THE EUCLID–MULLIN GRAPH

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ABSTRACT. We introduce the *Euclid–Mullin graph*, which encodes all instances of Euclid's proof of the infinitude of primes. We investigate structural properties of the graph both theoretically and numerically; in particular, we prove that it is not a tree.

1. INTRODUCTION

The Euclid-Mullin sequence begins [1,] 2, 3, 7, 43, 13, 53, 5, 6221671, where each term is the least prime factor of 1 plus the product of all the preceding terms. As such it can be viewed as a computational form of Euclid's proof that the number of primes is infinite. A companion sequence, sometimes referred to as the second Euclid-Mullin sequence takes the largest prime factor at each step. These sequences are A000945 and A000946 in the OEIS [14]. Both sequences were introduced by Mullin [11], who asked whether every prime occurs in these sequences. Mullin's question has been answered negatively for the second sequence and in fact the second sequence omits infinitely many primes [1, 12]. The question for the first sequence remains open.

Here a generalization is considered, where rather than choosing the least or largest prime factor at each stage, all prime factors are considered. Since there are now, in general, multiple choices for the next element, the result is not a single sequence, but a (directed) graph where each path from the root to a node corresponds to a particular sequence of primes. Questions asked about Mullin's sequence can now also be asked about the graph. In particular, does the graph contain every prime? If it were ever shown that Mullin's original sequence contains every prime, then the graph would also include every prime.

The graph admits other structural questions. While the graph is obviously infinite it would be interesting to know how the number of nodes grows at each level (or, indeed, to determine if it does grow!). As a first step in this direction, this paper establishes that the graph is not a tree.

The directed graph $G_n \subseteq (\mathbb{Z}, \mathbb{Z} \times \mathbb{Z})$ consists of a set of integer labelled nodes and edges defined by ordered pairs of nodes. G_n can be defined recursively by: n is a node in G_n . If m is a node in G_n , then so are all of mp_i where $m + 1 = \prod_{i=1}^k p_i^{e_i}$, $e_i > 0$, is the unique factorization of m + 1. Further, G_n has directed edges (m, mp_i) . It is sometimes convenient to think of the edge (m, mp_i) as being labelled p_i . We say, n is the root of the graph and has level 0. Any node adjacent to n is said to be level 1. In general, any node reachable by a directed path of r edges is said to be level r. In fact, a path of length r represents a product of r distinct primes. We call G_1 the Euclid–Mullin graph; its first few levels are shown in Figure 1.

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FIGURE 1. G_1 .

Theorem 1.1. The Euclid–Mullin graph G_1 is not a tree. In particular, each of the following nodes is connected to 1 by two distinct paths:

 $2 \cdot 3 \cdot 7 \cdot 43 \cdot 139 \cdot 50207 \cdot 1607 \cdot 38891 \cdot 71609249149971437 \cdot 104851$

 $\cdot 5914302068415095755097398828253214149923$

 $\cdot 103 \cdot 1750880132687750604376675981842334069$

 $+103451 \cdot 193 \cdot 22133 \cdot 5587528960270206397663051$

 $\cdot \ 73 \cdot 5 \cdot 13 \cdot 593$

and

 $2 \cdot 3 \cdot 7 \cdot 43 \cdot 139 \cdot 50207 \cdot 23 \cdot 217733 \cdot 4024572619121$

- $\cdot \ 539402497343 \cdot 72208156847017648587223 \cdot 79$
- $\cdot \ 7269452239696911635939429787229069136737446558564286318153183$
- $\cdot \ 8689 \cdot 107 \cdot 2895777621755988962510175673615781760909999040975810951$
- $\cdot 531543631 \cdot 73 \cdot 5 \cdot 13 \cdot 593.$

In each case, the order of the prime factors indicates one path, and the other path is obtained by swapping 73 and 593.

Note that the numbers given in the theorem both have level 21. Based on some probabilistic considerations presented in §3, we suspect that any node of lower level is connected to 1 by a unique path, but answering this definitively is likely to remain infeasible for the foreseeable future. Acknowledgements. The numerical computations for this paper were primarily carried out using the resources of Real Time Genomics at Hamilton, New Zealand, with additional factorizations performed on BlueCrystal phases 2 and 3 at the University of Bristol, UK, Michael Rubinstein's riemann cluster at the University of Waterloo, Canada, and by NFS@home. We thank these organizations for their support, without which this project would not have been possible. We also thank the users of mersenneforum.org, especially Oscar Östlin, who helped with many hard factorizations.

2. Multiple k-tuples of edges

Given a positive integer n, a path in G_n between n and $m = p_1 \cdots p_k n$ can be identified with the k-tuple of edge primes (p_1, \ldots, p_k) . In this section, we formalize this notion and formulate conditions under which nodes may be connected by more than one path. We also establish several theoretical results, including the following:

- For k = 3, we obtain a complete classification of the triples (p_1, p_2, p_3) that form one side of a loop in some G_n , given as the prime values of certain polynomials; see Theorem 2.5.
- We prove that there is a $k \leq 13$ such that, for any $q \in \mathbb{Z}_{>0}$, there are infinitely many k-tuples (p_1, \ldots, p_k) that form one side of a loop in some G_n and satisfy $(p_1 \cdots p_k, q) = 1$. Moreover, any given prime occurs as an edge of a loop of height at most 13 in some G_n ; see Theorem 2.15.

First, let \mathcal{P}_k denote the set of k-tuples (p_1, \ldots, p_k) , where each p_i is a prime number and $p_i \neq p_j$ for $i \neq j$. The symmetric group S_k acts on \mathcal{P}_k by permuting the indices; precisely, for $\pi \in S_k$ we write $\pi.(p_1, \ldots, p_k) = (q_1, \ldots, q_k)$, where $p_i = q_{\pi(i)}$ for $i = 1, \ldots, k$.

Definition 2.1. Let $P = (p_1, ..., p_k), Q = (q_1, ..., q_k) \in \mathcal{P}_k$.

(1) We say that P and Q are equivalent, and write $P \sim Q$, if there exists $\pi \in S_k$ such that $Q = \pi P$ and

$$p_1 \cdots p_{i-1} \equiv q_1 \cdots q_{\pi(i)-1} \pmod{p_i} \quad \text{for } i = 1, \dots, k.$$

- (2) The multiplicity of P, denoted m(P), is the number of $\pi \in S_k$ such that $P \sim \pi P$.
- (3) We say that P is multiple if m(P) > 1.
- (4) We call $p_1 \cdots p_k$ the modulus of P, and denote it by |P|.

It is straightforward to verify that ~ defines an equivalence relation on \mathcal{P}_k . Its relevance to the graphs G_n is described by the following key lemma.

Lemma 2.2. For $P = (p_1, \ldots, p_k) \in \mathcal{P}_k$, let N(P) denote the set of positive integers n such that n and |P|n are connected in G_n via edges p_1, \ldots, p_k , i.e.

$$p_1 | n+1, p_2 | p_1n+1, \dots, p_k | p_1 \cdots p_{k-1}n+1.$$

Then:

(1) N(P) is an arithmetic progression modulo |P|, i.e.

$$N(P) = \{ n \in \mathbb{Z}_{>0} : n \equiv a \pmod{|P|} \}$$

for some $a = a(P) \in \mathbb{Z}$ relatively prime to |P|.

(2) $Q \in \mathcal{P}_k$ is equivalent to P if and only if N(Q) = N(P).

(3) For any $n \in N(P)$, the paths in G_n between n and |P|n are in one-to-one correspondence with the equivalence class of P. In particular, the number of such paths is the multiplicity m(P).

Proof.

(2.

(1) The conditions on n can be rephrased as the system of congruences

$$n \equiv -1 \pmod{p_1}$$

$$n \equiv -p_1^{-1} \pmod{p_2}$$

$$\vdots$$

$$n \equiv -(p_1 \cdots p_{k-1})^{-1} \pmod{p_k},$$

and the solutions form an arithmetic progression, by the Chinese remainder theorem. Since none of the numbers on the right-hand side can be congruent to 0, the elements of N(P) lie in an invertible residue class modulo |P|.

(2) Suppose that $P = (p_1, \ldots, p_k)$ and $Q = (q_1, \ldots, q_k)$ are equivalent. Then there is a permutation $\pi \in S_k$ such that $Q = \pi P$. Choose $n \in N(P)$, $j \in \{1, \ldots, k\}$, and set $i = \pi^{-1}(j)$, so that $p_i = q_j$. Then,

1)
$$0 \equiv p_1 \cdots p_{i-1}n + 1 \equiv q_1 \cdots q_{j-1}n + 1 \pmod{p_i = q_j}.$$

Since this holds for every j, n is contained in N(Q). Since n was an arbitrary element of N(P), this shows that $N(P) \subseteq N(Q)$. Applying the argument again with the roles of P and Q reversed, we also get $N(Q) \subseteq N(P)$, and hence N(P) = N(Q).

Conversely, suppose that N(P) = N(Q). By part (1), we must have |P| = |Q|, and hence there is a permutation $\pi \in S_k$ such that $Q = \pi P$. Let $n \in N(P) = N(Q)$, $i \in \{1, \ldots, k\}$, and set $j = \pi(i)$, so that $p_i = q_j$. Then again we obtain (2.1), and since n is invertible modulo |P| = |Q|, it follows that

$$p_1 \cdots p_{i-1} \equiv q_1 \cdots q_{j-1} \pmod{p_i = q_j}.$$

Since this holds for all i, P and Q are equivalent.

(3) Let $P = (p_1, \ldots, p_k)$, $n \in N(P)$, and m = |P|n. Suppose that there is a path in G_n between n and m via edges q_1, \ldots, q_l . Then we have $m = q_1 \ldots q_l n$, so that $p_1 \cdots p_k = q_1 \ldots q_l$. By unique factorization, we have l = k and $Q = (q_1, \ldots, q_k) \in \mathcal{P}_k$. By part (1), N(P) and N(Q) are arithmetic progressions with the same modulus. Since they also have a common element $n \in N(P) \cap N(Q)$, they must be equal. By part (2), P and Q are therefore equivalent. Conversely, if P and Q are equivalent then N(P) = N(Q), so there is a path in G_n between n and |Q|n = m.

