

ON THE NUMBER OF PLANAR EULERIAN ORIENTATIONS

NICOLAS BONICHON, MIREILLE BOUSQUET-MÉLOU, PAUL DORBEC, AND CLAIRE PENNARUN

ABSTRACT. The number of planar Eulerian maps with n edges is well-known to have a simple expression. But what is the number of planar Eulerian *orientations* with n edges? This problem appears to be difficult. To approach it, we define and count families of subsets and supersets of planar Eulerian orientations, indexed by an integer k , that converge to the set of all planar Eulerian orientations as k increases. The generating functions of our subsets can be characterized by systems of polynomial equations, and are thus algebraic. The generating functions of our supersets are characterized by polynomial systems involving divided differences, as often occurs in map enumeration. We prove that these series are algebraic as well. We obtain in this way lower and upper bounds on the growth rate of planar Eulerian orientations, which appears to be around 12.5.

1. INTRODUCTION

The enumeration of planar maps (graphs embedded on the sphere) has received a lot of attention since the sixties. Many remarkable counting results have been discovered, which were often illuminated later by beautiful bijective constructions. For instance, it has been known¹ since 1963 that the number of rooted planar *Eulerian* maps (i.e., planar maps in which every vertex has even degree) with n edges is [54]:

$$m_n = \frac{3 \cdot 2^{n-1}}{(n+1)(n+2)} \binom{2n}{n}. \quad (1)$$

A bijective explanation involving plane trees can be found in [15]. The associated generating function $M(t) = \sum_{n \geq 0} m_n t^n$ is known to be *algebraic*, that is, to satisfy a polynomial equation. More precisely:

$$t^2 + 11t - 1 - (8t^2 + 12t - 1)M(t) + 16t^2M(t)^2 = 0.$$

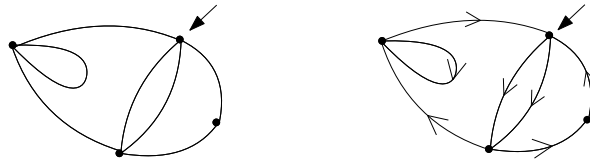


FIGURE 1. A rooted Eulerian map and a rooted Eulerian orientation.

Beyond their enumerative implications, bijections involving maps have been applied to encode, sample and draw maps efficiently [11, 21, 31, 51]. More recently, they have played a key role in the study of large random planar maps, culminating with the existence of a universal scaling limit known as the *Brownian map* [44].

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¹in disguise! The 1963 result involves *bicubic* maps, which are in one-to-one correspondence with Eulerian maps. See e.g. [15, Cor. 2.4] for the dual bijection between face-bicoloured triangulations and bipartite maps.

Planar maps equipped with an additional structure (e.g. a spanning tree [45], a proper colouring [56, 57], an Ising or Potts configuration [4, 12, 13, 16, 18, 22, 25, 40]...) are also much studied, both in combinatorics and in theoretical physics, where maps are considered as a model for two-dimensional quantum gravity [23]. However, for many of these structures, we are still in the early days of the study, as even their enumeration remains elusive, not to mention bijections and asymptotic properties.

Recent progresses in this direction include the enumeration of planar maps weighted by their Tutte polynomial, or equivalently, maps equipped with a Potts configuration. The associated generating function $P(t)$ is known to be *differentially algebraic*. That is, there exists a polynomial equation relating $P(t)$ and its derivatives [6, 7]. The Tutte polynomial has many interesting specializations (in particular, it counts all structures cited above, like spanning trees and colourings) and several special cases had been solved earlier. One key tool in the solution is that the Tutte polynomial of a map can be computed inductively, by deleting and contracting edges.

Another solved example, which does not seem to belong to the Tutte/Potts realm, consists of maps (in fact, triangulations) equipped with certain orientations called *Schnyder orientations*. The results obtained there have analogies with those obtained for another class of orientations, called *bipolar* (which *do* belong to the Tutte realm). Indeed, for both classes of oriented maps:

- oriented maps are counted by simple numbers, which are also known to count other combinatorial objects (various lattice paths and permutations, among others);
- there exist nice bijections explaining these equi-enumeration results [9, 10, 28, 33];
- for a fixed map M , the set of Schnyder/bipolar orientations of M has a lattice structure [52, 27, 46]. The above bijections, once specialized to maps equipped with their (unique) minimal orientation, coincide with attractive bijections designed earlier for (un-oriented) maps [5, 10];
- specializing the bijections further to maps that have only one Schnyder/bipolar orientation also yields interesting combinatorial results [5, 10].

These observations led us to wonder about another natural class of orientations, namely those in which every vertex has equal in- and out-degree, known as *Eulerian orientations* (Figure 1). Clearly, a map needs to be Eulerian to admit an Eulerian orientation. The condition is in fact sufficient (such maps even admit an Eulerian circuit [37]). One analogy with the above two classes is that the set of Eulerian orientations of a given planar map can be equipped with a lattice structure [52, 27]. Moreover, Eulerian maps (equivalently, Eulerian maps equipped with their minimal Eulerian orientation) have rich combinatorial properties: not only are they counted by simple numbers (see (1)), but they are equinumerous with several other families of objects, like certain trees [15] and permutations [8, 32]. And they are often related to them by beautiful bijections.

Hence our plan to count Eulerian orientations with n edges. However, this appears to be a difficult problem. In fact, we even lack a way to compute the corresponding numbers in, say, polynomial time. This leads us to resort to approximation methods that are ubiquitous when studying hard counting problems, like the enumeration of self-avoiding walks [1, 29, 36, 49], or polyominoes [2, 39, 41]: denoting by \mathcal{O} the set of Eulerian orientations, we construct subsets and supersets of \mathcal{O} , indexed by an integer parameter k , which converge to \mathcal{O} as k increases. And we count the elements of these sets.

One difference between our study and those dealing with tricky objects on regular lattices (like the above mentioned self-avoiding walks and polyominoes) is worth noting. The subsets and supersets that are defined to approximate lattice objects often have a one-dimensional structure, and *rational generating functions* that can be obtained using a transfer matrix approach. A typical example is provided by self-avoiding walks confined to a strip of fixed width. But our subsets and supersets of orientations belong to the world of maps (or *random lattices* in the physics terminology), and have *algebraic generating functions*. More precisely, our subsets have a branching, tree-like structure, which yields a system of algebraic equations for their generating functions, and a universal asymptotic behaviour in $\lambda^n n^{-3/2}$ (for a growth rate λ depending on

the index k). The generating functions of our supersets are more mysterious. They are bivariate series given by systems of equations involving *divided differences* of the form

$$\frac{F(t; x) - F(t; 1)}{x - 1},$$

and we have to resort to a deep theorem in algebra, due to Popescu [50], to prove their algebraicity for all k (we also solve these systems for small values of k). We conjecture that their asymptotic behaviour is also universal, this time in $\lambda^n n^{-5/2}$, as for planar maps (again, for varying λ).

