and
{	$\} \Longrightarrow \{$	(b)	¬0	=	1	(BOUNDARY CONDITION)	and
(INTUITIONISTIC NEGATION		(c)	¬ is	a F	UZZ	Y NEGATION	J

Theorem 2.19 ⁷⁸ Let $\neg \in X^X$ be a function on a BOUNDED LATTICE $L \triangleq (X, \lor, \land, 0, 1; \le)$. $\begin{cases} \neg is a \\ minimal \\ negation \end{cases} \implies \begin{cases} \neg x \lor \neg y \le \neg (x \land y) \quad \forall x, y \in X \text{ (conjunctive de Morgan inequality) and} \\ \neg (x \lor y) \le \neg x \land \neg y \quad \forall x, y \in X \text{ (disjunctive de Morgan inequality)} \end{cases}$

Theorem 2.20 ⁷⁹ Let $\neg \in X^X$ be a function on a BOUNDED LATTICE $L \triangleq (X, \lor, \land, 0, 1; \le)$. \neg is a de Morgan negation $\rbrace \implies \begin{cases} \neg(x \lor y) = \neg x \land \neg y \quad \forall x, y \in X \text{ (DISJUNCTIVE DE MORGAN)} and \\ \neg(x \land y) = \neg x \lor \neg y \quad \forall x, y \in X \text{ (CONJUNCTIVE DE MORGAN)} \end{cases}$

Theorem 2.21 ⁸⁰ Let $\neg \in X^X$ be a function on a BOUNDED LATTICE $L \triangleq (X, \lor, \land, 0, 1; \le)$.

	1	1.	$\neg 0$	=	1		(BOUNDARY CONDITION)	and	
$\left\{\begin{array}{l}\neg \text{ is an}\\ \text{ortho}\\ \text{negation}\end{array}\right\} =$		2.	-1	=	0		(BOUNDARY CONDITION)	and	
		З.	$\neg(x \lor y)$	=	$\neg x \land \neg y$	$\forall x, y \in X$	(disjunctive de Morgan) (conjunctive de Morgan)	and	l
	r — J	4.	$\neg(x \land y)$	=	$\neg x \lor \neg y$	$\forall x, y \in X$	(CONJUNCTIVE DE MORGAN)	and	ſ
		5.	$x \lor \neg x$	=	1	$\forall x \in X$	(EXCLUDED MIDDLE)	and	
		6.	$x \wedge \neg x$	\leq	$y \vee \neg y$	$\forall x, y \in X$	(KLEENE CONDITION).	J	

75 🝋 [26], page 78

⁷⁶ [77], $\langle \$2.1.2 \rangle$, [78], $\langle \$11.2 \rangle$ ⁷⁷ [77], $\langle \$2.1.2 \rangle$, [78], $\langle \$11.2 \rangle$ ⁷⁸ [77], $\langle \$2.1.2 \rangle$, [78], $\langle \$11.2 \rangle$ ⁷⁹ [77], $\langle \$2.1.2 \rangle$, [78], $\langle \$11.2 \rangle$ ⁷⁹ [77], $\langle \$2.1.2 \rangle$, [78], $\langle \$11.2 \rangle$

80 [77], ⟨\$2.1.2⟩, [78], ⟨\$11.2⟩

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2.3 Projections

Definition 2.22 ⁸¹ Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be an *orthocomplemented lattice* (Definition 1.72 page 20). A function $\phi_x \in X^X$ is a **Sasaki projection** on $x \in X$ if $\phi_x(y) \triangleq (y \lor x^{\perp}) \land x$. The *Sasaki projections* ϕ_x and ϕ_y are **permutable** if $\phi_x \circ \phi_y(u) = \phi_y \circ \phi_x(u) \quad \forall u \in X$.

Proposition 2.23 Let $\phi_x(y)$ be the SASAKI PROJECTION OF y ONTO x (Definition 2.22 page 31) in an ORTHOCOMPLEMENTED LATTICE $L \triangleq (X, \lor, \land, 0, 1; \le)$.

(1). $x \leq y$	\Rightarrow		$\phi_x(y)$	=	x	$\forall x, y \in X$
$(2). y \le x$	$\implies y$	\leq	$\phi_x(y)$	\leq	x	$\forall x, y \in X$
(3). $y \le x$ and L is BOOLEAN	\implies		$\phi_{x}(y)$	=	y	$\forall x, y \in X$

[®]Proof:

$$\begin{array}{ll} (1) \implies \phi_x(y) \triangleq \left(y \lor x^{\perp} \right) \land x & \text{by definition of } Sasaki \ projection \ (Definition 2.22 \ page 31)} \\ &= 1 \land x & \text{by } x \le y \ \text{hypothesis and Proposition } 3.1 \ page 34 \\ &= x & \text{by property of bounded lattices (Proposition 1.41 \ page 12)} \\ (2) \implies \boxed{y} = y \land x & \text{by } y \le x \ \text{hypothesis} \\ &\leq (y \lor x^{\perp}) \land x & \text{by definition of } \lor (Definition 1.27 \ page 9)} \\ &= \boxed{\phi_x(y)} & \text{by definition of } Sasaki \ projection \ (Definition 2.22 \ page 31)} \\ &\leq (y \lor x^{\perp}) \land x & \text{by definition of } Sasaki \ projection \ (Definition 2.22 \ page 31)} \\ &\leq (y \lor x^{\perp}) \land x & \text{by definition of } Sasaki \ projection \ (Definition 2.22 \ page 31)} \\ &\leq (x) & \text{by definition of } Sasaki \ projection \ (Definition 2.22 \ page 31)} \\ &\leq [x] & \text{by definition of } \land (Definition 1.28 \ page 9)} \\ &(3) \implies \phi_x(y) = (y \lor x^{\perp}) \land x & \text{by definition of } Sasaki \ projection \ (Definition 2.22 \ page 31)} \\ &= (y \land x) \lor (x^{\perp} \land x) & \text{by definition of } Sasaki \ projection \ (Definition 2.22 \ page 31)} \\ &= (y \land x) \lor (x^{\perp} \land x) & \text{by definition of } Sasaki \ projection \ (Definition 2.22 \ page 31)} \\ &= (y \land x) \lor (x^{\perp} \land x) & \text{by definition of } Sasaki \ projection \ (Definition 2.22 \ page 31)} \\ &= (y \land x) \lor (x^{\perp} \land x) & \text{by definition of } Sasaki \ projection \ (Definition 2.22 \ page 31)} \\ &= (y \land x) \lor (x^{\perp} \land x) & \text{by distributive property of Boolean lattices \ (Theorem 1.70 \ page 19)} \\ &= (y \land x) \lor 0 & \text{by non-contradiction of } Boolean \ lattices \ (Proposition 1.41 \ page 12)} \\ &= y & \text{by } y \le x \ \text{hypothesis and definition of } \land (Definition 1.28 \ page 9)} \end{aligned}$$

Proposition 2.24 Let $\phi_x(y)$ be the SASAKI PROJECTION OF y ONTO x (Definition 2.22 page 31) in an ORTHOCOMPLEMENTED LATTICE $(X, \lor, \land, 0, 1; \le)$.

(1). $\phi_0(y) = 0 \quad \forall y \in X$ (2). $\phi_x(0) = 0 \quad \forall x \in X$ (3). $\phi_1(y) = 1 \quad \forall y \in X$ (4). $\phi_x(1) = x \quad \forall x \in X$ (5). $\phi_x(x^{\perp}) = 0 \quad \forall x \in X$

⁸¹ S [127], pages 158–159, (equation (S)), S [152], page 300, (Def.5.1, cf Foulis 1962), S [98], page 117

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[∞]Proof:

$\phi_0(y) = 0$	because $0 \le y$ and by Proposition 2.23 page 31
$\phi_{\boldsymbol{x}}(0) \triangleq \left(\boldsymbol{0} \vee \boldsymbol{x}^{\perp} \right) \wedge \boldsymbol{x}$	by definition of Sasaki projection (Definition 2.22 page 31)
$= x^{\perp} \wedge x$	by property of bounded lattices (Proposition 1.41 page 12)
= 0	by definition of <i>orthocomplemented</i> (Definition 1.72 page 20)
$\phi_1(y) \triangleq \left(y \vee 1^{\perp} \right) \wedge 1$	by definition of Sasaki projection (Definition 2.22 page 31)
$= (y \lor 0) \land 1$	by <i>boundary condition</i> (Theorem 2.21 page 30)
$= y \wedge 1$	by property of bounded lattices (Proposition 1.41 page 12)
= 1	by property of bounded lattices (Proposition 1.41 page 12)
$\phi_x(1) = x$	because $x \le 1$ and by Proposition 2.23 page 31
$\phi_{\boldsymbol{x}} \big(\boldsymbol{x}^{\perp} \big) \triangleq \big(\boldsymbol{x}^{\perp} \lor \boldsymbol{x}^{\perp} \big) \land \boldsymbol{x}$	by definition of Sasaki projection (Definition 2.22 page 31)
$= x^{\perp} \wedge x$	by <i>idempotency</i> of lattices (Theorem 1.32 page 10)
= 0	by non-contradiction prop. of orthocomplemented lattice (Definition 1.72 page 20)

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Example 2.25 Here are some examples of projections in the O_6 *lattice* onto the element *x*:

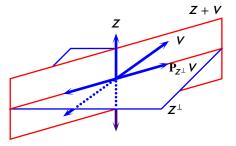
$\phi_p(q)$	≜	(q	V	p^{\perp})	Λ	р	=	p^{\perp}	Λ	р	=	0	(because $p \perp q$)	01
$\phi_p(p^{\perp})$	≜	(p^{\perp})	V	p^{\perp})	Λ	р	=	p^{\perp}	Λ	р	=	0	(because $p \perp p^{\perp}$)	
$\phi_p(q^{\perp})$	≜	$(q^{\perp}$	V	p^{\perp})	\wedge	р	=	1	Λ	р	=	р	(because $p \le q^{\perp}$)	$q^{\perp} \mathbf{q} \mathbf{q} \mathbf{p}^{\perp}$
$\phi_{q^{\perp}}(p)$	≜	(<i>p</i>	V	q)	\wedge	q^{\perp}	=	1	Λ	q^{\perp}	=	q^{\perp}	(because $q^{\perp} \leq 1$)	$p \diamond \phi q$
$\phi_p(1)$	≜	(1	V	p^{\perp})	\wedge	р	=	1	Λ	р	=	р	(because $p \le 1$)	
$\dot{\phi_p}(0)$	≜	(0	V	p^{\perp})	\wedge	р	=	p^{\perp}	Λ	р	=	0	(because $p \perp 0$)	

Example 2.26

Let \mathbb{R}^3 be the 3-dimensional Euclidean space (Example 1.75 page 22) with subspaces Z and V. Then the projection operator $P_{Z^{\perp}}$ onto Z^{\perp} is a sasaki projection $\phi_{Z^{\perp}}$. In particular

$$P_{Z^{\perp}}V \triangleq \phi_{Z^{\perp}}(V) \\ \triangleq (V + Z^{\perp \perp}) \cap Z^{\perp} \\ = (V + Z) \cap Z^{\perp}$$

as illustrated to the right.



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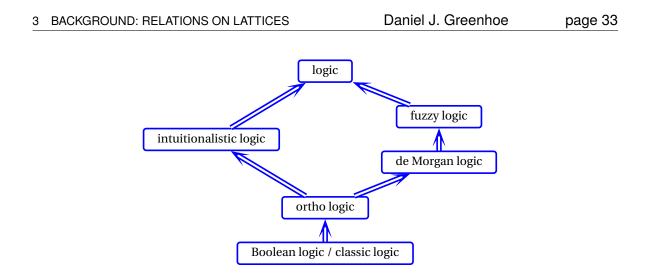


Figure 6: lattice of logics

2.4 Logics

Definition 2.27 ⁸² Let \rightarrow be an *implication* function defined on a *lattice with negation* $L \triangleq (X, \lor, \land, \neg, 0, 1; \leq)$ (Definition 2.16 page 30).

if ¬ is a <i>minimal negation</i> .
if ¬ is a <i>fuzzy negation</i> .
if \neg is an <i>intuitionalistic negation</i> .
if ¬ is a <i>de Morgan negation</i> .
if ¬ is a <i>Kleene negation</i> .
if ¬ is an <i>ortho negation</i> .
if ¬ is an <i>ortho negation</i> and
L is Boolean.

For examples and a definition of *implication*, see [77], (\$3.1).

3 Background: relations on lattices

The relations in this section are typically defined on an *orthocomplemented lattice* (Definition 1.72 page 20). Here, some relations are generalized to a *lattice with negation* (Definition 2.16 page 30). A *lattice* (Definition 1.31 page 10) with an *ortho negation* successfully defined on it is an *orthocomplemented lattice* (Definition 1.72 page 20). In many cases, these relations only work

⁸² [[159], page 136, (Definition 2.1), [[162], page 11, (Definition 16), [[77], (§3.1)

well on an *orthocomplemented lattice*, and thus many results are restricted to orthocomplemented lattices.

3.1 Orthogonality

Proposition 3.1 Let $(X, \lor, \land, 0, 1; \le)$ be an ORTHOCOMPLEMENTED LATTICE (Definition 1.72 page 20).

x < y	\rightarrow	ſ	x^{\perp}	V	у	=	1	and	
$x \leq y$		J	x	\wedge	y^{\perp}	=	0	,	$\forall x, y \in X$

[®]Proof:

$$x \le y \implies x \lor x^{\perp} \le y \lor x^{\perp}$$
by monotone property of lattices (Proposition 1.34 page 11) $\implies 1 \le y \lor x^{\perp}$ by excluded middle property (Definition 1.72 page 20) $\implies x^{\perp} \lor y = 1$ by upper bounded property of bounded lattices (Definition 1.39 page 12) $x \le y \implies x \land y^{\perp} \le y \land y^{\perp}$ by monotone property of lattices (Proposition 1.34 page 11) $\implies x \land y^{\perp} \le 0$ by monotone property of lattices (Proposition 1.34 page 11) $\implies x \land y^{\perp} \le 0$ by non-contradiction property (Definition 1.72 page 20) $\implies x \land y^{\perp} = 0$ by lower bounded property of bounded lattices (Definition 1.39 page 12)

Definition 3.2 ⁸³ Let $(X, \lor, \land, \neg, 0, 1; \le)$ be a *lattice with negation* (Definition 2.16 page 30). The **orthogonality** relation $\bot \in 2^{XX}$ is defined as

 $x \perp y \qquad \stackrel{\text{def}}{\iff} \qquad x \leq \neg y$ If $x \perp y$, we say that x is **orthogonal** to y.

Lemma 3.3 Let $(X, \lor, \land, \neg, 0, 1; \le)$ be a LATTICE WITH NEGATION (Definition 2.16 page 30). $\begin{cases} x \perp y \quad (\text{ORTHOGONAL Definition 3.2 page 34}) \end{cases} \implies \begin{cases} y \perp x \quad (\text{SYMMETRIC}) \end{cases}$

[®]Proof:

$x \perp y \implies x \le \neg y$	by definition of \perp (Definition 3.2 page 34)
\implies $(\neg \neg y) \leq \neg x$	by <i>antitone</i> property (Definition 1.72 page 20)
$\implies y \le \neg x$	by <i>weak double negation</i> property of <i>negation</i> (Definition 2.13 page 28)
$\implies y \perp x$	by definition of \perp (Definition 3.2 page 34)

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⁸³ 🖱 [157], page 12, 🖱 [112], page 3

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Lemma 3.4 ⁸⁴ Let $(X, \lor, \land, 0, 1; \leq)$ be an ORTHOCOMPLEMENTED LATTICE (Definition 1.72 page 20).

 $\underbrace{x \perp y}_{\text{ORTHOGONAL (Definition 3.2 page 34)}} \implies \left\{ \begin{array}{ll} 1. & x \wedge y &= 0 & and \\ 2. & x^{\perp} \vee y^{\perp} &= 1 \end{array} \right\}$

Remark 3.5 In an *orthocomplemented lattice* L, the *orthogonality* relation \perp is in general *non-associative*. That is,

 $\left\{\begin{array}{cc} x \perp y & \text{and} \\ y \perp z \end{array}\right\} \implies x \perp z$

SPROOF: Consider the L_2^4 Boolean lattice in Example 1.74 (page 21).

- $\Rightarrow a^{\perp} \perp p$ because $a^{\perp} \leq p^{\perp}$.
- $\Rightarrow p \perp r$ because $p \leq r^{\perp}$.
- ↔ But yet a^{\perp} is *not* orthogonal to *r* because $a^{\perp} \leq r^{\perp}$.

Example 3.6 In the O_6 lattice (Definition 1.73 page 20), there are a total of $\binom{6}{2} = \frac{6!}{(6-2)!2!} = \frac{6\times5}{2} = 15$ distinct unordered (the \perp relation is *symmetric* by Lemma 3.3 page 34 so the order doesn't matter) pairs of elements.

