

**COMPUTER ASSISTED PROOF FOR APWENIAN
SEQUENCES RELATED TO HANKEL
DETERMINANTS**

HAO FU AND GUO-NIU HAN

ABSTRACT. An infinite ± 1 -sequence is called *Apwenian* if its Hankel determinant of order n divided by 2^{n-1} is an odd number for every positive integer n . In 1998, Allouche, Peyrière, Wen and Wen discovered and proved that the Thue–Morse sequence is an Apwenian sequence by direct determinant manipulations. Recently, Bugeaud and Han re-proved the latter result by means of an appropriate combinatorial method. By significantly improving the combinatorial method, we prove that several other Apwenian sequences related to the Hankel determinants with Computer Assistance.

1. INTRODUCTION

For each infinite sequence $\mathbf{c} = (c_k)_{k \geq 0}$ and each nonnegative integer n the *Hankel determinant* of order n of the sequence \mathbf{c} is defined by

$$(1.1) \quad H_n(\mathbf{c}) := \begin{vmatrix} c_0 & c_1 & \cdots & c_{n-1} \\ c_1 & c_2 & \cdots & c_n \\ \vdots & \vdots & \ddots & \vdots \\ c_{n-1} & c_n & \cdots & c_{2n-2} \end{vmatrix}.$$

We also speak of the Hankel determinants of the power series $\tilde{\mathbf{c}}(x) = \sum_{k \geq 0} c_k x^k$ and write $H_n(\tilde{\mathbf{c}}(x)) = H_n(\mathbf{c})$. The Hankel determinants are widely studied in Mathematics and, in several cases, can be evaluated by basic determinant manipulation, *LU*-decomposition, or Jacobi continued fraction (see, e.g., [15, 16, 7, 18, 17]). However, the Hankel determinants studied in the present paper apparently have no closed-form expressions, and require additional efforts to obtain specific arithmetical properties.

An infinite ± 1 -sequence $\mathbf{c} = (c_k)_{k \geq 0}$ is called *Apwenian* if its Hankel determinant of order n divided by 2^{n-1} is an odd number, i.e., $H_n(\mathbf{c})/2^{n-1} \equiv 1 \pmod{2}$, for all positive integer n . The corresponding

Date: May 21, 2015.

2010 Mathematics Subject Classification. 05A05, 05A15, 11B50, 11B85, 11C20.

Key words and phrases. Hankel determinant, Thue–Morse sequence, Apwenian sequence, permutation, computer assisted proof.

The first author was supported by the National Basic Research Program of China Grant 2011CBA00300, 2011CBA00301, the National Natural Science Foundation of China Grant 61033001, 61361136003.

generating function or the power series $\tilde{\mathbf{c}}(x)$ is also said to be *Apwenian*. Recall that the *Thue–Morse sequence*, denoted by

$$\mathbf{e} = (e_k)_{k \geq 0} = (1, -1, -1, 1, -1, 1, 1, -1, -1, 1, 1, -1 \dots),$$

is a special ± 1 -sequence [10], defined by the generating function

$$(1.2) \quad \tilde{\mathbf{e}}(x) = \sum_{k=0}^{\infty} e_k x^k = \prod_{k=0}^{\infty} (1 - x^{2^k}),$$

or equivalently, by the recurrence relations

$$(1.3) \quad e_0 = 1, \quad e_{2k} = e_k \text{ and } e_{2k+1} = -e_k \text{ for } k \geq 0.$$

The Thue–Morse sequence is also called *Prouhet–Thue–Morse sequence*. For other equivalent definitions and properties related to the sequence, see [2, 3, 11, 9, 8]. In 1998, Allouche, Peyrière, Wen and Wen established a congruence relation concerning the Hankel determinants of the Thue–Morse sequence [1].

Theorem 1.1 (APWW). *The Thue–Morse sequence on $\{1, -1\}$ is Apwenian.*

Theorem 1.1 has an important application to Number Theory. As a consequence of Theorem 1.1, all the Hankel determinants of the Thue–Morse sequence are nonzero. This property allowed Bugeaud [5] to prove that the irrationality exponents of the Thue–Morse–Mahler numbers are exactly 2.

The goal of the paper is to find more Apwenian sequences. Let d be a positive integer and $\mathbf{v} = (v_0, v_1, v_2, \dots, v_{d-1})$ a finite ± 1 -sequence of length d such that $v_0 = 1$. The generating polynomial of \mathbf{v} is denoted by $\tilde{\mathbf{v}}(x) = \sum_{i=0}^{d-1} v_i x^i$. It is clear that the following power series

$$(1.4) \quad \Phi(\tilde{\mathbf{v}}(x)) = \prod_{k=0}^{\infty} \tilde{\mathbf{v}}(x^{d^k})$$

defines a ± 1 -sequence. Thus, the power series displayed in (1.2) is equal to $\Phi(1 - x)$. Our main result is stated next.

Theorem 1.2. *The following power series are all Apwenian:*

$$\begin{aligned} F_2(x) &= \Phi(1 - x), \\ F_3(x) &= \Phi(1 - x - x^2), \\ F_5(x) &= \Phi(1 - x - x^2 - x^3 + x^4), \\ F_{11}(x) &= \Phi(1 - x - x^2 + x^3 - x^4 + x^5 + x^6 + x^7 + x^8 - x^9 \\ &\quad - x^{10}), \\ F_{13}(x) &= \Phi(1 - x - x^2 + x^3 - x^4 - x^5 - x^6 - x^7 - x^8 + x^9 \\ &\quad - x^{10} - x^{11} + x^{12}), \end{aligned}$$

$$\begin{aligned}
F_{17a}(x) &= \Phi(1 - x - x^2 + x^3 - x^4 + x^5 + x^6 + x^7 + x^8 + x^9 \\
&\quad + x^{10} + x^{11} - x^{12} + x^{13} - x^{14} - x^{15} + x^{16}), \\
F_{17b}(x) &= \Phi(1 - x - x^2 - x^3 + x^4 + x^5 - x^6 + x^7 + x^8 + x^9 \\
&\quad - x^{10} + x^{11} + x^{12} - x^{13} - x^{14} - x^{15} + x^{16}).
\end{aligned}$$

Remarks. Let us make some useful comments about the above Theorem.

- (1) The fact that the generating function $F_2(x)$ for the Thue–Morse sequence is Apwenian has already been proved in [1].
- (2) By using the Jacobi continued fraction expansion of a power series $F(x)$, we know that $H_n(F(x)) = H_n(F(-x))$. See, for example, [15, 7, 12, 13]. Hence, Theorem 1.2 implies that $F_3(-x) = \Phi(1 + x - x^2)$, $F_5(-x) = \Phi(1 + x - x^2 + x^3 + x^4)$, etc. are all Apwenian.
- (3) There is no F_7 in Theorem 1.2, but two F_{17} (we mean F_{17a} and F_{17b}).
- (4) $\Phi(1 - x - x^2 + x^3)$ is Apwenian since it is equal to $\Phi(1 - x)$.

Actually, Theorem 1.1 has three proofs. The original proof of Theorem 1.1 is based on determinant manipulation by using the so-called *sudoku method* [1, 14]. The second one is a combinatorial proof derived by Bugeaud and Han [6]. The third proof is very short and makes use of Jacobi continued fraction algebra [13]. Unfortunately, the method developed in the short proof cannot be used for proving our main theorem, because the underlying Jacobi continued fractions are not ultimately periodic [13, 12]. However, another analogous result for the sequence $F_3(x)$ when dealing with modulo 3 (instead of modulo 2) is established using the short method, as stated in the next theorem [12].

Theorem 1.3. *For every positive integer n the Hankel determinant $H_n(F_3(x))$ of the sequence $F_3(x)$ verifies the following relation*

$$(1.5) \quad H_n(F_3(x)) \equiv \begin{cases} 1 \pmod{3} & \text{if } n \equiv 1, 2 \pmod{4}; \\ 2 \pmod{3} & \text{if } n \equiv 3, 0 \pmod{4}. \end{cases}$$

Combining Theorem 1.3 and Theorem 1.2 yields the following result.

Corollary 1.4. *For every positive integer n the Hankel determinant $H_n(F_3(x))$ verifies the following relation*

$$(1.6) \quad \frac{H_n(F_3(x))}{2^{n-1}} \equiv \begin{cases} 1 \pmod{6} & \text{if } n \equiv 0, 1 \pmod{4}; \\ 5 \pmod{6} & \text{if } n \equiv 2, 3 \pmod{4}. \end{cases}$$

In the following table we reproduce the first few values of the Hankel determinants of the sequence $F_3(x)$ for illustrating Theorems 1.2, 1.3 and Corollary 1.4.

| n | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
|--|---|----|----|---|----|-----|-----|-----|------|-------|
| $H_n(\mathbf{f})$ | 1 | -2 | -4 | 8 | 16 | -32 | -64 | 128 | 4864 | -9728 |
| $H_n(\mathbf{f}) \pmod{3}$ | 1 | 1 | 2 | 2 | 1 | 1 | 2 | 2 | 1 | 1 |
| $H_n(\mathbf{f})/2^{n-1}$ | 1 | -1 | -1 | 1 | 1 | -1 | -1 | 1 | 19 | -19 |
| $\frac{H_n(\mathbf{f})}{2^{n-1}} \pmod{2}$ | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 | 1 |
| $\frac{H_n(\mathbf{f})}{2^{n-1}} \pmod{6}$ | 1 | 5 | 5 | 1 | 1 | 5 | 5 | 1 | 1 | 5 |

Recently, Bugeaud and Han re-proved Theorem 1.1 by means of an appropriate combinatorial method [6]. The latter method has been significantly upgraded to prove that $F_3(x)$ is Apwenian. As can be seen, in Section 3 Step 2, a family of cases (called *types*) is considered for proving the various recurrence relations. Roughly speaking, the types are indexed by words $s_0s_1s_2 \cdots s_d$ of length $d+1$ over a d -letter alphabet. Comparing to the original combinatorial method, the upgrading does not provide a shorter proof; however, it involves of a systematic proof *by exhaustion* that only consists of checking all the types. The proof of Theorem 1.2 is then achieved with *Computer Assistance*.

In practice, the number of types is very large. For example, as described in §2.3 for the study of $F_{11}(x)$, there are 2274558 types! Fortunately, the set of permutations of each type can be decomposed into the Cartesian product of so-called *atoms* (see Substep 3(d) in the sequel), and moreover, the cardinality of each atom can be rapidly evaluated by a sequence of tests (see Definition 4.1 and Table 4.1).

Problem 1.5. *Is the following power series Apwenian:*

$$F_{19}(x) = \Phi(1 - x - x^2 - x^3 + x^4 - x^5 + x^6 - x^7 - x^8 + x^9 \\ + x^{10} - x^{11} - x^{12} - x^{13} - x^{14} - x^{15} + x^{16} - x^{17} - x^{18}) ?$$

Find a fast computer assisted proof for Theorem 1.2 to answer the above question.

For proving that $F_{17a}(x)$ is Apwenian, our *C* program has taken about one week by using 24 CPU cores. No hope for $F_{19}(x)$.

Problem 1.6. *Find a human proof of Theorem 1.2 without computer assistance.*

Problem 1.7. *Characterize all the finite ± 1 -sequences \mathbf{v} such that $\Phi(\tilde{\mathbf{v}}(x))$ is Apwenian.*

As an application of Theorem 1.2 in Number Theory, the irrationality exponents of $F_5(1/b)$, $F_{11}(1/b)$, $F_{17a}(1/b)$, $F_{17b}(1/b)$ are proved to be equal to 2 (see [4]).

2. PROOF OF THEOREM 1.2

Let d be a positive integer and $\mathbf{v} = (v_0, v_1, v_2, \dots, v_{d-1})$ be a finite ± 1 -sequence of length d with $v_0 = 1$. Let $\mathbf{f} = (f_k)_{k \geq 0}$ be the ± 1 -sequence defined by the following generating function

$$(2.1) \quad \tilde{\mathbf{f}}(x) = \Phi(\tilde{\mathbf{v}}(x)) = \prod_{k=0}^{\infty} \tilde{\mathbf{v}}(x^{d^k}),$$

where $\tilde{\mathbf{v}}(x) = \sum_{i=0}^{d-1} v_i x^i$. The above power series satisfies the following functional equation

$$(2.2) \quad \tilde{\mathbf{f}}(x) = \tilde{\mathbf{v}}(x) \prod_{k=1}^{\infty} \tilde{\mathbf{v}}(x^{d^k}) = \tilde{\mathbf{v}}(x) \tilde{\mathbf{f}}(x^d).$$

The sequence \mathbf{f} can also be defined by the recurrence relations

$$(2.3) \quad f_0 = 1, \quad f_{dn+i} = v_i f_n \text{ for } n \geq 0 \text{ and } 0 \leq i \leq d-1.$$

We divide the set $\{1, 2, \dots, d-1\}$ into two disjoint subsets

$$P = \{1 \leq i \leq d-1 \mid v_{i-1} \neq v_i\},$$

$$Q = \{1 \leq i \leq d-1 \mid v_{i-1} = v_i\}.$$

Two disjoint infinite sets of integers J and K play an important role in the proof of Theorem 1.2.

Definition 2.1. If $v_{d-1} = -1$, define

$$J = \{(dn + p)d^{2k} - 1 \mid n, k \in \mathbb{N}, p \in P\}$$

$$\bigcup \{(dn + q)d^{2k+1} - 1 \mid n, k \in \mathbb{N}, q \in Q\},$$

$$K = \{(dn + q)d^{2k} - 1 \mid n, k \in \mathbb{N}, q \in Q\}$$

$$\bigcup \{(dn + p)d^{2k+1} - 1 \mid n, k \in \mathbb{N}, p \in P\}.$$

If $v_{d-1} = 1$, define

$$J = \{(dn + p)d^k - 1 \mid n, k \in \mathbb{N}, p \in P\},$$

$$K = \{(dn + q)d^k - 1 \mid n, k \in \mathbb{N}, q \in Q\}.$$

From the above definition it is easy to see that $N = J \cup K$.

Lemma 2.1. For each $t \geq 0$ the integer $\delta_t := |(f_t - f_{t+1})/2|$ is equal to 1 if and only if t is in J .