Here is now an outline of the paper. In Section 2 we first present a simple recursive decomposition of (rooted) Eulerian orientations, based on the contraction of the root edge, and then a variant of this decomposition. Thanks to this variant, we can compute the number o_n of Eulerian orientations having n edges for $n \leq 15$ (Figure 2). By attaching two orientations at their root vertex, we see that the sequence $(o_n)_{n \geq 0}$ is super-multiplicative:

$$o_{m+n} \geq o_m o_n.$$

This classically implies that the limit μ of $o_n^{1/n}$ exists and satisfies

$$\mu = \sup_n o_n^{1/n}. \quad (2)$$

(see Fekete's Lemma in [59, p. 103]). We call μ the *growth rate* of Eulerian orientations. It is bounded from below by the growth rate 8 of Eulerian maps, and from above by the growth rate 16 of Eulerian maps equipped with an arbitrary orientation. Our data for $n \leq 15$ suggest that μ is around 12.5 (Figure 2, right). Using differential approximants [35], Tony Guttmann predicts $\mu = 12.568\dots$, and an asymptotic behaviour $o_n \sim c\mu^n n^{-\gamma}$ with $\gamma = 2.23\dots$

n	o_n	n	o_n	n	o_n
0	1	6	37 548	12	37 003 723 200
1	2	7	350 090	13	393 856 445 664
2	10	8	3 380 520	14	4 240 313 009 272
3	66	9	33 558 024	15	46 109 094 112 170
4	504	10	340 670 720		
5	4 216	11	3 522 993 656		

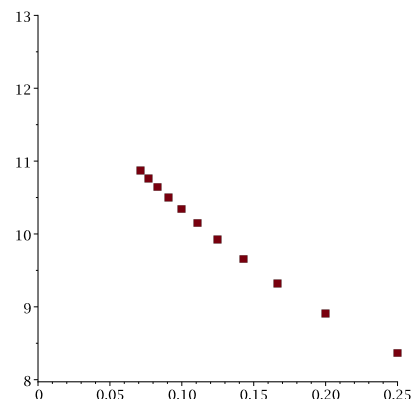


FIGURE 2. *Left*: First values of o_n , for n from 0 to 15 (entry A277493 of the OEIS [38]). *Right*: A plot of o_{n+1}/o_n vs. $1/n$, for $n = 4, \dots, 14$, suggests that the growth rate of Eulerian orientations, located at the intercept of the curve and the y -axis, is around 12.5.

Sections 3 and 4 deal with two families of subsets of Eulerian orientations. The first family uses our first recursive decomposition of Eulerian orientations, and should be considered as a warm up. The second family uses the variant of the standard decomposition of orientations. Its study is a bit more involved, but it gives better bounds on the growth constant. Both families are proved to have algebraic generating functions and a tree-like asymptotic behaviour in $\lambda^n n^{-3/2}$. The next two sections deal with two families of supersets of Eulerian orientations. Both are proved to have algebraic generating functions, and we conjecture a map-like asymptotic behaviour in $\lambda^n n^{-5/2}$. We solve our systems of equations explicitly for small values of k , and thus obtain Table 1. All calculations are supported by MAPLE sessions available on our web pages. We gather in Section 7 a few final comments and questions.

TABLE 1. Growth rates and cardinalities of subsets ($\mathcal{L}^{(k)}$ and $\mathbb{L}^{(k)}$) and supersets ($\mathcal{U}^{(k)}$ and $\mathbb{U}^{(k)}$) of Eulerian orientations. The table also records the degrees of the associated generating functions, which are systematically algebraic. The symbol \simeq refers to a numerical estimate. The other growth rates are algebraic numbers known exactly via their minimal polynomial.

	degree	growth	1	2	3	4	5	6	7
Eulerian maps	2	8	1	3	12	56	288	1 584	9 152
$\mathcal{L}^{(1)}$	2	9.68...	2	10	66	466	3 458	26 650	211 458
$\mathcal{L}^{(2)}$	4	10.16...	2	10	66	504	4 008	32 834	275 608
$\mathbb{L}^{(1)}$	3	10.60...	2	10	66	490	3 898	32 482	279 882
$\mathbb{L}^{(2)}$	6	10.97...	2	10	66	504	4 148	35 794	319 384
$\mathbb{L}^{(3)}$	20	11.22...	2	10	66	504	4 216	37 172	339 406
$\mathbb{L}^{(4)}$	258	$\simeq 11.41$	2	10	66	504	4 216	37 548	347 850
$\mathbb{L}^{(5)}$?	$\simeq 11.56$	2	10	66	504	4 216	37 548	350 090
Eulerian orientations		$\simeq 12.5$	2	10	66	504	4 216	37 548	350 090
$\mathbb{U}^{(5)}$?	$\simeq 13.005$	2	10	66	504	4 216	37 548	350 090
$\mathbb{U}^{(4)}$?	$\simeq 13.017$	2	10	66	504	4 216	37 548	350 538
$\mathbb{U}^{(3)}$?	$\simeq 13.031$	2	10	66	504	4 216	37 620	352 242
$\mathbb{U}^{(2)}$	28	13.047...	2	10	66	504	4 228	37 878	356 252
$\mathbb{U}^{(1)}$	3	13.065...	2	10	66	506	4 266	38 418	363 194
$\mathcal{U}^{(2)}$	27	13.057...	2	10	66	504	4 232	37 970	357 744
$\mathcal{U}^{(1)}$	3	13.065...	2	10	66	506	4 266	38 418	363 194
Oriented Eulerian maps	2	16	2	12	96	896	9 216	101 376	1 171 456

Let us mention that counting Eulerian orientations of 4-valent (rather than Eulerian) maps might be simpler: in this case, the number of Eulerian orientations is a specialization of the Tutte polynomial [60], and in fact some results exist in the physics literature [42, 61]. In the final section, we discuss further this problem, which seems to deserve more attention.