Of these 15 pairs, 8 are orthogonal to each other, and 0 is orthogonal to itself, making a total of 9 orthogonal pairs:

x	\bot	У	x	\bot	0	y^{\perp}	\bot	0
x	\bot	x^{\perp}	у	\bot	0	1	\bot	0
у	\bot	$y \\ x^{\perp} \\ y^{\perp}$	x^{\perp}	\bot	0	0	\bot	0

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Example 3.7 In lattice 5 of Example 1.74 (page 21), there are a total of $\binom{10}{2} = \frac{10!}{(10-2)!2!} = \frac{10\times9}{2} = 45$ distinct unordered pairs of elements.

Of these 45 pairs, 18 are orthogonal to each other, and 0 is orthogonal to itself, making a total of 19 orthogonal pairs:

р	T	p^{\perp}	x	T	x^{\perp}	У	\bot	Z.	x^{\perp}	T	0
р	\bot	x^{\perp}	x	\bot	у	y	\bot	0	y^{\perp}	\bot	0
р	\bot	У	x	\bot	z	z	\bot	z^{\perp}	z^{\perp}	\bot	0
р	\bot	z	x	\bot	0	z	\bot	0	0	\bot	0
р	\bot	0	y	\bot	y^{\perp}	p^{\perp}	\bot	0			

Example 3.8 In the \mathbb{R}^3 Euclidean space illustrated in Example 1.75 (page 22),

$$\begin{array}{cccc} X \subseteq Y^{\perp} & \Longrightarrow & X \perp Y & Y \subseteq X^{\perp} & \Longrightarrow & Y \perp X \\ X \subseteq Z^{\perp} & \Longrightarrow & X \perp Z & Y \subseteq Z^{\perp} & \Longrightarrow & Y \perp Z \\ X \wedge Y = X \wedge Z = Y \wedge Z = 0 \end{array}$$

⁸⁴ 🖱 [87], page 67, 🖱 [78], (Lemma 13.2)

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page 36

3.2 Commutativity

The *commutes* relation is defined next. Motivation for the name "commutes" is provided by Proposition 3.14 (page 36) which shows that if *x* commutes with *y* in a lattice *L*, then *x* and *y* commute in the *Sasaki projection* $\phi_x(y)$ on *L*.

Definition 3.9 ⁸⁵ Let $L \triangleq (X, \lor, \land, \neg, 0, 1; \le)$ be a *lattice with negation* (Definition 2.16 page 30). The **commutes** relation © is defined as

 $x \odot y \iff x = (x \land y) \lor (x \land \neg y) \quad \forall x, y \in X,$ in which case we say, "*x* **commutes** with *y* in *L*". That is, \odot is a relation in 2^{XX} such that $\odot \triangleq \{(x, y) \in X^2 | x = (x \land y) \lor (x \land \neg y)\}$

Proposition 3.10 ⁸⁶ Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be an ORTHOCOMPLEMENTED LATTICE.

x©0	and	0©x	$\forall x \in X$	x©y	\Leftrightarrow	$x @ y^{\perp}$	$\forall x, y \in X$
x©1	and	1@x	$\forall x \in X$	$x \leq y$	\Rightarrow	x©y	$\forall x, y \in X$
x©x			$\forall x \in X$	$x \perp y$	\Rightarrow	x @ y	$\forall x, y \in X$

Definition 3.11 Let © be the *commutes* relation (Definition 3.9 page 36) on a *lattice with negation* $L \triangleq (X, \lor, \land, \neg, 0, 1; \le)$ (Definition 2.16 page 30). *L* is **symmetric** if $x © y \implies y © x \quad \forall x, y \in X$

In general, the commutes relation is not *symmetric*. But Proposition 3.12 (next) describes some conditions under which it *is* symmetric.

Proposition 3.12 ⁸⁷ Let $(X, \lor, \land, 0, 1; \le)$ be an ORTHOCOMPLEMENTED LATTICE (Definition 1.72 page 20).

$\{x @ y \implies y @ x\}$	\Leftrightarrow	$\left\{x \le y \implies y = x \lor \left(x^{\perp} \land y\right)\right\}$	(ORTHOMODULAR IDENTITY)	(2)
© is symmetric at (x, y) (1)	\Leftrightarrow	$\left\{x \le y \implies x = y \land \left(x \lor y^{\perp}\right)\right\}$	$(x = \phi_y(x) \text{ (Sasaki projection) })$	(3)
(x, y) (1)		$\left\{ y = (x \land y) \lor \left[y \land (x \land y)^{\perp} \right] \right\}$		(4)
	\Leftrightarrow	$\left\{x = (x \lor y) \land \left[x \lor (x \lor y)^{\perp}\right]\right\}$		(5)

Theorem 3.13 ⁸⁸ Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be an ORTHOCOMPLEMENTED LATTICE (Definition 1.72 page 20).

 $\{ x \widehat{\otimes} c \mid \forall x \in X \} \iff \{ L \text{ is ISOMORPHIC to } [0, c] \times [0, c^{\perp}] \}$ with isomorphism $\theta(x) \triangleq ([0, c], [0, c^{\perp}]).$

Proposition 3.14 ⁸⁹ Let $(X, \lor, \land, 0, 1; \leq)$ be an ORTHOMODULAR lattice.

⁸⁵ **●** [98], page 20, **●** [88], page 79, (A. Commutativity), **●** [115], page 227, (Hilfssatz (Lemma) XII.1.2), **■** [152], page 301, (Def.5.2, cf Foulis 1962), **■** [15], page 833, (" $a = (a \cap x) \cup (a \cap x')$ ") ⁸⁶ **●** [87], page 67, **●** [78], (Proposition 13.2) ⁸⁷ **●** [87], page 68, **■** [127], page 158, **●** [78], (Proposition 13.3) ⁸⁸ **●** [98], page 20, **■** [114] ⁸⁹ **■** [62], page 66, **■** [152], (cf Foulis 1962)

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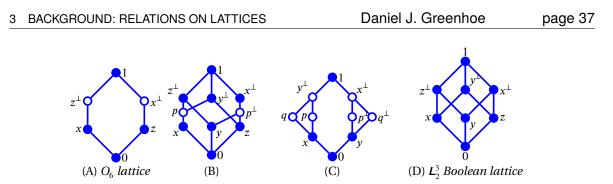


Figure 7: Lattices with centers marked with solid dots (see Example 3.17 page 37)

 $x \odot y \quad \iff \quad \phi_x(y) = \phi_y(x) = x \land y \quad \forall x, y \in X$

3.3 Center

An element in an *orthocomplemented lattice* (Definition 1.72 page 20) is in the *center* of the lattice if that element *commutes* (Definition 3.9 page 36) with every other element in the lattice (next definition). *All* the elements of an *orthocomplemented lattice* are in the *center* if and only if that lattice is *Boolean* (Proposition 1.81 page 24).

Definition 3.15 ⁹⁰ Let © be the *commutes* relation (Definition 3.9 page 36) on a *lattice with negation* $L \triangleq (X, \lor, \land, \neg, 0, 1; \le)$ (Definition 2.16 page 30). The **center** of L is defined as $\{x \in X | x © y \quad \forall y \in X\}$

Proposition 3.16 Let $L \triangleq (X, \lor, \land, 0, 1; \le)$ be an ORTHOCOMPLEMENTED LATTICE (Definition 1.72 page 20). The elements 0 and 1 are in the center of L.

PROOF: This follows directly from Definition 3.9 (page 36) and Proposition 3.10 (page 36).

Example 3.17 The **centers** of the lattices in Figure 7 (page 37) are illustrated with solid dots. Note that in the case of the Boolean lattice in (D), every dot is in the center (Proposition 1.81 page 24).

3.4 D-Posets

Definition 3.18 ⁹¹ Let 1 be the *upper bound* of an *ordered set* (X, \leq) . An operation \ is a **difference** on (X, \leq) if

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⁹⁰ 🖲 [<mark>88</mark>], page 80

⁹¹ \blacksquare [104], page 22,24, (Definitions 1,2)

page 38

1.	$x \leq y$	\Rightarrow	$y \setminus x \le y$	$\forall x, y \in X$	and		
2.	$x \leq y$	\Rightarrow	$y \setminus (y \setminus x) = x$	$\forall x, y \in X$	and		
3.	$x \le y \le z$	\Rightarrow	$z \setminus y \le z \setminus x$	$\forall x, y, z \in X$	and		
4.	$x \le y \le z$	\Rightarrow	$(z \setminus x) \setminus (z \setminus y) = y \setminus x$	$\forall x, y, z \in X$			
The structure $(X, \leq, \backslash, 1)$ is called a D-poset .							

Proposition 3.19 ⁹² Let X be a SET.

$$\left\{ \begin{array}{l} (X,\leq,\backslash,1) \text{ is a} \\ \text{D-POSET} \\ (\text{Definition 3.18 page 37)} \end{array} \right\} \implies \left\{ \begin{array}{l} 1. \quad x \leq y \leq z \quad \Longrightarrow \quad y \setminus x \leq z \setminus x \qquad \forall x,y,z \in X \quad and \\ 2. \quad x \leq y \leq z \quad \Longrightarrow \quad x \leq z \setminus (y \setminus x) \qquad \forall x,y,z \in X \quad and \\ 3. \quad x \leq y \leq z \quad \Longrightarrow \quad (z \setminus x) \setminus (y \setminus x) = z \setminus y \quad \forall x,y,z \in X \quad and \\ 4. \quad x \leq y \leq z \quad \Longrightarrow \quad \left[z \setminus (y \setminus x) \right] \setminus x = z \setminus y \quad \forall x,y,z \in X \quad . \end{array} \right\}$$

Example 3.20 ⁹³ The structure $(\mathbb{R}^+, -, \leq)$ is a *D*-poset where \mathbb{R}^+ is the set of positive real numbers, - is the standard subtraction operation on \mathbb{R} , and \leq is the standard ordering relation on \mathbb{R}^+ .

Example 3.21 ⁹⁴ The structure $(2^X, \backslash, \subseteq)$ is a *D*-poset where 2^X is the power set of a set *X*, \backslash is the set difference operator, and \subseteq is the set inclusion relation.

4 Background: MRA-wavelet analysis

4.1 Transversal Operators

Definition 4.1 95 1. **T** is the **translation operator** on $\mathbb{C}^{\mathbb{C}}$ defined as $\mathbf{T}_{\mathbf{r}}\mathbf{f}(x) \triangleq \mathbf{f}(x-\tau) \text{ and } \mathbf{T} \triangleq \mathbf{T}_{1}$ $\forall f \in \mathbb{C}^{\mathbb{C}}$ 2. **D** is the **dilation operator** on $\mathbb{C}^{\mathbb{C}}$ defined as and $\mathbf{D} \triangleq \sqrt{2}\mathbf{D}_2 \quad \forall f \in \mathbb{C}^{\mathbb{C}}$ $\mathbf{D}_{\alpha}\mathbf{f}(x) \triangleq \mathbf{f}(\alpha x)$ $\mathbf{D}\mathbf{f}(x)$ $\mathbf{T}^{-1}\mathbf{f}(x)$ f(x) = Tf(x)f(x)-2-10 2 -2-1 0 ⁹² [104], page 23, (PROPOSITION 1.) ⁹³ [104], page 22, (Example 1) ⁹⁴ [104], page 24, (Example 4) 95 🖻 [163], pages 79–80, (Definition 3.39), 🖻 [27], pages 41–42, 🖱 [168], page 18, (Definitions 2.3,2.4⟩, ► [100], page A-21, ► [8], page 473, ► [135], page 260, ► [11], page , ► [83], page 250, ⟨Notation 9.4⟩, <a>[25], page 74, [69], page 639, <a>[34], page 81, <a>[33], page 2, <a>[75], page 2

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Proposition 4.2 ⁹⁶ Let **T** be the TRANSLATION OPERATOR (Definition 4.1 page 38). $\sum_{n \in \mathbb{Z}} \mathbf{T}^n \mathbf{f}(x) = \sum_{n \in \mathbb{Z}} \mathbf{T}^n \mathbf{f}(x+1) \qquad \forall \mathbf{f} \in \mathbb{R}^{\mathbb{R}} \qquad \left(\sum_{n \in \mathbb{Z}} \mathbf{T}^n \mathbf{f}(x) \text{ is PERIODIC with period } 1\right)$

Proposition 4.3 ⁹⁷ Let **T** and **D** be as defined in Definition 4.1 page 38. **T** has an inverse \mathbf{T}^{-1} in $\mathbb{C}^{\mathbb{C}}$ expressed by the relation $\mathbf{T}^{-1}\mathbf{f}(x) = \mathbf{f}(x+1) \quad \forall \mathbf{f} \in \mathbb{C}^{\mathbb{C}}$ (translation operator inverse). **D** has an inverse \mathbf{D}^{-1} in $\mathbb{C}^{\mathbb{C}}$ expressed by the relation $\mathbf{D}^{-1}\mathbf{f}(x) = \frac{\sqrt{2}}{2}\mathbf{f}(\frac{1}{2}x) \quad \forall \mathbf{f} \in \mathbb{C}^{\mathbb{C}}$ (dilation operator inverse).

Proposition 4.4 ⁹⁸ Let **T** and **D** be as defined in Definition 4.1 page 38. Let $\mathbf{D}^0 = \mathbf{T}^0 \triangleq \mathbf{I}$ be

the identity operator.

 $\mathbf{D}^{j}\mathbf{T}^{n}\mathbf{f}(x) = 2^{j/2}\mathbf{f}\left(2^{j}x - n\right) \qquad \forall j, n \in \mathbb{Z}, \mathbf{f} \in \mathbb{C}^{\mathbb{C}}$

Example 4.5 (linear functions) ⁹⁹ Let **T** be the *translation operator* (Definition 4.1 page 38). Let $\mathcal{L}(\mathbb{C}, \mathbb{C})$ be the set of all *linear* functions in $L^2_{\mathbb{D}}$.

1. {x, Tx} is a *basis* for $\mathcal{L}(\mathbb{C}, \mathbb{C})$ and 2. f(x) = f(1)x - f(0)Tx $\forall f \in \mathcal{L}(\mathbb{C}, \mathbb{C})$

SPROOF: By left hypothesis, f is *linear*; so let $f(x) \triangleq ax + b$

 $f(1)x - f(0)\mathbf{T}x = f(1)x - f(0)(x - 1)$ by Definition 4.1 page 38 $= (ax + b)|_{x=1}x - (ax + b)|_{x=0}(x - 1)$ by left hypothesis and definition of f = (a + b)x - b(x - 1) = ax + bx - bx + b= f(x) by left hypothesis and definition of f

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Example 4.6 (Cardinal Series) Let **T** be the *translation operator* (Definition 4.1 page 38). The *Paley-Wiener* class of functions PW_{σ}^2 are those functions which are "*bandlimited*" with respect to their Fourier transform. The cardinal series forms an orthogonal basis for such a space. The *Fourier coefficients* for a projection of a function f onto the Cardinal series basis elements is particularly simple—these coefficients are samples of f(x) taken at regular intervals. In fact, one could represent the coefficients using inner product notation with

⁹⁶ 🖱 [75], page 3

⁹⁷ 🖱 [75], page 3

⁹⁸ 🖱 [75], page 4

^{🤊 🖱 [86],} page 2

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the *Dirac delta distribution* δ as follows:

$$\langle \mathbf{f}(x) \mid \mathbf{T}^{n} \delta(x) \rangle \triangleq \int_{\mathbb{R}} \mathbf{f}(x) \delta(x-n) \, \mathrm{d}x \triangleq \mathbf{f}(n)$$

$$1. \quad \left\{ \mathbf{T}^{n} \frac{\sin(\pi x)}{\pi x} \middle|_{n \in \mathbb{N}} \right\} \text{ is a basis for } \mathbf{PW}_{\sigma}^{2} \quad \text{and}$$

$$2. \quad \mathbf{f}(x) = \underbrace{\sum_{n=1}^{\infty} \mathbf{f}(n) \mathbf{T}^{n} \frac{\sin(\pi x)}{\pi x}}_{Cardinal \ series} \quad \forall \mathbf{f} \in \mathbf{PW}_{\sigma}^{2}, \sigma \leq \frac{1}{2}$$

Example 4.7 (Fourier Series)

1.
$$\{\mathbf{D}_{n}e^{ix}|_{n\in\mathbb{Z}}\}$$
 is a *basis* for $\mathbf{L}(0, 2\pi)$ and
2. $f(x) = \frac{1}{\sqrt{2\pi}}\sum_{n\in\mathbb{Z}}\alpha_{n}\mathbf{D}_{n}e^{ix}$ $\forall x\in(0,2\pi), f\in\mathbf{L}(0,2\pi)$ where
3. $\alpha_{n} \triangleq \frac{1}{\sqrt{2\pi}}\int_{0}^{2\pi}f(x)\mathbf{D}_{n}e^{-ix} dx$ $\forall f\in\mathbf{L}(0,2\pi)$

Example 4.8 (Fourier Transform)

1.
$$\{\mathbf{D}_{\omega}e^{ix} | \omega \in \mathbb{R}\}\$$
 is a basis for $\mathcal{L}^{2}_{\mathbb{R}}$ and
2. $f(x) = \frac{1}{\sqrt{2\pi}} \int_{\mathbb{R}} \tilde{f}(\omega)\mathbf{D}_{x}e^{i\omega} d\omega \quad \forall f \in \mathcal{L}^{2}_{\mathbb{R}}$ where
3. $\tilde{f}(\omega) \triangleq \frac{1}{\sqrt{2\pi}} \int_{\mathbb{R}} f(x)\mathbf{D}_{\omega}e^{-ix} dx \quad \forall f \in \mathcal{L}^{2}_{\mathbb{R}}$

Example 4.9 (Gabor Transform) ¹⁰⁰

1.
$$\left\{ \left(\mathbf{T}_{\tau} e^{-\pi x^{2}} \right) \left(\mathbf{D}_{\omega} e^{ix} \right) \middle|_{\tau,\omega \in \mathbb{R}} \right\}$$
 is a *basis* for $\mathcal{L}_{\mathbb{R}}^{2}$ and
2. $f(x) = \int_{\mathbb{R}} G(\tau, \omega) \mathbf{D}_{x} e^{i\omega} d\omega \quad \forall x \in \mathbb{R}, f \in \mathcal{L}_{\mathbb{R}}^{2}$ where
3. $G(\tau, \omega) \triangleq \int_{\mathbb{R}} f(x) \left(\mathbf{T}_{\tau} e^{-\pi x^{2}} \right) \left(\mathbf{D}_{\omega} e^{-ix} \right) dx \quad \forall x \in \mathbb{R}, f \in \mathcal{L}_{\mathbb{R}}^{2}$

Example 4.10 (wavelets) Let $\psi(x)$ be a *mother wavelet*.