Proof. Let $t = (dn + \ell)d^k - 1$. By (2.3) we have

$$f_{(dn+\ell)d^k-1} = f_{d[(dn+\ell)d^{k-1}-1]+(d-1)} = v_{d-1} f_{(dn+\ell)d^{k-1}-1},$$

and

$$f_{(dn+\ell)d^k-1} = v_{d-1} f_{(dn+\ell)d^{k-1}-1} = \dots = v_{d-1}^k f_{dn+\ell-1} = v_{d-1}^k v_{\ell-1} f_n.$$

In the same manner,

$$f_{(dn+\ell)d^k} = f_{(dn+\ell)d^{k-1}} = \dots = f_{dn+\ell} = v_{\ell} f_n.$$

Hence,

$$\delta_t = \left| \frac{1}{2}(f_{(dn+\ell)d^{k-1}} - f_{(dn+\ell)d^k}) \right| = \left| \frac{1}{2}(v_\ell - v_{d-1}^k v_{\ell-1}) \right|$$

is odd if and only if

$$(2.4) \quad v_{d-1}^k \cdot v_{\ell-1} v_\ell = -1.$$

There are two cases are to be considered: (i) If $v_{d-1} = 1$, condition (2.4) is equivalent to $v_{\ell-1} v_\ell = -1$, or $\ell \in P$, or $t \in J$ by Definition 2.1. (ii) If $v_{d-1} = -1$, condition (2.4) becomes $v_{\ell-1} v_\ell = (-1)^{k+1}$, which is equivalent to $\ell \in P$ when k is even and $\ell \in Q$ when k is odd. In other words, $t \in J$. \square

Let $\mathfrak{S}_m = \mathfrak{S}_{\{0,1,\dots,m-1\}}$ be the set of all permutations on $\{0, 1, \dots, m-1\}$. The following Theorem may be viewed as the combinatorial interpretation of Theorem 1.2.

Theorem 2.2. *Let \mathbf{v} be a ± 1 -sequence of length d with $v_0 = 1$. The sequence \mathbf{f} and the set J associated with \mathbf{v} are defined by (2.1) and Definition 2.1 respectively. Then, the sequence \mathbf{f} is Apwenian if, and only if, the number of permutations $\sigma \in \mathfrak{S}_m$ such that $i + \sigma(i) \in J$ for $i = 0, 1, \dots, m-2$ (no constraint on $m-1 + \sigma(m-1) \in N$) is an odd integer for every integer $m \geq 1$.*

Proof. Let m be a positive integer. By means of elementary transformations the Hankel determinant $H_m(\mathbf{f})$ is equal to

$$\begin{aligned} H_m(\mathbf{f}) &= \begin{vmatrix} f_0 & f_1 & \cdots & f_{m-1} \\ f_1 & f_2 & \cdots & f_m \\ \vdots & \vdots & \ddots & \vdots \\ f_{m-1} & f_m & \cdots & f_{2m-2} \end{vmatrix} \\ &= \begin{vmatrix} f_0 - f_1 & f_1 - f_2 & \cdots & f_{m-2} - f_{m-1} & f_{m-1} \\ f_1 - f_2 & f_2 - f_3 & \cdots & f_{m-1} - f_m & f_m \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ f_{m-1} - f_m & f_m - f_{m+1} & \cdots & f_{2m-3} - f_{2m-2} & f_{2m-2} \end{vmatrix} \\ &= 2^{m-1} \times \begin{vmatrix} \frac{f_0-f_1}{2} & \frac{f_1-f_2}{2} & \cdots & \frac{f_{m-2}-f_{m-1}}{2} & f_{m-1} \\ \frac{f_1-f_2}{2} & \frac{f_2-f_3}{2} & \cdots & \frac{f_{m-1}-f_m}{2} & f_m \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ \frac{f_{m-1}-f_m}{2} & \frac{f_m-f_{m+1}}{2} & \cdots & \frac{f_{2m-3}-f_{2m-2}}{2} & f_{2m-2} \end{vmatrix}. \end{aligned}$$

By Lemma 2.1, we have

$$(2.5) \quad \frac{H_m(\mathbf{f})}{2^{m-1}} \equiv \begin{vmatrix} \delta_0 & \delta_1 & \cdots & \delta_{m-2} & 1 \\ \delta_1 & \delta_2 & \cdots & \delta_{m-1} & 1 \\ \vdots & \vdots & \ddots & \vdots & \vdots \\ \delta_{m-1} & \delta_m & \cdots & \delta_{2m-3} & 1 \end{vmatrix} \pmod{2}.$$

By the very definition of a determinant or the Leibniz formula, the determinant occurring on the right-hand side of the congruence (2.5) is equal to

$$(2.6) \quad S_m := \sum_{\sigma \in \mathfrak{S}_m} (-1)^{\text{inv}(\sigma)} \delta_{0+\sigma_0} \delta_{1+\sigma_1} \cdots \delta_{m-2+\sigma_{m-2}},$$

where $\text{inv}(\sigma)$ is the number of inversions of the permutation σ . By Lemma 2.1 the product $\delta_{0+\sigma_0} \delta_{1+\sigma_1} \cdots \delta_{m-2+\sigma_{m-2}}$ is equal to 1 if $i + \sigma_i \in J$ for $i = 0, 1, \dots, m-2$, and to 0 otherwise. Hence, the summation S_m is congruent modulo 2 to the number of permutations $\sigma \in \mathfrak{S}_m$ such that $i + \sigma_i \in J$ for all $i = 0, 1, \dots, m-2$. Hence, \mathbf{f} is Apwenian if and only if the number of permutations $\sigma \in \mathfrak{S}_m$ such that $i + \sigma(i) \in J$ for $i = 0, 1, \dots, m-2$ (no constraint on $m-1 + \sigma(m-1) \in N$) is an odd integer for every integer $m \geq 1$. \square

For proving that the sequence \mathbf{f} is Apwenian by means of Theorem 2.2, it is convenient to introduce the following notations.

Definition 2.2. For $m \geq \ell \geq 0$ let $\mathfrak{J}_{m,\ell}$ (resp. $\mathfrak{K}_{m,\ell}$) be the set of all permutations $\sigma = \sigma_0 \sigma_1 \cdots \sigma_{m-1} \in \mathfrak{S}_m$ such that $i + \sigma_i \in J$ (resp. $i + \sigma_i \in K$) for $i \in \{0, 1, \dots, m-1\} \setminus \{\ell\}$. Let $n \geq 1$; for simplicity, write:

$$\begin{aligned} j_{m,\ell} &:= \#\mathfrak{J}_{m,\ell}, & k_{m,\ell} &:= \#\mathfrak{K}_{m,\ell}, \\ X_n &:= \sum_{i=0}^{n-1} j_{n,i}, & Y_n &:= j_{n,n}, & Z_n &:= j_{n,n-1}, \\ U_n &:= \sum_{i=0}^{n-1} k_{n,i}, & V_n &:= k_{n,n}, & W_n &:= k_{n,n-1}, \\ T_n &:= X_n + X_n Y_n + Y_n, \\ R_n &:= U_n + U_n V_n + V_n. \end{aligned}$$

Notice that if $\ell = m$, then $\{0, 1, \dots, m-1\} \setminus \{\ell\} = \{0, 1, \dots, m-1\}$, so that $j_{m,m}$ (resp. $k_{m,m}$) is the number of permutations $\sigma \in \mathfrak{S}_m$ such that $i + \sigma(i) \in J$ (resp. $\in K$) for all i .

By Theorem 2.2 and Definition 2.2 the sequence \mathbf{f} is Apwenian if and only if $Z_n \equiv 1 \pmod{2}$. In Section 4 we describe an algorithm enabling us to *find* and also *prove* a list of recurrence relations between $X_n, Y_n, X_n, U_n, V_n, W_n$. Then, it is routine to check whether $Z_n \equiv 1 \pmod{2}$ or not. Our program `Apwen.py` is an implementation of the latter algorithm in Python.

We now produce the proof of Theorem 1.2 by means of the program `Apwen.py`. Since $F_2(x)$ has been proved to be Apwenian in [1], only the three power series $F_3(x), F_5(x)$ and $F_{11}(x)$ require our attention. We can also prove that $F_{13}(x), F_{17a}(x), F_{17b}(x)$ are Apwenian in the same

manner. However, the full proofs are lengthy and are not reproduced in the paper.

2.1. $F_3(x)$ is Apwenian. Take $\mathbf{v} = (1, -1, -1)$ with $d = 3$ and $v_{d-1} = -1$. Then, the corresponding infinite ± 1 -sequence \mathbf{f} is equal to $F_3(x)$. We have $P = \{1\}$, $Q = \{2\}$ and

$$\begin{aligned} J &= \{(3n+1)3^{2k} - 1 \mid n, k \in \mathbb{N}\} \cup \{(3n+2)3^{2k+1} - 1 \mid n, k \in \mathbb{N}\} \\ &= \{0, 3, 5, 6, 8, 9, 12, 14, 15, 18, \dots\}, \end{aligned}$$

$$\begin{aligned} K &= \{(3n+2)3^{2k} - 1 \mid n, k \in \mathbb{N}\} \cup \{(3n+1)3^{2k+1} - 1 \mid n, k \in \mathbb{N}\} \\ &= \{1, 2, 4, 7, 10, 11, 13, 16, 17, \dots\} = \mathbb{N} \setminus J. \end{aligned}$$

By enumerating a list of 24 *types* of permutations (see Section 3), the program `Apwen.py` finds and proves the following recurrences.

Lemma 2.3. *For each $n \geq 1$ we have*

$$\begin{aligned} X_{3n+0} &\equiv U_n, & Y_{3n+0} &\equiv U_n + V_n, \\ X_{3n+1} &\equiv W_{n+1}(U_n + V_n), & Y_{3n+1} &\equiv W_{n+1}V_n, \\ X_{3n+2} &\equiv W_{n+1}(U_{n+1} + V_{n+1}), & Y_{3n+2} &\equiv W_{n+1}V_{n+1}, \\ \\ Z_{3n+0} &\equiv W_n(U_n + U_nV_n + V_n), \\ Z_{3n+1} &\equiv W_{n+1}(U_n + U_nV_n + V_n), \\ Z_{3n+2} &\equiv W_{n+1}. \end{aligned}$$

As explained in Section 3, the above relations express X, Y, Z in function of U, V, W since $v_{d-1} = -1$. By exchanging the values of P and Q , J and K , the program `Apwen.py` yields other relations which express U, V, W in terms of X, Y, Z by enumerating a list of 26 types of permutations.

Lemma 2.4. *For each $n \geq 1$ we have*

$$\begin{aligned} U_{3n+0} &\equiv X_n, & V_{3n+0} &\equiv X_n + Y_n, \\ U_{3n+1} &\equiv Z_{n+1}Y_n, & V_{3n+1} &\equiv Z_{n+1}X_n, \\ U_{3n+2} &\equiv Z_{n+1}Y_{n+1}, & V_{3n+2} &\equiv Z_{n+1}X_{n+1}, \\ \\ W_{3n+0} &\equiv Z_n(X_n + X_nY_n + Y_n), \\ W_{3n+1} &\equiv Z_{n+1}(X_n + X_nY_n + Y_n), \\ W_{3n+2} &\equiv Z_{n+1}. \end{aligned}$$

From Lemmas 2.3 and 2.4 we obtain the following ‘‘simplified’’ recurrence relations based on some elementary calculations.

Corollary 2.5. *For each positive integer n we have*

$$Z_{3n+0} \equiv W_n R_n, \quad W_{3n+0} \equiv Z_n T_n,$$

$$\begin{aligned}
Z_{3n+1} &\equiv W_{n+1}R_n, & W_{3n+1} &\equiv Z_{n+1}T_n, \\
Z_{3n+2} &\equiv W_{n+1}, & W_{3n+2} &\equiv Z_{n+1}, \\
T_{3n+0} &\equiv R_n, & R_{3n+0} &\equiv T_n, \\
T_{3n+1} &\equiv W_{n+1}R_n, & R_{3n+1} &\equiv Z_{n+1}T_n, \\
T_{3n+2} &\equiv W_{n+1}R_{n+1}, & R_{3n+2} &\equiv Z_{n+1}T_{n+1}.
\end{aligned}$$

Since $Z_1 = 1, T_1 = 3, W_1 = 1, R_1 = 1, Z_2 = 1, T_2 = 1, W_2 = 1$ and $R_2 = 7$, Corollary 2.5 yields $Z_m \equiv T_m \equiv W_m \equiv R_m \equiv 1 \pmod{2}$ for every positive integer m by induction. Hence, $F_3(x)$ is Apwenian.

2.2. $F_5(x)$ is Apwenian. Take $\mathbf{v} = (1, -1, -1, -1, 1)$ with $d = 5$ and $v_{d-1} = 1$. Then, the corresponding infinite ± 1 -sequence \mathbf{f} is equal to $F_5(x)$. We have

$$\begin{aligned}
P &= \{1, 4\}, \\
Q &= \{2, 3\}, \\
J &= \{(5n+1)5^k - 1 \mid n, k \in \mathbb{N}\} \cup \{(5n+4)5^k - 1 \mid n, k \in \mathbb{N}\} \\
&= \{0, 3, 4, 5, 8, 10, 13, 15, 18, 19, 20, 23, 24, 25, 28, 29, 30, 33, \dots\}, \\
K &= \{(5n+2)5^k - 1 \mid n, k \in \mathbb{N}\} \cup \{(5n+3)5^k - 1 \mid n, k \in \mathbb{N}\} \\
&= \{1, 2, 6, 7, 9, 11, 12, 14, 16, 17, 21, 22, 26, 27, 31, 32, 34, 36, \dots\}.
\end{aligned}$$

By enumerating a list of 225 *types* of permutations, the Python program `Apwen.py` finds and proves the following recurrences.

Lemma 2.6. *For each $n \geq 1$ we have*

$$\begin{aligned}
X_{5n+0} &\equiv X_n, & Y_{5n+0} &\equiv Y_n, \\
X_{5n+1} &\equiv Z_{n+1}Y_n, & Y_{5n+1} &\equiv Z_{n+1}(X_n + Y_n), \\
X_{5n+2} &\equiv Z_{n+1}(X_n + Y_n), & Y_{5n+2} &\equiv Z_{n+1}X_n, \\
X_{5n+3} &\equiv Z_{n+1}(X_{n+1} + Y_{n+1}), & Y_{5n+3} &\equiv Z_{n+1}X_{n+1}, \\
X_{5n+4} &\equiv Z_{n+1}Y_{n+1}, & Y_{5n+4} &\equiv Z_{n+1}(X_{n+1} + Y_{n+1}), \\
Z_{5n+0} &\equiv Z_n(X_n + X_nY_n + Y_n), \\
Z_{5n+1} &\equiv Z_{n+1}(X_n + X_nY_n + Y_n), \\
Z_{5n+2} &\equiv Z_{n+1}(X_n + X_nY_n + Y_n), \\
Z_{5n+3} &\equiv Z_{n+1}, \\
Z_{5n+4} &\equiv Z_{n+1}(X_{n+1} + X_{n+1}Y_{n+1} + Y_{n+1}).
\end{aligned}$$

As explained in Section 4, the above relations are between X, Y, Z without involving U, V, W , since $v_{d-1} = 1$. We obtain the following “simplified” recurrence relations based on some elementary calculations.