2. RECURSIVE DECOMPOSITIONS OF EULERIAN ORIENTATIONS

In this section, after a few basic definitions, we recall the standard recursive decomposition of Eulerian *maps* based on the contraction of the root edge, which can be traced back to the early papers of Tutte (e.g. [55]). We then adapt it to decompose Eulerian *orientations*. We also introduce a variant of the standard decomposition of Eulerian maps, based on a notion of *prime* maps, and adapt it again to Eulerian orientations. This variant is slightly more involved, but turns out to be more effective: it allows us to compute the numbers o_n for larger values of n , and it also leads to better lower and upper bounds on these numbers (Sections 4 and 6).

2.1. DEFINITIONS

A *planar map* is a proper embedding of a connected planar graph in the oriented sphere, considered up to orientation preserving homeomorphism. Loops and multiple edges are allowed (Figure 1). The *faces* of a map are the connected components of its complement. The number of edges of a planar map M is denoted by $e(M)$. The *degree* of a vertex is the number of edges incident to it, counted with multiplicity (e.g., a loop counts twice). A *corner* is a sector delimited by two consecutive edges around a vertex; hence a vertex of degree k defines k corners. Alternatively, a corner can be described as an incidence between a vertex and a face.

For counting purposes it is convenient to consider *rooted* maps. A map is rooted by choosing a corner, called the *root corner*. The vertex and face that are incident at this corner are respectively the *root vertex* and the *root face*. The *root edge* is the edge that follows the root corner in

counterclockwise order around the root vertex. In figures, we indicate the rooting by an arrow pointing to the root corner, and take the root face as the infinite face (Figure 1).

From now on, every map is *planar* and *rooted*, and these precisions will often be omitted. By convention, we include among rooted planar maps the *atomic map* having one vertex and no edge.

A map M is *Eulerian* if every vertex has even degree. In this case, we denote by $\text{dv}(M)$ the half degree of the root vertex. An *Eulerian orientation* is a map with oriented edges, in which the in- and out-degrees of every vertex are equal. Note that the underlying map must be Eulerian. We denote by \mathcal{M} the set of Eulerian maps, and by \mathcal{O} the set of Eulerian orientations.

2.2. EULERIAN MAPS: STANDARD DECOMPOSITION

Consider an Eulerian map M , not reduced to the atomic map, and its root edge e . If e is a loop, then M is obtained from two smaller maps M_1 and M_2 by joining M_1 and M_2 at their root vertices and adding a loop surrounding M_1 (Figure 3, left). The maps M_1 and M_2 are themselves Eulerian (because the sum of vertex degrees in a map is even, so that one cannot have a single odd vertex in M_1 or M_2). We call this operation the *merge* of M_1 and M_2 .

If the root edge e is not a loop, then we contract it, which gives a smaller Eulerian map M' . Note however that several maps give M' after contracting their root edge. All such maps can be obtained from M' as follows (see Figure 3, right): we split the root vertex v of M' into two vertices v and v' joined by an edge (which will be the root edge), and distribute the edges adjacent to v between v and v' in such a way the degrees of v and v' remain even. This operation is called a *split* of M' , and more precisely an *i-split* if v has degree $2i$ in the larger map. Note that if v has degree $2d$ in M' , then i can be chosen arbitrarily between 1 and d .

Let $M(t; x)$ be the generating function of Eulerian maps, counted by edges (variable t) and by the half degree of the root vertex (variable x):

$$M(t; x) = \sum_{M \in \mathcal{M}} t^{e(M)} x^{\text{dv}(M)} = \sum_{d \geq 0} x^d M_d(t),$$

where $M_d(t)$ denotes the edge generating function of Eulerian maps with root vertex degree $2d$. The above construction translates into the following functional equation, which we explain below:

$$\begin{aligned} M(t; x) &= 1 + txM(t; x)^2 + t \sum_{d \geq 0} M_d(t)(x + x^2 + \cdots + x^d) \\ &= 1 + txM(t; x)^2 + t \sum_{d \geq 0} M_d(t) \frac{x^{d+1} - x}{x - 1} \\ &= 1 + txM(t; x)^2 + \frac{tx}{x - 1} (M(t; x) - M(t; 1)). \end{aligned} \quad (3)$$

On the first line, the term 1 accounts for the atomic map, the next term for maps obtained by merging two smaller maps, and the third term for maps obtained from a split.

2.3. EULERIAN ORIENTATIONS: STANDARD DECOMPOSITION

Our recursive decomposition for Eulerian orientations is essentially the same as for Eulerian maps: if the root edge is a loop, we delete it and obtain two orientations, which are both Eulerian (in any oriented map, the sum over all vertices of in-degrees equals the sum of out-degrees, hence one cannot have a single vertex with distinct in- and out-degrees); otherwise we contract the root edge, which gives a smaller Eulerian orientation.

However, care must be taken when going in the opposite direction, that is, when constructing large orientations from smaller ones. The first type of orientations, obtained by a merge, do not raise any difficulty; one can orient the new root edge (the loop) in two different ways (Figure 4, left). But consider now an Eulerian orientation O' , with root vertex v of degree at least $2i$, and perform an *i-split* on O' : is there a way to orient the new edge so as to obtain an orientation O that is still Eulerian? The answer is yes if and only if the numbers of in- and out-edges in the

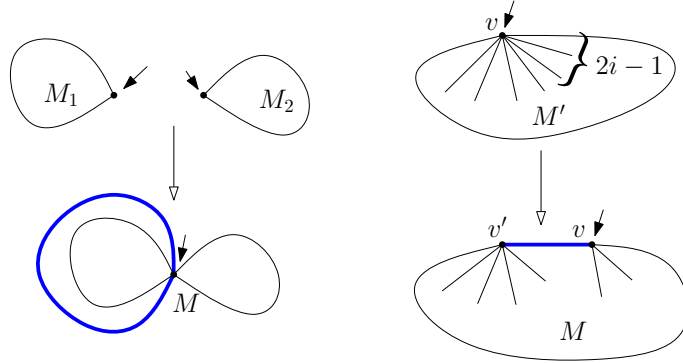


FIGURE 3. Construction of an Eulerian map with n edges: merge an ordered pair of Eulerian maps M_1, M_2 with n_1 and n_2 edges ($n_1 + n_2 = n - 1$) and add a loop, or make a split on an Eulerian map with $n - 1$ edges. The new edge (here thicker) is the root edge of M .

last $2i - 1$ edges incident to v in O' differ by ± 1 (edges are visited in counterclockwise order, starting from the root corner). The orientation of the root edge of O is then forced (Figure 4, right). In this case, we say that the i -split, performed on O' , is *legal*. Note that the 1-split and the d -split are always legal, where $2d$ is the degree of the root vertex of O' .