1.
$$\left\{ \mathbf{D}^{k}\mathbf{T}^{n}\psi(x)\big|_{k,n\in\mathbb{Z}} \right\}$$
 is a basis for $\mathcal{L}_{\mathbb{R}}^{2}$ and
2. $f(x) = \sum_{k\in\mathbb{Z}}\sum_{n\in\mathbb{Z}}\alpha_{k,n}\mathbf{D}^{k}\mathbf{T}^{n}\psi(x) \quad \forall f\in\mathcal{L}_{\mathbb{R}}^{2}$ where
3. $\alpha_{n} \triangleq \int_{\mathbb{R}}f(x)\mathbf{D}^{k}\mathbf{T}^{n}\psi^{*}(x) \, dx \quad \forall f\in\mathcal{L}_{\mathbb{R}}^{2}$

¹⁰⁰ 🖻 [143], (Chapter 3)

[61], page 32, (Definition 1.69)

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The Structure of Wavelets 4.2

In Fourier analysis, continuous dilations (Definition 4.1 page 38) of the complex exponential form a basis for the space of square integrable functions $L^2_{\mathbb{R}}$ such that

 $\boldsymbol{L}_{\mathbb{R}}^{2} = \operatorname{span} \left\{ \mathbf{D}_{\omega} e^{ix} \mid \omega \in \mathbb{R} \right\}.$

In Fourier series analysis, *discrete* dilations of the complex exponential form a basis for $\begin{aligned} \boldsymbol{L}^2_{\mathbb{R}}(0, \, 2\pi) & \text{such that} \\ \boldsymbol{L}^2_{\mathbb{R}}(0, \, 2\pi) = \operatorname{span} \left\{ \left. \mathbf{D}_j e^{ix} \right|_{j \in \mathbb{Z}} \right\}. \end{aligned}$

In Wavelet analysis, for some *mother wavelet* (Definition 4.18 page 47) $\psi(x)$, $\boldsymbol{L}_{\mathbb{D}}^{2} = \operatorname{span} \left\{ \mathbf{D}_{\omega} \mathbf{T}_{\tau} \boldsymbol{\psi}(x) \, | \, \boldsymbol{\omega}, \tau \in \mathbb{R} \right\}.$

However, the ranges of parameters ω and τ can be much reduced to the countable set \mathbb{Z} resulting in a *dyadic* wavelet basis such that for some mother wavelet $\psi(x)$,

 $\boldsymbol{L}_{\mathbb{R}}^{2} = \operatorname{span}\left\{ \mathbf{D}^{j}\mathbf{T}^{n}\boldsymbol{\psi}(x) \left| j, n \in \mathbb{Z} \right. \right\}.$

Wavelets that are both dyadic and compactly supported have the attractive feature that they can be easily implemented in hardware or software by use of the Fast Wavelet Transform (Figure 10 page 49).

In 1989, Stéphane G. Mallat introduced the *Multiresolution Analysis* (MRA, Definition 4.12 page 43) method for wavelet construction. The MRA has since become the dominate wavelet construction method. Moreover, P.G. Lemarié has proved that all wavelets with *compact sup*port are generated by an MRA.¹⁰¹

The MRA is an **analysis** of the linear space $L^2_{\mathbb{R}}$. An analysis of a linear space X is any sequence $(V_j)_{j \in \mathbb{Z}}$ of linear subspaces of X. The partial or complete reconstruction of Xfrom $(V_j)_{j \in \mathbb{Z}}$ is a **synthesis**.¹⁰² Some analyses are completely *characterized* by a *trans*form. For example, a Fourier analysis is a sequence of subspaces with sinusoidal bases. Examples of subspaces in a Fourier analysis include $V_1 = \text{span}\{e^{ix}\}, V_{2,3} = \text{span}\{e^{i2.3x}\}, V_{\sqrt{2}} =$ spen $\{e^{i\sqrt{2}x}\}$, etc. A transform is loosely defined as a function that maps a family of functions into an analysis. A very useful transform (a "Fourier transform") for Fourier Analysis is

$$\left[\tilde{\mathbf{F}}\mathbf{f}\right](\omega) \triangleq \frac{1}{\sqrt{2\pi}} \int_{\mathbb{R}} \mathbf{f}(x) e^{-i\omega x} \, \mathrm{d}x$$

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¹⁰¹ [[109], 🖱 [119], page 240

¹⁰²The word analysis comes from the Greek word ἀνάλυσις, meaning "dissolution" (🖻 [140], page 23, (entry 359)), which in turn means "the resolution or separation into component parts" [21], http://dictionary.reference.com/browse/dissolution)

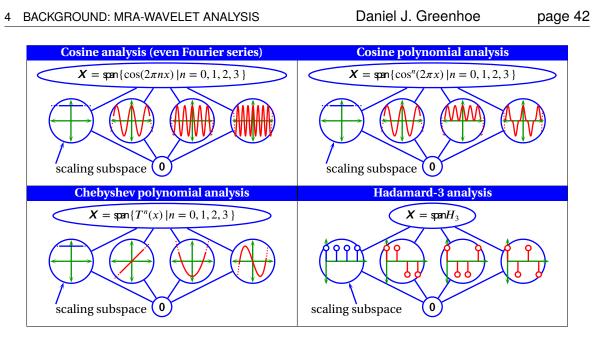
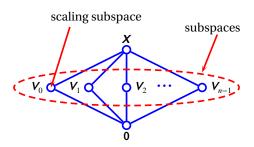


Figure 8: Examples of order structures for selected analyses (Example 4.11 page 42)

An analysis can be partially characterized by its order structure with respect to an order relation such as the set inclusion relation \subseteq . Most transforms have a very simple M-*n* order structure, as illustrated to the right.¹⁰³The M-*n* lattices for $n \ge 3$ are *modular* (Lemma 1.56 page 16) but not *distributive* (Theorem 1.57 page 16). Analyses typically have one subspace that is a *scaling* subspace; and this subspace is often simply a family of constants (as is the case with Fourier Analysis).



An analysis can be represented using three different structures:

- ① sequence of subspaces
- ② sequence of basis vectors
- ③ sequence of basis coefficients

These structures are isomorphic to each other, and can therefore be used interchangeably.

Example 4.11 ¹⁰⁴ Some examples of the order structures of some analyses are illustrated in Figure 8 (page 42).

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<sup>103</sup> [75], page 29, ⟨§2.2⟩
<sup>104</sup> [75], pages 30–31
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4.3 Multiresolution analysis

A multiresolution analysis provides "coarse" approximations of a function in a linear space $L^2_{\mathbb{R}}$ at multiple "scales" or "resolutions". Key to this process is a sequence of *scaling functions*. Most traditional transforms feature a single *scaling function* $\phi(x)$ set equal to one $(\phi(x) = 1)$. This allows for convenient representation of the most basic functions, such as constants.¹⁰⁵ A multiresolution system, on the other hand, uses a generalized form of the scaling concept:¹⁰⁶

- (1) Instead of the scaling function simply being set *equal to unity* ($\phi(x) = 1$), a multiresolution analysis (Definition 4.12 page 43) is often constructed in such a way that the scaling function $\phi(x)$ forms a *partition of unity* such that $\sum_{n \in \mathbb{Z}} \mathbf{T}^n \phi(x) = 1$.
- (2) Instead of there being *just one* scaling function, there is an entire sequence of scaling functions $(\mathbf{D}^{j}\phi(x))_{i\in\mathbb{Z}}$, each corresponding to a different "*resolution*".

Definition 4.12 ¹⁰⁷ Let $(V_j)_{j \in \mathbb{Z}}$ be a sequence of subspaces on $L^2_{\mathbb{R}}$. Let A^- be the *closure* of a set *A*. The sequence $(V_j)_{i \in \mathbb{Z}}$ is a **multiresolution analysis** on $L^2_{\mathbb{R}}$ if

1.
$$V_{j} = V_{j}^{-}$$
 $\forall j \in \mathbb{Z}$ (closed) and
2. $V_{j} \subset V_{j+1}$ $\forall j \in \mathbb{Z}$ (linearly ordered) and
3. $\left(\bigcup_{j \in \mathbb{Z}} V_{j}\right)^{-} = L_{\mathbb{R}}^{2}$ (dense in $L_{\mathbb{R}}^{2}$) and
4. $f \in V_{j} \iff Df \in V_{j+1}$ $\forall j \in \mathbb{Z}, f \in L_{\mathbb{R}}^{2}$ (self-similar) and
5. $\exists \phi$ such that $\left\{\mathbf{T}^{n} \phi | n \in \mathbb{Z}\right\}$ is a Riesz basis for V_{0} .

A multiresolution analysis is also called an MRA. An element V_j of $(V_j)_{j \in \mathbb{Z}}$ is a scaling subspace of the space $L^2_{\mathbb{R}}$. The pair $(L^2_{\mathbb{R}}, (V_j))$ is a multiresolution analysis space, or MRA space. The function ϕ is the scaling function of the *MRA space*.

The traditional definition of the MRA also includes the following:

6. $\mathbf{f} \in \mathbf{V}_j \iff \mathbf{T}^n \mathbf{f} \in \mathbf{V}_j \quad \forall n, j \in \mathbb{Z}, \mathbf{f} \in L^2_{\mathbb{R}} \quad (translation invariant)$ 7. $\bigcap_{i \in \mathbb{Z}} \mathbf{V}_j = \{\mathbb{O}\}$ (greatest lower bound is **0**)

However, these follow from the *MRA* as defined in Definition 4.12 (Proposition 4.13 page 44, Proposition 4.14 page 44).

¹⁰⁵ [95], page 8

¹⁰⁶ The concept of a scaling space was perhaps first introduced by Taizo Iijima in 1959 in Japan, and later as the *Gaussian Pyramid* by Burt and Adelson in the 1980s in the West. Solution [118], page 70, Solution [24], Solution [4], Solution [6], Solution [80], Solution [166], ⟨historical survey⟩

¹⁰⁷ S [85], page 44, S [119], page 221, (Definition 7.1), S [118], page 70, S [122], page 21, (Definition 2.2.1), S [27], page 284, (Definition 13.1.1), S [8], pages 451–452, (Definition 7.7.6), [163], pages 300–301, (Definition 10.16), S [35], pages 129–140, (Riesz basis: page 139)

Proposition 4.13 ¹⁰⁸ Let MRA be defined as in Definition 4.12 page 43. $\left\{ \left(\left(V_{j} \right) \right)_{j \in \mathbb{Z}} \text{ is an MRA} \right\} \implies \underbrace{\left\{ f \in V_{j} \iff \mathbf{T}^{n} \mathbf{f} \in V_{j} \quad \forall n, j \in \mathbb{Z}, \mathbf{f} \in \mathcal{L}_{\mathbb{R}}^{2} \right\}}_{\text{TRANSLATION INVARIANT}}$

Proposition 4.14 ¹⁰⁹ Let MRA be defined as in Definition 4.12 page 43.

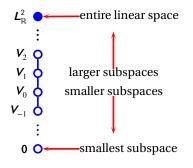
$$\left\{\left(\left(V_{j}\right)\right)_{j\in\mathbb{Z}}\ is\ an\ \mathrm{MRA}\right\}\qquad\Longrightarrow\qquad\left\{\bigcap_{j\in\mathbb{Z}}V_{j}=\left\{0\right\}\qquad(\text{greatest lower bound}\ is\ \mathbf{0})\right\}$$

The MRA (Definition 4.12 page 43) is more than just an interesting mathematical toy. Under some very "reasonable" conditions (next proposition), as $j \to \infty$, the *scaling subspace* V_j is *dense* in $L^2_{\mathbb{R}}$...meaning that with the MRA we can represent any "reasonable" function to within an arbitrary accuracy.

Proposition 4.15¹¹⁰

$$\left\{ \begin{array}{ll} (1). & (\mathbf{T}^{n}\phi) \text{ is a Riesz sequence} & and \\ (2). & \tilde{\phi}(\omega) \text{ is continuous at } 0 & and \\ (3). & \tilde{\phi}(0) \neq 0 \end{array} \right\} \implies \left\{ \left(\bigcup_{j \in \mathbb{Z}} \mathbf{V}_{j} \right)^{-} = \mathbf{L}_{\mathbb{R}}^{2} \quad (\text{dense in } \mathbf{L}_{\mathbb{R}}^{2}) \right\}$$

A multiresolution analysis (Definition 4.12 page 43) together with the set inclusion relation \subseteq form the *linearly ordered* set (Definition 1.4 page 4) (((V_j)), \subseteq), illustrated to the right by a *Hasse diagram* (Definition 1.6 page 4). Subspaces V_j increase in "size" with increasing *j*. That is, they contain more and more vectors (functions) for larger and larger *j*—with the upper limit of this sequence being $L^2_{\mathbb{R}}$. Alternatively, we can say that approximation within a subspace V_j yields greater "*resolution*" for increasing *j*.¹¹¹



Remark 4.16 ¹¹²Note that the *greatest lower bound* (g.*l.b.*) of the linearly ordered set $((V_j), \subseteq)$ is **0** (Proposition 4.14 page 44): All linear subspaces contain the zero vector. So the intersection of any two subspaces must at least contain 0. If the intersection of any two linear subspaces X and Y is exactly $\{0\}$, then for any vector in the sum of those subspaces $(u \in X + Y)$ there are **unique** vectors $f \in X$ and $g \in Y$ such that u = f + g. This is *not* necessarily true if the intersection contains more than just $\{0\}$.