Lemma 2.3. There are no multiple k-tuples for k < 3.

Proof. This is obvious for k = 1. For k = 2, the only non-trivial possibility is that (p_1, p_2) is equivalent to $(q_1, q_2) = (p_2, p_1)$. Then by Definition 2.1 we have

)

$$1 \equiv q_1 = p_2 \pmod{p_1}$$
$$p_1 \equiv 1 \pmod{p_2},$$

so that $p_1 < p_2 < p_1$, which is impossible.

2.1. Multiple triples.

Proposition 2.4. Let $P = (p_1, p_2, p_3) \in \mathcal{P}_3$. Then m(P) > 1 if and only if

(2.2)
$$p_2(p_1 + p_3) \equiv 1 \pmod{p_1 p_3}$$
 and $p_1 \equiv p_3 \pmod{p_2}$

In this case, m(P) = 2 and the equivalence class of P is $\{(p_1, p_2, p_3), (p_3, p_2, p_1)\}$.

Proof. Suppose that $P = (p_1, p_2, p_3)$ is equivalent to $Q = (q_1, q_2, q_3) = \pi P$ for some non-trivial $\pi \in S_3$. Since there are no multiple pairs, we must have $p_1 \neq q_1$ and $p_3 \neq q_3$, so $\pi \in \{(13), (123), (132)\}$.

First suppose that π is a 3-cycle. By reversing the roles of P and Q if necessary, we may assume that $\pi = (123)$. Then $(p_1, p_2, p_3) = (q_2, q_3, q_1)$, so by Definition 2.1 we have

$$1 \equiv q_1 = p_3 \pmod{p_1}$$
$$p_1 \equiv q_1 q_2 = p_1 p_3 \pmod{p_2} \Longrightarrow 1 \equiv p_3 \pmod{p_2}$$
$$p_1 p_2 \equiv 1 \pmod{p_3}.$$

Thus, $p_3 \equiv 1 \pmod{p_1 p_2}$ and $p_1 p_2 \equiv 1 \pmod{p_3}$, which is impossible.

The only remaining choice is $\pi = (13)$. Then $(p_1, p_2, p_3) = (q_3, q_2, q_1)$, and we have

$$1 \equiv q_1 q_2 = p_2 p_3 \pmod{p_1}$$
$$p_1 \equiv q_1 = p_3 \pmod{p_2}$$
$$p_1 p_2 \equiv 1 \pmod{p_3},$$

which is equivalent to the system (2.2). Conversely, the steps above are clearly reversible, so that any (p_1, p_2, p_3) satisfying (2.2) is equivalent to (p_3, p_2, p_1) .

Finally, since (13) is the only non-trivial permutation that can relate equivalent triples, any multiple $P \in \mathcal{P}_3$ must have m(P) = 2 and equivalence class $\{P, (13), P\}$.

Table 2.1 shows the first few solutions to (2.2) with $p_1 < p_3$, ordered by modulus.

Р	P	a(P)			
(2, 3, 5)	30	19			
(3,2,5)	30	29			
(7, 5, 17)	595	237			
(211, 197, 2969)	123412423	114015537			
(601, 577, 14449)	5010580873	4793484647			
(8191, 8101, 737281)	48922495303771	48372940054709			
(22921, 21169, 276949)	134379711825901	123251758931063			
TABLE 2.1. Multiple triples $P = (p_1, p_2, p_3)$ with $p_1 < p_3$					

2.1.1. Integer triples. Let us temporarily drop the restriction that p_1 , p_2 and p_3 be prime, and consider all solutions to (2.2) in integers. Then it turns out that we can give a complete classification. In order to state it, we recall that the *Fibonacci polynomials* $F_n(x)$ are defined by the recurrence

$$F_0(x) = 0$$
, $F_1(x) = 1$, and $F_n(x) = xF_{n-1}(x) + F_{n-2}(x)$ for $n \ge 2$,
₅

generalizing the usual Fibonacci numbers $F_n = F_n(1)$. By convention we extend the definition to negative indices by defining $F_{-n}(x) = F_n(-x) = (-1)^{n-1}F_n(x)$.

Theorem 2.5. Let $(p_1, p_2, p_3) \in \mathbb{Z}^3$. Then (p_1, p_2, p_3) satisfies (2.2) if and only if one of the following holds for some $n, x \in \mathbb{Z}$ and $\delta \in \{\pm 1\}$:

(2.3)
$$(p_1, p_2, p_3) = \begin{cases} \delta(F_{n-1}(x) + F_n(x), F_{-n}(x), F_n(x) + F_{n+1}(x)) \\ \delta(F_n(x), F_{-n}(x) + F_{-(n+1)}(x), F_{n+1}(x)) \\ \delta(1, x, 1) \\ \delta(x, 1, 1 - x). \end{cases}$$

Proof. The Fibonacci polynomials are given by the following explicit formula:

(2.4)
$$F_n(x) = \frac{\left(\frac{x+\sqrt{x^2+4}}{2}\right)^n - \left(\frac{x-\sqrt{x^2+4}}{2}\right)^n}{\sqrt{x^2+4}}.$$

Using this one can verify that

$$F_{n+1}(x)F_{n-1}(x) = F_n(x)^2 + (-1)^n,$$

and combined with the recurrence identity $F_{n+1}(x) - F_{n-1}(x) = xF_n(x)$ we see that if

$$(p_1, p_2, p_3) = \delta(F_{n-1}(x) + F_n(x), F_{-n}(x), F_n(x) + F_{n+1}(x))$$

then

$$p_2(p_1 + p_3) = 1 + (-1)^{n-1} p_1 p_3$$
 and $p_3 - p_1 = (-1)^{n-1} x p_2$.

Similarly, we obtain the identity

$$F_{n+1}(x)^2 - F_n(x)^2 = xF_n(x)F_{n+1}(x) + (-1)^n,$$

from which it follows that if

$$(p_1, p_2, p_3) = \delta(F_n(x), F_{-n}(x) + F_{-(n+1)}(x), F_{n+1}(x))$$

then

$$p_2(p_1 + p_3) = 1 + (-1)^n x p_1 p_3$$
 and $p_3 - p_1 = (-1)^n p_2$.

Thus, in either case, (p_1, p_2, p_3) is a solution to (2.2). The final two solutions are straightforward to verify directly.

Now suppose that $(p_1, p_2, p_3) \in \mathbb{Z}^3$ satisfies (2.2), and write

(2.5)
$$p_3 - p_1 = qp_2, \quad p_2(p_1 + p_3) = 1 + rp_1p_3$$

for some $q, r \in \mathbb{Z}$. If $p_1 p_2 p_3 qr = 0$ then it is easy to see that either $p_1 p_3 = 1$ or $p_2(p_1 + p_3) = 1$, and all such solutions are described by the third and fourth lines of (2.3). Otherwise q and r are uniquely determined and non-zero.

Next, set

(2.6)
$$s = r(p_1 + p_3) - 2p_2$$
 and $d = (qr)^2 + 4$.

Then d is not a square, and a computation shows that s and p_2 are related by the Pell-type equation

(2.7)
$$s^2 - dp_2^2 = -4r.$$

In other words, $\frac{s+p_2\sqrt{d}}{2}$ is an element of norm -r in the quadratic order $\mathcal{O} = \mathbb{Z}\left[\frac{d+\sqrt{d}}{2}\right]$. (Note that \mathcal{O} need not be the maximal order in $\mathbb{Q}(\sqrt{d})$.)

If $r = \pm 1$ then (2.7) is just the unit equation for \mathcal{O} . It is easy to see that $\frac{q+\sqrt{d}}{2}$ is a fundamental unit (of norm -1), so the general solution of (2.7) in this case is given by

$$\frac{s+p_2\sqrt{d}}{2} = \delta\left(\frac{q+\sqrt{d}}{2}\right)^r$$

for $\delta \in \{\pm 1\}$ and $n \in \mathbb{Z}$ with $(-1)^{n-1} = r$. Thus,

$$p_2 = \delta \frac{\left(\frac{q+\sqrt{d}}{2}\right)^n - \left(\frac{q-\sqrt{d}}{2}\right)^n}{\sqrt{d}} = \delta F_n(q) \quad \text{and} \quad s = \delta \left[\left(\frac{q+\sqrt{d}}{2}\right)^n + \left(\frac{q-\sqrt{d}}{2}\right)^n\right] = \delta L_n(q),$$

where $L_n(x) = F_{n+1}(x) + F_{n-1}(x)$ is the Lucas polynomial. Recalling the definition of s, we have

$$p_1 + p_3 = \delta'(L_n(q) + 2F_n(q))$$

where $\delta' = (-1)^{n-1}\delta$. Together with $p_3 - p_1 = qp_2 = (-1)^{n-1}\delta' qF_n(q)$, this yields $p_1 = \delta' \frac{L_n(q) + 2F_n(q) - (-1)^{n-1}qF_n(q)}{2}, \quad p_3 = \delta' \frac{L_n(q) + 2F_n(q) + (-1)^{n-1}qF_n(q)}{2}.$

From the identities

 $L_n(x) - xF_n(x) = 2F_{n-1}(x), \quad L_n(x) + xF_n(x) = 2F_{n+1}(x) \text{ and } (-1)^{n-1}F_n(x) = F_{-n}(x),$ we get

$$(p_1, p_2, p_3) = \delta'(F_n(q) + F_{n-1}(q), F_{-n}(q), F_n(q) + F_{n+1}(q))$$

if n is odd, and

$$(p_1, p_2, p_3) = \delta'(F_n(q) + F_{n+1}(q), F_{-n}(q), F_n(q) + F_{n-1}(q)) = \delta'(F_{-n}(-q) + F_{-n-1}(-q), F_n(-q), F_{-n}(-q) + F_{-n+1}(-q))$$

if n is even. In either case, this is in the form of the first line of (2.3).