The fact that not all splits are legal makes it difficult to write a single functional equation for the generating function of Eulerian orientations. However, we can write an *infinite system* of equations relating the generating functions of orientations with prescribed orientations at the root.

Let us be more precise. Given an Eulerian orientation O with root vertex v of degree $2d$, the *root word* $w(O)$ of O is a word of length $2d$ on the alphabet $\{0, 1\}$ describing the orientation of the edges around v (in counterclockwise order, starting from the root corner): the k -th letter of $w(O)$ is 0 (resp. 1) if the k -th edge around v is in-going (resp. out-going). Note that this word is always *balanced*, meaning that it contains as many 0's as 1's. We call a word \mathbf{w} *quasi-balanced* if the number of 0's and 1's in \mathbf{w} differ by ± 1 . The length (number of letters) of \mathbf{w} is denoted by $|\mathbf{w}|$, while the number of occurrences of the letter a in it is denoted by $|\mathbf{w}|_a$. We define the *balance* of \mathbf{w} to be $b(\mathbf{w}) := ||\mathbf{w}|_1 - |\mathbf{w}|_0|$. The empty word is denoted by ε .

Now we can decide from the root word of O' if the i -split of O' is legal: this holds if and only if the last $2i - 1$ letters of $w(O')$ form a quasi-balanced word.

For \mathbf{w} a word on $\{0, 1\}$, let $O_{\mathbf{w}}(t) \equiv O_{\mathbf{w}}$ be the generating function of Eulerian orientations having \mathbf{w} as root word, counted by their edge number. Clearly, $O_{\mathbf{w}} = 0$ if \mathbf{w} is not balanced and $O_{\varepsilon} = 1$ (accounting for the atomic map). Now if \mathbf{w} is non-empty and balanced,

$$O_{\mathbf{w}} = t \sum_{a\bar{u}\bar{a}\mathbf{v}=\mathbf{w}} O_{\mathbf{u}}O_{\mathbf{v}} + t \sum_{\mathbf{u}} O_{\mathbf{u}\mathbf{w}_s}. \quad (4)$$

This identity is illustrated in Figure 4. Here, a stands for any of the letters 0, 1, and the first sum runs over all factorisations of \mathbf{w} of the form $a\bar{u}\bar{a}\mathbf{v}$, with $\bar{a} := 1 - a$. This sum counts orientations obtained by a merge. The second sum runs over all possible words \mathbf{u} , and \mathbf{w}_s denotes the suffix of \mathbf{w} of length $|\mathbf{w}| - 1$. This sum counts orientations obtained by a (legal) split of an orientation having root word $\mathbf{u}\mathbf{w}_s$. Now the generating function O of Eulerian orientations is

$$\sum_{\mathbf{w}} O_{\mathbf{w}},$$

where the sum runs over all (balanced) words \mathbf{w} .

We do not know how to solve this system. But a map with n edges has a root word of length at most $2n$, and hence we can use our system to compute the numbers o_n for n small. We obtain in this way the first 11 values of Figure 2.

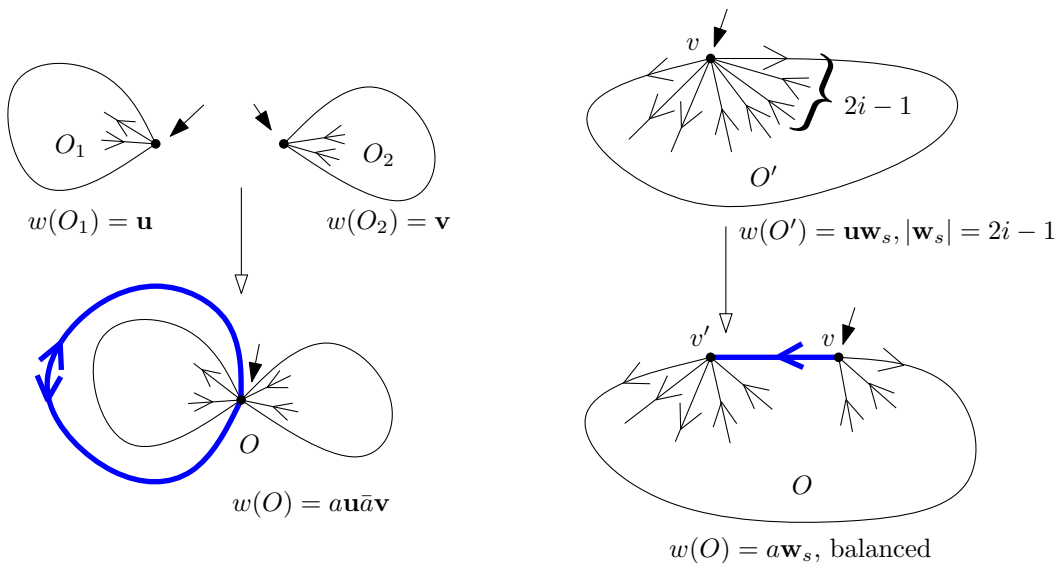


FIGURE 4. Construction of an Eulerian orientation: merge two Eulerian orientations (the loop can be oriented in either way), or split (legally) an Eulerian orientation. Observe how the root word changes.

In Sections 3 to 6, we define subsets and supersets of \mathcal{O} that we can generate by just looking at the last $2k - 1$ letters of the root word (for k fixed). This allows us to write finitely many equations for the generating functions of these subsets and supersets. Solving them gives lower and upper bounds on the growth rate of Eulerian orientations. However, we obtain more precise bounds by using a variant of the standard decomposition of maps and orientations. We now present this variant.

2.4. PRIME DECOMPOSITION OF MAPS AND ORIENTATIONS

A (non-atomic) map is said to be *prime* if the root vertex appears only once when walking around the root face. A planar map M can be seen as a sequence of prime maps M_1, \dots, M_ℓ (Figure 5). We say that the M_i are the *prime submaps* of M , and denote $M = M_1 \cdots M_\ell$. Note that if M is Eulerian, then each M_i is Eulerian too.