Corollary 2.7. *For each positive integer n we have*

$$Z_{5n+0} \equiv Z_n T_n, \quad T_{5n+0} \equiv T_n,$$

$$\begin{aligned}
Z_{5n+1} &\equiv Z_{n+1}T_n, & T_{5n+1} &\equiv Z_{n+1}T_n, \\
Z_{5n+2} &\equiv Z_{n+1}T_n, & T_{5n+2} &\equiv Z_{n+1}T_n, \\
Z_{5n+3} &\equiv Z_{n+1}, & T_{5n+3} &\equiv Z_{n+1}T_{n+1}, \\
Z_{5n+4} &\equiv Z_{n+1}T_{n+1}, & T_{5n+4} &\equiv Z_{n+1}T_{n+1}.
\end{aligned}$$

Since $Z_1 = 1, T_1 = 3, Z_2 = 1, T_2 = 1, Z_3 = 1, T_3 = 9, Z_4 = 5, T_4 = 129$, Corollary 2.7 yields $Z_n \equiv T_n \equiv 1 \pmod{2}$ for each $n \geq 1$ by induction. Hence, $F_5(x)$ is Apwenian.

2.3. $F_{11}(x)$ is Apwenian. Take

$$\mathbf{v} = (1, -1, -1, 1, -1, 1, 1, 1, 1, -1, -1)$$

with $d = 11$ and $v_{d-1} = -1$. Then, the corresponding infinite ± 1 -sequence \mathbf{f} is equal to $F_{11}(x)$. We have

$$\begin{aligned}
P &= \{1, 3, 4, 5, 9\}, \\
Q &= \{2, 6, 7, 8, 10\}, \\
J &= \{0, 2, 3, 4, 8, 11, 13, 14, 15, 19, 21, 22, 24, 25, 26, 30, 33, 35, \dots\}, \\
K &= \{1, 5, 6, 7, 9, 10, 12, 16, 17, 18, 20, 23, 27, 28, 29, 31, 32, 34, \dots\}.
\end{aligned}$$

By enumerating a list of 2274558 *types* of permutations, the program `Apwen.py` finds and proves the following recurrences.

Lemma 2.8. *For each $n \geq 1$ we have*

$$\begin{aligned}
X_{11n+0} &\equiv U_n, & Y_{11n+0} &\equiv U_n + V_n, \\
X_{11n+1} &\equiv W_{n+1}(V_n + U_n), & Y_{11n+1} &\equiv V_n W_{n+1}, \\
X_{11n+2} &\equiv U_n W_{n+1}, & Y_{11n+2} &\equiv W_{n+1}(V_n + U_n), \\
X_{11n+3} &\equiv W_{n+1}(V_n + U_n), & Y_{11n+3} &\equiv V_n W_{n+1}, \\
X_{11n+4} &\equiv V_n W_{n+1}, & Y_{11n+4} &\equiv U_n W_{n+1}, \\
X_{11n+5} &\equiv U_n W_{n+1}, & Y_{11n+5} &\equiv W_{n+1}(V_n + U_n), \\
X_{11n+6} &\equiv U_{n+1} W_{n+1}, & Y_{11n+6} &\equiv W_{n+1}(U_{n+1} + V_{n+1}), \\
X_{11n+7} &\equiv V_{n+1} W_{n+1}, & Y_{11n+7} &\equiv U_{n+1} W_{n+1}, \\
X_{11n+8} &\equiv W_{n+1}(U_{n+1} + V_{n+1}), & Y_{11n+8} &\equiv V_{n+1} W_{n+1}, \\
X_{11n+9} &\equiv U_{n+1} W_{n+1}, & Y_{11n+9} &\equiv W_{n+1}(U_{n+1} + V_{n+1}), \\
X_{11n+10} &\equiv W_{n+1}(U_{n+1} + V_{n+1}), & Y_{11n+10} &\equiv V_{n+1} W_{n+1}, \\
\\
Z_{11n+0} &\equiv W_n(V_n + U_n + U_n V_n), \\
Z_{11n+i} &\equiv W_{n+1}(U_n V_n + V_n + U_n), & (i = 1, 2, 3, 4, 5) \\
Z_{11n+6} &\equiv W_{n+1}, \\
Z_{11n+i} &\equiv W_{n+1}(U_{n+1} + U_{n+1} V_{n+1} + V_{n+1}). & (i = 7, 8, 9, 10)
\end{aligned}$$

As explained in Section 4, the above relations express X, Y, Z in function of U, V, W , since $v_{d-1} = -1$. By exchanging the values of P

and Q , J and K , the program `Apwen.py` yields another list of relations which express U, V, W in terms of X, Y, Z . For this purpose, a long list of 2350964 *types* of permutations are enumerated.

Lemma 2.9. *For each $n \geq 1$ we have*

$$\begin{aligned}
 U_{11n+0} &\equiv X_n, & V_{11n+0} &\equiv X_n + Y_n, \\
 U_{11n+1} &\equiv Y_n Z_{n+1}, & V_{11n+1} &\equiv X_n Z_{n+1}, \\
 U_{11n+2} &\equiv X_n Z_{n+1}, & V_{11n+2} &\equiv Z_{n+1}(X_n + Y_n), \\
 U_{11n+3} &\equiv Y_n Z_{n+1}, & V_{11n+3} &\equiv X_n Z_{n+1}, \\
 U_{11n+4} &\equiv Z_{n+1}(X_n + Y_n), & V_{11n+4} &\equiv Y_n Z_{n+1}, \\
 U_{11n+5} &\equiv X_n Z_{n+1}, & V_{11n+5} &\equiv Z_{n+1}(X_n + Y_n), \\
 U_{11n+6} &\equiv X_{n+1} Z_{n+1}, & V_{11n+6} &\equiv Z_{n+1}(Y_{n+1} + X_{n+1}), \\
 U_{11n+7} &\equiv Z_{n+1}(Y_{n+1} + X_{n+1}), & V_{11n+7} &\equiv Y_{n+1} Z_{n+1}, \\
 U_{11n+8} &\equiv Y_{n+1} Z_{n+1}, & V_{11n+8} &\equiv X_{n+1} Z_{n+1}, \\
 U_{11n+9} &\equiv X_{n+1} Z_{n+1}, & V_{11n+9} &\equiv Z_{n+1}(Y_{n+1} + X_{n+1}), \\
 U_{11n+10} &\equiv Y_{n+1} Z_{n+1}, & V_{11n+10} &\equiv X_{n+1} Z_{n+1}, \\
 \\
 W_{11n+0} &\equiv Z_n(Y_n + X_n + X_n Y_n), \\
 W_{11n+i} &\equiv Z_{n+1}(X_n Y_n + X_n + Y_n), & (i = 1, 2, 3, 4, 5) \\
 W_{11n+6} &\equiv Z_{n+1}, \\
 W_{11n+i} &\equiv Z_{n+1}(Y_{n+1} + X_{n+1} Y_{n+1} + X_{n+1}). & (i = 7, 8, 9, 10)
 \end{aligned}$$

From Lemmas 2.8 and 2.9 we obtain the following “simplified” recurrence relations based on some elementary calculations.

Corollary 2.10. *For each positive integer n we have*

$$\begin{aligned}
 T_{11n+0} &\equiv R_n, & R_{11n+0} &\equiv T_n, \\
 T_{11n+i} &\equiv R_n W_{n+1}, & R_{11n+i} &\equiv T_n Z_{n+1}, & (1 \leq i \leq 5) \\
 T_{11n+i} &\equiv R_{n+1} W_{n+1}, & R_{11n+i} &\equiv T_{n+1} Z_{n+1}, & (6 \leq i \leq 10) \\
 Z_{11n+0} &\equiv R_n W_n, & W_{11n+0} &\equiv T_n Z_n, \\
 Z_{11n+i} &\equiv R_n W_{n+1}, & W_{11n+i} &\equiv T_n Z_{n+1}, & (1 \leq i \leq 5) \\
 Z_{11n+6} &\equiv W_{n+1}, & W_{11n+6} &\equiv Z_{n+1}, \\
 Z_{11n+i} &\equiv R_{n+1} W_{n+1}, & W_{11n+i} &\equiv T_{n+1} Z_{n+1}. & (7 \leq i \leq 10)
 \end{aligned}$$

The first values of Z_m, T_m, W_m, R_m are reproduced in the following table.

| m | 1 | 2 | 3 | 4 | 5 | 6 | 7 | 8 | 9 | 10 |
|-------|---|---|----|-----|-----|------|--------|---------|----------|------------|
| Z_m | 1 | 1 | 3 | 11 | 13 | 25 | 39 | 117 | 739 | 4431 |
| T_m | 3 | 5 | 47 | 237 | 487 | 419 | 3503 | 66905 | 3527039 | 82080975 |
| W_m | 1 | 1 | 1 | 1 | 5 | 25 | 177 | 1091 | 3839 | 19791 |
| R_m | 1 | 5 | 1 | 11 | 107 | 5151 | 198769 | 4802755 | 56576127 | 2717644635 |

Corollary 2.10 yields $Z_m \equiv T_m \equiv W_m \equiv R_m \equiv 1 \pmod{2}$ for every positive integer m by induction. Hence, $F_{11}(x)$ is Apwenian.

3. ALGORITHM FOR FINDING THE RECURRENCES

Keep the same *notations* as in Section 2. We will show how to *find* and also *prove* a list of recurrence relations between the quantities $X_n, Y_n, Z_n, U_n, V_n, W_n$. The set N of nonnegative integers is partitioned into d disjoint subsets A_0, A_1, \dots, A_{d-1} according to the value modulo d :

$$(3.1) \quad A_i = \{dn + i \mid n \in N\} \quad (i = 0, 1, \dots, d-1).$$

For an infinite set S let $S|_m$ be the set composed of the m smallest integers in S . Let $\beta : N \rightarrow N$ denote the transformation $k \mapsto \lfloor \frac{k}{d} \rfloor$. In other words,

$$(3.2) \quad \beta(k) = (k - i)/d \quad \text{if } k \in A_i.$$

For simplicity, write

$$\bar{J} = \begin{cases} J & \text{if } v_{d-1} = 1, \\ K & \text{if } v_{d-1} = -1, \end{cases} \quad \text{and} \quad \bar{K} = \begin{cases} K & \text{if } v_{d-1} = 1, \\ J & \text{if } v_{d-1} = -1. \end{cases}$$

Then $\bar{\mathfrak{J}}_{m,\ell}, \bar{X}_n, \bar{Y}_n, \bar{Z}_n$ mean $\mathfrak{J}_{m,\ell}, X_n, Y_n, Z_n$ (resp. $\mathfrak{K}_{m,\ell}, U_n, V_n, W_n$) if $v_{d-1} = 1$ (resp. $v_{d-1} = -1$).

Lemma 3.1. *For each $p \in P$ and $q \in Q$ we have*

- (i) $A_{p-1} \subset J$ and $A_{q-1} \subset K$;
- (ii) $A_{q-1} \cap J = \emptyset$ and $A_{p-1} \cap K = \emptyset$;
- (iii) $\beta(A_{d-1} \cap J) = \bar{J}$ and $\beta(A_{d-1} \cap K) = \bar{K}$.

Proof. (i) By the definition of the sets J and K we have

$$J \supset \{(dn + p)d^{2k} - 1 \mid n, k \in N\} \supset \{(dn + p)d^0 - 1 \mid n \in N\} = A_{p-1};$$

$$K \supset \{(dn + q)d^{2k} - 1 \mid n, k \in N\} \supset \{(dn + q)d^0 - 1 \mid n \in N\} = A_{q-1}.$$

(ii) By (i) and the relation $K \cap J = \emptyset$.

(iii) By Definition 2.1 with $v_{d-1} = -1$, we have

$$A_{d-1} \cap J = \{(dn + p)d^{2k+2} - 1 \mid n, k \in N, p \in P\}$$

$$\bigcup \{(dn + q)d^{2k+1} - 1 \mid n, k \in N, q \in Q\}.$$

Thus,

$$\beta(A_{d-1} \cap J) = \{(dn + p)d^{2k+1} - 1 \mid n, k \in N, p \in P\}$$

$$\bigcup \{(dn + q)d^{2k} - 1 \mid n, k \in N, q \in Q\}.$$

$$= K = \bar{J}.$$

If $v_{d-1} = 1$, then $A_{d-1} \cap J = \{(dn + p)d^{k+1} - 1 \mid n, k \in N, p \in P\}$. Thus, $\beta(A_{d-1} \cap J) = \{(dn + p)d^k - 1 \mid n, k \in N, p \in P\} = J = \bar{J}$. The second part $\beta(A_{d-1} \cap K) = \bar{K}$ is proved in the same manner. \square

Let $0 \leq i, j \leq d-1$ and $x \in A_i, y \in A_j$. For determining the condition of i and j such that the sum $x + y$ belongs to J or K , there are three cases to be considered.

- (S1) If $i + j + 1 \pmod{d} \in P$, then, $x + y \in J$;
- (S2) If $i + j + 1 \pmod{d} \in Q$, then, $x + y \in K$;
- (S3) If $i + j + 1 \pmod{d} = 0$, then, $x + y \in A_{d-1}$. In this case, the sum $x + y$ may belong to J or K .

Let $m \geq \ell \geq 0$. We want to enumerate the permutations in $\mathfrak{J}_{m,\ell}$ modulo 2. Each permutation $\sigma = \sigma_0\sigma_1 \cdots \sigma_{m-1} \in \mathfrak{S}_m$ may be written as the *two-line representation*

$$\begin{pmatrix} 0 & 1 & 2 & \cdots & m-1 \\ \sigma_0 & \sigma_1 & \sigma_2 & \cdots & \sigma_{m-1} \end{pmatrix}.$$

The columns $\binom{i}{\sigma_i}$ are called *biletters*. For each $\sigma \in \mathfrak{J}_{m,\ell}$ a billetter $\binom{i}{\sigma_i}$ in σ is said to be of (*normal*) *form* $\binom{a_j}{a_k}$ (resp. *specific form* $\binom{\ell}{a_k}$) if $i \neq \ell$ and $(i, \sigma_i) \in A_j \times A_k$ (resp. $i = \ell$ and $\sigma_i \in A_k$). To count the permutations from $\mathfrak{J}_{m,\ell}$ modulo 2, we proceed in several steps. In most cases the calculations are illustrated with $d = 5$.