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¹⁰⁸ **(**85], page 45, (Theorem 1.6), **(**75], pages 32–33, (Proposition 2.1)

¹⁰⁹ \square [168], pages 19–28, (Proposition 2.14), \square [85], page 45, (Theorem 1.6), \square [141], pages 313–314, (Lemma 6.4.28), \square [75], pages 33–35, (Proposition 2.2)

¹¹⁰ Temposition 2.15 , Temposition 2.15 , Temposition 2.3 , Proposition 2.3 ,

¹¹¹ [123], page 83, ⟨Theorem 3.2.12⟩, ■ [106], page 67, ⟨Theorem 2.14⟩, ■ [76], ⟨Theorem 7.1⟩

¹¹² [75], page 38, (§2.3.2 Order structure)

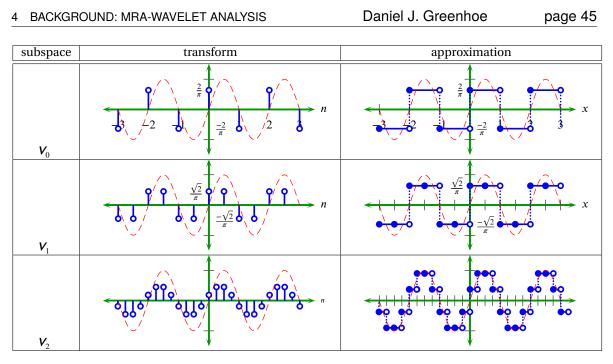
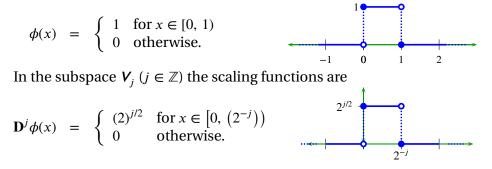


Figure 9: Example approximations of $sin(\pi x)$ in 3 Haar scaling subspaces (see Example 4.17 page 45)

Example 4.17

In the *Haar* MRA, the scaling function $\phi(x)$ is the *pulse function*

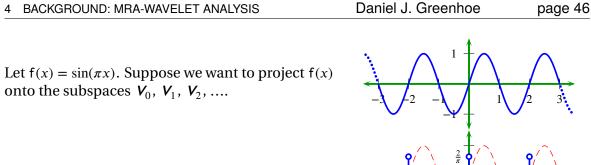


The scaling subspace V_0 is the span $V_0 \triangleq \operatorname{span} \{ \mathbf{T}^n \phi | n \in \mathbb{Z} \}$. The scaling subspace V_j is the span $V_j \triangleq \operatorname{span} \{ \mathbf{D}^j \mathbf{T}^n \phi | n \in \mathbb{Z} \}$. Note that $\| \mathbf{D}^j \mathbf{T}^n \phi \|$ for each resolution j and shift n is unity:

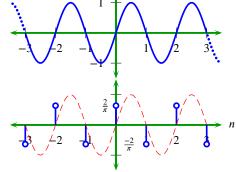
$$\|\mathbf{D}^{j}\mathbf{T}^{n}\phi\|^{2} = \|\phi\|^{2}$$

= $\int_{0}^{1} |1|^{2} dx$ by definition of $\|\cdot\|$ on $\mathcal{L}_{\mathbb{R}}^{2}$
= 1

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The values of the transform coefficients for the subspace V_i are given by



$$\begin{split} \left[\mathbf{R}_{j} \mathbf{f}(x) \right](n) &= \frac{1}{\|\mathbf{D}^{j} \mathbf{T}^{n} \phi\|^{2}} \left\langle \mathbf{f}(x) \mid \mathbf{D}^{j} \mathbf{T}^{n} \phi \right\rangle \\ &= \frac{1}{\|\phi\|^{2}} \left\langle \mathbf{f}(x) \mid 2^{j/2} \phi(2^{j} x - n) \right\rangle \\ &= 2^{j/2} \left\langle \mathbf{f}(x) \mid \phi(2^{j} x - n) \right\rangle \\ &= 2^{j/2} \int_{2^{-j}(n+1)}^{2^{-j}(n+1)} \mathbf{f}(x) \, \mathrm{d}x \\ &= 2^{j/2} \int_{2^{-j}(n+1)}^{2^{-j}(n+1)} \sin(\pi x) \, \mathrm{d}x \\ &= 2^{j/2} \left(-\frac{1}{\pi} \right) \cos(\pi x) \Big|_{2^{-j}(n+1)}^{2^{-j}(n+1)} \\ &= \frac{2^{j/2}}{\pi} \left[\cos\left(2^{-j} n\pi\right) - \cos\left(2^{-j}(n+1)\pi\right) \right] \end{split}$$

by Proposition 4.4 page 39

And the projection $\mathbf{A}_n f(x)$ of the function f(x) onto the subspace V_j is

$$\begin{split} \mathbf{A}_{j}\mathbf{f}(x) &= \sum_{n \in \mathbb{Z}} \left\langle \mathbf{f}(x) \mid \mathbf{D}^{j}\mathbf{T}^{n}\phi \right\rangle \mathbf{D}^{j}\mathbf{T}^{n}\phi \\ &= \frac{2^{j/2}}{\pi} \sum_{n \in \mathbb{Z}} \left[\cos\left(2^{-j}n\pi\right) - \cos\left(2^{-j}(n+1)\pi\right) \right] 2^{j/2}\phi(2^{j}x-n) \\ &= \frac{2^{j}}{\pi} \sum_{n \in \mathbb{Z}} \left[\cos\left(2^{-j}n\pi\right) - \cos\left(2^{-j}(n+1)\pi\right) \right] \phi(2^{j}x-n) \end{split}$$

The transforms into the subspaces V_0 , V_1 , and V_2 , as well as the approximations in those subspaces are as illustrated in Figure 9 (page 45).

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4.4 Wavelet analysis

The term "wavelet" comes from the French word "ondelette", meaning "small wave". And in essence, wavelets are "small waves" (as opposed to the "long waves" of Fourier analysis) that form a basis for the Hilbert space $L^2_{\mathbb{D}}$.¹¹³

Definition 4.18 ¹¹⁴ Let **T** and **D** be as defined in Definition 4.1 page 38. A function $\psi(x)$ in $L^2_{\mathbb{R}}$ is a wavelet function for $L^2_{\mathbb{R}}$ if

 $\left\{\mathbf{D}^{j}\mathbf{T}^{n}\psi|_{j,n\in\mathbb{Z}}\right\}$ is a *Riesz basis* for $\boldsymbol{L}_{\mathbb{R}}^{2}$.

In this case, ψ is also called the **mother wavelet** of the basis $\{\mathbf{D}^{j}\mathbf{T}^{n}\psi|_{j,n\in\mathbb{Z}}\}$. The sequence of subspaces $(W_i)_{i \in \mathbb{Z}}$ is the wavelet analysis induced by ψ , where each subspace W_i is defined as

 $\boldsymbol{W}_{i} \triangleq \operatorname{span} \left\{ \mathbf{D}^{j} \mathbf{T}^{n} \boldsymbol{\psi} \middle|_{n \in \mathbb{Z}} \right\} \,.$

A wavelet analysis (W_j) is often constructed from a multiresolution analysis (Definition 4.12 page 43) (V_i) under the relationship

where $\hat{+}$ is subspace addition (*Minkowski addition*). $\boldsymbol{V}_{i+1} = \boldsymbol{V}_i + \boldsymbol{W}_i,$ By this relationship alone, (W_i) is in no way uniquely defined in terms of a multiresolution analysis (V_i) . In general there are many possible complements of a subspace V_i . To uniquely define such a wavelet subspace, one or more additional constraints are required.

One of the most common additional constraints is *orthogonality*, such that V_i and W_j are orthogonal to each other.

Definition 4.19 Let $(L^2_{\mathbb{R}}, (V_j), \phi, (h_n))$ be a multiresolution system (Definition 4.12 page 43) and $(W_j)_{j\in\mathbb{Z}}$ a wavelet analysis (Definition 4.18 page 47) with respect to $(V_j)_{j\in\mathbb{Z}}$. Let $(g_n)_{n\in\mathbb{Z}}$ be a sequence of coefficients such that $\psi = \sum_{n \in \mathbb{Z}} g_n \mathbf{D} \mathbf{T}^n \phi$.

A wavelet system is the tuple

 $(L^2_{\mathbb{R}}, (V_j), (W_j), \phi, \psi, (h_n), (g_n))$ and the sequence $(g_n)_{n \in \mathbb{Z}}$ is the wavelet coefficient sequence.

Theorem 4.20 ¹¹⁵ Let $(L^2_{\mathbb{R}}, (V_i), (W_i), \phi, \psi, (h_n), (g_n))$ be a WAVELET SYSTEM (Definition 4.19) page 47). Let $V_1 + V_2$ represent MINKOWSKI ADDITION of two subspaces V_1 and V_2 of a Hilbert space H.

$$\begin{aligned} L_{\mathbb{R}}^{2} &= \lim_{j \to \infty} V_{j} \\ &= V_{j} \stackrel{?}{+} W_{j} \stackrel{?}{+} W_{j+1} \stackrel{?}{+} W_{j+2} \stackrel{?}{+} \cdots \\ &= \cdots \stackrel{?}{+} W_{-2} \stackrel{?}{+} W_{-1} \stackrel{?}{+} W_{0} \stackrel{?}{+} W_{1} \stackrel{?}{+} W_{2} \stackrel{?}{+} \cdots \\ \end{aligned}$$

$$\begin{aligned} (L_{\mathbb{R}}^{2} \text{ is equivalent to one very large scaling space} \\ and a sequence of wavelet subspaces}) \mathbb{I} \\ &= \cdots \stackrel{?}{+} W_{-2} \stackrel{?}{+} W_{-1} \stackrel{?}{+} W_{0} \stackrel{?}{+} W_{1} \stackrel{?}{+} W_{2} \stackrel{?}{+} \cdots \\ \end{aligned}$$

¹¹³ 🖱 [158], page ix, 🖱 [7], page 191

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¹¹⁴ [168], page 17, (Definition 2.1), [75], page 50, (Definition 2.4)

¹¹⁵ **[75]**, page 53, (Theorem 2.8)

[∞]Proof:

(1) Proof for (1):

$$\boldsymbol{L}_{\mathbb{R}}^2 = \lim_{j \to \infty} \boldsymbol{V}_j \qquad \qquad \text{by Definition 4.12 page 43}$$

(2) Proof for (2):

$$\underbrace{\underbrace{V_{j} + W_{j}}_{V_{j+1}} + W_{j+1} + W_{j+2} + \cdots}_{V_{j+2}} = \underbrace{\underbrace{V_{j+1} + W_{j+1}}_{V_{j+2}} + W_{j+2} + W_{j+3} + \cdots}_{V_{j+3}}$$

$$= \underbrace{\underbrace{V_{j+2} + W_{j+2}}_{V_{j+3}} + W_{j+3} + W_{j+4} + \cdots}_{V_{j+4}}$$

$$= \underbrace{\underbrace{V_{j+3} + W_{j+3}}_{V_{j+4}} + W_{j+5} + \cdots}_{V_{j+5}}$$

$$= \underbrace{\underbrace{V_{j+5} + W_{j+5}}_{V_{j+5}} + W_{j+6} + W_{j+6} + \cdots}_{J_{j+6}}$$

$$= \underbrace{\lim_{j \to \infty} V_{j+5} + W_{j+5} + W_{j+6} + W_{j+6} + \cdots}_{J_{R}}$$

(3) Proof for (3):

$$L_{\mathbb{R}}^{2} = \underbrace{V_{0}}_{V_{-1} + W_{-1}} + W_{0} + W_{1} + W_{2} + W_{3} + \cdots \qquad by (2)$$

$$= \underbrace{V_{-1}}_{V_{-2} + W_{-1}} W_{-1} + W_{0} + W_{1} + W_{2} + W_{3} + \cdots$$

$$= \underbrace{V_{-2}}_{V_{-3} + W_{-2}} W_{-2} + W_{-1} + W_{0} + W_{1} + W_{2} + W_{3} + \cdots$$

$$= \underbrace{V_{-3}}_{V_{-4} + W_{-3}} W_{-3} + W_{-2} + W_{-1} + W_{0} + W_{1} + W_{2} + W_{3} + \cdots$$

$$\vdots$$

$$= \cdots + W_{-3} + W_{-2} + W_{-1} + W_{0} + W_{1} + W_{2} + W_{3} + \cdots$$

Remark 4.21 In the special case that two subspaces W_1 and W_2 are *orthogonal* to each other, then the *subspace addition* operation $W_1 + W_2$ is frequently expressed as $W_1 \oplus W_2$. In the case of an *orthonormal wavelet system*, the expressions in Theorem 4.20 (page 47)

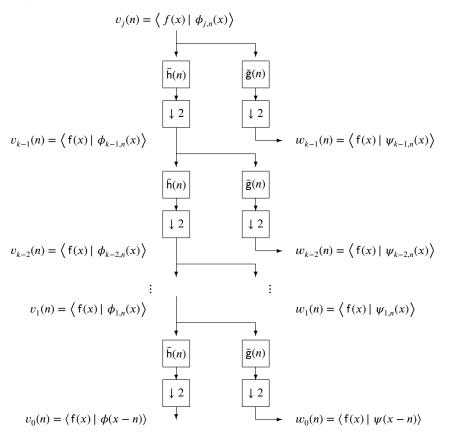
could be expressed as

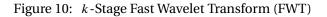
.

$$\begin{aligned} \mathcal{L}_{\mathbb{R}}^{2} &= \lim_{j \to \infty} \mathcal{V}_{j} \\ &= \mathcal{V}_{j} \oplus \mathcal{W}_{j} \oplus \mathcal{W}_{j+1} \oplus \mathcal{W}_{j+2} \oplus \cdots \\ &= \cdots \oplus \mathcal{W}_{-2} \oplus \mathcal{W}_{-1} \oplus \mathcal{W}_{0} \oplus \mathcal{W}_{1} \oplus \mathcal{W}_{2} \oplus \cdots \end{aligned}$$

4.5 Fast Wavelet Transform (FWT)

Filter banks can be used to implement a "*Fast Wavelet Transform*" (*FWT*). This is illustrated in Figure 10 page 49.¹¹⁶





¹¹⁶ S [119], page 257, (Figure 7.12), S [75], pages 371–372, (Figure L.1)

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and

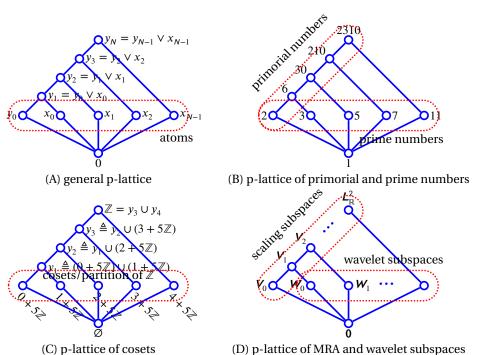


Figure 11: Some selected *primorial lattices* (see Example 5.2 page 50–Example 5.5 page 51)

5 Main Results

5.1 Primorial Lattices

Definition 5.1 Let $X \triangleq \{0, x_0, x_1, \dots, x_N, y_0, y_1, \dots, y_N\}$ be a set. A *lattice* $L \triangleq (X, \lor, \land; \le)$ is **primorial** if

- 1. 0 is the *least element* of *L*
- 2. L is *atomic* (Definition 1.44 page 13) and $\{y_0, x_0, x_1, \dots, x_N\}$ are *atoms* of L and 3. $y_{n+1} = y_n \lor x_n$.

A lattice that is *primorial* is a **primorial lattice**, or simply a **p-lattice**.

Example 5.2 A general *primorial lattice* is illustrated to in Figure 11 page 50 (A).

Example 5.3 ¹¹⁷ The set of *primorial numbers* and *prime numbers* ordered by the *divides* ("|") relation forms a *primorial lattice*, as illustrated in Figure 11 page 50 (B).

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^{117 [75],} page 30, 🖳 [2] (http://oeis.org/A002110)

5 MAIN RESULTS

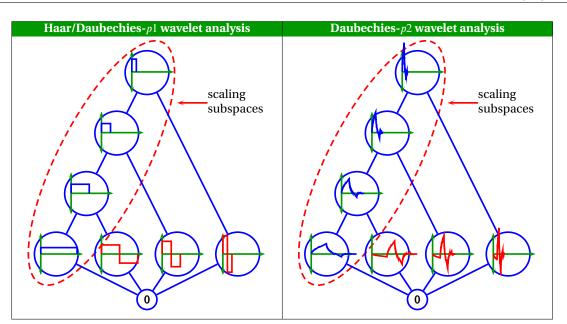


Figure 12: some MRA-wavelet systems

Example 5.4 Any *partition*, along with successive unions of the partition elements, generates a *primorial lattice*. One example of this is the *cosets* of \mathbb{Z} , which generate a *finite* primorial lattice, as illustrated in Figure 11 page 50 (C).

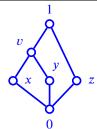
Example 5.5 A special characteristic of MRA-wavelet analysis is that it's order structure with respect to the \subseteq relation is not a simple M_n *lattice* (as is with the case of Fourier and several other analyses). Rather, it is a *primorial lattice* as illustrated in Figure 11 page 50 (D) and in Figure 12 page 51.