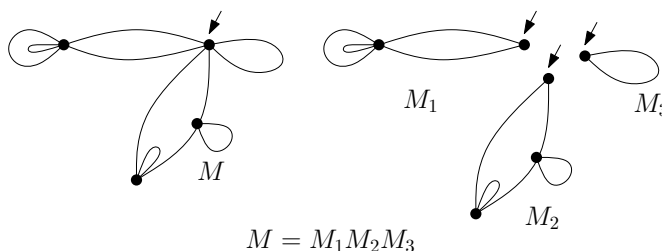


FIGURE 5. Decomposition of an Eulerian map M into prime Eulerian maps M_1, M_2, M_3 .

Now take a *prime* Eulerian map M , and apply the standard decomposition of Section 2.2, illustrated in Figure 3: either M is an (arbitrary) Eulerian map M_1 surrounded by a loop, or M is obtained by an i -split in another Eulerian map M' , provided the last prime submap of M' (in counterclockwise order) has root degree at least $2i$ (otherwise, the resulting map would not be prime). Alternatively, if $M' = M'_1 \cdots M'_\ell$, we can obtain M by performing an i -split in the

prime map M'_ℓ , and attaching the map $M'' := M'_1 \cdots M'_{\ell-1}$ at the new vertex v' created by this split (Figure 6).

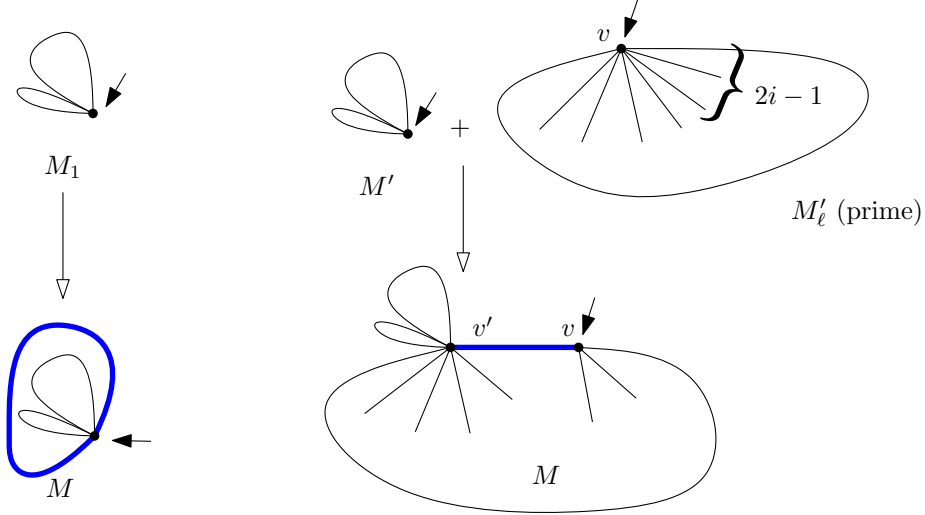


FIGURE 6. Construction of a prime Eulerian map: add a loop around any Eulerian map, or split a prime Eulerian map, and attach an arbitrary Eulerian map at the end of its root edge.

This alternative decomposition of Eulerian maps gives a system of two equations defining the generating function $M(t; x)$ of Eulerian maps (still counted by edges and root vertex degree) and its counterpart $M'(t; x)$ for prime maps:

$$\begin{aligned} M(t; x) &= 1 + M(t; x)M'(t; x), \\ M'(t; x) &= txM(t; x) + txM(t; 1) \frac{M'(t; x) - M'(t; 1)}{x - 1}. \end{aligned}$$

In the first equation, the term $M'(t; x)$ accounts for the last prime submap attached at the root vertex (denoted M_ℓ above). In the second equation, the divided difference $(M'(t; x) - M'(t; 1))/(x - 1)$ has the same explanation as in (3). This equation is easily recovered by eliminating $M'(t; x)$ from the above system.

This decomposition can also be applied to Eulerian orientations: an Eulerian orientation is a sequence of prime Eulerian orientations, and a prime orientation is either obtained by adding an oriented loop around another orientation, or by performing a legal split in a prime orientation, and attaching another orientation at the vertex v' created by the split.

Thus, denoting again by $O_{\mathbf{w}}$ the generating function of orientations with root word \mathbf{w} , and by $O'_{\mathbf{w}}$ its counterpart for prime orientations, we now have $O_{\mathbf{w}} = O'_{\mathbf{w}} = 0$ if \mathbf{w} is not balanced, $O_\varepsilon = 1$, $O'_\varepsilon = 0$ and finally for \mathbf{w} balanced and non-empty,

$$\begin{aligned} O_{\mathbf{w}} &= \sum_{\mathbf{u}\mathbf{v}=\mathbf{w}} O_{\mathbf{u}}O'_{\mathbf{v}}, \\ O'_{\mathbf{w}} &= tO_{\mathbf{w}_c} + tO \sum_{\mathbf{u}} O'_{\mathbf{u}\mathbf{w}_s}. \end{aligned}$$

In the second equation, \mathbf{w}_c denotes the central factor or \mathbf{w} of length $|\mathbf{w}| - 2$, and $O = \sum_{\mathbf{w}} O_{\mathbf{w}}$ is the generating function of all Eulerian orientations. Recall that \mathbf{w}_s is the suffix of \mathbf{w} of length $|\mathbf{w}| - 1$.

Using these equations, we have been able to push further the enumeration of Eulerian orientations of small size, thus obtaining the values of Figure 2.

3. SUBSETS OF EULERIAN ORIENTATIONS, VIA THE STANDARD DECOMPOSITION

In this section and the three following ones, we define certain subsets and supersets of Eulerian orientations, indexed by an integer k , which converge (monotonously) to the set \mathcal{O} of all Eulerian orientations as k tends to infinity. Those sets are respectively denoted by $\mathcal{L}^{(k)}$ (and $\mathbb{L}^{(k)}$), $\mathcal{U}^{(k)}$ (and $\mathbb{U}^{(k)}$), as they give *lower* and *upper* bounds on the numbers o_n and their growth rate. For each set, we give a system of functional equations defining its generating function: for the subsets $\mathcal{L}^{(k)}$ and $\mathbb{L}^{(k)}$, these systems are algebraic, so that the associated generating functions are algebraic series. For the supersets $\mathcal{U}^{(k)}$ and $\mathbb{U}^{(k)}$, the systems define bivariate series and involve divided differences as in (3). However, we prove that the resulting series are also algebraic.