Step 1. Occurrences of biletters. Since we want to enumerate permutations modulo 2, we can delete suitable pairs of the permutations and the result will not be changed. Let $i \in N|_d$, if a permutation $\sigma \in \mathfrak{J}_{m,\ell}$ contains more than two biletters of form $\binom{a_i}{a_j}$ such that $i + j + 1 \pmod{d} \in P$, select the first two such biletters $\binom{i_1}{j_1}$ and $\binom{i_2}{j_2}$. We define another permutation τ obtained from σ by exchanging j_1 and j_2 in the bottom line. This procedure is reversible. By (S1), it is easy to verify that τ is also in $\mathfrak{J}_{m,\ell}$, so that we can delete the pair σ and τ . Then, there only remain the permutations containing 0 or 1 billetter of form $\binom{a_i}{a_j}$ such that $i + j + 1 \pmod{d} \in P$.

Let $\mathfrak{J}'_{m,\ell}$ be the set of permutations $\sigma \in \mathfrak{J}_{m,\ell}$ which, for each $i \in N|_d$, contains 0 or 1 billetter of form $\binom{a_i}{a_j}$ such that $i + j + 1 \pmod{d} \in P$. We have $j_{m,\ell} = \#\mathfrak{J}_{m,\ell} \equiv \#\mathfrak{J}'_{m,\ell} \pmod{2}$. By (S2), each permutation $\sigma \in \mathfrak{J}'_{m,\ell}$ does not contain any billetter of form $\binom{a_i}{a_j}$ such that $i + j + 1 \pmod{d} \in Q$. Thus, most of the biletters are of form $\binom{a_i}{a_{d-i-1}}$. In conclusion, the number of occurrences of each form is summarized in Table 3.1. A billetter of form $\binom{a_i}{a_j}$ such that $i + j + 1 \pmod{d} \in P$ is said to be *unsociable*. A billetter of form $\binom{a_j}{a_{d-j-1}}$ is said to be *friendly*. By Table 3.1, each permutation in $\mathfrak{J}'_{m,\ell}$ contains only a few unsociable biletters. *Step 2. Form and type.* The two-line representation of a permutation can be seen as a word of biletters. In fact, the order of the biletters does not matter. Let $m \geq 2d$. The *form* $f(\sigma)$ of a permutation $\sigma \in \mathfrak{J}'_{m,\ell}$ is obtained from σ by replacing each billetter of σ by its (normal or specific) form. From Table 3.1, the form $f(\sigma)$ of a

| form | total times |
|---|-----------------|
| $\left\{ \binom{a_k}{a_j} \mid k + j + 1 \pmod{d} \in Q \right\}$ | 0 |
| $\left\{ \binom{a_i}{a_j} \mid i + j + 1 \pmod{d} \in P \right\}$ for each $i \in N _d$ | 0, 1 |
| $\left\{ \binom{a_j}{a_{d-j-1}} \mid j \in N _d \right\}$ | 0, 1, 2, 3, ... |
| $\left\{ \binom{\ell}{a_j} \mid j \in N _d \right\}$ ($\ell = m$) | 0 |
| $\left\{ \binom{\ell}{a_j} \mid j \in N _d \right\}$ ($0 \leq \ell \leq m - 1$) | 1 |

TABLE 3.1. Number of occurrences of biletters

permutation $\sigma \in \mathfrak{J}'_{m,\ell}$ is

$$(3.3) \quad \left(\begin{array}{ccc|ccc|ccc} a_0 & a_0 & a_0 & a_1 & a_1 & a_1 & \dots & a_{d-1} & a_{d-1} & a_{d-1} \\ a_{d-1} & a_{d-1} & s_0 & a_{d-2} & a_{d-2} & s_1 & \dots & a_0 & a_0 & s_{d-1} \end{array} \right),$$

for $\ell = m$, or

$$(3.4) \quad \left(\begin{array}{ccc|ccc|ccc|c} a_0 & a_0 & a_0 & a_1 & a_1 & a_1 & \dots & a_{d-1} & a_{d-1} & a_{d-1} & \ell \\ a_{d-1} & a_{d-1} & s_0 & a_{d-2} & a_{d-2} & s_1 & \dots & a_0 & a_0 & s_{d-1} & s_d \end{array} \right),$$

for $0 \leq \ell \leq m - 1$, where

$$(3.5) \quad \begin{cases} s_i & \in \{a_j \mid i + j + 1 \pmod{d} \in P \cup \{0\}\}, & (i \in N|_d) \\ s_d & \in \{a_0, a_1, \dots, a_{d-1}\}. \end{cases}$$

Consequently, it can be characterized by a word $t(\sigma) = s_0 s_1 \dots s_{d-1}$ or $s_0 s_1 \dots s_{d-1} s_d$, of length d or $d + 1$ respectively. The word $t(\sigma)$ is called the *type* of the permutation σ . We classify the permutations from the set $\mathfrak{J}'_{m,\ell}$ according to the type $t = s_0 s_1 \dots s_{d-1}$ (resp. $t = s_0 s_1 \dots s_{d-1} s_d$) by defining

$$(3.6) \quad \mathfrak{J}^t_{m,\ell} = \{\sigma \in \mathfrak{J}'_{m,\ell} \mid t(\sigma) = t\}.$$

Hence,

$$(3.7) \quad j_{m,\ell} \equiv \sum_t \#\mathfrak{J}^t_{m,\ell} \pmod{2}.$$

Some types do not have any contribution for counting the permutations modulo 2, as stated in the following two lemmas.

Lemma 3.2. *Let $\ell = m$ and $t = s_0 s_1 s_2 \dots s_{d-1}$ (resp. $\ell \in N|_m$ and $t = s_0 s_1 s_2 \dots s_{d-1} s_d$). If there are $0 \leq i < j \leq d - 1$ (resp. $0 \leq i < j \leq d$), such that $s_i = s_j$, $s_i \neq a_{d-i-1}$ and $s_j \neq a_{d-j-1}$, then*

$$(3.8) \quad \#\mathfrak{J}^t_{m,\ell} \equiv 0 \pmod{2}.$$

Proof. If $\mathfrak{J}^t_{m,\ell} = \emptyset$, then (3.8) holds. Otherwise, each permutation $\sigma \in \mathfrak{J}^t_{m,\ell}$ has two biletters $\binom{i_1}{i_2}$ and $\binom{j_1}{j_2}$ of forms $\binom{a_i}{a_k}$ and $\binom{a_j}{a_k}$, respectively, where $a_k = s_i = s_j$. We define another permutation τ obtains from σ by exchanging i_2 and j_2 in the bottom line. This procedure is reversible.

By Lemma 3.1(i) or Table 3.1, it is easy to verify that τ is also in $\mathfrak{J}_{m,\ell}^t$. Thus, the transformation $\sigma \leftrightarrow \tau$ is an *involution* on $\mathfrak{J}_{m,\ell}^t$. Hence, $\#\mathfrak{J}_{m,\ell}^t \equiv 0 \pmod{2}$. \square

Lemma 3.3. *Let $\ell \in N|_m$ and $t = s_0s_1s_2 \dots s_d$. If there is $i \in N|_m$ such that $s_i \neq a_{d-i-1}$, then*

$$(3.9) \quad \sum_{\ell \in N|_m \cap A_i} \#\mathfrak{J}_{m,\ell}^t \equiv 0 \pmod{2}.$$

Proof. For any $\ell \in N|_m \cap A_i$, each permutation $\sigma \in \mathfrak{J}_{m,\ell}^t$ contains two biletters $\binom{i_1}{i_2}$ and $\binom{\ell}{\sigma_\ell}$ of forms $\binom{a_i}{a_j}$ and $\binom{\ell}{a_k}$, respectively. We define another permutation τ by exchanging i_2 and σ_ℓ in the bottom line. This procedure is reversible. By Lemma 3.1(i) it is easy to verify that $\tau \in \mathfrak{J}_{m,\ell'}^t$, where $\ell' = i_1 \in N|_m \cap A_i$. Thus, the transformation $\sigma \leftrightarrow \tau$ is an *involution* on $\sum_{\ell \in N|_m \cap A_i} \mathfrak{J}_{m,\ell}^t$. Hence, (3.9) holds. \square

Let $m = dn + h$ ($n \geq 2$, $h \in N|_d$), $k \in N|_d$ and

$$\mathfrak{P}_Y := \mathfrak{J}_{m,m}^t, \quad \mathfrak{P}_Z := \mathfrak{J}_{m,m-1}^t, \quad \mathfrak{P}_X := \sum_{\ell \in N|_m \cap A_k} \mathfrak{J}_{m,\ell}^t.$$

The recurrence relations listed in Lemmas 2.3, 2.4, 2.6, 2.8, 2.9 can be generated by Algorithm 1. The procedure `EvalAtoms(P,t,h,k)` appearing in Algorithm 1 evaluates the cardinality of the set $\mathfrak{P} := \mathfrak{P}_Y$, \mathfrak{P}_Z or \mathfrak{P}_X for each type t , and will be discussed in Section 4 (see Algorithm 2).

Remark. By Step 2 the recurrence relations generated by Algorithms 1 and 2 are valid for $n \geq 2$. However, we can certify that they are also true for $n = 1$ by using the method described in Sections 3 and 4.

Algorithm 1 Finding the recurrences

```

for P in ['PX', 'PY', 'PZ']:
  for h in range(d):
    Val=0
    for k in range(d) if P=='PX' else range(1):
      for t in PossibleTypes(P,h,k):
        Val=Val+EvalAtoms(P,t,h,k)
    print P,h,k,Val

```

Step 3. Counting permutations. Throughout this step we fix $m = dn + h$ ($h \in N|_d$). Counting permutations from $\mathfrak{J}_{m,\ell}^t$ is lengthy; it is made in several substeps. We illustrate the entire calculations by means of four well-selected examples, using some *compressed and intuitive notations*. Then, we explain what those compressed notations mean in full detail. The examples are given for $d = 5$. We write A, B, C, D, E

instead of A_0, A_1, A_2, A_3, A_4 and a, b, c, d, e instead of a_0, a_1, a_2, a_3, a_4 , respectively.

Example 3.1. Consider $m = \ell = 5n + 1$ and the type ‘adbca’ which satisfies condition (3.5). We have

$$\begin{aligned}
\mathfrak{J}_{5n+1,5n+1}^{adbca} &\stackrel{w}{=} \left(\begin{array}{ccc|ccc|ccc|ccc} 0 & \tilde{5} & 10 & 15 & 1 & 6 & 11 & & 2 & \tilde{7} & 12 & & \tilde{3} & 8 & 13 & & 4 & 9 & 14 \\ e & a & e & e & d & d & d & & c & b & c & & c & b & b & & a & a & a \end{array} \right) \\
&\stackrel{a}{=} \left(\begin{array}{ccc|ccc|ccc|ccc} 0 & \tilde{5} & 10 & 15 & 1 & 6 & 11 & & 2 & \tilde{7} & 12 & & \tilde{3} & 8 & 13 & & 4 & 9 & 14 & 19 \\ e & a & e & e & d & d & d & & c & b & c & & c & b & b & & a & a & a & \underline{19} \end{array} \right) \\
&\stackrel{e}{=} \left(\begin{array}{ccc|ccc|ccc|ccc} 0 & \tilde{5} & 10 & 15 & 1 & 6 & 11 & & 2 & \tilde{7} & 12 & & \tilde{3} & 8 & 13 & & 4 & 9 & 14 & 19 \\ e & \underline{19} & e & e & d & d & d & & c & \underline{c} & c & & \underline{b} & b & b & & a & a & a & \underline{a} \end{array} \right) \\
&\stackrel{d}{=} \left(\begin{array}{ccc} 0 & \tilde{5} & 10 & 15 \\ e & \underline{19} & e & e \end{array} \right) \left(\begin{array}{ccc} 1 & 6 & 11 \\ d & d & d \end{array} \right) \left(\begin{array}{ccc} 2 & \tilde{7} & 12 \\ c & \underline{c} & c \end{array} \right) \left(\begin{array}{ccc} \tilde{3} & 8 & 13 \\ \underline{b} & b & b \end{array} \right) \left(\begin{array}{ccc} 4 & 9 & 14 & 19 \\ a & a & a & \underline{a} \end{array} \right) \\
&\stackrel{b}{=} Z_{n+1} \times Y_n \times X_n \times X_n \times Z_{n+1}.
\end{aligned}$$

Example 3.2. Consider $m = 5n + 2, \ell = 5n + 1$ and the type ‘dcbbaa’ which satisfies condition (3.5). We have

$$\begin{aligned}
\mathfrak{J}_{5n+2,5n+1}^{dcbbaa} &\stackrel{w}{=} \left(\begin{array}{ccc|ccc|ccc|ccc} 0 & \tilde{5} & 10 & 15 & \tilde{1} & 6 & 11 & 16 & 2 & \tilde{7} & 12 & & 3 & 8 & 13 & & 4 & 9 & 14 \\ e & d & e & e & c & d & d & \underline{a} & c & b & c & & b & b & b & & a & a & a \end{array} \right) \\
&\stackrel{a}{=} \left(\begin{array}{ccc|ccc|ccc|ccc} 0 & \tilde{5} & 10 & 15 & \tilde{1} & 6 & 11 & 16 & 2 & \tilde{7} & 12 & & 3 & 8 & 13 & 18 & 4 & 9 & 14 & 19 \\ e & d & e & e & c & d & d & \underline{a} & c & b & c & & b & b & b & \underline{18} & a & a & a & \underline{19} \end{array} \right) \\
&\stackrel{e}{=} \left(\begin{array}{ccc|ccc|ccc|ccc} 0 & \tilde{5} & 10 & 15 & \tilde{1} & 6 & 11 & 16 & 2 & \tilde{7} & 12 & & 3 & 8 & 13 & 18 & 4 & 9 & 14 & 19 \\ e & \underline{19} & e & e & \underline{d} & d & d & \underline{18} & c & \underline{c} & c & & b & b & b & \underline{b} & a & a & a & \underline{a} \end{array} \right) \\
&\stackrel{d}{=} \left(\begin{array}{ccc} 0 & \tilde{5} & 10 & 15 \\ e & \underline{19} & e & e \end{array} \right) \left(\begin{array}{ccc} \tilde{1} & 6 & 11 & 16 \\ \underline{d} & d & d & \underline{18} \end{array} \right) \left(\begin{array}{ccc} 2 & \tilde{7} & 12 \\ c & \underline{c} & c \end{array} \right) \left(\begin{array}{ccc} 3 & 8 & 13 & 18 \\ b & b & b & \underline{b} \end{array} \right) \left(\begin{array}{ccc} 4 & 9 & 14 & 19 \\ a & a & a & \underline{a} \end{array} \right) \\
&\stackrel{b}{=} Z_{n+1} \times X_n \times X_n \times Z_{n+1} \times Z_{n+1}.
\end{aligned}$$