Next suppose that $q = \pm 1$. Since $\frac{r-2+\sqrt{d}}{2} \in \mathcal{O}$ has norm -r, we get a family of solutions defined by

(2.8)
$$\frac{s+p_2\sqrt{d}}{2} = \delta \frac{r-2+\sqrt{d}}{2} \left(\frac{r+\sqrt{d}}{2}\right)^n$$

for $\delta \in \{\pm 1\}$ and $n \in 2\mathbb{Z}$. Thus,

$$p_{2} = \delta \frac{\frac{r-2+\sqrt{d}}{2} \left(\frac{r+\sqrt{d}}{2}\right)^{n} - \frac{r-2-\sqrt{d}}{2} \left(\frac{r-\sqrt{d}}{2}\right)^{n}}{\sqrt{d}} = \delta \frac{(r-2)F_{n}(r) + L_{n}(r)}{2}$$
$$= \delta (F_{n+1}(r) - F_{n}(r)) = \delta (F_{-(n+1)}(r) + F_{-n}(r)),$$
$$s = \delta \left[\frac{r-2+\sqrt{d}}{2} \left(\frac{r+\sqrt{d}}{2}\right)^{n} + \frac{r-2-\sqrt{d}}{2} \left(\frac{r-\sqrt{d}}{2}\right)^{n}\right]$$
$$= \delta \frac{(r-2)L_{n}(r) + dF_{n}(r)}{2}$$

and

$$p_1 + p_3 = \frac{s + 2p_2}{r} = \delta \frac{(r+2)F_n(r) + L_n(r)}{\frac{2}{7}} = \delta(F_{n+1}(r) + F_n(r)).$$

Combining this with $p_3 - p_1 = qp_2$, we obtain

$$(p_1, p_2, p_3) = \delta(F_n(r), F_{-n}(r) + F_{-(n+1)}(r), F_{n+1}(r))$$

if q = 1 and

$$(p_1, p_2, p_3) = \delta(F_{n+1}(r), F_{-n}(r) + F_{-(n+1)}(r), F_n(r)) = \delta(F_{-n-1}(-r), F_{n+1}(-r) + F_n(-r), F_{-n}(-r))$$

if q = -1. In either case, this is in the form of the second line of (2.3).

In the case just presented, it is not obvious that we obtain all solutions in this manner, but we now proceed to show that this is indeed the case. Let us assume first that $4 \nmid r$, and let $\mathfrak{a} = \left(\frac{s+p_2\sqrt{d}}{2}\right) \mathcal{O}$ be the \mathcal{O} -ideal associated to the pair (s, p_2) . Then $s+p_2\sqrt{d} \equiv 0 \pmod{\mathfrak{a}}$, and by (2.6) we have $s+2p_2 \equiv 0 \pmod{\mathfrak{a}}$. It follows from (2.5) that p_2 is invertible modulo r, so we conclude that $\sqrt{d} \equiv 2 \pmod{\mathfrak{a}}$.

Now if p is an odd prime factor of r, then from (2.7) we see that $\left(\frac{d}{p}\right) = 1$. Thus, $p\mathcal{O}$ splits as a product of two prime ideals that are distinguished by the reduction of \sqrt{d} , i.e. there is a unique prime ideal $\mathfrak{p} \subseteq \mathcal{O}$ with norm p such that $\sqrt{d} \equiv 2 \pmod{\mathfrak{p}}$.

If r is even then $r \equiv 2 \pmod{4}$, and from (2.6) we see that $4 \mid s$. If q is also even then $d \equiv 4 \pmod{16}$, so that $s^2 - dp_2^2 \equiv 12 \pmod{16}$, in contradiction to (2.7). Hence, q must be odd and $d \equiv 8 \pmod{16}$. It follows that the conductor of \mathcal{O} is odd and 2 is ramified in $\mathbb{Q}(\sqrt{d})$, so there is anyway a unique prime ideal $\mathfrak{p} \subseteq \mathcal{O}$ lying above 2.

In summary, provided that $4 \nmid r$, we have shown that r is co-prime to the conductor of \mathcal{O} and that the prime factors of \mathfrak{a} are uniquely determined. Therefore, any solution of (2.6) and (2.7) generates the same ideal as the solution noted above, viz. $\left(\frac{r-2+\sqrt{d}}{2}\right)\mathcal{O}$. Hence, (2.8) describes all solutions.

Next, to handle the case when $4 \mid r$ we need to modify the above argument since the conductor of \mathcal{O} is even. In this case we set

$$d' = d/4, \quad r' = r/4, \quad s' = s/2 \text{ and } \mathcal{O}' = \mathbb{Z}\left[\frac{1+\sqrt{d'}}{2}\right],$$

and we work over \mathcal{O}' instead of \mathcal{O} . Then

$$d' = 4(qr')^2 + 1$$
 and $(s')^2 - d'p_2^2 = -4r'$,

and if $\mathfrak{a}' = \left(\frac{s'+p_2\sqrt{d'}}{2}\right)\mathcal{O}'$ then $N(\mathfrak{a}') = |r'|$ and $\frac{1+\sqrt{d'}}{2} \equiv 1 \pmod{\mathfrak{a}'}$. Note that if r' is even then $d' \equiv 1 \pmod{8}$, so that $\left(\frac{d'}{2}\right) = 1$. Hence, proceeding as above, for each prime $p \mid r'$, we find that there is a unique prime $\mathfrak{p} \subseteq \mathcal{O}'$ such that $N(\mathfrak{p}) = p$ and $\frac{1+\sqrt{d'}}{2} \equiv 1 \pmod{\mathfrak{p}}$. Thus, the ideal \mathfrak{a}' is again uniquely determined, so (2.8) describes all solutions.

It remains only to show that (2.7) admits no solutions if $\min(|q|, |r|) > 1$. For this we appeal to the reduction theory of primitive ideals in quadratic orders; see, for instance, [3, Chapters 8 and 9] for terminology and fundamental results. When $4 \nmid r$, we apply the reduction algorithm to see that the cycle of \mathcal{O} has length 1; in other words, \mathcal{O} is the only reduced principal \mathcal{O} -ideal. On the other hand, by [3, Prop. 9.1.8], any primitive \mathcal{O} -ideal of norm less than $\sqrt{d}/2$ is reduced. Note that if $|q| \geq 2$ then

$$|r| \le \frac{|qr|}{2} < \frac{\sqrt{(qr)^2 + 4}}{8} = \frac{\sqrt{d}}{2}.$$

Together these imply that if $|q|, |r| \neq 1$ then there is no primitive, principal \mathcal{O} -ideal of norm |r|, so (2.7) is not solvable.

For r divisible by 4, we similarly apply the reduction algorithm to \mathcal{O}' and find that its cycle consists of \mathcal{O}' together with the ideals $\left(\frac{qr'-2\pm\sqrt{d'}}{2}\right)\mathcal{O}'$ of norm |qr'|. In this case we have $|r'| < \frac{1}{2}\sqrt{d'}$ for every value of q, so there are no primitive, principal \mathcal{O}' -ideals of norm |r'| if $|q|, |r'| \neq 1$.

2.1.2. Prime triples. We now return to the prime case. Clearly the third and fourth lines of (2.3) never yield primes, and since the sum of the entries of the second line is even, the only (positive) prime solutions that it yields are permutations of (2,3,5). As for the first line, note that $F_n(x)$ is irreducible only if |n| is prime [10]. If we take $p_1 < p_3$, then we may assume that n is an odd prime, x is positive, and $\delta = 1$.

In particular, with n = 3 we get the solutions

$$(p_1, p_2, p_3) = (x^2 + x + 1, x^2 + 1, x^3 + x^2 + 2x + 1).$$

By standard conjectures (Schinzel's Hypothesis), we expect that these polynomials are simultaneously prime for infinitely many values of x > 0, and that motivates the following conjecture.

Conjecture 2.6. There are infinitely many $P \in \mathcal{P}_3$ with m(P) > 1.

In fact, it is natural to expect triples of primes to occur with probability proportional to $(\log x)^{-3}$, so there should be a constant c > 0 such that

$$\#\{P \in \mathcal{P}_3 : m(P) > 1 \text{ and } |P| < X\} = (c + o(1)) \frac{X^{1/7}}{(\log X)^3} \text{ as } X \to \infty.$$

Such a statement seems far from what can be proven with present technology, but we are able to obtain somewhat weaker results in Section 2.3 below.

2.2. Multiple quadruples. In this section we compute the systems of congruences giving rise to multiple quadruples of edge primes, analogous to Proposition 2.4 in the case of triples. Note first that if $(p_1, p_2, p_3) \sim (p_3, p_2, p_1) \in \mathcal{P}_3$ is a multiple triple, then clearly $(p_0, p_1, p_2, p_3) \sim (p_0, p_3, p_2, p_1)$ and $(p_1, p_2, p_3, p_4) \sim (p_3, p_2, p_1, p_4)$ are multiple quadruples for any suitable choice of p_0 or p_4 . More interesting are the solutions giving rise to loops of height 4 in the graph. More generally, we will be interested in pairs $P = (p_1, \ldots, p_k), Q =$ $(q_1, \ldots, q_k) \in \mathcal{P}_k$ defining paths in G_n that meet only at n and |P|n = |Q|n, so that they form a loop of height k; that is the content of the following definition.

Definition 2.7. Let $P = (p_1, \ldots, p_k), Q = (q_1, \ldots, q_k) \in \mathcal{P}_k$. We say that the pair $(P, Q) \in \mathcal{P}_k^2$ is irreducible if $P \neq Q, P \sim Q$ and

 $p_1 \cdots p_i \neq q_1 \cdots q_i \quad for \ 0 < i < k.$

Remark 2.8. Note that (P,Q) is irreducible if and only if (Q,P) is irreducible, so we may regard the pair as unordered.

Next, we observe that any equivalence $P \sim Q$ gives rise to another equivalence, as follows.