Recall from Figure 4 that planar Eulerian orientations can be obtained recursively from the atomic map by either:

- the merge of two orientations $O_1, O_2 \in \mathcal{O}$ (with the root loop oriented in either way),
- or a legal split on an orientation $O' \in \mathcal{O}$.

Definition 1. Let $k \geq 1$. Let $\mathcal{L}^{(k)}$ be the set of planar orientations obtained recursively from the atomic map by either:

- the merge of two orientations $O_1, O_2 \in \mathcal{L}^{(k)}$ (with the root loop oriented in either way),
- or a legal i -split on an orientation $O' \in \mathcal{L}^{(k)}$ such that $i \leq k$ or $i = \text{dv}(O')$.

In other words, the only allowed splits are the small splits ($i \leq k$) and the maximal split ($i = \text{dv}(O')$).

Obviously, all orientations of $\mathcal{L}^{(k)}$ are Eulerian. Moreover, the sets $\mathcal{L}^{(k)}$ form an increasing sequence since more and more (legal) splits are allowed as k grows. Finally, all Eulerian orientations of size n belong to $\mathcal{L}^{(n)}$ (and even to $\mathcal{L}^{(n-2)}$). Hence the limit of the sets $\mathcal{L}^{(k)}$ is the set \mathcal{O} of all Eulerian orientations.

Figure 7 shows a (random) orientation of $\mathcal{L}^{(1)}$.

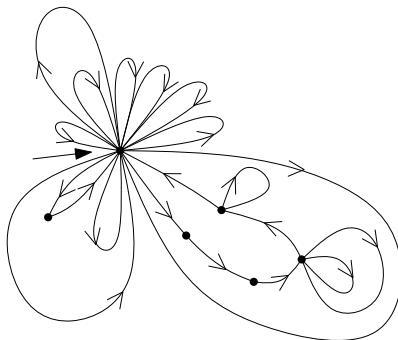


FIGURE 7. An Eulerian orientation in $\mathcal{L}^{(1)}$, taken uniformly at random among those with 20 edges.

3.1. AN ALGEBRAIC SYSTEM FOR $\mathcal{L}^{(k)}$

In this section, k is a fixed integer.

Definition 2. A word \mathbf{w} on $\{0, 1\}$ is valid (for k) if there exists a balanced word of length $2k$ having \mathbf{w} as a factor. Equivalently, the balance of \mathbf{w} satisfies $b(\mathbf{w}) \leq 2k - |\mathbf{w}|$. This holds automatically if $|\mathbf{w}| \leq k$.

Given a word \mathbf{w} , it will be convenient to have notation for several words that differ from \mathbf{w} by one or two letters. We have already defined \mathbf{w}_c , the central factor of \mathbf{w} of length $|\mathbf{w}| - 2$, and \mathbf{w}_s , the suffix of \mathbf{w} of length $|\mathbf{w}| - 1$. We similarly define \mathbf{w}_p as the prefix of \mathbf{w} of length $|\mathbf{w}| - 1$.

Finally, if \mathbf{w} is quasi-balanced, then $\overleftarrow{\mathbf{w}}$ stands for the unique balanced word of the form $a\mathbf{w}$, for $a \in \{0, 1\}$.

For any word \mathbf{w} , we denote by $L_{\mathbf{w}}^{(k)}(t)$ the generating function of orientations of $\mathcal{L}^{(k)}$ whose root word ends with \mathbf{w} , counted by edges. In particular, the generating function counting all orientations of $\mathcal{L}^{(k)}$ is $L_{\varepsilon}^{(k)}(t)$. We denote by $K_{\mathbf{w}}^{(k)}(t)$ the generating function of orientations of $\mathcal{L}^{(k)}$ having root word exactly \mathbf{w} . In order to lighten notation, we often omit the dependence of our series in t and the superscript (k) .

We now give equations defining the series $L_{\mathbf{w}}$ (for $|\mathbf{w}| \leq 2k - 1$) and $K_{\mathbf{w}}$ (for $|\mathbf{w}| \leq 2k$). First, we note that $K_{\mathbf{w}} = 0$ if \mathbf{w} is not balanced, and that $K_{\varepsilon} = 1$. Now for \mathbf{w} balanced of length between 2 and $2k$, we have:

$$K_{\mathbf{w}} = t \sum_{\mathbf{w}=\mathbf{au}\bar{\mathbf{a}}\mathbf{v}} K_{\mathbf{u}}K_{\mathbf{v}} + tL_{\mathbf{w}_s}, \quad (5)$$

where, as before, a is any of the letters 0, 1. This equation is analogous to (4): the first term counts orientations obtained from a merge, the second orientations obtained from a split. Now for $L_{\mathbf{w}}$, with \mathbf{w} of length at most $2k - 1$, we have:

$$L_{\mathbf{w}} = \mathbf{1}_{\mathbf{w}=\varepsilon} + 2tL_{\varepsilon}L_{\mathbf{w}} + t \sum_{\mathbf{w}=\mathbf{u}\mathbf{a}\mathbf{v}} L_{\mathbf{u}}K_{\mathbf{v}} + t \sum_{\mathbf{w}=\mathbf{au}\bar{\mathbf{a}}\mathbf{v}} K_{\mathbf{u}}K_{\mathbf{v}} + t(L_{\mathbf{w}} - \mathbf{1}_{\mathbf{w}=\varepsilon}) + t \sum_{\substack{\mathbf{u}=\mathbf{v}\mathbf{w} \\ 2 \leq |\mathbf{u}| \leq 2k \\ \mathbf{u} \text{ balanced}}} (L_{\mathbf{u}_s} - K_{\mathbf{u}}). \quad (6)$$

This equation deserves some explanations. The first line counts the atomic map (if $\mathbf{w} = \varepsilon$), and the orientations obtained by a merge. The second (resp. third, fourth) term of this line counts orientations such that no (resp. one, both) half-edge(s) of the root loop is/are involved in the suffix \mathbf{w} of the root word. Equivalently, denoting by O_1 and O_2 the merged orientations, those three terms respectively correspond to $|\mathbf{w}| \leq 2 \operatorname{dv}(O_2)$, $2 \operatorname{dv}(O_2) < |\mathbf{w}| < 2 + 2 \operatorname{dv}(O_1) + 2 \operatorname{dv}(O_2)$ and $|\mathbf{w}| = 2 + 2 \operatorname{dv}(O_1) + 2 \operatorname{dv}(O_2)$.