Proposition 5.6 ¹¹⁸ Let $L \triangleq (X, \lor, \land; \leq)$ be a LATTICE.

$$\left\{\begin{array}{ll} L \text{ is} \\ primorial \end{array}\right\} \implies \left\{\begin{array}{ll} 1. \quad L \text{ is NONDISTRIBUTIVE} & (Definition 1.53 page 15) & and \\ 2. \quad L \text{ is NONMODULAR} & (Definition 1.47 page 14) & and \\ 3. \quad L \text{ is COMPLEMENTED} \iff L \text{ is FINITE} & (Definition 1.63 page 17) & and \\ 4. \quad L \text{ is NOT UNIQUELY COMPLEMENTED} & (Definition 1.63 page 17) & and \\ 5. \quad L \text{ is NONORTHOCOMPLEMENTED} & (Definition 1.72 page 20) & and \\ 6. \quad L \text{ is NONBOOLEAN} & (Definition 1.69 page 18) & . \end{array}\right\}$$

¹¹⁷ ● [75], page 72, 〈Section 2.4.3 Order structure〉
 ¹¹⁸ ● [75], page 52, 〈Proposition 2.6〉

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[®]Proof:

- (1) Proof that *L* is *nondistributive*:
 - (a) *L* contains the *N5 lattice* (Definition 1.49 page 14).
 - (b) Because *L* contains the *N5* lattice, *L* is *nondistributive* (Theorem 1.57 page 16).
- (2) Proof that *L* is *nonmodular* and *nondistributive*:
 - (a) L contains the N5 lattice (Definition 1.49 page 14).
 - (b) Because *L* contains the *N*5 lattice, *L* is *nonmodular* (Theorem 1.50 page 14).
- (3) Proof that *L* is *noncomplemented*:

$$x' = y' = v' = z$$

$$z' = \{x, y, v\}$$

$$x'' = (x')'$$

$$= z'$$

$$= \{x, y, v\}$$

$$\neq x$$

- (4) Proof that *L* is *nonBoolean*:
 - (a) *L* is *nondistributive* (item 1 page 52).
 - (b) Because *L* is *nondistributive*, it is *nonBoolean* (Definition 1.69 page 18).

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5.2 Reduction operator on boolean lattices

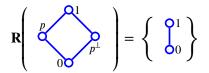
Definition 5.7 Let \mathbb{B} be the set of all *bounded lattices* (Definition 1.39 page 12). Let $L_2^N \triangleq (X, \lor, \land, 0, 1; \le)$ be a *Boolean lattice* (Definition 1.69 page 18) with 2^N elements and $N \in \mathbb{N}$ (N is a positive integer). The operator **R** is the **lattice reduction operator** of L_2^N and $\mathbf{R}L_2^N$ is the **reduction of** L_2^N if

$$\mathbf{RL}_{2}^{N} \triangleq \left\{ \boldsymbol{L} \in \mathbb{B} \middle| \begin{array}{c} 1. \quad \boldsymbol{L} \text{ is a } 2^{N-1} \text{ element } Boolean \ lattice & \text{and} \\ 2. \quad \boldsymbol{L} \subseteq \boldsymbol{L}_{2}^{N} & \text{and} \\ 3. \quad \{0,1\} \in \boldsymbol{L} & \text{and} \\ 4. \quad \{x,y\} \text{ is an orthocomplemented pair in } \boldsymbol{L} \Longrightarrow \\ \quad \{x,y\} \text{ is an orthocomplemented pair in } \boldsymbol{L}_{2}^{N} & \end{array} \right\}$$

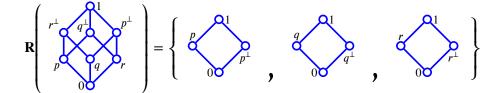
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Note that in Definition 5.7, the *order relation* \leq is the same for both L_2^N and any L in $\mathbb{R}L_2^N$. That is, if $x \leq y$ in L_2^N , then $x \leq y$ in L as well.

Example 5.8 Let L_2^2 be a *Boolean lattice* (Definition 1.69 page 18) of order 2. Let **R** be the *lattice* reduction operator **R** and $\mathbf{R}L_2^2$ be the reduction of L_2^2 (Definition 5.7 page 52). Then $\mathbf{R}L_2^2$ yields a set of exactly one 2^{2-1} value Boolean lattice, as illustrated next:



Example 5.9 Let L_2^3 be a *Boolean lattice* (Definition 1.69 page 18) of order 3. Let **R** be the **lattice** reduction operator **R** and **R** L_2^3 be the reduction of L_2^3 (Definition 5.7 page 52). The operation **R** L_2^3 yields a set of three 2^2 value Boolean lattices, as illustrated next:



Example 5.10 Let L_2^4 be a *Boolean lattice* (Definition 1.69 page 18) of order 4. Let **R** be the **lattice** reduction operator **R** and **R** L_2^4 be the reduction of L_2^4 (Definition 5.7 page 52). The operation **R** L_2^4 yields a set of ten 2^3 value Boolean lattices, as illustrated in Figure 13 (page 54).

Remark 5.11 In a *boolean lattice* L_2^N (Definition 1.69 page 18), besides the pair $\{0, 1\}$, there are a total of $2^{N-1} - 1$ orthocomplemented (Definition 1.72 page 20) pairs of elements. But note that any arbitrary $2^{N-1} - 2$ pairs of orthocomplemented pairs does not in general generate a *boolean lattice*. The lattice L_2^4 , for example, has $2^{4-1} - 1 = 7$ orthocomplemented pairs besides $\{0, 1\}$. To generate an L_2^3 lattice, we need 3 orthocomplemented pairs. There are $\binom{7}{3} = \frac{7!}{3!4!} = 35$ ways of selecting 3 pairs from L_2^4 , but only 10 of these ways generate a *boolean lattice* (Example 5.10 page 53). All other ways fail.

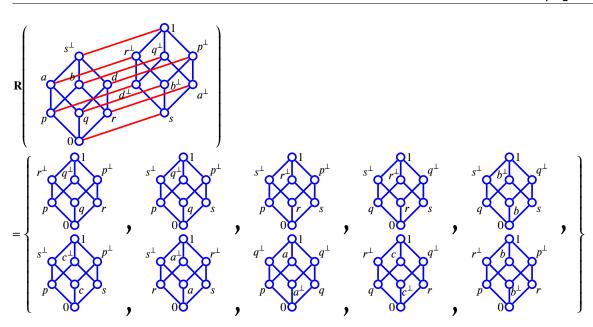
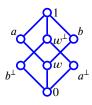


Figure 13: reduction of L_2^4 (Example 5.10 page 53)

For example, if we were to select the pairs $\{0, w, w^{\perp}, a, a^{\perp}, b, b^{\perp}, 1\}$, we would get the *orthocomplemented*, but **non-boolean** (Definition 1.69 page 18) lattice illustrated to the right; In particular, it is *complemented*, but *non-distributive*. For example, $w^{\perp} \wedge (a \vee b) = w^{\perp} \neq 0 = 0 \vee 0 = (w^{\perp} \wedge a) \vee (w^{\perp} \wedge b)$. Alternatively, note that the set $\{1, a, w, 0, b^{\perp}, w^{\perp}\}$ together with the ordering relation \leq form an O_6 sublattice (Definition 1.73 page 20), which contains an N_5 sublattice, which implies that the lattice to the right is *non-distributive* (by the *Birkhoff distributivity criterion* Theorem 1.57 page 16).



Example 5.12 Let L_2^5 be a *Boolean lattice* (Definition 1.69 page 18) of order 5. Let **R** be the **lattice** reduction operator **R** and **R** L_2^5 be the reduction of L_2^5 (Definition 5.7 page 52). The result of the operation **R** L_2^5 is partially illustrated in Figure 14 (page 55).

5.3 Difference operator on bounded lattices

Definition 5.13 Let $X \setminus Y$ be the standard *set difference* of a set X and a set Y. Let $L_x \triangleq (X, \lor, \land, 0, 1; \le)$ and $L_y \triangleq (Y, \lor, \land, 0, 1; \le)$ be *bounded lattices* (Definition 1.39 page 12). The **bounded lattice difference** $L_x \otimes L_y$ of L_x and L_y is the lattice L such that $L \triangleq ((X \setminus Y) \cup \{0, 1\}, \lor, \land, 0, 1; \le)$

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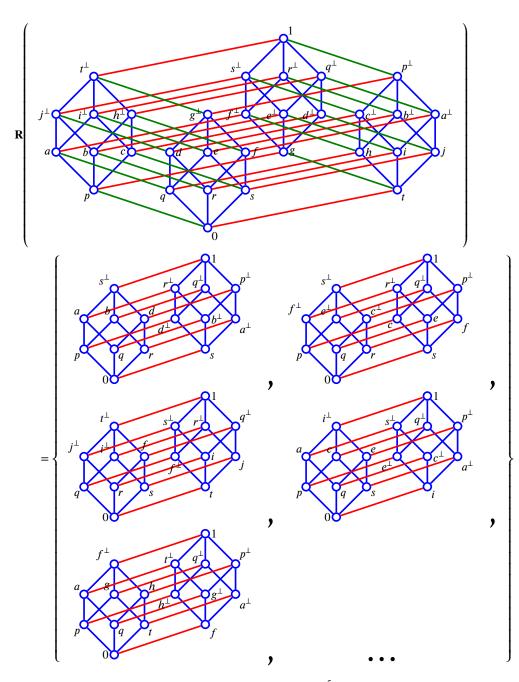
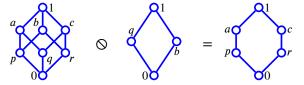


Figure 14: reduction of L_2^5 (Example 5.12 page 54)

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Example 5.14 Let \otimes be the *bounded lattice difference* operator (Definition 5.13 page 54).



Proposition 5.15 Let B be the set of all BOUNDED LATTICES (Definition 1.39 page 12). Let \bigcirc be the BOUNDED LATTICE DIFFERENCE Operator (Definition 5.13 page 54). $(\mathbb{B}, \mathbb{Q}, \subseteq)$ is a D-POSET (Definition 3.18 page 37).

Theorem 5.16 Let $L \triangleq L_2^N \otimes L_2^{N-1}$ be the BOUNDED LATTICE DIFFERENCE (Definition 5.13 page 54) of a BOOLEAN LATTICE L_2^N (Definition 1.69 page 18) and a BOOLEAN LATTICE L_2^{N-1} selected from the set \mathbf{RL}_2^N (Definition 5.7 page 52). Let $X \triangleq \{L_2^n | n = 1, 2, ...\} \cup \{L_2^n \otimes L_2^{n-1} | n = 2, 3, ...\}$. 1. $L_2^N \otimes L_2^{N-1}$ is an **orthocomplemented lattice** (Definition 1.72 page 20) and 2. The structure $\mathbb{P} \triangleq (X, \lor, \land; \subseteq)$ is a **primorial lattice** (Definition 5.1 page 50).

Proof:

- (1) Proof that $L_2^N \otimes L_2^{N-1}$ is an **orthocomplemented lattice**:
 - (a) L_2^N is a *Boolean lattice* by definition.
 - (b) L_2^{N-1} is also a *Boolean lattice* (Definition 5.7 page 52).
 - (c) Every lattice that is *Boolean* is also *orthocomplemented* (Proposition 1.80 page 23).
 - (d) By definition of $L_2^N \otimes L_2^{N-1}$, orthocomplemented pairs are removed from L_2^N and the orthocomplemented pair $\{0, 1\}$ is put back in.
 - (e) What remains in $L_2^N \otimes L_2^{N-1}$ is a set of *orthocomplemented pairs*, ordered with the same ordering relation \leq that orders L_2^N .
 - (f) All remaining orthocomplemented pairs are still involutory: $x = x^{\perp \perp}$ $\forall x \in X$
 - (g) All remaining *orthocomplemented pairs* are still *antitone* because the *ordering relation* $\leq \text{ in } L_2^N \otimes L_2^{N-1}$ is the same.
 - (h) All remaining *orthocomplemented pairs* still have the *non-contradiction* property be-cause suppose that in $L_2^N \otimes L_2^{N-1}$, there is an element *x* such that $x \wedge x^{\perp} = m \neq 0$. Then in L_2^N , it would also be true that $x \wedge x^{\perp} \neq 0$. This cannot be true (is a contradiction); so therefore for all *x* in $L_2^N \otimes L_2^{N-1}$, $x \wedge x^{\perp} = 0$ (*non-contradiction* property).
 - (i) So $L_2^N \otimes L_2^{N-1}$ is an orthocomplemented lattice (Definition 1.72 page 20).
- (2) Proof that $(X \triangleq \{L_2^n | n = 1, 2, ...\} \cup \{L_2^n \otimes L_2^{n-1} | n = 2, 3, ...\}, \subseteq)$ is a **primorial lattice**: This follows directly from the construction of the *bounded lattice difference* (Definition 5.13 page 54) and the definition of *primorial lattices* (Definition 5.1 page 50).

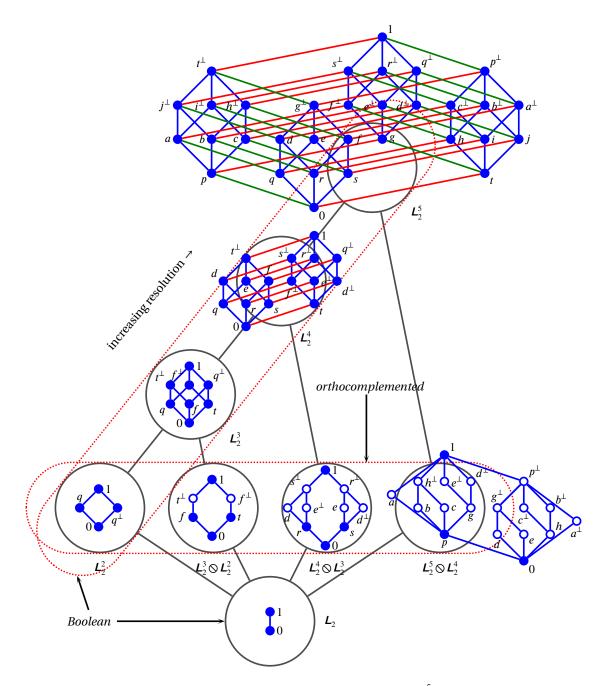


Figure 15: *a primorial lattice generated by* L_2^5

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Definition 5.17 Let L_2^N be a 2^N element *Boolean lattice* (Definition 1.69 page 18). The lattice \mathbb{P} as described in Theorem 5.16 is a **primorial lattice generated by** L_2^N .

Example 5.18 Figure 15 (page 57) illustrates a *primorial lattice generated by* L_2^5 .

Projections on primorial lattices 5.4

This section introduces three lattice projections. When performing analysis in a *primorial lattice* (Definition 5.1 page 50), it is necessary to project a point that exists in a lattice of "high resolution" onto a lattice L of lower resolution that may or may not contain this point. The three projections introduced here are the

- 1. zero primorial projection (Definition 5.19 page 58) which assigns to 0 any point that does not exist in L
- 2. Sasaki primorial projection (Definition 5.20 page 58) which assigns a projection value using the Sasaki projection (Definition 2.22 page 31)
- 3. metric primorial projection (Definition 5.22 page 59) which assigns a projection value based on a *lattice metric* (Definition 2.7 page 27).

Definition 5.19 Let P be a primorial lattice (Definition 5.17 page 58) generated by a Boolean lat*tice* L_2^N (Definition 1.69 page 18). Let $L \triangleq (Y, \lor, \land, 0, 1; \le)$ be a lattice in \mathbb{P} . Let $x \triangleq (x_n)$ be a *sequence* over the set X. The **zero primorial projection** $\Phi_L(x)$ of x onto L is defined as $\Phi_L^z(x) \triangleq \bigvee_L [\{x, 0\} \cap Y] \quad \forall x \in X$

The **zero primorial projection** $\Phi_L^z(x)$ of x onto L is defined as $\Phi_L^z(x) \triangleq (y_n)$ where $y_n \triangleq \Phi_L^z(x_n) \quad \forall x_n \in (x_n), y_n \in (y_n)$.