Lemma 2.9. Let $P \in \mathcal{P}_k$, and suppose that P is equivalent to $Q = \pi . P$ for some $\pi \in S_k$. Let $\sigma = \begin{pmatrix} 1 & k \end{pmatrix} \begin{pmatrix} 2 & k-1 \end{pmatrix} \cdots \begin{pmatrix} \lfloor \frac{k}{2} \rfloor & k+1-\lfloor \frac{k}{2} \rfloor \end{pmatrix} \in S_k$ be the permutation that reverses the order of indices, and put $\widetilde{P} = \sigma . P$, $\widetilde{Q} = \sigma . Q$. Then:

- (1) \widetilde{P} is equivalent to $\widetilde{Q} = \sigma \pi \sigma . \widetilde{P}$;
- (2) P, Q, \widetilde{P} and \widetilde{Q} all have the same multiplicity.

Proof. Suppose that $P = (p_1, \ldots, p_k)$ and $Q = (q_1, \ldots, q_k)$. Then

 $p_1 \cdots p_{i-1} \equiv q_1 \cdots q_{j-1} \pmod{p_i}$

whenever $p_i = q_j$. Note that we also have |P| = |Q|, and cancelling the common factor of $p_i = q_j$ yields

$$(p_1 \cdots p_{i-1})(p_{i+1} \cdots p_k) = (q_1 \cdots q_{j-1})(q_{j+1} \cdots q_k).$$

Dividing this equality by the above congruence, we obtain

$$p_{i+1}\cdots p_k \equiv q_{j+1}\cdots q_k \pmod{p_i}.$$

Thus, (p_k, \ldots, p_1) is equivalent to (q_k, \ldots, q_1) , as desired.

For the second assertion, P and Q clearly have the same multiplicity since they are equivalent, and likewise for \tilde{P} and \tilde{Q} , so it is enough to show that $m(P) = m(\tilde{P})$. But by the first assertion, P is equivalent to Q if and only if $\tilde{P} = \sigma P$ is equivalent to $\tilde{Q} = \sigma Q$, so σ defines a bijection between the equivalence classes of P and \tilde{P} .

Proposition 2.10. Let $(p_1, p_2, p_3, p_4) \in \mathcal{P}_4$. If the conditions listed in the middle column of the following table are satisfied in any one case, then each of the corresponding quadruples in the right column has multiplicity 2, with equivalence classes as indicated. Conversely, every multiple quadruple has multiplicity 2, and every irreducible pair of multiple quadruples occurs in the table for a unique choice of (p_1, p_2, p_3, p_4) .

	$p_4 \equiv 1 \pmod{p_1}$	$\{(p_1, p_2, p_3, p_4), (p_4, p_1, p_3, p_2)\}\$
Case I	$p_3(p_1p_2 + p_4) \equiv 1 \pmod{p_2p_4}$	$\{(p_4, p_3, p_2, p_1), (p_2, p_3, p_1, p_4)\}\$
	$p_2 \equiv p_4 \pmod{p_3}$	
	$p_1 < p_2$	$\{(p_1, p_2, p_3, p_4), (p_4, p_3, p_1, p_2)\}\$
Case II	$p_3(p_1p_2 + p_4) \equiv 1 \pmod{p_1p_2p_4}$	$\{(p_4, p_3, p_2, p_1), (p_2, p_1, p_3, p_4)\}\$
	$p_1p_2 \equiv p_4 \pmod{p_3}$	
	$p_1 < p_4, p_2 < p_3$	$\{(p_1, p_2, p_3, p_4), (p_4, p_2, p_3, p_1)\}\$
Case III	$(p_1 + p_4)p_2p_3 \equiv 1 \pmod{p_1p_4}$	$\{(p_4, p_3, p_2, p_1), (p_1, p_3, p_2, p_4)\}\$
	$p_1 \equiv p_4 \pmod{p_2 p_3}$	
	$p_1 < p_4$	$\{(p_1, p_2, p_3, p_4), (p_4, p_3, p_2, p_1)\}\$
Case IV	$(p_1 + p_4)p_2p_3 \equiv 1 \pmod{p_1p_4}$	
	$p_1 \equiv p_3 p_4 \pmod{p_2}$	
	$p_1p_2 \equiv p_4 \pmod{p_3}$	

Remarks 2.11.

- (1) Note that the non-trivial permutations of $P = (p_1, p_2, p_3, p_4)$ appearing in the table are those labelled Q, \tilde{P} and \tilde{Q} in Lemma 2.9; they are all distinct except in Case IV, where we have $Q = \tilde{P}$ and $\tilde{Q} = P$.
- (2) The proposition asserts that a given quadruple cannot appear on the right-hand side of the table more than once, and that there are never more than two paths in G_n between n and $p_1p_2p_3p_4n$. However, it can happen that different permutations of (p_1, p_2, p_3, p_4) arise from different cases in the table or from the same case multiple times; for instance, eight permutations of (2, 3, 11, 13) give rise to quadruples with

multiplicity 2, and they arise once in Case I and twice in Case IV. This is not a contradiction because the sets N(P) and N(P') are disjoint for inequivalent permutations P and P', and thus the corresponding paths cannot emerge together from the same node.

(3) We will see below that solutions exist in each of the Cases I–IV.

Proof. Let $P = (p_1, p_2, p_3, p_4)$, $Q = (q_1, q_2, q_3, q_4)$, and suppose that $(P, Q) \in \mathcal{P}_4^2$ form an irreducible pair. Then there is a non-trivial permutation $\pi \in S_4$ such that P is equivalent to $Q = \pi . P$. Since (P, Q) is irreducible, π cannot stabilize any of the sets $\{1\}, \{1, 2\}$ or $\{1, 2, 3\}$. Moreover, by Lemma 2.9, the solutions for a given π are in one-to-one correspondence with those for π^{-1} , $\sigma\pi\sigma$ and $\sigma\pi^{-1}\sigma$, where $\sigma = (14)(23)$, so we may group those permutations together into classes and consider the solutions for only one permutation from each class.

With some straightforward computations in S_4 , we find that there are seven classes:

$$(2.9) \qquad \{(1234), (1432)\}, \{(1243), (1342)\}, \{(13)(24)\}, \\ \{(124), (142), (134), (143)\}, \{(1324), (1423)\}, \{(14)\}, \{(14)(23)\}. \}$$

The first three turn out to yield no solutions, while the last four correspond to the four cases in the table. We consider each class in turn and take π to be the first element listed in each case.

$$\pi = (1234): \text{ Then } (p_1, p_2, p_3, p_4) = (q_2, q_3, q_4, q_1), \text{ and we have}$$

$$1 \equiv q_1 = p_4 \pmod{p_1}$$

$$p_1 \equiv q_1 q_2 = p_1 p_4 \pmod{p_2} \Longrightarrow 1 \equiv p_4 \pmod{p_2}$$

$$p_1 p_2 \equiv q_1 q_2 q_3 = p_1 p_2 p_4 \pmod{p_3} \Longrightarrow 1 \equiv p_4 \pmod{p_3}$$

$$p_1 p_2 p_3 \equiv 1 \pmod{p_4}.$$

Thus, we have both $p_4 \equiv 1 \pmod{p_1 p_2 p_3}$ and $p_1 p_2 p_3 \equiv 1 \pmod{p_4}$, which is impossible.

 $\pi = (1243)$: Then $(p_1, p_2, p_3, p_4) = (q_2, q_4, q_1, q_3)$, and we have $1 = q_1 - p_2 \pmod{p_1}$

$$p_1 \equiv q_1 - p_3 \pmod{p_1}$$

$$p_1 \equiv q_1 q_2 q_3 = p_1 p_3 p_4 \pmod{p_2} \Longrightarrow 1 \equiv p_3 p_4 \pmod{p_2}$$

$$p_1 p_2 \equiv 1 \pmod{p_3}$$

$$p_1 p_2 p_3 \equiv q_1 q_2 = p_1 p_3 \pmod{p_4} \Longrightarrow p_2 \equiv 1 \pmod{p_4}.$$

Thus, p_1 divides $1-p_3$, and $p_2 \equiv \frac{1-p_3}{p_1} \pmod{p_3}$. Note that applying the permutation (14)(23) to the indices leaves the system unchanged, so we may assume without loss of generality that $p_2 < p_3$. Therefore, $p_2 = p_3 + \frac{1-p_3}{p_1}$, whence

$$p_3 \equiv \frac{p_3 - 1}{p_1} \pmod{p_2} \Longrightarrow p_1 p_3 \equiv p_3 - 1 \pmod{p_2}.$$

Since we also have $p_3p_4 \equiv 1 \pmod{p_2}$, this implies that $p_1 + p_4 \equiv 1 \pmod{p_2}$.

Now, if $p_2 < 5$ then we must have $p_2 = 3$, $p_4 = 2$, so $p_1 > 3$ and $p_2 < 2p_1 - 3$. On the other hand, if $p_2 \ge 5$ then $p_2 \ge 1 + 2p_4$, and

$$1 + p_2 \le p_1 + p_4 \le p_1 + \frac{p_2 - 1}{2} \Longrightarrow p_2 \le 2p_1 - 3.$$

Since $p_2 = p_3 + \frac{1-p_3}{p_1}$, this implies that $p_3 < \frac{2(p_1 - \frac{3}{2})p_1}{p_1 - 1} < 2p_1$. Hence, $p_3 = p_1 + 1$, so that $p_1 = 2, p_3 = 3$. But then $p_2 \le 2p_1 - 3 = 1$, which is impossible.