Example 3.3. Consider $m = 5n + 4, \ell \in C|_{n+1}$ and the type ‘adcbac’ which satisfies condition (3.5). We have

$$\begin{aligned}
&\sum_{\ell \in C|_{n+1}} \mathfrak{J}_{5n+4,\ell}^{adcbac} \\
&\stackrel{w}{=} \left(\begin{array}{ccc|ccc|ccc|ccc} 0 & \tilde{5} & 10 & 15 & 1 & 6 & 11 & 16 & 2 & \tilde{7} & 12 & 17 & 3 & 8 & 13 & 18 & 4 & 9 & 14 \\ e & a & e & e & d & d & d & d & c & \underline{c} & c & c & b & b & b & b & a & a & a \end{array} \right) \\
&\stackrel{a}{=} \left(\begin{array}{ccc|ccc|ccc|ccc} 0 & \tilde{5} & 10 & 15 & 1 & 6 & 11 & 16 & 2 & \tilde{7} & 12 & 17 & 3 & 8 & 13 & 18 & 4 & 9 & 14 & 19 \\ e & a & e & e & d & d & d & d & c & \underline{c} & c & c & b & b & b & b & a & a & a & \underline{19} \end{array} \right) \\
&\stackrel{e}{=} \left(\begin{array}{ccc|ccc|ccc|ccc} 0 & \tilde{5} & 10 & 15 & 1 & 6 & 11 & 16 & 2 & \tilde{7} & 12 & 17 & 3 & 8 & 13 & 18 & 4 & 9 & 14 & 19 \\ e & \underline{19} & e & e & d & d & d & d & c & \underline{c} & c & c & b & b & b & b & a & a & a & \underline{a} \end{array} \right) \\
&\stackrel{d}{=} \left(\begin{array}{ccc} 0 & \tilde{5} & 10 & 15 \\ e & \underline{19} & e & e \end{array} \right) \left(\begin{array}{ccc} 1 & 6 & 11 & 16 \\ d & d & d & d \end{array} \right) \left(\begin{array}{ccc} 2 & \tilde{7} & 12 & 17 \\ c & \underline{c} & c & c \end{array} \right) \left(\begin{array}{ccc} 3 & 8 & 13 & 18 \\ b & b & b & b \end{array} \right) \left(\begin{array}{ccc} 4 & 9 & 14 & 19 \\ a & a & a & \underline{a} \end{array} \right)
\end{aligned}$$

$$\stackrel{b}{=} Z_{n+1} \times Y_{n+1} \times X_{n+1} \times Y_{n+1} \times Z_{n+1}.$$

Example 3.4. Consider $m = 5n + 1$, $\ell = 5n$ and the type ‘ $edcaab$ ’ which satisfies condition (3.5). We have

$$\begin{aligned} & \mathfrak{J}_{5n+1,5n}^{edcaab} \\ & \stackrel{w}{=} \left(\begin{array}{c|c|c|c|c} 0 & 5 & 10 & 15 & \\ e & e & e & \underline{b} & \\ \hline 1 & 6 & 11 & & \\ d & d & d & & \\ \hline 2 & 7 & 12 & & \\ c & c & c & & \\ \hline \tilde{3} & 8 & 13 & & \\ a & b & b & & \\ \hline 4 & 9 & 14 & & \\ a & a & a & & \end{array} \right) \\ & \stackrel{a}{=} \left(\begin{array}{c|c|c|c|c} 0 & 5 & 10 & 15 & \\ e & e & e & \underline{b} & \\ \hline 1 & 6 & 11 & & \\ d & d & d & & \\ \hline 2 & 7 & 12 & & \\ c & c & c & & \\ \hline \tilde{3} & 8 & 13 & & \\ a & b & b & & \\ \hline 4 & 9 & 14 & 19 & \\ a & a & a & \underline{19} & \end{array} \right) \\ & \stackrel{e}{=} \left(\begin{array}{c|c|c|c|c} 0 & 5 & 10 & 15 & \\ e & e & e & \underline{19} & \\ \hline 1 & 6 & 11 & & \\ d & d & d & & \\ \hline 2 & 7 & 12 & & \\ c & c & c & & \\ \hline \tilde{3} & 8 & 13 & & \\ \underline{b} & b & b & & \\ \hline 4 & 9 & 14 & 19 & \\ a & a & a & \underline{a} & \end{array} \right) \\ & \stackrel{d}{=} \left(\begin{array}{c} 0 & 5 & 10 & 15 \\ e & e & e & \underline{19} \end{array} \right) \left(\begin{array}{c} 1 & 6 & 11 \\ d & d & d \end{array} \right) \left(\begin{array}{c} 2 & 7 & 12 \\ c & c & c \end{array} \right) \left(\begin{array}{c} \tilde{3} & 8 & 13 \\ \underline{b} & b & b \end{array} \right) \left(\begin{array}{c} 4 & 9 & 14 & 19 \\ a & a & a & \underline{a} \end{array} \right) \\ & \stackrel{b}{=} Y_n \times Y_n \times Y_n \times X_n \times Z_{n+1}. \end{aligned}$$

Notation 1. In the above compressed writing, the letter w, a, e, d, b over the symbol “=” means that the equality is obtained by substep $3(w), 3(a), 3(e), 3(d), 3(b)$ respectively.

Notation 2. In the compressed writing the integer n is represented by the explicit value 3. Hence, the second block in the first equality in Example 3.1 has the following meaning:

$$\left| \begin{array}{c} 1 & 6 & 11 \\ d & d & d \end{array} \right| := \left| \begin{array}{c} 1 & 6 & 11 & 16 & \dots & 5n-4 \\ d & d & d & d & \dots & d \end{array} \right|.$$

Also, the added biletter $\binom{19}{19}$ (see Substep 3(a)) in the second equality in Example 3.1 means $\binom{5n+4}{5n+4}$.

Substep 3(w). Rewrite the set. For each permutation σ from $\mathfrak{J}_{m,\ell}^t$, we reorder the biletters of σ such that $\binom{i}{\sigma_i}$ is on the left of $\binom{j}{\sigma_j}$ if $i \bmod d < j \bmod d$, or if $i \equiv j \bmod d$ and $i < j$. Then, we replace each letter $y \in a_k$ in the bottom line by a_k . To facilitate readability, vertical bars are inserted between the biletters $\binom{i}{\sigma_i}$ and $\binom{j}{\sigma_j}$ such that $i \not\equiv j \pmod{d}$. We get a biword w , denoted by $\rho(\sigma) = w$, called *shape* of σ .

Applying this operation on the following permutation $\sigma \in \mathfrak{J}_{5n+1,5n+1}^{adbca}$ considered in Example 3.1

$$(3.10) \quad \sigma = \left(\begin{array}{cccccccccccccccc} 0 & 1 & 2 & 3 & 4 & 5 & 6 & 7 & 8 & 9 & 10 & 11 & 12 & 13 & 14 & 15 \\ 4 & 3 & 2 & 7 & 15 & 0 & 13 & 1 & 11 & 10 & 14 & 8 & 12 & 6 & 5 & 9 \end{array} \right),$$

we get the shape $\rho(\sigma) = w$, where

$$(3.11) \quad w = \left(\begin{array}{c|c|c|c|c} 0 & 5 & 10 & 15 & \\ e & a & e & e & \\ \hline 1 & 6 & 11 & & \\ d & d & d & & \\ \hline 2 & 7 & 12 & & \\ c & b & c & & \\ \hline 3 & 8 & 13 & & \\ c & b & b & & \\ \hline 4 & 9 & 14 & & \\ a & a & a & & \end{array} \right).$$

Notation 3. In the compressed writing, the above shape w represents also the set $\rho^{-1}(w)$ of all the permutations σ such that $\rho(\sigma) = w$.

Each permutation $\sigma \in \mathfrak{J}_{5n+1, 5n+1}^{adbca}$ contains exactly three unsociable biletters of form $\binom{a}{a}$, $\binom{c}{b}$, $\binom{d}{c}$, denoted by $\binom{i_0}{j_0}$, $\binom{i_1}{j_1}$, $\binom{i_2}{j_2}$, respectively. So that, for example, in the block

$$\begin{vmatrix} 2 & 7 & 12 \\ c & b & c \end{vmatrix}$$

there is exactly one letter ‘ b ’ in the bottom line. All other letters are ‘ c ’. However, the position of the letter ‘ b ’ is not fixed. The shape of another permutation may contain the block

$$\begin{vmatrix} 2 & 7 & 12 \\ b & c & c \end{vmatrix} \quad \text{or} \quad \begin{vmatrix} 2 & 7 & 12 \\ c & c & b \end{vmatrix}.$$

Notation 4. The underlined biletters $\binom{i}{\underline{a}_j}$ in the shape of a permutation σ means that there is no constraint $i + \sigma_i \in J$ for the corresponding biletters $\binom{i}{\sigma_i}$ of σ . All other biletters of σ must satisfy the latter constraint.

In the first equality of each calculation, there is no underlined billetter if $m = \ell$ (Example 3.1) or exactly one underlined billetter if $0 \leq \ell \leq m - 1$ (Examples 3.2 and 3.3). In the latter case, the underscore sign indicates the position of ℓ .

Notation 5. The shape w , with a tilde sign $\tilde{}$ over a billetter $\binom{i}{a_j}$, represents the sum of all shapes w' which are obtained from w by moving the letter a_j , including the underscore sign if it is underlined, to other non-underlined position in the block. For examples, we write (see Example 3.1)

$$\begin{vmatrix} 2 & \tilde{7} & 12 \\ c & b & c \end{vmatrix} := \begin{vmatrix} 2 & 7 & 12 \\ b & c & c \end{vmatrix} + \begin{vmatrix} 2 & 7 & 12 \\ c & b & c \end{vmatrix} + \begin{vmatrix} 2 & 7 & 12 \\ c & c & b \end{vmatrix},$$

and (see Example 3.2)

$$\begin{vmatrix} \tilde{1} & 6 & 11 & 16 \\ c & d & d & \underline{a} \end{vmatrix} := \begin{vmatrix} 1 & 6 & 11 & 16 \\ c & d & d & \underline{a} \end{vmatrix} + \begin{vmatrix} 1 & 6 & 11 & 16 \\ d & c & d & \underline{a} \end{vmatrix} + \begin{vmatrix} 1 & 6 & 11 & 16 \\ d & d & c & \underline{a} \end{vmatrix},$$

$$\begin{vmatrix} \tilde{1} & 6 & 11 & 16 \\ \underline{d} & d & d & \underline{18} \end{vmatrix} := \begin{vmatrix} 1 & 6 & 11 & 16 \\ \underline{d} & d & d & \underline{18} \end{vmatrix} + \begin{vmatrix} 1 & 6 & 11 & 16 \\ d & \underline{d} & d & \underline{18} \end{vmatrix} + \begin{vmatrix} 1 & 6 & 11 & 16 \\ d & d & \underline{d} & \underline{18} \end{vmatrix}.$$

Notice that there is at most one tilde in each block by Lemma 3.3

Substep 3(a). Add biletters. For each $\sigma \in \mathfrak{J}_{dn+h, \ell}^t$, we add all biletters $\binom{i}{}$ such that $\max\{dn + h, dn + d - h\} \leq i \leq dn + d - 1$. Thus, the number of occurrences of a_j in the bottom row becomes the same as the number of occurrences of a_{d-j-1} for any $j \in N|_d$.

For instance, the bottom row of the right-hand side of $\overset{w}{}$ in Example 3.1 contains $4 \times a, 3 \times b, 3 \times c, 3 \times d, 3 \times e$. By adding the billetter $\binom{19}{19}$ to the shape the number of occurrences of a in the bottom row

becomes the same as the number of occurrences of e (since 19 is also an ‘ e ’). The added biletter in the shape is still represented by $\binom{19}{19}$, instead of $\binom{19}{e}$. Notice that it is underlined (see Notation 4).

Substep 3(e). Exchange. Consider all the biletters of the permutation σ , which are unsocial, or which were added in Substep 3(a), or still which have the specific form $\binom{\ell}{a_k}$ with $0 \leq \ell \leq m - 1$. Exchange the bottom letters of those biletters in such a way that all the biletters will become friendly. In most of the cases, each block contains zero or one bad biletter. The only exception is the block containing the specific form $\binom{\ell}{a_k}$ with $\ell = m - 1$, and another unsocial biletter $\binom{i}{a_j}$. In such a case we put the appropriate *explicit letter*, which was added in Substep 3(a), under the letter ℓ when the exchange was made. The whole procedure is reversible.

In Examples 3.1 and 3.2, the exchanges of the bad biletters are realized respectively as follows:

$$\begin{aligned} \left(\cdots \begin{array}{cccc} \tilde{5} & \tilde{7} & \tilde{3} & 19 \\ a & b & c & \underline{19} \end{array} \cdots \right) & \mapsto \left(\cdots \begin{array}{cccc} \tilde{5} & \tilde{7} & \tilde{3} & 19 \\ \underline{19} & \underline{c} & \underline{b} & \underline{a} \end{array} \cdots \right) \\ \left(\cdots \begin{array}{cccc} \tilde{5} & \tilde{1} & 16 & \tilde{7} & 18 & 19 \\ d & c & \underline{a} & b & \underline{18} & \underline{19} \end{array} \cdots \right) & \mapsto \left(\cdots \begin{array}{cccc} \tilde{5} & \tilde{1} & 16 & \tilde{7} & 18 & 19 \\ \underline{19} & \underline{d} & \underline{18} & \underline{c} & \underline{b} & \underline{a} \end{array} \cdots \right) \end{aligned}$$

In the second example, the block $\left| \begin{array}{c} \tilde{1} & 16 \\ c & \underline{a} \end{array} \right|$ contains two bad biletters. We put the *explicit letter* 18 instead of the symbol ‘ d ’ under the letter $\ell = 16$.