Definition 5.20 Let \mathbb{P} and x be defined as in Definition 5.19 (page 58). Let \mathbb{P} be a *primorial lattice* (Definition 5.17 page 58) generated by a *Boolean lattice* L_2^N (Definition 1.69 page 18). Let $L \triangleq (Y, \lor, \land, 0, 1; \le)$ be a lattice in \mathbb{P} . Let $x \triangleq ((x_n))$ be a sequence over the set X. The **Sasaki primorial projection** $\Phi_L^s(x)$ of x onto L is defined as

 $\Phi_{L}^{s}(x) \triangleq \bigvee_{I} \left[\left\{ \phi_{y}(x) \mid y \in Y \right\} \cap Y \right]$ ∀*x*∈*L*

where $\phi_y(x)$ is the Sasaki projection of x onto y (Definition 2.22 page 31) in the smallest Boolean *lattice* L_2^{M} that contains both x and L. The Sasaki primorial projection $\Phi_L^s(x)$ of x onto L is defined as

 $\boldsymbol{\Phi}_{\boldsymbol{I}}^{s}(\mathbf{x}) \triangleq (y_{n}) \text{ where } y_{n} \triangleq \boldsymbol{\Phi}_{\boldsymbol{L}}^{s}(x_{n}) \quad \forall x_{n} \in (x_{n}).$

The Sasaki primorial projection yields a kind of maxmini (Theorem 1.35 page 11) result:

MONDAY 13TH OCTOBER, 2014 MRA-Wa^{V^el}et subspace architecture for logic, probability, and symbolic sequence processing VERSION 0.65 **Proposition 5.21** Let $\Phi_L(x)$ be the SASAKI PRIMORIAL PROJECTION of x onto L in a PRIMO-RIAL LATTICE \mathbb{P} .

$$\Phi_{L}^{s}(x) = \bigvee_{L} \left[\left\{ x \land y \, | \, y \in Y \right\} \cap Y \right] \quad \forall x \in X$$

PROOF:

$$\begin{split} \varPhi_{L}^{s}(x) &\triangleq \bigvee \left[\left\{ \phi_{y}(x) \mid y \in Y \right\} \cap Y \right] & \text{by def. of } Sasaki \, primorial \, projection \, (Definition 5.20 \, page 58) \\ &\triangleq \bigvee \left[\left\{ (x \lor y^{\perp}) \land y \mid y \in Y \right\} \cap Y \right] & \text{by definition of } Sasaki \, projection \, (Definition 2.22 \, page 31) \\ &= \bigvee \left[\left\{ (x \land y) \lor (y^{\perp} \land y) \mid y \in Y \right\} \cap Y \right] & \text{by distributive prop. (Theorem 1.70 \, page 19)} \\ &= \bigvee \left[\left\{ (x \land y) \lor (0) \mid y \in Y \right\} \cap Y \right] & \text{by noncontradiction property (Theorem 1.70 \, page 19)} \\ &= \bigvee \left[\left\{ (x \land y \mid y \in Y \right\} \cap Y \right] & \text{by bounded property (Theorem 1.70 \, page 19)} \\ \end{split}$$

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and

Definition 5.22 Let \mathbb{P} and \times be defined as in Definition 5.19. The *metric primorial projection* $\Phi_L^m(x)$ of x onto **L** is defined as

 $\Phi_{L}^{m}(x) \triangleq \bigwedge_{r} \left[\overline{\mathsf{B}}(x,r) \cap Y\right]$ where

- 1. $\overline{B}(x, r)$ is the *closed ball* in (L_2^M, d) with the smallest radius *r* that contains *x* 2. (L_2^M, d) is a *metric lattice* (Definition 2.7 page 27) and
- 3. L_2^{M} is the smallest *Boolean lattice* (Definition 1.69 page 18) containing x and
- 4. the *valuation* function defining d is the *height* function on L_2^M .
- The metric primorial projection $\Phi_L(\mathbf{x})$ of \mathbf{x} onto \mathbf{L} is defined as

 $\Phi_L(\mathbf{x}) \triangleq (y_n)$ such that $y_n \triangleq \Phi_L(x_n)$.

Example 5.23 Here are examples of the *primorial projections* $\Phi_{O_6}^z(x)$ (Definition 5.19 page 58), $\Phi_{O_{\epsilon}}^{s}(x)$ (Definition 5.20 page 58), and $\Phi_{O_{\epsilon}}^{m}(x)$ (Definition 5.22 page 59) in the primorial lattice (Definition 5.1 page 50) generated by the *Boolean lattice* (Definition 1.69 page 18) $L_2^5 \triangleq (X, \lor, \land, 0, 1; \le)$ as illustrated in Figure 15 page 57 onto the lattice $O_6 \triangleq L_2^3 \otimes L_2^2 \triangleq (Y, \vee, \wedge, 0, 1; \leq)$.

projection)	$x \text{ in } \boldsymbol{O}_6 \triangleq \boldsymbol{L}_2^3 \otimes \boldsymbol{L}_2^2$					x i	$\frac{1}{n}L_2^3$		$\frac{2}{x \text{ ir}}$	L_2^4				x i	n L_2^5					
<i>x</i> =																					
$\Phi_{O_6}^z(x) =$	0	f	t	t^{\perp}	f^{\perp}	1	0	0	0	0	0	0	0	0	0	0	0	0			
$\Phi^{s}_{O_6}(x) =$	0	f	t	t^{\perp}	f^{\perp}	1	0	1	0	f^{\perp}	0	f^{\perp}	t	f	0	1	0	t			
	0	f	t	t^{\perp}	f^{\perp}	1	0	0	0	f^{\perp}	0	f^{\perp}	t	f	0	1	0	t			

PROOF:

MRA-Wa^{V^{Cl}et} subspace architecture for logic, probability, and symbolic sequence processing 4 Monday 13TH October, 2014 VERSION 0.65 (1) Proof for zero primorial projection values:

$$\begin{array}{lll} \varPhi_{\mathcal{O}_{6}}^{z}(0) = \bigvee \left[(\{0\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, f^{\perp}, 1\} \right] &= \bigvee \left[\{0\} \right] &= 0 \\ \varPhi_{\mathcal{O}_{6}}^{z}(f) = \bigvee \left[(\{f\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, f^{\perp}, 1\} \right] &= \bigvee \left[\{0, f\} \right] &= f \\ \varPhi_{\mathcal{O}_{6}}^{z}(t) = \bigvee \left[(\{t\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, f^{\perp}, 1\} \right] &= \bigvee \left[\{0, t\} \right] &= t \\ \varPhi_{\mathcal{O}_{6}}^{z}(t^{\perp}) = \bigvee \left[(\{t^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, f^{\perp}, 1\} \right] &= \bigvee \left[\{0, t^{\perp}\} \right] &= t^{\perp} \\ \varPhi_{\mathcal{O}_{6}}^{z}(f^{\perp}) = \bigvee \left[(\{f^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, f^{\perp}, 1\} \right] &= \bigvee \left[\{0, t^{\perp}\} \right] &= f^{\perp} \\ \varPhi_{\mathcal{O}_{6}}^{z}(f^{\perp}) = \bigvee \left[(\{f^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, f^{\perp}, 1\} \right] &= \bigvee \left[\{0, t^{\perp}\} \right] &= f \\ \varPhi_{\mathcal{O}_{6}}^{z}(f^{\perp}) = \bigvee \left[(\{f^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, f^{\perp}, 1\} \right] &= \bigvee \left[\{0, t^{\perp}\} \right] &= 0 \\ \varPhi_{\mathcal{O}_{6}}^{z}(g) = \bigvee \left[(\{q^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, f^{\perp}, 1\} \right] &= \bigvee \left[\{0\} \right] &= 0 \\ \varPhi_{\mathcal{O}_{6}}^{z}(g^{\perp}) = \bigvee \left[(\{f^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, f^{\perp}, 1\} \right] &= \bigvee \left[\{0\} \right] &= 0 \\ \varPhi_{\mathcal{O}_{6}}^{z}(g^{\perp}) = \bigvee \left[(\{r^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, r^{\perp}, 1\} \right] &= \bigvee \left[\{0, r^{\perp}\} \right] \\ = 0 \\ \varPhi_{\mathcal{O}_{6}}^{z}(g^{\perp}) = \bigvee \left[(\{g^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, r^{\perp}, 1\} \right] &= \bigvee \left[\{0\} \right] \\ = 0 \\ \varPhi_{\mathcal{O}_{6}}^{z}(g^{\perp}) = \bigvee \left[(\{g^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, r^{\perp}, 1\} \right] \\ = \bigvee \left[\{0\} \right] \\ = 0 \\ \varPhi_{\mathcal{O}_{6}}^{z}(g^{\perp}) = \bigvee \left[(\{g^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, r^{\perp}, 1\} \right] \\ = \bigvee \left[\{0\} \right] \\ = 0 \\ \varPhi_{\mathcal{O}_{6}}^{z}(g^{\perp}) = \bigvee \left[(\{g^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, r^{\perp}, 1\} \right] \\ = \bigvee \left[\{0\} \right] \\ = 0 \\ \varPhi_{\mathcal{O}_{6}}^{z}(g^{\perp}) = \bigvee \left[(\{g^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, r^{\perp}, 1\} \right] \\ = \bigvee \left[\{0\} \right] \\ = 0 \\ \varPhi_{\mathcal{O}_{6}}^{z}(g^{\perp}) = \bigvee \left[(\{g^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, r^{\perp}, 1\} \right] \\ = \bigvee \left[\{0\} \right] \\ = 0 \\ \varPhi_{\mathcal{O}_{6}}^{z}(g^{\perp}) = \bigvee \left[(\{g^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, r^{\perp}, 1\} \right] \\ = \bigvee \left[\{0\} \right] \\ = 0 \\ \varPhi_{\mathcal{O}_{6}}^{z}(g^{\perp}) = \bigvee \left[(\{g^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, r^{\perp}, 1\} \right] \\ = \bigvee \left[\{0\} \right] \\ = 0 \\ \liminf_{\mathcal{O}_{6}}^{z}(g^{\perp}) = \bigvee \left[(\{g^{\perp}\} \cup \{0\}) \cap \{0, f, t, t^{\perp}, r^{\perp}, 1\} \right] \\ = \bigvee \left[\{0\} \right] \\ = 0 \\ \liminf_{\mathcal{O}_{6}}^{z}(g^{\perp}) = 0 \\ \liminf_{\mathcal{O}_{6}}^{z}(g^{\perp}) = 0 \\ \liminf_{\mathcal{O}_{6}}^{z}(g^$$

(2) Proof for Sasaki primorial projection (Definition 5.20 page 58):

$$\begin{split} \varPhi_{O_{6}}^{s}(0) &= \bigvee \left[\{0 \land y \mid y \in Y \} \cap Y \right] &= \bigvee \left[\{0, 0, 0, 0, 0\} \cap Y \right] &= \bigvee \{0\} &= 0 \\ \varPhi_{O_{6}}^{s}(f) &= \bigvee \left[\{f \land y \mid y \in Y \} \cap Y \right] &= \bigvee \left[\{0, f, 0, f, 0, f\} \cap Y \right] &= \bigvee \{0, f\} &= f \\ \varPhi_{O_{6}}^{s}(t) &= \bigvee \left[\{t \land y \mid y \in Y \} \cap Y \right] &= \bigvee \left[\{0, t, 0, t, t\} \cap Y \right] &= \bigvee \{0, t\} &= t \\ \varPhi_{O_{6}}^{s}(t^{\perp}) &= \bigvee \left[\{t^{\perp} \land y \mid y \in Y \} \cap Y \right] &= \bigvee \left[\{0, f, 0, t^{\perp}, q, t^{\perp} \} \cap Y \right] &= \bigvee \left\{ 0, f, t^{\perp} \right\} &= t^{\perp} \\ \varPhi_{O_{6}}^{s}(f^{\perp}) &= \bigvee \left[\{f^{\perp} \land y \mid y \in Y \} \cap Y \right] &= \bigvee \left[\{0, 0, t, q, f^{\perp}, f^{\perp} \} \cap Y \right] &= \bigvee \left\{ 0, t, f^{\perp} \right\} &= f^{\perp} \\ \varPhi_{O_{6}}^{s}(1) &= \bigvee \left[\{1 \land y \mid y \in Y \} \cap Y \right] &= \bigvee \left[\{0, f, t, f^{\perp}, t^{\perp}, 1\} \cap Y \right] &= \bigvee Y &= 1 \\ \varPhi_{O_{6}}^{s}(q) &= \bigvee \left[\{q \land y \mid y \in Y \} \cap Y \right] &= \bigvee \left[\{0, 0, 0, q, 0, q\} \cap Y \right] &= \bigvee \{0\} &= 0 \\ \varPhi_{O_{6}}^{s}(q^{\perp}) &= \bigvee \left[\{q^{\perp} \land y \mid y \in Y \} \cap Y \right] &= \bigvee \left[\{0, f, t, f, t, q^{\perp} \} \cap Y \right] &= \bigvee \{0, f, t\} &= 1 \end{split}$$

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$\boldsymbol{\varPhi}_{\boldsymbol{O}_{\boldsymbol{b}}}^{s}(r) = \bigvee \left[\left\{ r \land y y \in Y \right\} \cap Y \right]$	$=\bigvee [\{0,r,0,r,0,r\}\cap Y]$	$= \bigvee \{0\}$	= 0			
$\Phi^{s}_{\mathcal{O}_{6}}(r^{\perp}) = \bigvee \left[\left\{ r^{\perp} \land y \mid y \in Y \right\} \cap Y \right]$	$= \bigvee \left[\left\{ 0, s, t, e, f^{\perp}, r^{\perp} \right\} \cap Y \right]$	$=\bigvee \left\{ 0,t,f^{\perp }\right\}$	$= f^{\perp}$			
$\boldsymbol{\Phi}^{s}_{\boldsymbol{\mathcal{O}}_{6}}(s) = \bigvee \left[\left\{ s \land y \mid y \in Y \right\} \cap Y \right]$	$= \bigvee [\{0, s, 0, s, 0, s\} \cap Y]$	$=\bigvee \{0\}$	= 0			
$\boldsymbol{\varPhi}_{\boldsymbol{\mathcal{O}}_{6}}^{s}(s^{\perp}) = \bigvee \left[\left\{ s^{\perp} \land y y \in Y \right\} \cap Y \right]$	$=\bigvee\left[\left\{0,0,t,d,f^{\perp},s^{\perp}\right\}\cap Y\right]$	$=\bigvee\left\{ 0,t,f^{\bot}\right\}$	$= f^{\perp}$			
$\boldsymbol{\Phi}_{\boldsymbol{O}_{6}}^{s}(g) = \bigvee \left[\left\{ g \land y y \in Y \right\} \cap Y \right]$	$=\bigvee \left[\left\{ 0,0,t,p,g,g \right\} \cap Y \right]$	$= \bigvee \{0, t\}$	= t			
$\boldsymbol{\varPhi}_{\boldsymbol{\mathcal{O}}_{6}}^{s}(g^{\perp}) = \bigvee \left[\left\{ g^{\perp} \land y y \in Y \right\} \cap Y \right]$	$=\bigvee\left[\left\{0,f,0,g^{\bot},0,g^{\bot}\right\}\cap Y\right]$	$= \bigvee \{0, f\}$	= f			
$\boldsymbol{\Phi}_{\boldsymbol{O}_{6}}^{s}(p) = \bigvee \left[\left\{ p \land y \mid y \in Y \right\} \cap Y \right]$	$=\bigvee \left[\{0,0,0,p,p,p\} \cap Y \right]$	$= \bigvee \{0\}$	= 0			
$\boldsymbol{\varPhi}_{\boldsymbol{\mathcal{O}}_{6}}^{s}(p^{\perp}) = \bigvee \left[\left\{ p^{\perp} \land y y \in Y \right\} \cap Y \right]$	$=\bigvee\left[\left\{0,f,t,g^{\bot},t,p^{\bot}\right\}\cap Y\right]$	$=\bigvee \left\{ 0,f,t\right\}$	= 1			
$\varPhi_{\mathcal{O}_{6}}^{s}(d) = \bigvee \left[\left\{ d \land y \middle y \in Y \right\} \cap Y \right]$	$= \bigvee \left[\left\{ 0, r, 0, d, 0, d \right\} \cap Y \right]$	$= \bigvee \{0\}$	= 0			
$\varPhi^{s}_{\mathcal{O}_{6}}(d^{\perp}) = \bigvee \left[\left\{ d^{\perp} \wedge y y \in Y \right\} \cap Y \right]$	$=\bigvee\left[\left\{0,s,t,0,g,d^{\perp}\right\}\cap Y\right]$	$=\bigvee \{0,t\}$	= t			