 $\pi = (13)(24)$: Then $(p_1, p_2, p_3, p_4) = (q_3, q_4, q_1, q_2)$ and we have

$$1 \equiv q_1 q_2 = p_3 p_4 \pmod{p_1}$$
$$p_1 \equiv q_1 q_2 q_3 = p_1 p_3 p_4 \pmod{p_2} \Longrightarrow 1 \equiv p_3 p_4 \pmod{p_2}$$
$$p_1 p_2 \equiv 1 \pmod{p_3}$$
$$p_1 p_2 p_3 \equiv q_1 = p_3 \pmod{p_4} \Longrightarrow p_1 p_2 \equiv 1 \pmod{p_4}.$$

Thus, we have both $p_3p_4 \equiv 1 \pmod{p_1p_2}$ and $p_1p_2 \equiv 1 \pmod{p_3p_4}$, which is impossible.

$$\pi = (124): \text{ Then } (p_1, p_2, p_3, p_4) = (q_2, q_4, q_3, q_1), \text{ and we have}$$

$$1 \equiv q_1 = p_4 \pmod{p_1}$$

$$p_1 \equiv q_1 q_2 q_3 = p_1 p_3 p_4 \pmod{p_2} \Longrightarrow 1 \equiv p_3 p_4 \pmod{p_2}$$

$$p_1 p_2 \equiv q_1 q_2 = p_1 p_4 \pmod{p_3} \Longrightarrow p_2 \equiv p_4 \pmod{p_3}$$

$$p_1 p_2 p_3 \equiv 1 \pmod{p_4},$$

which is equivalent to the set of conditions in Case I. The equivalence classes in the right-hand column are $\{P, \pi.P\}, \{\sigma.P, \sigma\pi.P\}$.

$$\pi = (1324): \text{ Then } (p_1, p_2, p_3, p_4) = (q_3, q_4, q_2, q_1), \text{ and we have}$$

$$1 \equiv q_1 q_2 = p_3 p_4 \pmod{p_1}$$

$$p_1 \equiv q_1 q_2 q_3 = p_1 p_3 p_4 \pmod{p_2} \Longrightarrow 1 \equiv p_3 p_4 \pmod{p_2}$$

$$p_1 p_2 \equiv q_1 = p_4 \pmod{p_3}$$

$$p_1 p_2 p_3 \equiv 1 \pmod{p_4},$$

which is equivalent to the system of congruences in Case II. In this case, the system is invariant under the action of $(12) = \sigma \pi$, but the normalization condition $p_1 < p_2$ ensures that each set of solutions $\{P, \pi.P\}, \{\sigma.P, \sigma \pi.P\}$ is counted only once.

$$\pi = (14): \text{ Then } (p_1, p_2, p_3, p_4) = (q_4, q_2, q_3, q_1), \text{ and we have}$$

$$1 \equiv q_1 q_2 q_3 = p_2 p_3 p_4 \pmod{p_1}$$

$$p_1 \equiv q_1 = p_4 \pmod{p_2}$$

$$p_1 p_2 \equiv q_1 q_2 = p_2 p_4 \pmod{p_3} \Longrightarrow p_2 \equiv p_4 \pmod{p_3}$$

$$p_1 p_2 p_3 \equiv 1 \pmod{p_4},$$

which is equivalent to the system of congruences in Case III. In this case, the system is invariant under both σ and π , but the normalization conditions $p_1 < p_4$ and $p_2 < p_3$ ensure that each set of solutions $\{P, \pi.P\}, \{\sigma.P, \sigma\pi.P\}$ is counted only once.

$$\pi = (14)(23): \text{ Then } (p_1, p_2, p_3, p_4) = (q_4, q_3, q_2, q_1), \text{ and we have}$$

$$1 \equiv q_1 q_2 q_3 = p_2 p_3 p_4 \pmod{p_1}$$

$$p_1 \equiv q_1 q_2 = p_3 p_4 \pmod{p_2}$$

$$p_1 p_2 \equiv q_1 = p_4 \pmod{p_3}$$

$$p_1 p_2 p_3 \equiv 1 \pmod{p_4},$$

which is equivalent to the system of congruences in Case IV. In this case, we have $\pi = \sigma$, so we get only one equivalence class of solutions. The system is also invariant under $\pi = \sigma$, but the normalization condition $p_1 < p_4$ ensures that each set of solutions $\{P, \pi.P\}$ is counted only once.

Conversely, it is easy to see that the logic is reversible in the last four cases considered, so any $(p_1, p_2, p_3, p_4) \in \mathcal{P}_4$ satisfying one of the given sets of conditions gives rise to multiple quadruples as indicated.

It remains to prove the assertion that the multiplicity is 2 in each case. Suppose that P is equivalent to both $Q = \pi . P$ and $Q' = \pi' . P$ for some non-trivial $\pi \neq \pi'$. Then Q is equivalent to $Q' = \pi' \pi^{-1} . Q$. Hence, π, π' and $\pi' \pi^{-1}$ are all contained in the union

 $\{(124), (142), (134), (143), (1324), (1423), (14), (14)(23), (13), (24)\}$

of the last four classes in (2.9), together with the permutations giving rise to multiple triples (p_1, p_2, p_3) or (p_2, p_3, p_4) . Note that we are free to replace P, Q, Q' by $\sigma.P, \sigma.Q, \sigma.Q'$ or to permute them arbitrarily, which is to say that we can replace (π, π') by any of the pairs

 $(\pi,\pi'), (\pi',\pi), (\pi^{-1},\pi'\pi^{-1}), (\pi'\pi^{-1},\pi^{-1}), (\pi''^{-1},\pi\pi'^{-1}) \text{ or } (\pi\pi'^{-1},\pi'^{-1}),$

or their conjugates by σ . Going through all possibilities, we find that we may assume that

$$(\pi, \pi') \in \{((124), (142)), ((13), (124)), ((13), (134))\}.$$

We consider these three cases in turn.

 $\pi = (124), \pi' = (142)$: Recall that $\pi = (124)$ leads to the system in Case I. For $\pi' = (142)$ and $Q' = (q'_1, q'_2, q'_3, q'_4)$, we have $(p_1, p_2, p_3, p_4) = (q'_4, q'_1, q'_3, q'_2)$, so that $p_1 \equiv 1 \pmod{p_2}$ and $p_1 p_2 p_3 \equiv q'_1 = p_2 \pmod{p_4}$. Hence, $p_2 \equiv 1 \pmod{p_4}$, and we also have $p_4 \equiv 1 \pmod{p_1}$, so that $p_4 < p_2 < p_1 < p_4$, which is impossible.

 $\pi = (13), \pi' = (124)$: Then (p_1, p_2, p_3, p_4) satisfies the system in Case I as well as (2.2). Thus we have

$$1 \equiv p_4 \equiv p_2 p_3 p_4 \pmod{p_1}$$

$$1 \equiv p_3 p_4 \equiv p_1 p_4 \pmod{p_2}$$

$$1 \equiv p_1 p_2 \equiv p_1 p_4 \pmod{p_3},$$

so that $p_4(p_1 + p_2p_3) \equiv 1 \pmod{p_1p_2p_3}$. Also, $p_1p_2p_3 \equiv 1 \pmod{p_4}$, so that $p_4 = \frac{p_1p_2p_3-1}{t}$ for some $t \in (0, p_1p_2p_3) \cap \mathbb{Z}$. Substituting for p_4 , we have $t \equiv -p_1 - p_2p_3 \pmod{p_1p_2p_3}$, whence $t = p_1p_2p_3 - p_1 - p_2p_3$. Thus,

$$p_4(p_1p_2p_3 - p_1 - p_2p_3) = p_1p_2p_3 - 1,$$

which implies

$$p_4p_1 - 1 = ((p_4 - 1)p_1 - p_4)p_2p_3 \ge 6[(p_4 - 1)p_1 - p_4]$$

Hence $p_1 \leq \frac{6p_4-1}{5p_4-6}$. If $p_4 \geq 3$ then this gives $p_1 < 2$, while if $p_4 = 2$ then $2 < p_1 < 3$, but both of these are impossible.

 $\pi = (13), \pi' = (134)$: We have $Q' = (q'_1, q'_2, q'_3, q'_4) = (p_4, p_2, p_1, p_3)$, and in view of (2.2) we get

$$1 \equiv q_1'q_2' = p_2p_4 \pmod{p_1} \Longrightarrow p_4 \equiv p_2^{-1} \equiv p_3 \pmod{p_1}$$
$$p_1 \equiv q_1' = p_4 \pmod{p_2} \Longrightarrow p_4 \equiv p_3 \pmod{p_2}$$
$$p_1p_2 \equiv q_1'q_2'q_3' = p_1p_2p_4 \pmod{p_3} \Longrightarrow p_4 \equiv 1 \equiv p_1p_2 \pmod{p_3}$$
$$p_1p_2p_3 \equiv 1 \pmod{p_4}.$$

Hence $p_4 \equiv p_3 + p_1 p_2 \pmod{p_1 p_2 p_3}$ and $p_4 < p_1 p_2 p_3$, so that $p_4 = p_3 + p_1 p_2$. By parity considerations we see that at least one of p_1 , p_2 and p_3 must be 2, and it follows from Theorem 2.5 that (p_1, p_2, p_3) is a permutation of (2, 3, 5). Therefore, $p_1 p_2 p_3 - 1 = 29$ is prime, so that $p_4 = p_1 p_2 p_3 - 1 > p_3 + p_1 p_2$, which is a contradiction.

Finally, suppose that a quadruple P occurs in the table for two different choices of (p_1, p_2, p_3, p_4) . Then, by the above argument, in both instances P must be related to the other element of its equivalence class by the same permutation. Thus, either P appears once in each equivalence class in Case II or Case III, or twice in Case IV. However, the normalization conditions rule out all of these possibilities.

Table 2.2 shows the first several solutions to the conditions in Proposition 2.10, ordered by modulus.

2.3. Multiple k-tuples for large k. The alert reader will note that the congruence constraints in Cases II and III of Proposition 2.10 are nothing but (2.2) with (p_1, p_2, p_3) replaced by (p_1p_2, p_3, p_4) or (p_1, p_2p_3, p_4) ; in particular, the solutions are parametrized by Theorem 2.5. This turns out to be a general phenomenon, in the sense that the system of congruences arising from a given element of S_k can be embedded in a system for any K > k by grouping the primes into products, as the following lemma shows.