The second line counts orientations O obtained by a legal i -split in a smaller orientation O' . The first term accounts for maximal splits ($i = \operatorname{dv}(O')$), which, we recall, do not change the root word (note also that no split is possible on the atomic map). The second term counts orientations O obtained from a non-maximal split. The word \mathbf{u} stands for the root word of O . The subtraction of $K_{\mathbf{u}}$ comes from the condition that the split is not maximal.

Proposition 3. *Consider the collection of equations consisting of:*

- Equation (5), written for all balanced words \mathbf{w} of length between 2 and $2k$,
- Equation (6), written for all valid words \mathbf{w} of length at most $2k - 1$.

In this collection, replace all trivial K -series by their value: $K_{\mathbf{w}} = 0$ when \mathbf{w} is not balanced, $K_{\varepsilon} = 1$. Let S_0 denote the resulting system. The number of series it involves is

$$f(k) = \binom{2k+2}{k+1} - 1 + \sum_{i=1}^{k-1} \binom{2i}{i}. \quad (7)$$

The system S_0 defines uniquely these $f(k)$ series. Its size can be (roughly) divided by two upon noticing that replacing all 0's by 1's, and vice-versa, in a word \mathbf{w} , does not change the series $L_{\mathbf{w}}$ nor $K_{\mathbf{w}}$.

Proof. To see that S_0 defines all the series that it involves, it suffices to note the factor t in the right-hand sides of (5) and (6), and to check that each series occurring in the right-hand side of some equation also occurs as the left-hand side of another. This is readily done, as any factor of a valid word is still valid.

Let us now count the equations of the system. The number of non-empty balanced words of length at most $2k$ is

$$\sum_{i=1}^k \binom{2i}{i}.$$

Then, all words of length at most k are valid, while the number of valid words of length $k+i$, for $1 \leq i \leq k-1$, is

$$\sum_{j=i}^k \binom{k+i}{j}.$$

(One can interpret j as the number of occurrences of 0 in the word.) Hence the number of equations in the system is

$$f(k) = \sum_{i=1}^k \binom{2i}{i} + \sum_{i=0}^k 2^i + \sum_{i=1}^{k-1} \sum_{j=i}^k \binom{k+i}{j}.$$

The second sum evaluates to $2^{k+1} - 1$. The third one is

$$\begin{aligned} \sum_{j=1}^k \sum_{i=1}^{\min(j, k-1)} \binom{k+i}{j} &= \sum_{j=1}^{k-1} \sum_{i=1}^j \binom{k+i}{j} + \sum_{i=1}^{k-1} \binom{k+i}{k} \\ &= \left(1 + \binom{2k+1}{k} - 2^{k+1}\right) + \left(\frac{k}{k+1} \binom{2k}{k} - 1\right) \\ &= \frac{3k+1}{k+1} \binom{2k}{k} - 2^{k+1}. \end{aligned}$$

The sums are evaluated using classical summation identities, or Gosper's algorithm [47]. The expression of $f(k)$ given in the proposition then follows after elementary manipulations. \square

Remark 4. If \mathbf{w} is such that $0\mathbf{w}$ and $1\mathbf{w}$ are both valid of length less than $2k$, we can define $L_{\mathbf{w}}$ by a simpler "forward" equation, without increasing the size of the system:

$$L_{\mathbf{w}} = K_{\mathbf{w}} + L_{0\mathbf{w}} + L_{1\mathbf{w}}. \quad (8)$$

This is obviously smaller than (6), and possibly better suited to feed a computer algebra system. However, mixing equations of type (6) and (8) makes some proofs of Section 3.3 heavier.

3.2. EXAMPLES

3.2.1. **When $k = 1$,** the system S_0 contains $f(1) = 5$ equations and reads

$$\begin{cases} K_{01} = tK_{\varepsilon}K_{\varepsilon} + tL_1, \\ K_{10} = tK_{\varepsilon}K_{\varepsilon} + tL_0, \\ L_{\varepsilon} = 1 + 2tL_{\varepsilon}L_{\varepsilon} + t(L_{\varepsilon} - 1) + t(L_0 - K_{10} + L_1 - K_{01}), \\ L_0 = 2tL_{\varepsilon}L_0 + tL_{\varepsilon}K_{\varepsilon} + tL_0 + t(L_0 - K_{10}), \\ L_1 = 2tL_{\varepsilon}L_1 + tL_{\varepsilon}K_{\varepsilon} + tL_1 + t(L_1 - K_{01}), \end{cases} \quad (9)$$

with $K_{\varepsilon} = 1$. Using the 0/1 symmetry, this system can be compacted into

$$\begin{cases} K_{01} = t + tL_0, \\ L_{\varepsilon} = 1 + 2tL_{\varepsilon}L_{\varepsilon} + t(L_{\varepsilon} - 1) + 2t(L_0 - K_{01}), \\ L_0 = 2tL_{\varepsilon}L_0 + tL_{\varepsilon} + tL_0 + t(L_0 - K_{01}). \end{cases}$$

The variant mentioned in Remark 4 consists in replacing the second equation by $L_{\varepsilon} = 1 + 2L_0$. The reader may check that this is consistent with the above system.

This leads us to define, for \mathbf{w} quasi-balanced of length less than $2k$, the series $L_{\mathbf{w}}^+ := L_{\mathbf{w}} - K_{\overleftarrow{\mathbf{w}}}$. We will also need to define, for \mathbf{w} balanced, $L_{\mathbf{w}}^+ := L_{\mathbf{w}} - K_{\mathbf{w}}$. These series have natural combinatorial interpretations in terms of orientations whose root word ends *strictly* with \mathbf{w} (if \mathbf{w} is balanced) or $\overleftarrow{\mathbf{w}}$ (if \mathbf{w} is quasi-balanced). Then we alter the original system S_0 as follows.