Substep 3(d). Decomposition. After Substep 3(e) Exchange, the set $\mathfrak{J}_{dn+h,\ell}^t$ is decomposed, in a natural way, into the Cartesian product of d sets of biwords, which are called *atoms* in the sequel. According to the situation of the tilde and underscore signs, the atoms are classified into six families:

$$\begin{aligned} (i) : & \quad \left(\begin{array}{ccc} 1 & 6 & 11 \\ d & d & d \end{array} \right), \left(\begin{array}{cccc} 1 & 6 & 11 & 16 \\ d & d & d & d \end{array} \right), \left(\begin{array}{cccc} 3 & 8 & 13 & 18 \\ b & b & b & b \end{array} \right), \left(\begin{array}{ccc} 2 & 7 & 12 \\ c & c & c \end{array} \right); \\ (i') : & \quad \left(\begin{array}{cccc} 0 & 5 & 10 & 15 \\ e & e & e & \underline{19} \end{array} \right); \\ (ii) : & \quad \left(\begin{array}{cccc} 4 & 9 & 14 & 19 \\ a & a & a & \underline{a} \end{array} \right), \left(\begin{array}{cccc} 3 & 8 & 13 & 18 \\ b & b & b & \underline{b} \end{array} \right); \\ (ii') : & \quad \left(\begin{array}{cccc} 0 & \tilde{5} & 10 & 15 \\ e & \underline{19} & e & e \end{array} \right); \\ (iii) : & \quad \left(\begin{array}{ccc} 2 & \tilde{7} & 12 \\ c & \underline{c} & c \end{array} \right), \left(\begin{array}{ccc} \tilde{3} & 8 & 13 \\ \underline{b} & b & b \end{array} \right), \left(\begin{array}{cccc} 2 & \tilde{7} & 12 & 17 \\ c & \underline{c} & c & c \end{array} \right); \\ (iii') : & \quad \left(\begin{array}{cccc} \tilde{1} & 6 & 11 & 16 \\ \underline{d} & d & d & \underline{18} \end{array} \right). \end{aligned}$$

It suffices to count the permutations in each atom.

Substep 3(b). Beta transformation. The cardinalities of the atoms can be derived by means of the transformation β defined in (3.2). We discuss the method according to the classification given in Substep 3(d).

(i) The atom

$$(3.12) \quad \mathfrak{A}_0 = \begin{pmatrix} 1 & 6 & 11 \\ d & d & d \end{pmatrix}$$

represents the set

$$\left\{ \begin{pmatrix} 1 & 6 & 11 \\ \tau_1 & \tau_6 & \tau_{11} \end{pmatrix} \mid \begin{array}{l} \{\tau_1, \tau_6, \tau_{11}\} = D|_n \\ i + \tau_i \in J \text{ for } i \in B|_n \end{array} \right\}.$$

If $i + \sigma_i \in A_{d-1}$, then

$$i + \sigma_i \in J \iff \beta(i) + \beta(\sigma_i) = \beta(i + \sigma_i) \in \bar{J}$$

by Lemma 3.1 (iii). Applying the transformation β to each letter in the top and bottom rows of each element τ of \mathfrak{A}_0 , we get a permutation λ from $\tilde{\mathfrak{J}}_{n,n}$:

$$\left\{ \beta \begin{pmatrix} 1 & 6 & 11 \\ \tau_1 & \tau_6 & \tau_{11} \end{pmatrix} = \begin{pmatrix} 0 & 1 & 2 \\ \lambda_1 & \lambda_2 & \lambda_3 \end{pmatrix} \mid \begin{array}{l} \{\lambda_0, \lambda_1, \lambda_2\} = N|_n \\ i + \lambda_i \in \bar{J} \text{ for } i \in N|_n \end{array} \right\}.$$

The above transformation is reversible and the atom \mathfrak{A}_0 is in bijection with $\tilde{\mathfrak{J}}_{n,n}$. Thus $\#\mathfrak{A}_0 = \#\tilde{\mathfrak{J}}_{n,n} = \bar{Y}_n$. For example, the second factor appearing in the right-hand side of the equality $\stackrel{b}{=}$ in Example 3.1, is equal to $\bar{Y}_n = Y_n$.

Notation 6. In the compressed writing, a set symbol may designate also the cardinality of the set, if necessary. For example, we may write $\mathfrak{S}_4 = 24$.

(i') The atom

$$(3.13) \quad \mathfrak{A}_1 = \begin{pmatrix} 0 & 5 & 10 & 15 \\ e & e & e & \underline{19} \end{pmatrix}$$

represents the set

$$\left\{ \begin{pmatrix} 0 & 5 & 10 & 15 \\ \tau_0 & \tau_5 & \tau_{10} & 19 \end{pmatrix} \mid \begin{array}{l} \{\tau_0, \tau_5, \tau_{10}\} = E|_n \\ i + \tau_i \in J \text{ for } i \in A|_n \end{array} \right\}.$$

Thus, it has the same cardinality of the atom \mathfrak{A}_0 defined in (3.12).

(ii) The atom

$$(3.14) \quad \mathfrak{A}_2 = \begin{pmatrix} 3 & 8 & 13 & 18 \\ b & b & b & \underline{b} \end{pmatrix}$$

is meant to be the set

$$\left\{ \begin{pmatrix} 3 & 8 & 13 & 18 \\ \tau_3 & \tau_8 & \tau_{13} & \underline{\tau_{18}} \end{pmatrix} \mid \begin{array}{l} \{\tau_3, \tau_8, \tau_{13}, \tau_{18}\} = B|_{n+1} \\ i + \tau_i \in J \text{ for } i \in D|_n \end{array} \right\}.$$

Applying the transformation β to each letter in each element σ in the atom \mathfrak{A}_2 , we get a permutation λ from $\tilde{\mathfrak{J}}_{n+1,n}$:

$$\left\{ \left(\begin{array}{cccc} 0 & 1 & 2 & 3 \\ \lambda_0 & \lambda_1 & \lambda_2 & \lambda_3 \end{array} \right) \middle| \begin{array}{l} \{\lambda_0, \lambda_1, \lambda_2, \lambda_3\} = N|_{n+1} \\ i + \lambda_i \in \bar{J} \text{ for } i \in N|_n \end{array} \right\}.$$

The transformation is reversible, so that \mathfrak{A}_2 is in bijection with $\tilde{\mathfrak{J}}_{n+1,n}$. Hence, $\#\mathfrak{A}_2 = \#\tilde{\mathfrak{J}}_{n+1,n} = \bar{Z}_{n+1}$.

(ii') The atom

$$(3.15) \quad \mathfrak{A}_3 = \left(\begin{array}{cccc} 0 & \tilde{5} & 10 & 15 \\ e & \underline{19} & e & e \end{array} \right)$$

represents the set

$$\left\{ \left(\begin{array}{cccc} 0 & 5 & 10 & 15 \\ \tau_0 & \tau_5 & \tau_{10} & \tau_{15} \end{array} \right) \middle| \begin{array}{l} \{\tau_0, \tau_5, \tau_{10}, \tau_{15}\} = E|_{n+1} \\ i + \tau_i \in J \text{ for } i \in A|_{n+1} \text{ such that } \tau_i \neq 5n+4 \end{array} \right\}.$$

By inverting the top and bottom rows of each biword, the above set becomes

$$\left\{ \left(\begin{array}{cccc} 4 & 9 & 14 & 19 \\ \rho_4 & \rho_9 & \rho_{14} & \underline{\rho_{19}} \end{array} \right) \middle| \begin{array}{l} \{\rho_4, \rho_9, \rho_{14}, \rho_{19}\} = A|_{n+1} \\ i + \rho_i \in J \text{ for } i \in E|_n \end{array} \right\},$$

which is equal to the atom

$$(3.16) \quad \left(\begin{array}{cccc} 4 & 9 & 14 & 19 \\ a & a & a & \underline{a} \end{array} \right)$$

already studied in (ii).

(iii') The atom

$$(3.17) \quad \mathfrak{A}_4 = \left(\begin{array}{ccc} \tilde{3} & 8 & 13 \\ \underline{b} & b & b \end{array} \right)$$

represents the set

$$\sum_{r \in D|_n} \left\{ \left(\begin{array}{ccc} 3 & 8 & 13 \\ \tau_3 & \tau_8 & \tau_{13} \end{array} \right) \middle| \begin{array}{l} \{\tau_3, \tau_8, \tau_{13}\} = B|_n \\ i + \tau_i \in J \text{ for } i \in D|_n \text{ such that } i \neq r \end{array} \right\}.$$

Applying the transformation β , the latter set becomes

$$\sum_{r=0}^{n-1} \left\{ \left(\begin{array}{ccc} 0 & 1 & 2 \\ \lambda_0 & \lambda_1 & \lambda_2 \end{array} \right) \middle| \begin{array}{l} \{\lambda_0, \lambda_1, \lambda_2\} = N|_n \\ i + \tau_i \in \bar{J} \text{ for } i \in N|_n \text{ such that } i \neq r \end{array} \right\}.$$

Hence,

$$(3.18) \quad \mathfrak{A}_4 = \sum_{r=0}^{n-1} \#\tilde{\mathfrak{J}}_{n,i} = \bar{X}_n.$$

(iii') Similar to (i'), we have

$$\left(\begin{array}{cccc} \tilde{1} & 6 & 11 & 16 \\ \underline{d} & d & d & \underline{18} \end{array} \right) = \left(\begin{array}{cccc} \tilde{1} & 6 & 11 \\ \underline{d} & d & d \end{array} \right),$$

which was already studied in (iii).

4. ALGORITHM FOR EVALUATING THE ATOMS

Keep the same notations as in Section 3, in particular, $m = dn + h$ ($h \in N|_d$). Let $k \in N|_d$. For simplicity, we write

$$\mathfrak{P}_Y := \mathfrak{J}_{m,m}^t, \quad \mathfrak{P}_Z := \mathfrak{J}_{m,m-1}^t, \quad \mathfrak{P}_X := \sum_{\ell \in N|_m \cap A_k} \mathfrak{J}_{m,\ell}^t.$$

For each type t the cardinality of the set $\mathfrak{P} := \mathfrak{P}_Y, \mathfrak{P}_Z$ or \mathfrak{P}_X , is evaluated by the substeps 3(w), 3(a), 3(e), 3(d), 3(b), which are fully described in Section 3. As a consequence, the latter cardinality is equal to the product of d factors (see Examples 3.1–3.4) corresponding to the d atoms respectively. In this section, we show that the substeps in Step 3 can be combined onto one super-step. In fact, each factor can be evaluated directly by using a prefabricated dictionary.

Definition 4.1. Let $i \in N|_d$ be a fixed integer. We define several parameters depending on i, k, d, m, t , where $t = s_0 s_1 \dots s_{d-1}$ (if $\mathfrak{P} = \mathfrak{P}_Y$) or $s_0 s_1 \dots s_{d-1} s_d$ (if $\mathfrak{P} = \mathfrak{P}_X$ or \mathfrak{P}_Z):

$$\begin{aligned} \eta_0 &= \begin{cases} 1, & \text{if } i+1 \leq h, \\ 0, & \text{otherwise;} \end{cases} \\ \eta_1 &= \begin{cases} 1, & \text{if } d-i \leq h, \\ 0, & \text{otherwise;} \end{cases} \\ \eta_2 &= \begin{cases} 1, & \text{if } s_i = a_{d-i-1}, \\ 0, & \text{otherwise;} \end{cases} \\ \eta_3 &= \begin{cases} 1, & \text{if } \mathfrak{P} \neq \mathfrak{P}_Y \text{ and } s_d = a_{d-i-1}, \\ 0, & \text{otherwise.} \end{cases} \\ \nu &= \begin{cases} \text{'Z'}, & \text{if } \mathfrak{P} = \mathfrak{P}_Z \text{ and } m-1 \in A_i, \\ \text{'X'}, & \text{if } \mathfrak{P} = \mathfrak{P}_X \text{ and } k = i, \\ \text{'G'}, & \text{otherwise;} \end{cases} \\ \mu_i &= \begin{cases} \Psi_Z(\eta_0, \eta_1, \eta_2, \eta_3), & \text{if } \nu = \text{'Z'}, \\ \Psi_X(\eta_0, \eta_1, \eta_2, \eta_3), & \text{if } \nu = \text{'X'}, \\ \Psi_G(\eta_0, \eta_1, \eta_2), & \text{if } \nu = \text{'G'}, \end{cases} \end{aligned}$$

where the explicit values of the functions Ψ_Z, Ψ_X, Ψ_G are given in Table 4.1.

Notice that each permutation contains $n + \eta_0$ (resp. $n + \eta_1$) letters in A_i (resp. in A_{d-i-1}).

Example 4.1. Consider $\mathfrak{P} = \mathfrak{J}_{5n+1, 5n+1}^{adbca}$, studied in Example 3.1. In this case, $d = 5, m = 5n + 1, h = 1, \ell = m = 5n + 1, t = s_0 s_1 s_2 s_3 s_4 = \text{'adbca'}$. For $i = 1$ we have $\eta_0 = 0, \eta_1 = 0, \eta_2 = 1$. Hence, $\mu_1 = \Psi_G(0, 0, 1) = \bar{Y}_n$.

| | | | | | | | | |
|----------------|-------------|-------------|------|-----------------|-----------------|-----------------|-----------------|-----------------|
| η | 000 | 001 | 010 | 011 | 100 | 101 | 110 | 111 |
| $\Psi_G(\eta)$ | \bar{X}_n | \bar{Y}_n | 0 | \bar{Z}_{n+1} | \bar{Z}_{n+1} | 0 | \bar{X}_{n+1} | \bar{Y}_{n+1} |
| η | 0000 | 0010 | 0100 | 0110 | 1000 | 1010 | 1100 | 1110 |
| $\Psi_Z(\eta)$ | 0 | \bar{Z}_n | 0 | 0 | \bar{X}_n | \bar{Y}_n | 0 | \bar{Z}_{n+1} |
| η | 0001 | 0011 | 0101 | 0111 | 1001 | 1011 | 1101 | 1111 |
| $\Psi_Z(\eta)$ | 0 | \bar{Z}_n | 0 | 0 | \bar{X}_n | 0 | 0 | \bar{Z}_{n+1} |
| η | 0000 | 0010 | 0100 | 0110 | 1000 | 1010 | 1100 | 1110 |
| $\Psi_X(\eta)$ | 0 | \bar{X}_n | 0 | 0 | 0 | \bar{Z}_{n+1} | 0 | \bar{X}_{n+1} |
| η | 0001 | 0011 | 0101 | 0111 | 1001 | 1011 | 1101 | 1111 |
| $\Psi_X(\eta)$ | 0 | \bar{X}_n | 0 | 0 | 0 | 0 | 0 | \bar{X}_{n+1} |

TABLE 4.1. Explicit values of the functions Ψ_Z, Ψ_X, Ψ_G

Example 4.2. Consider $\mathfrak{P} = \mathfrak{J}_{5n+2,5n+1}^{dcbbaa}$, studied in Example 3.2. In this case, $d = 5, m = 5n + 2, h = 2, \ell = m - 1 = 5n + 1 \in A_1, t = s_0s_1s_2s_3s_4 = 'dcbbaa'$. For $i = 1$ we have $\eta_0 = 1, \eta_1 = 0, \eta_2 = 0, \eta_3 = 0$. So that $\mu_1 = \Psi_Z(1, 0, 0, 0) = \bar{X}_n$.