Lemma 2.12. For i = 1, ..., k, let $P_i > 1$ be an integer with prime factors p_{ij} for $j = 1, ..., r_i$, and assume that $P_1 \cdots P_k$ is squarefree. Put $K = r_1 + ... + r_k$, and set

$$P = (p_{11}, \ldots, p_{1r_1}, \ldots, p_{k1}, \ldots, p_{kr_k}) \in \mathcal{P}_K.$$

Suppose that $\pi \in S_k$ is a non-trivial permutation such that

$$P_1 \cdots P_{i-1} \equiv Q_1 \cdots Q_{\pi(i)-1} \pmod{P_i} \quad for \ i = 1, \dots, k,$$

where $(Q_1, \ldots, Q_k) = \pi.(P_1, \ldots, P_k)$. Then there is a non-trivial permutation $\Pi \in S_K$ such that $P \sim \Pi.P$. Further, the pair $(P, \Pi.P)$ is irreducible if and only if

$$P_1 \cdots P_i \neq Q_1 \cdots Q_i \quad for \ 0 < i < k.$$

Remark 2.13. Note that the order of the prime factors of P_i is not specified, so each solution (P_1, \ldots, P_k) gives rise to $\prod_{i=1}^k r_i!$ multiple K-tuples.

Proof. The main idea is to apply π to the blocks of indices of length r_i . More formally, for $i = 1, \ldots, k+1$, let $s_i = r_1 + \ldots + r_{i-1}$ and $t_i = r_{\pi^{-1}(1)} + \ldots + r_{\pi^{-1}(i-1)}$. Note that $s_i + j$ is the index of the *j*th prime factor of P_i in P. Given $I \in \{1, \ldots, K\}$ we define $\Pi(I) = t_{\pi(i)} + j$, where $i \in \{1, \ldots, k\}$ and $j \in \{1, \ldots, r_i\}$ are the unique indices for which $I = s_i + j$.

(p_1, p_2, p_3, p_4)	P	a(P)	case			
(2, 5, 7, 3)	210	107, 149	II			
(3, 13, 2, 7)	546	181, 251	Ι			
(3, 2, 11, 13)	858	467, 779	Ι			
(11, 3, 2, 13)	858	571	IV			
(13, 3, 2, 11)	858	857	IV			
(3, 19, 11, 2)	1254	1127	IV			
(7, 3, 2, 41)	1722	1721	IV			
(41, 3, 2, 7)	1722	1147	IV			
(41, 7, 2, 3)	1722	491	IV			
(41, 7, 3, 2)	1722	1639	IV			
(5, 29, 2, 17)	4930	3909	IV			
(13, 2, 5, 43)	5590	3353, 5589	III			
(2, 3, 31, 37)	6882	1183, 5771	II			
(3, 7, 17, 89)	31773	22427, 26966	II			
(103, 31, 2, 5)	31930	5149	IV			
(7, 23, 2, 107)	34454	29959	IV			
(3, 17, 31, 79)	124899	81764, 81922	Ι			
(41, 17, 2, 199)	277406	32635	IV			
(5, 53, 37, 43)	421615	39559, 173203	II			
(73, 5, 13, 593)	2813785	1125513, 1861426	III			
(449, 67, 2, 191)	11491706	6517683	IV			
(241, 2, 113, 3631)	197766046	183764909, 42003407	III			
(2, 3541, 997, 103)	727257662	714062125	IV			
(23, 367, 401, 421)	1425018061	418499259, 1226476565	II			
TABLE 2.2. Multiple quadruples of small modulus						

Note that $\Pi(I) = t_{\pi(i)} + j \leq t_{\pi(i)} + r_i \leq K$, so Π maps $\{1, \ldots, K\}$ to itself. To see that it defines an element of S_K , it suffices to show that it is surjective. To that end, given any $I \in \{1, \ldots, K\}$, choose *i* to be the largest positive integer such that $t_i < I$, and set $j = I - t_i > 0$. Then $t_i + r_{\pi^{-1}(i)} = t_{i+1} \geq I$, so $j \leq r_{\pi^{-1}(i)}$. Hence $I = \Pi(s_{\pi^{-1}(i)} + j)$, as required.

We must show that P is equivalent to $\Pi.P$. Let u_1, \ldots, u_K denote the entries of P and v_1, \ldots, v_K the entries of $\Pi.P$. Given $I \in \{1, \ldots, K\}$, let $I = s_i + j$ for $i \in \{1, \ldots, k\}$ and $j \in \{1, \ldots, r_i\}$. Then

$$u_1 \cdots u_{I-1} = \left(\prod_{i'=1}^{i-1} P_{i'}\right) \left(\prod_{j'=1}^{j-1} u_{s_i+j'}\right).$$

Since $u_I \mid P_i$ and $u_{s_i+j'} = v_{t_{\pi(i)}+j'}$ for $j' = 1, \ldots, j$, this is congruent modulo $u_I = v_{\Pi(I)}$ to

$$\left(\prod_{i'=1}^{\pi(i)-1} Q_{i'}\right) \left(\prod_{j'=1}^{j-1} v_{t_{\pi(i)}+j'}\right) = v_1 \cdots v_{\Pi(I)-1}.$$

Since I was arbitrary, $P \sim \Pi P$.

As for the final claim, if (P, Π, P) is not irreducible then $u_1 \cdots u_I = v_1 \ldots v_I$ for some $I \in (0, K) \cap \mathbb{Z}$. If $I < r_1$ then by definition we have $\Pi(I) = t_{\pi(1)} + I$. Since u_I divides $v_1 \cdots v_I$, we also have $\Pi(I) \leq I$. Thus $t_{\pi(1)} = 0$, which implies $\pi(1) = 1$ and $P_1 = Q_1$. Hence we may assume that $I \geq r_1$.

Let i < k be the largest positive integer such that $I \ge s_{i+1}$, and i' < k the largest nonnegative integer such that $I \ge t_{i'+1}$. It follows that $u_1 \cdots u_I$ is divisible by P_1, \ldots, P_i but not by P_j for any j > i. Similarly, $v_1 \cdots v_I$ is divisible by $Q_1 \ldots, Q_{i'}$, but not by Q_j for any j > i'. Since $P_j = Q_{\pi(j)}$ for every j, it follows that π is a bijection between $\{1, \ldots, i\}$ and $\{1, \ldots, i'\}$; hence i' = i and π stabilizes $\{1, \ldots, i\}$. In particular, $P_1 \cdots P_i = Q_1 \cdots Q_i$.

Conversely, suppose that $P_1 \cdots P_i = Q_1 \cdots Q_i$ for some $i \in (0, k) \cap \mathbb{Z}$. We have $P_1 \cdots P_i = u_1 \cdots u_I$ and $Q_1 \cdots Q_i = v_1 \cdots v_{I'}$ for some $I, I' \in (0, K) \cap \mathbb{Z}$. By unique factorization, I = I', and thus $(P, \Pi. P)$ is not irreducible.

In the following we let T_r denote the set of squarefree integers with at most r prime factors, and $T_{\infty} = \bigcup_{r=0}^{\infty} T_r$ the set of all squarefree integers.

Lemma 2.14. Let $f(x) = (x^2+x+1)(x^2+1)(x^3+x^2+2x+1)$ and $g(x) = x(x^2-x+1)(x^2+1)$. Then, for any $q \in \mathbb{Z}_{>0}$ and all sufficiently large X > 0 (with the meaning of "sufficiently large" possibly depending on q), we have

 $\begin{array}{l} (1) \ \#\{x \in \mathbb{Z} \cap [1, X] : (f(x), q) = 1 \ and \ f(x) \in T_{\infty}\} \gg_{q} X; \\ (2) \ \#\{x \in \mathbb{Z} \cap [1, X] : (f(x), q) = 1 \ and \ f(x) \in T_{13}\} \gg_{q} \frac{X}{(\log X)^{3}}; \\ (3) \ \#\{x \in \mathbb{Z} \cap [1, X] : (g(x), q^{2}) = q \ and \ q^{-1}g(x) \in T_{\infty}\} \gg_{q} X; \\ (4) \ \#\{x \in \mathbb{Z} \cap [1, X] : (g(x), q^{2}) = q \ and \ q^{-1}g(x) \in T_{12}\} \gg_{q} \frac{X}{(\log X)^{3}}. \end{array}$

Proof. Let $h \in \mathbb{Z}[x]$ be a squarefree polynomial with k irreducible factors and content 1, and suppose that there exists $a \in \mathbb{Z}$ such that $p \nmid h(a)$ for every prime $p \leq \deg h$. Then it was shown in [2] that if every irreducible factor of h has degree at most 3 then there are positive numbers c = c(h) and $r = r(k, \deg h)$ such that

$$#\{x \in \mathbb{Z} \cap [1, X] : h(x) \in T_{\infty}\} = (c + o(1))X \quad \text{as } X \to \infty,$$

and

$$#\{x \in \mathbb{Z} \cap [1, X] : h(x) \in T_r\} \gg_h \frac{X}{(\log X)^k} \text{ for } X \gg_h 1.$$

Further, for k = 3 and deg h = 7 we may take r = 13. Thus, (1) and (2) follow on applying these results to h(x) = f(qx).

For (3) and (4) we set $Q = \operatorname{lcm}(q, 2)$ and take $h(x) = Q^{-1}g(Q + Q^2x)$. Then $h \in \mathbb{Z}[x]$, and if $a \in \mathbb{Z}$ is such that

$$Qa \equiv -1 \; (\text{mod } \frac{15}{(q,15)}),$$

then (h(a), 30) = 1. From [2] we find that r = 11 is admissible for h, from which (3) and (4) follow.

Theorem 2.15.

- (1) For any $q \in \mathbb{Z}_{>0}$, there are infinitely many positive integers k such that \mathcal{P}_k^2 contains an irreducible pair of modulus co-prime to q.
- (2) There is a positive integer $k \leq 13$ such that, for any $q \in \mathbb{Z}_{>0}$, \mathcal{P}_k^2 contains infinitely many irreducible pairs of modulus co-prime to q.