(3) Proof for metric primorial projection (Definition 5.22 page 59):

$$\begin{split} & \varPhi_{O_{6}}^{m}(0) = \bigwedge \left[\overline{B}(0,0) \cap Y\right] &= \bigwedge \left[\{0\} \cap \{0,f,t,t^{\perp},f^{\perp},1\}\right] &= \bigwedge \{0\} &= 0 \\ & \varPhi_{O_{6}}^{m}(f) = \bigwedge \left[\overline{B}(f,0) \cap Y\right] &= \bigwedge \left[\{f\} \cap \{0,f,t,t^{\perp},f^{\perp},1\}\right] &= \bigwedge \{f\} &= f \\ & \varPhi_{O_{6}}^{m}(t) = \bigwedge \left[\overline{B}(t,0) \cap Y\right] &= \bigwedge \left[\{t\} \cap \{0,f,t,t^{\perp},f^{\perp},1\}\right] &= \bigwedge \{t\} &= t \\ & \varPhi_{O_{6}}^{m}(t^{\perp}) = \bigwedge \left[\overline{B}(t^{\perp},0) \cap Y\right] &= \bigwedge \left[\{t^{\perp}\} \cap \{0,f,t,t^{\perp},f^{\perp},1\}\right] &= \bigwedge \{t^{\perp}\} &= t^{\perp} \\ & \varPhi_{O_{6}}^{m}(f^{\perp}) = \bigwedge \left[\overline{B}(f^{\perp},0) \cap Y\right] &= \bigwedge \left[\{f^{\perp}\} \cap \{0,f,t,t^{\perp},f^{\perp},1\}\right] &= \bigwedge \{f^{\perp}\} &= f^{\perp} \\ & \varPhi_{O_{6}}^{m}(f^{\perp}) = \bigwedge \left[\overline{B}(f,0) \cap Y\right] &= \bigwedge \left[\{f^{\perp}\} \cap \{0,f,t,t^{\perp},f^{\perp},1\}\right] &= \bigwedge \{f^{\perp}\} &= f^{\perp} \\ & \varPhi_{O_{6}}^{m}(f) = \bigwedge \left[\overline{B}(f,0) \cap Y\right] &= \bigwedge \left[\{f^{\perp}\} \cap \{0,f,t,t^{\perp},f^{\perp},1\}\right] &= \bigwedge \{f^{\perp}\} &= f^{\perp} \\ & \varPhi_{O_{6}}^{m}(g) = \bigwedge \left[\overline{B}(g,1) \cap Y\right] &= \bigwedge \left[\{q,0,t^{\perp}\} \cap Y\right] &= \bigwedge \{\{0,t^{\perp}\} &= 0 \\ & \varPhi_{O_{6}}^{m}(g^{\perp}) = \bigwedge \left[\overline{B}(g,1) \cap Y\right] &= \bigwedge \left[\{q,0,t^{\perp}\} \cap Y\right] &= \bigwedge \{f,t,1\} &= t \\ & \varPhi_{O_{6}}^{m}(g) = \bigwedge \left[\overline{B}(g,1) \cap Y\right] &= \bigwedge \left[\{r,0,d,f\} \cap Y\right] &= \bigwedge \{f^{\perp},1\} &= f^{\perp} \\ & \varPhi_{O_{6}}^{m}(g) = \bigwedge \left[\overline{B}(g,1) \cap Y\right] &= \bigwedge \left[\{r^{\perp},d^{\perp},f^{\perp},1\} \cap Y\right] &= \bigwedge \{f^{\perp},1\} &= f^{\perp} \\ & \varPhi_{O_{6}}^{m}(g) = \bigwedge \left[\overline{B}(g,1) \cap Y\right] &= \bigwedge \left[\{s^{\perp},e^{\perp},f^{\perp},d^{\perp},f^{\perp},1\} \cap Y\right] &= \bigwedge \{f^{\perp},1\} &= f^{\perp} \\ & \varPhi_{O_{6}}^{m}(g) = \bigwedge \left[\overline{B}(g,1) \cap Y\right] &= \bigwedge \left[\{g^{\perp},d,e,f,p^{\perp},t^{\perp}\} \cap Y\right] &= \bigwedge \{f^{\perp},1\} &= f^{\perp} \\ & \varPhi_{O_{6}}^{m}(g) = \bigwedge \left[\overline{B}(g,1) \cap Y\right] &= \bigwedge \left[\{g^{\perp},d,e,f,p^{\perp},t^{\perp}\} \cap Y\right] &= \bigwedge \{f^{\perp},t\} &= f \\ & \varPhi_{O_{6}}^{m}(g) = \bigwedge \left[\overline{B}(g,1) \cap Y\right] &= \bigwedge \left[\{g^{\perp},d,e,f,p^{\perp},t^{\perp}\} \cap Y\right] &= \bigwedge \{f^{\perp},t\} &= f \\ & \varPhi_{O_{6}}^{m}(g) = \bigwedge \left[\overline{B}(g,1) \cap Y\right] &= \bigwedge \left[\{g^{\perp},d,e,f,p^{\perp},t^{\perp}\} \cap Y\right] &= \bigwedge \{f^{\perp},t\} &= f \\ & \varPhi_{O_{6}}^{m}(g) = \bigwedge \left[\overline{B}(g,1) \cap Y\right] &= \bigwedge \left[\{g^{\perp},d,e,f,g^{\perp},t^{\perp}\} \cap Y\right] &= \bigwedge \{f^{\perp},t\} &= f \\ & \varPhi_{O_{6}}^{m}(g) = \bigwedge \left[\overline{B}(g,1) \cap Y\right] &= \bigwedge \left\{g^{\perp}(g,e,g,g^{\perp}) \cap Y\right] &= \bigwedge \{g^{\perp}(g) = g \\ & = \bigwedge \{g^{\perp}(g) = (g^{\perp},g^{\perp}) \cap Y\right] &= \bigwedge \{g^{\perp}(g^{\perp},g^{\perp},g^{\perp}) \cap Y\right] &= \bigwedge \{g^{\perp}(g^{\perp},g^{\perp}) \cap Y\right] &= \bigwedge \{g^{\perp}(g^{\perp},g^{\perp},g^{\perp}) \cap Y\right] &= \bigwedge \{g^{\perp}(g^{\perp},g^{\perp},g^{\perp}) \cap Y\right] \\ & = \bigwedge \{g^{\perp}(g^{\perp},g^{\perp}) \cap Y\right] &= \bigwedge \{g^{$$

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$$\boldsymbol{\Phi}_{\boldsymbol{\mathcal{O}}_{6}}^{m}(p^{\perp}) = \bigwedge \left[\overline{\mathsf{B}}\left(p^{\perp},1\right) \cap Y \right] \qquad = \bigwedge \left[\left\{ p^{\perp},a^{\perp},b^{\perp},c^{\perp},g^{\perp},1 \right\} \cap Y \right] \qquad = \bigwedge \left\{ 1 \right\} \qquad = 1$$

$$\begin{split} \boldsymbol{\varPhi}_{\mathcal{O}_{6}}^{m}(d) &= \bigwedge \left[\overline{\mathsf{B}}\left(d,2\right) \cap \left\{0,f,t,t^{\perp},f^{\perp},1\right\} \right] \\ &= \bigwedge \left[\left\{0,a,b,d,e,f,h,i,q,r,c^{\perp},g^{\perp},j^{\perp},p^{\perp},s^{\perp},t^{\perp}\right\} \cap \left\{0,f,t,t^{\perp},f^{\perp},1\right\} \right] \\ &= \bigwedge \left\{0,f,t^{\perp}\right\} \\ &= 0 \\ \boldsymbol{\varPhi}_{\mathcal{O}_{6}}^{m}(d^{\perp}) &= \bigwedge \left[\overline{\mathsf{B}}\left(d^{\perp},2\right) \cap \left\{0,f,t,t^{\perp},f^{\perp},1\right\} \right] \\ &= \bigwedge \left[\left\{c,g,j,p,s,t,a^{\perp},b^{\perp},d^{\perp},e^{\perp},f^{\perp},h^{\perp},i^{\perp},q^{\perp},r^{\perp},1\right\} \cap \left\{0,f,t,t^{\perp},f^{\perp},1\right\} \right] \\ &= \bigwedge \left\{t,f^{\perp},1\right\} \\ &= t \end{split}$$

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5.5 A generalized probability function

This paper introduces a new definition for a lattice-valued probability function (next).

Definition 5.24 Let $L \triangleq (X, \lor, \land, \neg, 0, 1; \le)$ be a *lattice with negation* (Definition 2.16 page 30). Let D be the *distributivity* relation (Definition 1.52 page 15). A function p in \mathbb{R}^X is a **probability** on L if

1.				p(0)	=	0		(nondegenerate)	and
2.				p (1)	=	1		(normalized)	and
3.	$x \leq y$		\Rightarrow	p(x)	\leq	p (<i>y</i>)	$\forall x, y \in X$	(monotone)	and
4.	$\begin{cases} x \land y = 0\\ (z, x, y) \in \textcircled{D} \end{cases}$	$\left. \begin{array}{c} \text{and} \\ \forall z \in X \end{array} \right\}$	\Rightarrow	$p(x \lor y)$	=	p(x) + p(y)	$\forall x, y \in X$	(additive).	
ic	is a probability on a lattice with population I then $(I - n)$ is a probability space								

If p is a *probability* on a *lattice with negation L*, then (*L*, p) is a **probability space**.

Remark 5.25 Definition 5.24 page 62 (previous) is not any standard definition of the *probability function*. On a *Boolean lattice*, the **measure-theoretic probability** function, due to A. N. Kolmogorov, is defined as¹¹⁹

(1).
$$p(1) = 1 \qquad (normalized) \text{ and}$$

(2).
$$p(x) \ge 0 \qquad \forall x \in X \qquad (nonnegative) \text{ and}$$

(3).
$$\bigwedge_{n=1}^{\infty} x_n = 0 \implies p\left(\bigvee_{n=1}^{\infty} x_n\right) = \sum_{n=1}^{\infty} p(x_n) \quad \forall x_n \in X \qquad (\sigma\text{-additive}) \qquad .$$

¹¹⁹ \square [13], pages 22–23, (Probability Measures), \square [103], \square [102], page 16, (*field of probability*), \square [137], pages 8–9, (Definition 2.3(13)), \square [99], page 27

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The advantage of this definition is that p is a *measure*, and hence all the power of measure theory is subsequently at one's disposal in using p. However, it has often been argued that the requirement of σ -additivity is unnecessary for a probability function. Even as early as 1930, de Finetti argued against it, in what became a kind of polite running debate with Fréchet.¹²⁰ In fact, Kolmogorov himself provided some argument against σ -additivity when referring to the closely related Axiom of Continuity saying, "Since the new axiom is essential for infinite fields of probability only, it is almost impossible to elucidate its empirical meaning...For, in describing any observable random process we can obtain only finite fields of probability...." But in its support he added, "This limitation has been found expedient in researches of the most diverse sort."¹²¹

There are several other definitions of probability that only require *additivity* rather than σ -*additivity*. On a *Boolean lattice*, the **traditional probability** function is defined as¹²²

(1).	•	p (1)	=	1		(normalized)	and
(2).		p(x)	\geq	0	$\forall x \in X$	(nonnegative)	and
(3).	$x \wedge y = 0 \implies$	p(x	$(\lor y)$	=	p(x) + p(y)	$\forall x, y \in X$	(additive)	
This defi	nition implies (or	n a E	Booled	ın le	<i>ittice</i>) that			
(a).	p (0)	=	0				(nondegenerate	e) and
(b).	p(x)	\leq	1			$\forall x \in X$	(upper bounde	d) and
(c).	p(x)	=	1 — p	$\nabla(\neg x)$)	$\forall x \in X$		and
(d).	$p(x \lor y)$	\leq	p(x)	+ p(<i>y</i>)	$\forall x, y \in X$	(subadditive)	and
(e).	$p(x \lor y)$	=	p(x)	+ p($(y) - p(x \land y)$	$\forall x, y \in X$		and
(f).	$x \le y \implies p(x)$	\leq	p(y)			$\forall x, y \in X$	(monotone)	

On a *distributive pseudocomplemented lattice*, the **generalized probability** function has been defined as¹²³

(1).	p (0)	=	0	(nondegenerate)	and
(2).	p (1)	=	1	(normalized)	and
(3).	$0 \le p(1)$	\leq	1		and
(4).	$p(x \lor y)$	=	$p(x) + p(y) - p(x \land y)$	$\forall x, y \in X$	

On an *orthomodular lattice*, or a *finite modular lattice*, the **quantum probability** function is defined as¹²⁴

(1).			p (0)	=	0		(nondegenerate)	and
(2).			p (1)	=	1		(normalized)	and
(3).	$x \perp y$	\Rightarrow	$p(x \lor y)$	=	p(x) + p(y)	$\forall x, y \in X$	(additive)	

However, for lattices that are not *distributive*, *modular*, or *orthomodular*, none of these definitions work out so well. Take for example the O_6 *lattice* with the "very reasonable"

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¹²⁰ [[60], [[65], [[59], [[66], [[58], [[28], pages 258–260

¹²¹ 🖱 [102], page 15

¹²² [138], pages 21–22, [102], page 2, ⟨§1. Axioms I–V⟩

¹²³ [[129], page 118, 🖱 [128]

¹²⁴ [[74], page 126, (Definitions), [[129], page 118

probability function given in Example 5.31 (page 66). This probability space (O_6, p) fails to be any of the 4 probability functions defined in this Remark. It fails to be a *measure-theoretic* or *traditional probability* function because

 $a \wedge b = 0$ but $p(a \vee b) = p(1) = 1 \neq \frac{1}{3} + \frac{1}{2} = p(a) + p(b)$. It fails to be a *generalized probability* function because

 $p(a \lor b) = p(1) = 1 \neq \frac{1}{3} + \frac{1}{2} - 0 = p(a) + p(b) - p(0) = p(a) + p(b) - p(a \land b)$. It fails to be an *quantum probability* function because

 $a \perp b = 0$ but $p(a \lor b) = p(1) = 1 \neq \frac{1}{3} + \frac{1}{2} = p(a) + p(b)$.

In each of these cases, the function p fails to be *additive*. The solution of Definition 5.24 (page 62) is simply to "switch off" *additivity* when the lattice is not *distributive*. This method is a little "crude", but at least it allows us to define probability on a very wide class of lattices, while retaining compatibility with the *Boolean* case (Proposition 5.26 page 64, Proposition 5.27 page 64, Proposition 5.28 page 65).

Proposition 5.26 ¹²⁵ Let (L, p) be a PROBABILITY SPACE (Definition 5.24 page 62).

 $0 \leq \mathsf{p}(x) \leq 1 \quad \forall x \in X$

Proof:

0 = p(0)	by previous result
$\leq p(x)$	because $0 \le x$ and <i>monotone</i> property (Definition 5.24 page 62)
$p(x) \le p(1)$	because $x \le 1$ and <i>monotone</i> property (Definition 5.24 page 62)
= 1	by property of p (Definition 5.24 page 62)

Proposition 5.27 ¹²⁶ Let (L, p) be a PROBABILITY SPACE (Definition 5.24 page 62).

Į	L is		$\int \mathbf{p}(\mathbf{x})$	=	$1 - p(\neg x)$	$\forall x \in X$	}
J	ORTHOCOMPLEMENTED) í	(P(0))		- P()	VACA	J

Proof:

 $1 - p(\neg x) = p(1) - p(\neg x)$ by Definition 5.24 page 62 = $p(x \lor \neg x) - p(\neg x)$ by *excluded middle* property of *ortho negation* (Definition 2.14 page 29) = $p(x) + p(\neg x) - p(\neg x)$ because $(x)(\neg x) = 0$ and *additive* property (Definition 5.24 page 62) = p(x)

 \square

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¹²⁵ (138], page 21, ((2-11))¹²⁶ (138], page 21, ((2-12))

[®]Proof:

(1) lemma: Proof that $p((\neg x) \land y) = p(y) - p(x \land y)$:

$p(y) - p(xy) = p(1 \land y) - p(xy)$	by definition of 1 and \land (Definition 1.28 page 9)
$= p[(x \lor \neg x)y] - p(xy)$	by excluded middle property of Boolean lattices
$= p(xy \lor \neg xy) - p(xy)$	by distributive property of Boolean lattices
$= p(xy) + p(\neg xy) - p(xy)$	because $(xy)(\neg xy) = 0$ and by <i>additive</i> property
$= p(\neg xy)$	

(2) Proof that
$$p(x \lor y) = p(x) + p(y) - p(x \land y)$$
:

 $p(x \lor y) = p(x \lor \neg xy)$ by property of Boolean lattices $= p(x) + p(\neg xy)$ because $(x)(\neg xy) = 0$ and by additive property $= p(x) + p(y) - p(x \land y)$ by item 1 (page 65)

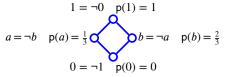
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 $1 = \neg 0 \mathbf{Op}(1) = 1$ $a = \neg a \mathbf{Op}(a) = \frac{1}{2}$

0 = -100(0) = 0

Example 5.29 The function \neg on the lattice *L* as illustrated to the right is a *Kleene negation* (Definition 2.14 page 29). Together with the probability function p, also illustrated to the right, the pair (*L*, p) is a *probability space* (Definition 5.24 page 62).

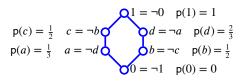
Example 5.30 The *lattice with negation* L (Definition 2.16 page 30) illustrated to the right is a *Boolean lattice*. Together with the probability function p, also illustrated to the right, the pair (L, p) is a *probability space* (Definition 5.24 page 62).



¹²⁷ \blacksquare [138], page 21, $\langle (2-13) \rangle$, \blacksquare [57], pages 22–23, $\langle (7.4), (7.6) \rangle$

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Example 5.31 The *lattice with negation L* (Definition 2.16 page 30) illustrated to the right is an *orthocom*-plemented O_6 lattice (Definition 1.73 page 20). Together $p(c) = \frac{1}{2}$ $c = \neg b$ with the probability function p, also illustrated to $p(a) = \frac{1}{3}$ $a = \neg d$ $b = \neg c$ $p(b) = \frac{1}{2}$ the right the pair (I, p) is a probability space (Definition 1.73 page 20). tion 2.16 page 30) illustrated to the right is an orthocomthe right, the pair (L, p) is a probability space (Definition 5.24 page 62).