Example 4.3. Consider $\mathfrak{P} = \sum_{\ell \in C|_{n+1}} \mathfrak{J}_{5n+4, \ell}^{adcbac}$, studied in Example 3.3. In this case, $d = 5, m = 5n + 4, h = 4, \ell \in A_2, t = s_0s_1s_2s_3s_4 = 'adcbac'$. For $i = 2$ we have $\eta_0 = 1, \eta_1 = 1, \eta_2 = 1, \eta_3 = 1$ and $\mu_2 = \Psi_X(1, 1, 1, 1) = \bar{X}_{n+1}$.

Theorem 4.1. *With the above notations, the cardinality of the set $\mathfrak{P} := \mathfrak{P}_Y, \mathfrak{P}_Z, \mathfrak{P}_X$ is equal to*

$$(4.1) \quad \#\mathfrak{P} = \mu_0 \times \mu_1 \times \mu_2 \times \cdots \times \mu_{d-1}.$$

For example, the set $\mathfrak{P} = \sum_{\ell \in C|_{n+1}} \mathfrak{J}_{5n+4, \ell}^{adcbac}$, studied in Example 3.3, is evaluated by means of Theorem 4.1 as follows:

$$\begin{aligned} & \sum_{\ell \in C|_{n+1}} \mathfrak{J}_{5n+4, \ell}^{adcbac} \\ &= \mu_0 \mu_1 \mu_2 \mu_3 \mu_4 \\ &= \Phi_G(1, 0, 0) \Phi_G(1, 1, 1) \Psi_X(1, 1, 1, 1) \Psi_G(1, 1, 1) \Phi_G(0, 1, 1) \\ &= Z_{n+1} Y_{n+1} X_{n+1} Y_{n+1} Z_{n+1}. \end{aligned}$$

By Theorem 4.1, the procedure `EvalAtoms(P, t, h, k)` figured in Algorithm 1, which evaluates the cardinality of the set $\mathfrak{P} := \mathfrak{P}_Y, \mathfrak{P}_Z$ or \mathfrak{P}_X for each type t , is described in Algorithm 2.

Proof of Theorem 4.1. When we speak of *case*, we refer to a tuple $(\nu = 'G', \eta_0, \eta_1, \eta_2)$, $(\nu = 'Z', \eta_0, \eta_1, \eta_2, \eta_3)$ or $(\nu = 'X', \eta_0, \eta_1, \eta_2, \eta_3)$, which depends on $i, k, d, m, t, \mathfrak{P}$ by Definition 4.1. The case is reproduced without non-significant symbols. For example, we write $X1000$ for the case $('X', 1, 0, 0, 0)$.

In fact, the cases $G101, Z1011, X1011$ do not appear in product (4.1) and can take any value, in particular, zero. In the cases $X1000$ and

Algorithm 2 Evaluating the atoms

```

def EvalAtoms(P,t,h,k):
  Prod=1
  for i in Ch:
    nu='G'
    if P=='PZ' and i==(h+d-1)%d: nu='Z'
    if P=='PX' and i==k: nu='X'
    eta=(i+1<=h, d-i<=h, t[i]==d-i-1)
    if nu=='X' or nu=='Z': eta=eta+(t[d]==d-i-1,)
  Prod=Prod*Psi(nu, eta)
return Prod

```

X_{1001} , we have $\#\mathfrak{P} = 0$ by Lemma 3.3, so that Identity (4.1) is true. In the cases $(G, 0, 1, 0)$ and (ν, η) for

$$\begin{aligned} \nu &= 'Z', 'X'; \\ \eta &= (0, 0, 0, 0), (0, 1, 0, 0), (0, 1, 1, 0), (1, 1, 0, 0), \\ &\quad (0, 0, 0, 1), (0, 1, 0, 1), (0, 1, 1, 1), (1, 1, 0, 1), \end{aligned}$$

Lemma 3.2 implies that $\#\mathfrak{P} = 0$. Hence, Identity (4.1) is true. All other cases are proved as follows.

The evaluations of product (4.1) are explained in Section 3, see Examples 3.1–3.4. The factors $\mu_0, \mu_1, \dots, \mu_{d-1}$ are obtained at the same time by proceeding with the substeps $3(w), 3(a), 3(e), 3(d), 3(b)$. In fact, we can evaluate each sole factor μ_i without keeping in mind the others. For this purpose, we extract all biletters such that either its top letter is in A_i or its bottom letter is a_{d-i-1} in the first two substeps $3(w)$ and $3(a)$.

Again, consider $i = 1$ and $\mathfrak{P} = \mathfrak{J}_{5n+2, 5n+1}^{dcbbaa}$. We extract all biletters such that either its top letter is in $\{1, 6, 11, 16, \dots\}$ or its bottom letter is d in the first two substeps $3(w)$ and $3(a)$ of Example 3.2. We have

$$\mathfrak{J}_{5n+2, 5n+1}^{dcbbaa} \stackrel{w}{=} \left(\begin{array}{c|ccc|c|c|c} \tilde{5} & \tilde{1} & 6 & 11 & 16 & ? & ? & ? \\ d & c & d & d & \underline{a} & & & \end{array} \right) \stackrel{a}{=} \left(\begin{array}{c|ccc|c|c|c} \tilde{5} & \tilde{1} & 6 & 11 & 16 & ? & \underline{18} & ? \\ d & c & d & d & \underline{a} & & \underline{18} & ? \end{array} \right),$$

and

$$\mathfrak{J}_{5n+2, 5n+1}^{dcbbaa} \stackrel{e}{=} \left(\begin{array}{c|ccc|c|c|c} \tilde{5} & \tilde{1} & 6 & 11 & 16 & ? & \underline{18} & ? \\ \underline{19} & \underline{d} & d & d & \underline{18} & ? & \underline{b} & ? \end{array} \right) \stackrel{d}{=} ? \left(\begin{array}{c|ccc} \tilde{1} & 6 & 11 & 16 \\ \underline{d} & d & d & \underline{18} \end{array} \right) ? \stackrel{b}{=} ? X_n ???.$$

It means that $\mu_1 = X_n = \bar{X}_n$. On the other hand, this case corresponds to the tuple $(Z, 1, 0, 0, 0)$ that takes the value \bar{X}_n , as shown in Example 4.2.

In the sequel, $i_0 = i, i_1 = d + i, i_2 = 2d + i, \dots, i_{n-1} = (n-1)d + i, i_n = nd + i$ are integers from A_i . Let $j = d - i - 1, j_n = dn + j$. We prove (4.1) case by case using the method described in the above

example. Without loss of generality, the proof is illustrated for $n = 3$.

$$\begin{aligned}
G000 &: \left(\begin{array}{ccc|c} i_0 & \tilde{i}_1 & i_2 & ? \\ a_j & ? & a_j & a_j \end{array} \right) \stackrel{e}{=} \left(\begin{array}{ccc} i_0 & \tilde{i}_1 & i_2 \\ a_j & \underline{a_j} & a_j \end{array} \right) \stackrel{b}{=} \bar{X}_n, \\
G001 &: \left(\begin{array}{ccc} i_0 & i_1 & i_2 \\ a_j & a_j & a_j \end{array} \right) \stackrel{b}{=} \bar{Y}_n, \\
G011 &: \left(\begin{array}{ccc|c} i_0 & i_1 & i_2 & ? \\ a_j & a_j & a_j & a_j \end{array} \right) \\
&\stackrel{a}{=} \left(\begin{array}{ccc|c} i_0 & i_1 & i_2 & ? \\ a_j & a_j & a_j & a_j \end{array} \right) \left(\begin{array}{c} i_n \\ ? \end{array} \right) \stackrel{e}{=} \left(\begin{array}{cccc} i_0 & i_1 & i_2 & i_n \\ a_j & a_j & a_j & \underline{a_j} \end{array} \right) \stackrel{b}{=} \bar{Z}_{n+1}, \\
G100 &: \left(\begin{array}{cccc} i_0 & \tilde{i}_1 & i_2 & i_3 \\ a_j & ? & a_j & a_j \end{array} \right) \\
&\stackrel{a}{=} \left(\begin{array}{cccc} i_0 & \tilde{i}_1 & i_2 & i_3 \\ a_j & ? & a_j & a_j \end{array} \right) \left(\begin{array}{c} ? \\ \underline{j_n} \end{array} \right) \stackrel{e}{=} \left(\begin{array}{cccc} i_0 & \tilde{i}_1 & i_2 & i_3 \\ a_j & \underline{j_n} & a_j & a_j \end{array} \right) \stackrel{b}{=} \bar{Z}_{n+1}, \\
G110 &: \left(\begin{array}{cccc|c} i_0 & \tilde{i}_1 & i_2 & i_3 & ? \\ a_j & ? & a_j & a_j & a_j \end{array} \right) \stackrel{e}{=} \left(\begin{array}{cccc} i_0 & \tilde{i}_1 & i_2 & i_3 \\ a_j & \underline{a_j} & a_j & a_j \end{array} \right) \stackrel{b}{=} \bar{X}_{n+1}, \\
G111 &: \left(\begin{array}{cccc} i_0 & i_1 & i_2 & i_3 \\ a_j & a_j & a_j & a_j \end{array} \right) \stackrel{b}{=} \bar{Y}_{n+1}, \\
Z0010 &: \left(\begin{array}{ccc|c} i_0 & i_1 & i_{n-1} & ? \\ a_j & a_j & ? & a_j \end{array} \right) \stackrel{e}{=} \left(\begin{array}{ccc} i_0 & i_1 & i_{n-1} \\ a_j & a_j & \underline{a_j} \end{array} \right) \stackrel{b}{=} \bar{Z}_n, \\
Z1110 &: \left(\begin{array}{cccc|c} i_0 & i_1 & i_2 & i_n & ? \\ a_j & a_j & a_j & ? & a_j \end{array} \right) \stackrel{e}{=} \left(\begin{array}{cccc} i_0 & i_1 & i_2 & i_n \\ a_j & a_j & a_j & \underline{a_j} \end{array} \right) \stackrel{b}{=} \bar{Z}_{n+1}, \\
Z0011 &: \left(\begin{array}{ccc} i_0 & i_1 & i_{n-1} \\ a_j & a_j & \underline{a_j} \end{array} \right) \stackrel{b}{=} \bar{Z}_n, \\
Z1111 &: \left(\begin{array}{cccc} i_0 & i_1 & i_2 & i_n \\ a_j & a_j & a_j & \underline{a_j} \end{array} \right) \stackrel{b}{=} \bar{Z}_{n+1}, \\
Z1000 &: \left(\begin{array}{ccc|c} i_0 & \tilde{i}_1 & i_2 & i_n \\ a_j & ? & a_j & ? \\ & & & a_j \end{array} \right) \\
&\stackrel{a}{=} \left(\begin{array}{ccc|c} i_0 & \tilde{i}_1 & i_2 & i_n \\ a_j & ? & a_j & ? \\ & & & a_j \end{array} \right) \left(\begin{array}{c} ? \\ \underline{j_n} \end{array} \right) \stackrel{e}{=} \left(\begin{array}{cccc} i_0 & \tilde{i}_1 & i_2 & i_n \\ a_j & \underline{a_j} & a_j & \underline{j_n} \end{array} \right) \stackrel{b}{=} \bar{X}_n, \\
Z1001 &: \left(\begin{array}{cccc} i_0 & \tilde{i}_1 & i_2 & i_n \\ a_j & ? & a_j & \underline{a_j} \end{array} \right) \\
&\stackrel{a}{=} \left(\begin{array}{cccc} i_0 & \tilde{i}_1 & i_2 & i_n \\ a_j & ? & a_j & \underline{a_j} \end{array} \right) \left(\begin{array}{c} ? \\ \underline{j_n} \end{array} \right) \stackrel{e}{=} \left(\begin{array}{cccc} i_0 & \tilde{i}_1 & i_2 & i_n \\ a_j & \underline{a_j} & a_j & \underline{j_n} \end{array} \right) \stackrel{b}{=} \bar{X}_n, \\
Z1010 &: \left(\begin{array}{cccc} i_0 & i_1 & i_2 & i_n \\ a_j & a_j & a_j & ? \end{array} \right) \\
&\stackrel{a}{=} \left(\begin{array}{cccc} i_0 & i_1 & i_2 & i_n \\ a_j & a_j & a_j & ? \end{array} \right) \left(\begin{array}{c} ? \\ \underline{j_n} \end{array} \right) \stackrel{e}{=} \left(\begin{array}{cccc} i_0 & i_1 & i_2 & i_n \\ a_j & a_j & a_j & \underline{j_n} \end{array} \right) \stackrel{b}{=} \bar{Y}_n,
\end{aligned}$$

$$\begin{aligned}
X0010 &: \left(\begin{array}{ccc|c} i_0 & \tilde{i}_1 & i_2 & ? \\ a_j & ? & a_j & a_j \end{array} \right) \stackrel{e}{=} \left(\begin{array}{ccc} i_0 & \tilde{i}_1 & i_2 \\ a_j & \underline{a_j} & a_j \end{array} \right) \stackrel{b}{=} \bar{X}_n, \\
X0011 &: \left(\begin{array}{ccc} i_0 & \tilde{i}_1 & i_2 \\ a_j & \underline{a_j} & a_j \end{array} \right) \stackrel{b}{=} \bar{X}_n, \\
X1110 &: \left(\begin{array}{cccc|c} i_0 & \tilde{i}_1 & i_2 & i_3 & ? \\ a_j & ? & a_j & a_j & a_j \end{array} \right) \stackrel{e}{=} \left(\begin{array}{cccc} i_0 & \tilde{i}_1 & i_2 & i_3 \\ a_j & \underline{a_j} & a_j & a_j \end{array} \right) \stackrel{b}{=} \bar{X}_{n+1}, \\
X1111 &: \left(\begin{array}{cccc} i_0 & \tilde{i}_1 & i_2 & i_3 \\ a_j & \underline{a_j} & a_j & a_j \end{array} \right) \stackrel{b}{=} \bar{X}_{n+1}, \\
X1010 &: \left(\begin{array}{cccc} i_0 & \tilde{i}_1 & i_2 & i_3 \\ a_j & ? & a_j & a_j \end{array} \right) \\
&\stackrel{a}{=} \left(\begin{array}{cccc} i_0 & \tilde{i}_1 & i_2 & i_3 \\ a_j & ? & a_j & a_j \end{array} \right) \left(\begin{array}{c} ? \\ \underline{j_n} \end{array} \right) \stackrel{e}{=} \left(\begin{array}{cccc} i_0 & \tilde{i}_1 & i_2 & i_3 \\ a_j & \underline{j_n} & a_j & a_j \end{array} \right) \stackrel{b}{=} \bar{Z}_{n+1}. \quad \square
\end{aligned}$$

5. IMPLEMENTATION AND OUTPUTS

Our program `Apwen.py` is an implementation of Algorithms 1 and 2 in Python. The proofs of Lemmas 2.3, 2.4 and 2.6 are achieved by the following Outputs 1–3 of the program `Apwen.py` respectively. For simplicity, the expression $n+1$ is reproduced by letter m . Thus, " $Y(3n+1) = V_n W_m$ " means the following recurrence relation

$$Y_{3n+1} = V_n W_{n+1},$$

which appeared in Lemma 2.3. The calculations made in Examples 3.1–3.4 in Section 3 can be found in Output 3, types 168 `adbca`, 219 `dcbaa`, 145 `adcbac`, 213 `edcaab` respectively.