(3) For any squarefree $q \in \mathbb{Z}_{>0}$, there are infinitely many positive integers k such that \mathcal{P}_k^2 contains an irreducible pair of modulus divisible by q, and the least such k is at most $\omega(q) + 12$.

Remark 2.16.

- Combining (1) and (2) with Lemma 2.2 and the Chinese remainder theorem, we see that if $a, q, k \in \mathbb{Z}_{>0}$, then for a positive proportion of the numbers $n \equiv a \pmod{q}$, G_n contains both a loop of height ≤ 13 and a loop of height $\geq k$. If (a, q) = 1 then the same assertion holds with n restricted to primes.
- Similarly, by (3), for any squarefree $q \in \mathbb{Z}_{>0}$ there is a prime *n* such that G_n contains a loop of height $\leq \omega(q) + 12$ that has every prime factor of *q* as an edge. In particular, every prime occurs as an edge of a loop in some G_n .

Proof. Let f(x) be as in Lemma 2.14. Suppose that f(x) is squarefree for some $x \in \mathbb{Z}_{>0}$, and put

$$(P_1, P_2, P_3) = (x^2 + x + 1, x^2 + 1, x^3 + x^2 + 2x + 1).$$

Then the P_i are squarefree and pairwise co-prime. By Theorem 2.5, (P_1, P_2, P_3) satisfies (2.2), and applying Lemma 2.12 with $\pi = (13)$, we obtain an irreducible pair $(P, \Pi.P) \in \mathcal{P}_K^2$, where |P| = f(x) and $K = \omega(f(x))$. (Recall that $\omega(n)$ denotes the number of distinct prime factors of n.)

Now, to prove (1), we construct a sequence of positive integers x_i as follows. Assume that x_1, \ldots, x_{i-1} have been chosen, and set

$$r = \begin{cases} 0 & \text{if } i = 1, \\ \omega(f(x_{i-1})) & \text{if } i > 1. \end{cases}$$

It was shown by Halberstam [8] that, for any irreducible polynomial $h \in \mathbb{Z}[x]$, $\frac{\omega(h(x)) - \log \log x}{\sqrt{\log \log x}}$ has a Gaussian distribution, as in the Erdős–Kac theorem. Taking h to be one of the irreducible factors of f, we have in particular that

$$\#\{x \in \mathbb{Z} \cap [1, X] : f(x) \in T_r\} \le \#\{x \in \mathbb{Z} \cap [1, X] : h(x) \in T_r\} = o(X) \text{ as } X \to \infty.$$

Thus, by part (1) of Lemma 2.14, we may choose $x_i \in \mathbb{Z}_{>0}$ such that $(f(x_i), q) = 1$, $f(x_i)$ is squarefree and $\omega(f(x_i)) > r$.

Hence, for the sequence of x_i thus constructed, $\omega(f(x_i))$ is strictly increasing. By the above, for each i, $\mathcal{P}^2_{\omega(f(x_i))}$ contains an irreducible pair of modulus $f(x_i)$, and (1) follows.

Turning to (2), suppose that there is no such k. Then for each k = 1, ..., 13, there exists $q_k \in \mathbb{Z}_{>0}$ such that \mathcal{P}_k^2 contains at most finitely many irreducible pairs of modulus co-prime to q_k , and replacing q_k by a suitable multiple if necessary, we may assume that there are no such pairs. Applying part (2) of Lemma 2.14 with $q = q_1 \cdots q_{13}$, there exists $x \in \mathbb{Z}_{>0}$ such that $f(x) \in T_{13}$ and (f(x), q) = 1. By the above construction, we obtain an irreducible pair $(P, \Pi.P) \in \mathcal{P}_K^2$ of modulus co-prime to q, where $K = \omega(f(x)) \leq 13$. This is a contradiction, and (2) follows.

Finally, (3) is proved in much the same way using the triple

$$(P_1, P_2, P_3) = (x, x^2 - x + 1, x^2 + 1),$$

corresponding to the second line of Theorem 2.5 with n = 2, and g(x) in place of f(x); we omit the details.

2.4. Multiple k-tuples with small modulus. One could continue as in Propositions 2.4 and 2.10 to classify the multiple k-tuples for $k = 5, 6, \ldots$, but as the proof of Proposition 2.10 shows, this quickly becomes cumbersome. A more practical means of identifying relatively dense arithmetic progressions N(P) of nodes giving rise to loops is to do a direct search for small values of |P|.

One procedure for finding all multiple k-tuples of a given modulus is as follows. Suppose that m is a squarefree positive integer (our candidate for |P|), and rewrite the system of congruences in Definition 2.1 as

$$(2.10) p_1 \cdots p_{i-1} \equiv d_i \pmod{p_i},$$

where d_1, \ldots, d_k are proper divisors of m satisfying

(2.11)
$$d_i \neq d_j \text{ and } \min(d_i, d_j) \mid \max(d_i, d_j)$$

for all $i \neq j$. (If we wish to find only irreducible pairs, then we impose the further constraint $d_i \neq p_1 \cdots p_{i-1}$.) We search for solutions to (2.10) recursively: suppose that p_1, \ldots, p_{i-1} and d_1, \ldots, d_{i-1} have been chosen, loop over all proper divisors d_i of m such that (2.11) holds for all j < i, and then over all primes $p_i \mid \frac{m}{p_1 \cdots p_{i-1}}$ such that (2.10) holds. Since (2.10) is a very restrictive condition, most branches of the search tree are pruned quickly, so this method is substantially more efficient than naively trying all permutations of the prime factors of m.

We coded this procedure and used it to find 195167 (unordered) irreducible pairs of modulus $|P| < 10^9$. The results reveal that for large k, topologies that are much more intricate than the simple loops observed in Propositions 2.4 and 2.10 can arise. For instance, for any $n \equiv 58183403 \pmod{635825190}$, G_n has a subgraph as shown in Figure 2, in which there are 7 paths between n and 635825190n, 12 out of the 21 pairs of paths are irreducible, and there are subloops of heights 3, 4, 5, 6 and 8.

Note that only pairs of modulus co-prime to $2 \cdot 3 \cdot 7 \cdot 43$ can possibly appear in G_1 . Imposing that restriction reduces the list to just 18 moduli $|P| < 10^9$ with 42 associated arithmetic progressions N(P), as shown in Table 2.3. Consider, for instance, the progressions with modulus $115908845 = 5 \cdot 13 \cdot 23 \cdot 31 \cdot 41 \cdot 61$. It is known (see the introduction of [1]) that none of these primes can occur as an edge of the right-most branch of G_1 (sequence A000946). Therefore, it seems natural to expect the nodes of the right-most branch to vary randomly among the invertible residue classes mod 115908845 as the level increases, in the sense that each residue class should occur with equal frequency. (This is the same heuristic reasoning as that supporting Shanks' conjecture [13] that the first Euclid–Mullin sequence contains every prime.) Thus, we would expect one of the four corresponding residue classes in Table 2.3 to occur with frequency $4/\varphi(115908845) = 1/19008000$. In particular, we are led to the following conjecture:

Conjecture 2.17. G_1 contains infinitely many loops.

More generally, it seems likely that each of the residue classes N(P) in Table 2.3 will be met infinitely often by the nodes of G_1 ; we provide some evidence towards this in the next section. It is difficult to compute the overall probability of a random node on the graph landing in one of the residue classes, since these events are not independent, i.e. the classes overlap in non-trivial ways. However, it is apparent from the first few lines of the table that the greatest chance of finding a loop comes from the progressions of modulus $2813785 = 5 \cdot 13 \cdot 73 \cdot 593$, with density $2/\varphi(2813785) = 1/1022976$. Thus, on the sub-graph



FIGURE 2. Some nodes of G_n for $n \equiv 58183403 \pmod{635825190}$

of nodes co-prime to 2813785, we expect roughly one out of every million nodes to be the base of a loop of height 4.

3. Numerical results

We used two methods for exploring G_1 numerically. First, we used freely available software implementations of the elliptic curve method (see GMP-ECM [7]) and general number field sieve (see YAFU [6], msieve [5] and GGNFS [4]) to compute as many nodes as was practical for levels up to 17. This was a community effort, with support from users of mersenneforum.org.

Table 3.1 lists the number of nodes that we have computed at each level of the graph G_1 . The final column is the number of remaining unfactored composites at that level. Factoring a composite at a given level will increase the number of nodes at that level (by at least 2) and all subsequent levels. The single remaining composite at level 13 is the 253-digit:

30741638041263757309600460000064107032998604910525153993522043894945654246227689310806-05652579832748915879865519993669161314951649593763245464995966627308199534468607184384-744257573685683221611440202806222725727083224756010635164700144499225512799343807.

We have been unable to factor this number despite running GMP-ECM on approximately 225,000 curves at $B_1 = 8.5 \times 10^8$ and 44,000 curves at $B_1 = 3 \times 10^9$ (and default B_2 values); this is a level of effort comparable to a "t70", meaning it has a reasonable chance of revealing any prime factors with up to 70 digits.