5.6 Applications

This section discusses some possible applications of *primorial lattices*.

5.6.1 Logic analysis

Let L_2^N be a 2^N -valued *Boolean logic* (Definition 2.27 page 33). Let \mathbb{P} be the *primorial lattice generated by* L_2^N (Definition 5.17 page 58). The *sequence* of lattices $((L_2^N, L_2^{N-1}, \dots, L_2^2, L_2))$ in \mathbb{P} are *Boolean logics* with decreasing "resolution" (higher values of n in L_2^n correspond to greater resolution). Thus, we can reduce a very complex logic in L_2^N to a simpler lower resolution logic.

Moreover, the sequence of *ortho logics* (Definition 2.27 page 33) in \mathbb{P} $((L_2^N \otimes L_2^{N-1}, L_2^{N-1} \otimes L_2^{N-2}, ..., L_2^3 \otimes L_2^2, L_2))$

represents the Boolean logic L_2^N at N - 1 progressively lower "frequencies". Alternatively, we could say that the *Boolean logic* at resolution N is "decomposed" into (or *analyzed* by) N-1 ortho logics. Moreover, a proposition p in a higher resolution space can be projected into a lower resolution space (including the two-value classic logic space) by a *projection* operator (Section 5.4 page 58).

5.6.2 Fuzzy logic analysis

Fuzzy logics (Definition 2.27 page 33) can be constructed on Boolean and orthocomplemented lattices¹²⁸ such that together with the subset ordering relation \subseteq , form of a *primorial lattice* \mathbb{P} (Definition 5.1 page 50). A Boolean fuzzy logic L_2^N can then be rendered at N-1 different "resolutions" using the Boolean lattices of \mathbb{P} and analyzed at N - 1 "frequencies" using the orthocomplemented lattices of \mathbb{P} , as described in Section 5.6.1 (page 66).

¹²⁸ [77], (§2.2)

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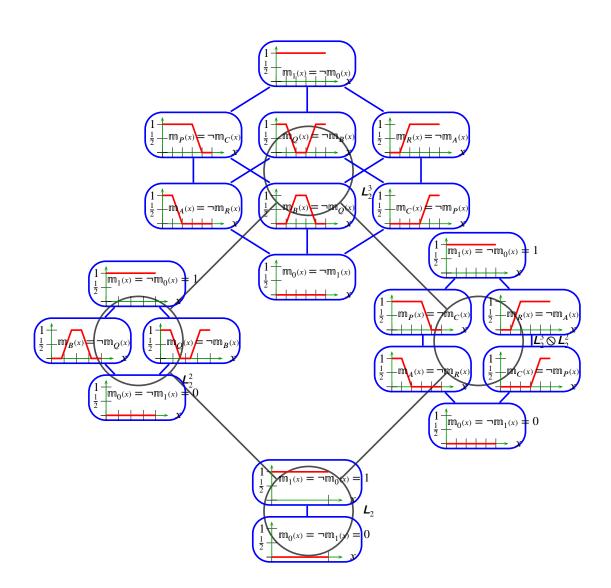


Figure 16: primorial lattice for *fuzzy subset logic* (Example 5.32 page 68)

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Example 5.32 Figure 16 (page 67) illustrates a *fuzzy subset logic*¹²⁹ on a primorial lattice. The lattice L_2^3 contains both *monotonic* and *non-monotonic membership functions*. These are separated into lower resolution spaces L_2^2 containing the *non-monotonic* membership functions (neglecting 1 and 0), $L_2^3 \otimes L_2^2$ containing the *monotonic* membership functions, and L_2 containing crisp set logic. A projection operator (Section 5.4 page 58) can be used to project a membership function onto any of these spaces as perhaps called for by a given application.

Probability analysis 5.6.3

A logic is a lattice with negation (Definition 2.16 page 30) and with an implication function defined on it. A probability is a lattice with negation and with a probability function (Definition 5.24 page 62) defined on it.

Let L_2^N be the 2^N -element Boolean lattice generated by an *N*-event Boolean probability *space* (Definition 5.24 page 62). Let \mathbb{P} be the *primorial lattice* (Definition 5.1 page 50) generated by \mathcal{L}_{2}^{N} . Then in \mathbb{P} , the probability space can be rendered at progressively lower resolutions using the Boolean lattices of \mathbb{P} , and can be analyzed at assorted "frequencies" using the orthocomplemented lattices of \mathbb{P} .

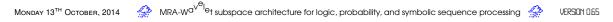
Example 5.33 A *primorial lattice* with a probability function is illustrated in Figure 17 (page 69).

5.6.4 Symbolic sequence analysis

Definitions. Finding some properties of a sequence \times that is constructed over a field \mathbb{F} may be referred to as sequence analysis or discrete-time signal analysis. If we somehow mathematically alter x with an operator A to produce a new sequence $y \triangleq Ax$, then this may be referred to as *sequence processing*, or more commonly as *discrete-time signal pro*cessing or digital signal processing (DSP).

Basis theory. Sequence analysis and sequence processing typically make use of basis theory. In basis theory in general (of which Fourier analysis and wavelet analysis are special cases), we represent some point x (x is a sequence) in a Banach space (a complete normed linear space) by a linear combination of a basis sequence (x_n) such that $x \stackrel{*}{=} \sum_{n \in \mathbb{Z}} a_n x_n$

¹²⁹ [77], (\$3.2)



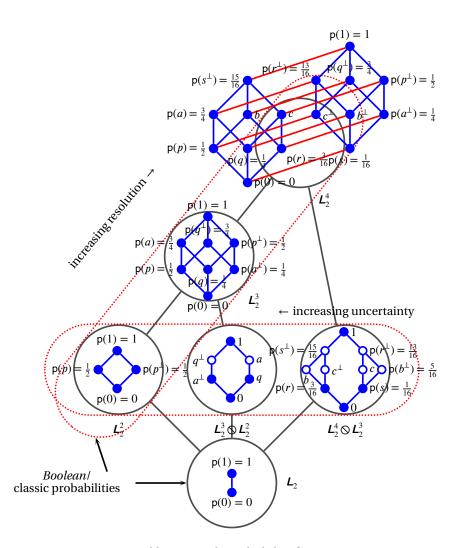


Figure 17: primorial lattice with probability function (Example 5.33 page 68)

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where \neq represents strong convergence with respect to the norm $\|\cdot\|$ of the Banach space. Each element a_n is a member of the field \mathbb{F} of the Banach space and the sequence (a_n) is often referred to as a "transform" (Fourier transform, discrete-time Fourier transform, wavelet transform, etc.)

In order to be able to successfully compute any transform (such as a Fourier transform or wavelet transform) in a Banach space or even a finite linear space, the sequence \times needs to be somehow related to the field \mathbb{F} over which the Banach space is constructed.

The problem. Let $\tilde{\mathbf{F}}$ be the discrete-time Fourier transform operator and \mathbf{W} be a discrete-time wavelet transform. Suppose we want to compute $\tilde{\mathbf{F}}_{\times}$ or \mathbf{W}_{\times} . This is a problem in *symbolic sequence analysis* and *symbolic signal processing* in general because of the following reasons:

- 1. The symbols in x have no field structure; so we can't even add them.
- 2. The symbols in x have no order structure; so if *A*, *B*, and *C* are symbols, we can't say, for example, *A* < *B* or *B* < *C*, etc.
- 3. The symbols in \times have no topology except for some arguably trivial topologies;¹³⁰ so we can't say, for example, that *A* is "closer" to *B* than it is to *C*, etc.

In fact, *symbol sequence analysis* does not just cause problems for Fourier or wavelet analysis only—it causes problems for basis theory in general because a basis is constructed in a Banach space, and symbolic sequences are in general not constructed in Banach spaces.

A kind of "hack" solution may be to map the symbols to points $(p_1, p_2, ..., p_N)$ in the *complex plane* \mathbb{C} . If these points are chosen such that they are distinct, not on either the real or imaginary axes, and $|p_1| = |p_2| = ... = |p_N|$, then that would seem to be a good start, because now the mapped symbols have a field structure, and they are arguably unordered (arguably we can't say any one of them is greater or less than any other, just as in the original symbol sequence).

But we still have the topology problem. If we map, say, 4 symbols to 4 points in \mathbb{C} as $p_1 = 1$, $p_2 = -1$, $p_3 = i$, and $p_4 = -i$, then " p_1 " is closer (with respect to the metric induced by the norm $|\cdot|$) to " p_3 " then it is to " p_2 ":

 $d(p_1, p_3) = |p_1 - p_3| = (p_1^2 - p_3^2)^{\frac{1}{2}} = (1^2 - i^2)^{\frac{1}{2}} = \sqrt{2} \leq 2 = (2^2 - 0^2)^{\frac{1}{2}} = d(p_1, p_2)$ This unwanted topological property is introduced by the mapping, will affect the transform, but yet is not a property of the original symbolic sequence.

¹³⁰ These topologies include the *indiscrete topology* { \emptyset , X} where $X \triangleq \{A, B, C\}$, *discrete topology* 2^X (references: \square [126], page 77, \square [107], page 107, (Example 3.J), \square [156], pages 42–43, (II.4), \square [44], page 18), and the topology induced by the *discrete metric* $d(x, y) \triangleq \{1 \text{ for } x \neq y, 0 \text{ for } x = y\}$ (references: \square [67], page 13, \square [31], page 24, \square [101], page 19, (Example 2.1)).

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"Frequency" properties may be useful in *symbolic sequence analysis* and *symbolic sequence processing*. But the point here is that any kind of basis theory technique (including Fourier or wavelet techniques) may result in a kind of imperfect "hack" solution.

Proposed solution. The solution proposed here is to perform symbolic sequence analysis using primorial lattices. Suppose we have a sequence \times over a set of N symbols (each element in the sequence can be any one of N different symbols). Let \mathbb{P} be the primorial lattice generated by L_2^N . The orthogonal N atoms of L_2^N represent the N symbols. The element $A \vee B$ in L_2^N , where A and B are 2 symbols, represents the event of a particular position in the sequence being A OR B (it is not possible for a particular position to be both A AND B).

Any symbol in L_2^N can be projected onto any other Boolean or orthocomplemented lattice in \mathbb{P} by use of a *lattice projection* (Section 5.4 page 58). The result of projecting an entire sequence onto a lattice in \mathbb{P} is another sequence (Definition 5.19 page 58). So after projection, a sequence on L_2^N results in N - 1 sequences of lower resolution and N - 1 sequences of assorted frequencies. This is similar in form to the *Fast Wavelet Transform*, as illustrated in Figure 10 (page 49).

5.6.5 Symbolic sequence processing (SSP)

Introduction. The previous section discusses symbolic sequence analysis—meaning we are not trying to change the properties of the sequence, we are only trying to understand its properties. This section discusses *symbolic sequence processing* (or *symbolic signal processing*)—meaning we *are* trying to change the properties of the sequence.

Digital signal processing (*DSP*) or *discrete-time signal processing* operates on a sequence constructed over a field \mathbb{F} , where \mathbb{F} is typically either \mathbb{R} or \mathbb{C} . Often by use of simple multiplication and addition operations on elements of the sequence, one can change the properties of the sequence. Often when the properties are related to Fourier analysis, the DSP operations are called "filtering".

The problem. Multiplication and addition operations commonly used in DSP require field properties. In symbolic sequence processing, we don't in general have a field.

Proposed solution. Sequence processing of, or "filtering" on, a symbolic sequence x can be performed by judicious selection and/or rejection of the various projections onto the logics in the primorial lattice \mathbb{P} .

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For example, if one wants x at a lower "resolution"s, then simply select the sequence from a projection onto the *Boolean logic* at resolution lower than *N*. If one wants to "filter out" the "high frequency" components of x, then simply discard the projections onto the higher frequency orthocomplemented lattices before synthesizing a new sequence from the "low frequency" component sequences.

Synthesis of two projection sequences y and z into a new sequence x' can be performed, for example, by pointwise join such that

$$y \bigoplus \mathbb{Z} \triangleq ((y_n)_{n \in \mathbb{Z}} \bigvee ((z_n)_{n \in \mathbb{Z}})$$
$$\triangleq ((y_n \lor z_n)_{n \in \mathbb{Z}})$$
$$\triangleq ((x_n)_{n \in \mathbb{Z}})$$
$$\triangleq \mathbb{X}$$

5.6.6 Genomic Signal Processing (GSP)

Genomic Signal Processing (GSP) is simply a special case of Symbolic Sequence Processing with N = 4. In GSP, the 4 symbols are commonly referred to as A, C, T, and G, each of which corresponds to a nucleobase (adenine, thymine, cytosine, and guanine, respectively).¹³¹ The sequence itself is called a *genome*. A typical genome sequence contains a large number of symbols (about 3 billion for humans, 29751 for the SARS virus).¹³²

Example 5.34 Traditionally in GSP, the symbols $(A \lor T)$ and $(C \lor G)$ are of special interest. Portions of a genome sequence high in $(A \lor T)$ content separate at lower temperatures than do those with high $(C \lor G)$ content.¹³³ Therefore, one could construct a primorial lattice induced by L_2^4 that allows for convenient analysis of $A \lor T$ and/or $C \lor G$ in some lower resolution space. An example is illustrated in Figure 18 (page 73).

Example 5.35 In some cases, genomic sequences with more than 4 symbols (N > 4) have been studied.¹³⁴ Figure 19 (page 74) illustrates a primorial lattice with an extra symbol X

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¹³¹ [121], (Mendel (1853): gene coding uses discrete symbols), [165], page 737, (Watson and Crick (1953): gene coding symbols are adenine, thymine, cytosine, and guanine), [164], page 965, (142], page 52

^{132 [1], &}lt;http://www.ncbi.nlm.nih.gov/genome/guide/human/>, <Homo sapiens, NC_000001-NC_000022 (22 chromosome pairs), NC_000023 (X chromosome), NC_000024 (Y chromosome), NC_012920 (mitochondria)>, [1], <http://www.ncbi.nlm.nih.gov/nuccore/ 30271926>, <SARS coronavirus, NC_004718.3> [150], <homo sapien chromosome 1>, [149], <SARS coronavirus>

¹³³ **●** [32], page 13, 〈Remark 1.2〉 ¹³⁴ **■** [30], **■** [53]

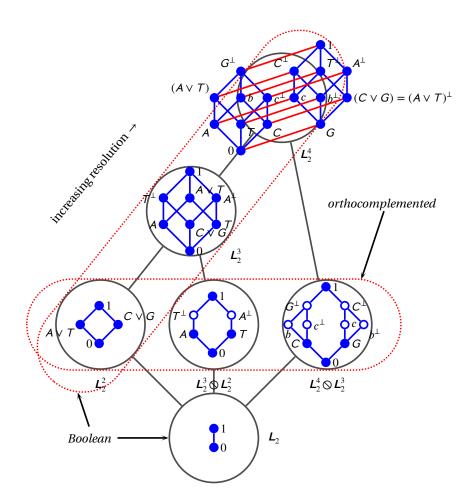


Figure 18: primorial lattice for genomic signal processing (GSP) with $A \lor T$ and $C \lor G$ analysis features (Example 5.34 page 72)

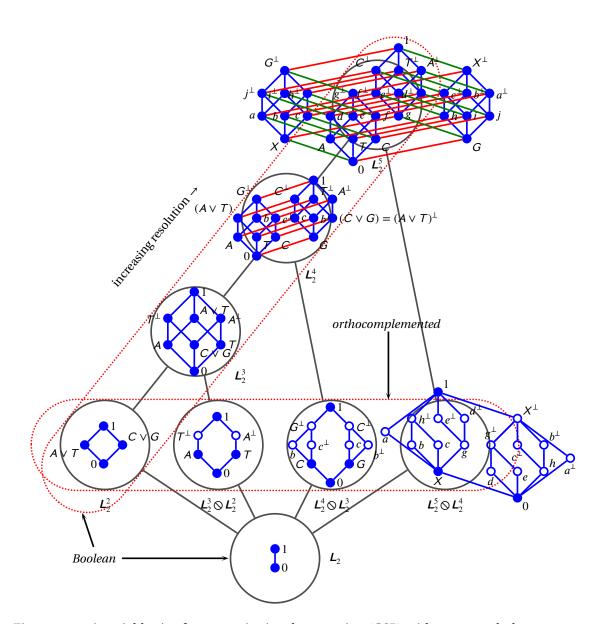


Figure 19: primorial lattice for genomic signal processing (GSP) with extra symbol X (Example 5.35 page 72)

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in the higher resolution L_2^5 Boolean lattice, but with only the symbols *A*, *C*, *G*, and *T* in the lower resolution L_2^4 Boolean lattice. The symbol *X* can be projected onto any of the lower resolution spaces using a *projection operator* (Section 5.4 page 58).

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