Output 1 "python Apwen.py 3"

```

v= [1, -1, -1]      direction = XYZ -> UVW
P= [1]
Q= [2]
J= [0, 3, 5, 6, 8, 9, 12, 14, 15, 18, 21, 23, 24, 27, ...]
K= [1, 2, 4, 7, 10, 11, 13, 16, 17, 19, 20, 22, 25, ...]

k : 3N+0
1 cbac: [Un:X0011] [Vn:G001] [Vn:G001]
2 ccab: [Un:X0010] [Un:G000] [Vn:G001]
3 ccba: [Un:X0010] [Un:G000] [Un:G000]
k : 3N+1
4 abbc: [Un:G000] [Un:X0010] [Un:G000]
5 cbab: [Vn:G001] [Un:X0011] [Vn:G001]
6 cbba: [Vn:G001] [Un:X0010] [Un:G000]
k : 3N+2
7 abac: [Un:G000] [Vn:G001] [Un:X0010]
8 acab: [Un:G000] [Un:G000] [Un:X0010]
9 cbaa: [Vn:G001] [Vn:G001] [Un:X0011]
X(3n+0) = Un

```

k : 3N+0
 10 cbaa: [Wm:X1010] [Vn:G001] [Wm:G011]
 k : 3N+1
 11 abab: [Wm:G100] [Un:X0011] [Wm:G011]
 k : 3N+2
 $X(3n+1) = U_n W_m + V_n W_m$

k : 3N+0
 12 cbaa: [Wm:X1010] [Vm:G111] [Wm:G011]
 k : 3N+1
 13 abab: [Wm:G100] [Um:X1111] [Wm:G011]
 k : 3N+2
 $X(3n+2) = U_m W_m + V_m W_m$

14 acb: [Un:G000] [Un:G000] [Un:G000]
 15 cba: [Vn:G001] [Vn:G001] [Vn:G001]
 $Y(3n+0) = U_n + V_n$

16 aba: [Wm:G100] [Vn:G001] [Wm:G011]
 $Y(3n+1) = V_n W_m$

17 aba: [Wm:G100] [Vm:G111] [Wm:G011]
 $Y(3n+2) = V_m W_m$

18 abac: [Un:G000] [Vn:G001] [Wn:Z0010]
 19 acab: [Un:G000] [Un:G000] [Wn:Z0010]
 20 cbaa: [Vn:G001] [Vn:G001] [Wn:Z0011]
 $Z(3n+0) = U_n V_n W_n + U_n W_n + V_n W_n$

21 abac: [Un:Z1001] [Vn:G001] [Wm:G011]
 22 acab: [Un:Z1000] [Un:G000] [Wm:G011]
 23 cbaa: [Vn:Z1010] [Vn:G001] [Wm:G011]
 $Z(3n+1) = U_n V_n W_m + U_n W_m + V_n W_m$

24 abab: [Wm:G100] [Wm:Z1111] [Wm:G011]
 $Z(3n+2) = W_m$

Output 2 "python Apwen.py -3"

```
v= [1, -1, -1]      direction = UVW -> XYZ
P= [2]
Q= [1]
J= [1, 2, 4, 7, 10, 11, 13, 16, 17, 19, 20, 22, 25, ...]
K= [0, 3, 5, 6, 8, 9, 12, 14, 15, 18, 21, 23, 24, 27, ...]
```

k : 3N+0
 1 cacb: [Xn:X0010] [Xn:G000] [Xn:G000]
 2 cbac: [Xn:X0011] [Yn:G001] [Yn:G001]
 3 cbca: [Xn:X0010] [Yn:G001] [Xn:G000]
 k : 3N+1
 4 bbac: [Xn:G000] [Xn:X0010] [Yn:G001]
 5 bbca: [Xn:G000] [Xn:X0010] [Xn:G000]
 6 cbab: [Yn:G001] [Xn:X0011] [Yn:G001]

k : 3N+2
 7 baac: [Xn:G000] [Xn:G000] [Xn:X0010]
 8 caab: [Yn:G001] [Xn:G000] [Xn:X0010]
 9 cbaa: [Yn:G001] [Yn:G001] [Xn:X0011]
 $U(3n+0) = Xn$

k : 3N+0
 10 caab: [Zm:X1010] [Xn:G000] [Zm:G011]
 11 cbaa: [Zm:X1010] [Yn:G001] [Zm:G011]
 k : 3N+1
 12 bbaa: [Zm:G100] [Xn:X0010] [Zm:G011]
 k : 3N+2
 $U(3n+1) = Yn Zm$

k : 3N+0
 13 caab: [Zm:X1010] [Xm:G110] [Zm:G011]
 14 cbaa: [Zm:X1010] [Ym:G111] [Zm:G011]
 k : 3N+1
 15 bbaa: [Zm:G100] [Xm:X1110] [Zm:G011]
 k : 3N+2
 $U(3n+2) = Ym Zm$

16 bac: [Xn:G000] [Xn:G000] [Xn:G000]
 17 cba: [Yn:G001] [Yn:G001] [Yn:G001]
 $V(3n+0) = Xn + Yn$

18 baa: [Zm:G100] [Xn:G000] [Zm:G011]
 $V(3n+1) = Xn Zm$

19 baa: [Zm:G100] [Xm:G110] [Zm:G011]
 $V(3n+2) = Xm Zm$

20 baac: [Xn:G000] [Xn:G000] [Zn:Z0010]
 21 caab: [Yn:G001] [Xn:G000] [Zn:Z0010]
 22 cbaa: [Yn:G001] [Yn:G001] [Zn:Z0011]
 $W(3n+0) = Xn Yn Zn + Xn Zn + Yn Zn$

23 baac: [Xn:Z1001] [Xn:G000] [Zm:G011]
 24 caab: [Yn:Z1010] [Xn:G000] [Zm:G011]
 25 cbaa: [Yn:Z1010] [Yn:G001] [Zm:G011]
 $W(3n+1) = Xn Yn Zm + Xn Zm + Yn Zm$

26 bbaa: [Zm:G100] [Zm:Z1110] [Zm:G011]
 $W(3n+2) = Zm$

Output 3 "python Apwen.py 5" (extract)

```

v= [1, -1, -1, -1, 1]    direction = XYZ -> XYZ
P= [1, 4]
Q= [2, 3]
J= [0, 3, 4, 5, 8, 10, 13, 15, 18, 19, 20, 23, 24, 25, 28, 29]
K= [1, 2, 6, 7, 9, 11, 12, 14, 16, 17, 21, 22, 26, 27]

```

```

...
k : 5N+2
144 accbad: [Zm:G100] [Xm:G110] [Xm:X1110] [Ym:G111] [Zm:G011]
145 adcbac: [Zm:G100] [Ym:G111] [Xm:X1111] [Ym:G111] [Zm:G011]
146 adccab: [Zm:G100] [Ym:G111] [Xm:X1110] [Xm:G110] [Zm:G011]
147 dccaab: [Zm:G100] [Xm:G110] [Xm:X1110] [Xm:G110] [Zm:G011]
148 dccbaa: [Zm:G100] [Xm:G110] [Xm:X1110] [Ym:G111] [Zm:G011]
k : 5N+3
149 acbbad: [Zm:G100] [Xm:G110] [Xm:G110] [Xm:X1110] [Zm:G011]
150 acdbab: [Zm:G100] [Xm:G110] [Xm:G110] [Xm:X1111] [Zm:G011]
151 adbbac: [Zm:G100] [Ym:G111] [Xm:G110] [Xm:X1110] [Zm:G011]
152 adcbab: [Zm:G100] [Ym:G111] [Ym:G111] [Xm:X1111] [Zm:G011]
153 dcbbaa: [Zm:G100] [Xm:G110] [Xm:G110] [Xm:X1110] [Zm:G011]
k : 5N+4
X(5n+4) = Ym Zm

...
167 acdba: [Zm:G100] [Xn:G000] [Xn:G000] [Yn:G001] [Zm:G011]
168 adbca: [Zm:G100] [Yn:G001] [Xn:G000] [Xn:G000] [Zm:G011]
169 adcba: [Zm:G100] [Yn:G001] [Yn:G001] [Yn:G001] [Zm:G011]
170 dcbaa: [Zm:G100] [Xn:G000] [Xn:G000] [Xn:G000] [Zm:G011]
Y(5n+1) = Xn Zm + Yn Zm

...
212 edbcaa: [Yn:Z1010] [Yn:G001] [Xn:G000] [Xn:G000] [Zm:G011]
213 edcaab: [Yn:Z1010] [Yn:G001] [Yn:G001] [Xn:G000] [Zm:G011]
214 edcbaa: [Yn:Z1010] [Yn:G001] [Yn:G001] [Yn:G001] [Zm:G011]
Z(5n+1) = Xn Yn Zm + Xn Zm + Yn Zm

215 acbbad: [Zm:G100] [Xn:Z1001] [Xn:G000] [Zm:G011] [Zm:G011]
216 acdbab: [Zm:G100] [Xn:Z1000] [Xn:G000] [Zm:G011] [Zm:G011]
217 adbbac: [Zm:G100] [Yn:Z1010] [Xn:G000] [Zm:G011] [Zm:G011]
218 adcbab: [Zm:G100] [Yn:Z1010] [Yn:G001] [Zm:G011] [Zm:G011]
219 dcbbaa: [Zm:G100] [Xn:Z1000] [Xn:G000] [Zm:G011] [Zm:G011]
Z(5n+2) = Xn Yn Zm + Xn Zm + Yn Zm

...

```

The proof of that F_{13} is Apwenian takes 11 hours by using the program `Apwen.py` on a modern personal computer. For proving that F_{17a} and F_{17b} are Apwenian, it was necessary to rewrite the program in the C language with some optimizations. The running times of the two programs are reproduced in the following table:

| f | F_3 | F_5 | F_{11} | F_{13} | F_{17a}, F_{17b} | F_{19} |
|--------|-------|-------|----------|----------|-------------------------|----------|
| Python | < 1s | < 1s | 11m | 11h | ∞ | ∞ |
| C | < 1s | < 1s | 16s | 29m | 7 days \times 24 CPUs | ∞ |

REFERENCES

- [1] J.-P. Allouche, J. Peyrière, Z.-X. Wen, and Z.-Y. Wen. Hankel determinants of the Thue–Morse sequence. *Ann. Inst. Fourier, Grenoble*, 48:1–27, 1998.

- [2] J.-P. Allouche and J. Shallit. *Automatic sequences*. Cambridge University Press, Cambridge, 2003. Theory, applications, generalizations.
- [3] Jean-Paul Allouche and Jeffrey Shallit. The ubiquitous Prouhet-Thue-Morse sequence. In *Sequences and their applications (Singapore, 1998)*, Springer Ser. Discrete Math. Theor. Comput. Sci., pages 1–16. Springer, London, 1999.
- [4] Y. Bugeaud, G.-N. Han, Wen Z.-Y., and Yao J.-Y. Hankel determinants, Padé approximations, and irrationality exponents. *to appear in IMRN*, 28 pages, 2015.
- [5] Yann Bugeaud. On the rational approximation to the Thue-Morse-Mahler numbers. *Ann. Inst. Fourier (Grenoble)*, 61(5):2065–2076 (2012), 2011.
- [6] Yann Bugeaud and Guo-Niu Han. A combinatorial proof of the non-vanishing of Hankel determinants of the Thue-Morse sequence. *Electron. J. Combin.*, 21(3):Paper 3.26, 17, 2014.
- [7] P. Flajolet. Combinatorial aspects of continued fractions. *Discrete Math.*, 32(2):125–161, 1980.
- [8] OEIS Foundation. Sequence A010060. *The On-Line Encyclopedia of Integer Sequences*, 2015.
- [9] OEIS Foundation. Sequence A106400. *The On-Line Encyclopedia of Integer Sequences*, 2015.
- [10] Wikimedia Foundation. ± 1 -sequence. *Wikipedia: The Free Encyclopedia*, 2015.
- [11] Wikimedia Foundation. Thue–Morse sequence. *Wikipedia: The Free Encyclopedia*, 2015.
- [12] Guo-Niu Han. Hankel continued fraction and its applications. *arXiv:1406.1593*, 2014.
- [13] Guo-Niu Han. Hankel determinant calculus for the Thue-Morse and related sequences. *J. Number Theory*, 147:374–395, 2015.
- [14] Guo-Niu Han and Wen Wu. Evaluations of the Hankel determinants of a Thue–Morse-like sequence. *International J. Number Theory*, 2015.
- [15] C. Krattenthaler. Advanced determinant calculus. *Sém. Lothar. Combin.*, 42:Art. B42q, 67 pp. (electronic), 1999. The Andrews Festschrift (Maratea, 1998).
- [16] C. Krattenthaler. Advanced determinant calculus: a complement. *Linear Algebra Appl.*, 411:68–166, 2005.
- [17] T. Muir. *The theory of determinants in the historical order of development*, volume 4. Macmillan, London, 1906-1923.
- [18] H. S. Wall. *Analytic Theory of Continued Fractions*. D. Van Nostrand Company, Inc., New York, N. Y., 1948.

INSTITUTE FOR INTERDISCIPLINARY INFORMATION SCIENCES, TSINGHUA UNIVERSITY, BEIJING, 100084, P.R.CHINA

E-mail address: fu-h13@mails.tsinghua.edu.cn

INSTITUT DE RECHERCHE MATHÉMATIQUE AVANCÉE, UNIVERSITÉ DE STRASBOURG ET CNRS, 7 RUE RENÉ DESCARTES, 67084 STRASBOURG, FRANCE

E-mail address: guoniu.han@unistra.fr