$\{p_1,\ldots,p_k\}$	P	a(P)	1/density		
$\{5, 13, 73, 593\}$	2813785	1125513, 1861426	1022976		
$\{5, 11, 13, 79, 523\}$	29541655	2913109, 19876614	9771840		
$\{5, 13, 17, 53, 563\}$	32972095	45473, 14501753, 15173846, 15665474	5611008		
$\{5, 11, 23, 31, 1307\}$	51254005	29374824, 37354844	17239200		
$\{5, 13, 23, 31, 41, 61\}$	115908845	30432518, 43262953, 74975328, 87805763	19008000		
{197, 211, 2969}	123412423	114015537	122162880		
$\{5, 11, 23, 67, 1831\}$	155186405	92870549	106286400		
$\{5, 13, 19, 23, 71, 89\}$	179491195	106001778, 120823468, 140339224, 145796156	29272320		
$\{5, 13, 29, 61, 1597\}$	183631045	16718992, 26947777, 40801752, 51030537	32175360		
$\{5, 11, 13, 41, 73, 113\}$	241819435	31106978, 108457851	77414400		
$\{5, 11, 13, 17, 97, 233\}$	274715155	161397329, 273388114	85524480		
$\{5, 11, 13, 733, 773\}$	405125435	26064013, 332551556	135624960		
$\{5, 11, 19, 41, 83, 127\}$	451629145	16717776, 363759119	148780800		
$\{5, 13, 19, 23, 37, 449\}$	471892265	331562178, 399028904	153280512		
$\{5, 19, 53, 337, 421\}$	714350695	171690041, 516232304	264176640		
$\{5, 11, 19, 53, 71, 199\}$	782534665	504298018, 599617009	259459200		
{11, 19, 29, 71, 1871}	805149301	159790883,664072158	329868000		
$\{5, 11, 23, 61, 67, 167\}$	863399185	132876677, 201396989	289238400		
TABLE 2.3. Multiple k-tuples of modulus $ P < 10^9$ with $(P , 2 \cdot 3 \cdot 7 \cdot 43) = 1$					

level	nodes	composites	level	nodes	$\operatorname{composites}$
≤ 4	1	0	11	555	0
5	2	0	12	2020	0
6	4	0	13	7948	1
7	9	0	14	32738	8
8	24	0	15	141619	636
9	52	0	16	622317	13445
10	165	0	17	2550301	186060

TABLE 3.1. Number of nodes by level in G_1 .

The numbers appearing in Theorem 1.1 were found by checking our data for the residue classes in Table 2.3. If there is a lower node with multiple paths to 1 then there must be a loop of height k starting from some node of level $\leq 20 - k$, viz. at most 17 if k = 3 and 16 if $k \geq 4$. Although we have not been able to compute the full graph up to level 17, we expect that there are no more than one million nodes remaining to be found up to that level, with at most 50 thousand of those at level 16 or below (and fewer still that are co-prime to 5 and 13). In view of Table 2.3, it seems unlikely that one of those will yield a loop. However, that cannot be established definitively until the full graph is computed up to level 18, which is out of reach with present technology.

Our second numerical method aimed to produce large quantities of nodes rather than a comprehensive list of all of them. We began with our list of nodes at level 17 and followed only the edges corresponding to primes below some bound B. Taking $B = 2^{24}$ and computing up to level 50, we found at least one match to every congruence class listed in Table 2.3;

k	$\frac{X_{k+1}/X_k}{\sqrt{2k}}$
9	0.748
10	0.752
11	0.776
12	0.803
13	0.808
14	0.825

TABLE 4.1. Estimated values of $\frac{X_{k+1}/X_k}{\sqrt{2k}}$

in particular we found loops of heights 3, 4, 5 and 6. This method was also helpful for investigating some other statistical questions, as we describe in the next section.

4. Related questions

In this final section, we record some numerical observations and heuristics on related questions:

- Does every prime occur as an edge in G_1 ? This seems very likely. With the second method described above, we verified that every prime below 10^9 occurs.
- How does the number of nodes at level k grow asymptotically as $k \to \infty$? Let X_k denote the number of nodes of G_1 of level k. Heuristically, based on the Erdős–Kac theorem, we expect that for a typical node n, n + 1 will have about $\log \log n$ prime factors, with the values of $\frac{\log \log p}{\log \log(n+1)}$ uniformly distributed on [0, 1] as p varies over the prime factors of n + 1.

Let n_k be the nodes of a typical path in G_1 , with $n_0 = 1$, and define θ_k so that

(4.1)
$$\frac{n_{k+1}}{n_k} = \exp([\log(n_k+1)]^{\theta_k}).$$

Then by the above heuristic, θ_k should vary uniformly over [0, 1] as $k \to \infty$. If we instead treat the θ_k formally as independent, uniform random variables on [0, 1] and define n_k by (4.1), then it is not hard to see that

$$\lim_{k \to \infty} \frac{\log \log n_k}{\sqrt{2k}} = 1$$

holds almost surely. Thus, we might expect the typical node of level k to be of size $\exp \exp([1+o(1)]\sqrt{2k})$. (This analysis ignores the fact that n_k+1 is co-prime to n_k and hence typically has no small prime factors; however, in the random model, the bulk of the contribution to $\log \log n_k$ comes from the values of θ_k close to 1, corresponding to the large prime factors, so this makes little difference.) In turn, this leads to the conjecture that $\frac{X_{k+1}}{X_k} = (1+o(1))\sqrt{2k}$, or equivalently $\log X_k = \frac{k}{2} \left(\log \frac{2k}{e} + o(1)\right)$. As far as we are aware, it is not even known that X_k is unbounded, so this remains

As far as we are aware, it is not even known that X_k is unbounded, so this remains largely guesswork. Table 4.1 shows estimated values of $\frac{X_{k+1}/X_k}{\sqrt{2k}}$ for $k \leq 14$, based on the data in Table 3.1. Although the data are very limited, they are at least consistent with the above guess, in that the ratio appears to grow slowly towards 1.

• Are there arbitrarily long chains of nodes with only one child each? This is related to the previous two questions. The basic heuristic underlying Shanks' conjecture is that

the nodes of a given path in G_1 should vary randomly among the invertible residue classes modulo a fixed prime p, until p occurs as an edge (beyond which every node is divisible by p). One (perhaps the only) conceivable way in which this heuristic might fail is if n + 1 is prime for every node n of sufficiently large level along the path. In fact, as discovered by Kurokawa and Satoh [9], that *can* happen for the analogous question over $\mathbb{F}_p[x]$.

All numerics to date indicate that this pathology does not occur over \mathbb{Z} , but it is an interesting question whether there are arbitrarily long chains in G_1 of nodes n such that n + 1 is prime. For a random node n, we can estimate the probability that n + 1 is prime as $\frac{n}{\varphi(n)\log n}$, so the chance that there is a unique path of length ℓ descending from n is roughly

$$\frac{n}{\varphi(n)\log n} \times \frac{n}{\varphi(n)\log(n^2)} \times \dots \times \frac{n}{\varphi(n)\log(n^{2^{\ell-1}})} = \left(\frac{n}{\varphi(n)2^{\frac{\ell-1}{2}}\log n}\right)^{\ell} \gg_{\ell} (\log n)^{-\ell}.$$

As above, we expect the *n* of level *k* typically satisfy $\log n = e^{O(\sqrt{k})}$. Hence, if our asymptotic guess for X_k holds then we should indeed expect chains of length ℓ to occur for sufficiently large *k*, and in fact we might expect ℓ as large as about $\sqrt{\frac{k \log k}{\log 2}}$. By our second method we found several examples of nodes followed by a unique path of length 4; the lowest (after the root node 1) is the following node at level 20:

- $\begin{array}{l} 2\cdot3\cdot7\cdot43\cdot139\cdot50207\cdot1607\cdot38891\cdot71609249149971437\cdot97272377313541\cdot318004829\\ \cdot1555110880896883\cdot39807662109343\cdot53437\cdot35251\cdot79\cdot2011283825921\cdot29\cdot17\cdot241.\end{array}$
 - Is G_1 planar? Our search for multiple k-tuples of small modulus uncovered several arithmetic progressions of n, e.g. $n \equiv 93397 \pmod{510510}$, such that G_n is not planar. As a generalization of Theorem 1.1, it is a natural question whether G_1 itself is planar. However, despite making an extended search, every progression that we found leading to non-planar graphs had modulus divisible by 6, and it is unclear whether or not that is a necessary condition. In any case, if G_1 is non-planar, that fact is likely not manifested until astronomically large level, so this question is unlikely to be settled in the near future.
 - How does the number of irreducible pairs of modulus $\leq X$ grow asymptotically as $X \to \infty$? The proof of Theorem 2.15 shows that, for large X,

(4.2)
$$\#\{q \in \mathbb{Z} \cap [1, X] : \exists \text{ an irreducible pair of modulus } q\} \gg X^{1/5},$$

and this gives a lower bound for the number of irreducible pairs of modulus up to X. However, a log-log fit of our data up to 10^9 suggests that this is too low, and that (4.2) is perhaps asymptotic to $cX^{5/8}$ for some c > 0. Note that the moduli exhibited in the lower bound in (4.2) are all even (for odd moduli the proof of Theorem 2.15 gives only a lower bound $\gg X^{1/7}$); our numerics also suggest that almost all irreducible pairs have even modulus.

Finally, we record the latest results on the computation of the Euclid–Mullin sequence and some of its relatives. Let M_n denote the first Euclid–Mullin sequence starting with the prime n, i.e. the edges of the left-most path in G_n . Wagstaff [15] computed M_2 up through the 43rd term (180 digits). Much computation effort, including several large GNFS world-wide

n	p	step	digits	OEIS	$\mid n$	p	step	digits	OEIS
2	41	52	335	A000945	47	23	36	194	A051319
5	31	58	347	A051308	53	71	92	526	A051320
11	29	56	313	A051309	59	37	79	1059	A051321
13	17	58	353	A051310	61	29	47	501	A051322
17	37	31	232	A051311	67	19	43	200	A051323
19	43	73	922	A051312	71	79	140	991	A051324
23	29	62	515	A051313	73	83	131	949	A051325
29	67	80	566	A051314	79	17	32	292	A051326
31	29	38	240	A051315	83	71	65	296	A051327
37	59	77	826	A051316	89	79	79	743	A051328
41	43	56	933	A051317	97	53	52	261	A051330
	TABLE 4.2. Summary of M_n for $n < 100$.								

distributed efforts, has since been expended on factoring the integers needed to extend the sequence. In 2012, Ryan Propper found a 75-digit factor using ECM; it remains the fifth largest factor ever produced by ECM.

Table 4.2 is a summary of known computational results for the distinct sequences with n < 100. The 'p' column is the smallest prime not yet confirmed as a member of the corresponding sequence. The 'step' column indicates the number of known terms and the 'digits' column the number of decimal digits in the unfactored composite needed for the next step. The final column is the corresponding entry number in the OEIS. It is unlikely that any of the blocking composites has a factor of less than 45 digits.

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