- (i) For \mathbf{w} balanced or quasi-balanced, we replace the equation (6) defining $L_{\mathbf{w}}$ by an equation defining $L_{\mathbf{w}}^+$:

$$L_{\mathbf{w}}^+ = 2tL_{\varepsilon}L_{\mathbf{w}} + t \sum_{\mathbf{w}=\mathbf{uav}} (L_{\mathbf{u}} - K_{\mathbf{u}})K_{\mathbf{v}} + tL_{\mathbf{w}}^+ + t \sum_{\substack{\mathbf{u}=\mathbf{vw}, \mathbf{u} \neq \mathbf{w} \\ |\mathbf{u}| \leq 2k-1 \\ \mathbf{u} \text{ quasi-balanced}}} L_{\mathbf{u}}^+. \quad (13)$$

To obtain it, either we get back to the explanation of (6) and remove from its right-hand side the terms that count orientations with root word exactly \mathbf{w} (if \mathbf{w} is balanced) or $\overleftarrow{\mathbf{w}}$ (if \mathbf{w} is quasi-balanced). Or we simply subtract from (6) Equation (5), written for \mathbf{w} if \mathbf{w} is balanced, for $\overleftarrow{\mathbf{w}}$ if \mathbf{w} is quasi-balanced.

- (ii) In the new system thus obtained, we replace every series $K_{\mathbf{w}}$ such that \mathbf{w} is not balanced by 0, every series $L_{\mathbf{w}}$ such that \mathbf{w} is balanced by $K_{\mathbf{w}} + L_{\mathbf{w}}^+$, and every series $L_{\mathbf{w}}$ such that \mathbf{w} is quasi-balanced by $K_{\overleftarrow{\mathbf{w}}} + L_{\mathbf{w}}^+$. In particular, the series $L_{\mathbf{u}} - K_{\mathbf{u}}$ occurring in (13) becomes $L_{\mathbf{u}}^+$ when \mathbf{u} is balanced, $L_{\mathbf{u}}$ otherwise. The only series $L_{\mathbf{w}}$ that remain in the system are such that the balance of \mathbf{w} is at least 2.

We thus obtain a positive system, denoted S_1 , defining the following series:

- $K_{\mathbf{w}}$, for \mathbf{w} balanced of length between 2 and $2k$,
- $L_{\mathbf{w}}^+$, for \mathbf{w} balanced or quasi-balanced of length less than $2k$,
- $L_{\mathbf{w}}$, for \mathbf{w} valid of length less than $2k$ and balance at least 2.

For instance, when $k = 1$, the system (9) becomes (after exploiting the 0/1 symmetry):

$$\begin{cases} K_{01} = t + t(K_{01} + L_0^+), \\ L_{\varepsilon}^+ = 2t(1 + L_{\varepsilon}^+)^2 + tL_{\varepsilon}^+ + 2tL_0^+, \\ L_0^+ = 2t(1 + L_{\varepsilon}^+)(K_{01} + L_0^+) + tL_{\varepsilon}^+ + tL_0^+. \end{cases}$$

Similarly, when $k = 2$, the system (11) is replaced by:

$$\begin{cases} K_{10} = K_{01} = t + t(K_{10} + L_0^+), \\ K_{1100} = tK_{10} + t(K_{1100} + L_{100}^+), \\ K_{1010} = t(K_{10} + K_{01}) + t(K_{1010} + L_{010}^+), \\ K_{0110} = tK_{01} + t(K_{0110} + L_{110}^+), \\ L_{\varepsilon}^+ = 2t(1 + L_{\varepsilon}^+)^2 + tL_{\varepsilon}^+ + 2t(L_0^+ + L_{100}^+ + L_{010}^+ + L_{110}^+), \\ L_0^+ = L_1^+ = 2t(1 + L_{\varepsilon}^+)(K_{10} + L_0^+) + tL_{\varepsilon}^+ + tL_0^+ + t(L_{100}^+ + L_{010}^+ + L_{110}^+), \\ L_{00} = L_{11} = 2t(1 + L_{\varepsilon}^+)L_{00} + t(K_{10} + L_0^+) + tL_{00} + tL_{100}^+, \\ L_{10}^+ = L_{01}^+ = 2t(1 + L_{\varepsilon}^+)(K_{10} + L_{10}^+) + t(K_{01} + L_1^+) + tL_{10}^+ + t(L_{110}^+ + L_{010}^+), \\ L_{100}^+ = 2t(1 + L_{\varepsilon}^+)(K_{1100} + L_{100}^+) + tL_{10}^+ + tL_{100}^+, \\ L_{010}^+ = 2t(1 + L_{\varepsilon}^+)(K_{1010} + L_{010}^+) + t(L_{01}^+ + L_{\varepsilon}^+K_{10}) + tL_{010}^+, \\ L_{110}^+ = 2t(1 + L_{\varepsilon}^+)(K_{0110} + L_{110}^+) + t(L_{11} + L_{\varepsilon}^+K_{10}) + tL_{110}^+. \end{cases}$$

The second condition that we need is *properness* (again, in the sense of [30, Sec. VII.6.3]). But the system S_1 that we have just obtained is proper, thanks to the factor t occurring in the right-hand side of (5), (6) and (13).

The next condition is aperiodicity. The coefficients of t^1 and t^2 in the series $L_{\varepsilon}^+(t)$ are both non-zero. This implies that this series is aperiodic. Consequently, if we prove that the system S_1 is *irreducible* (which will be our final objective below), then it will be aperiodic [30, p. 483].

So let us finally prove that S_1 is irreducible. Recall that in such a polynomial system, a series F depends on a series G if G occurs in the right-hand side of the equation defining F . Irreducibility means that the digraph of dependences is strongly connected. Recall that S_1 involves two families of series: the series $K_{\mathbf{w}}$, for \mathbf{w} balanced of length between 2 and $2k$, and $L_{\mathbf{w}}$ (or $L_{\mathbf{w}}^+$) for any valid \mathbf{w} of length at most $2k-1$. To lighten notation, for any \mathbf{w} we denote by $\tilde{L}_{\mathbf{w}}$ the corresponding L -series, be it $L_{\mathbf{w}}$ or $L_{\mathbf{w}}^+$.

Let us first prove that every series in S_1 depends on $\tilde{L}_{\varepsilon} = L_{\varepsilon}^+$. By (6) and (13), this holds for every $\tilde{L}_{\mathbf{w}}$. Now, each $K_{\mathbf{w}}$ depends on at least one \tilde{L} -series (see (5)), and thus by transitivity on L_{ε}^+ .

Conversely, let us prove that L_{ε}^+ depends on all other series occurring in S_1 .

- First, Equation (13) applied to $\mathbf{w} = \varepsilon$ shows that L_{ε}^+ depends on all series $\tilde{L}_{\mathbf{u}}$ such that \mathbf{u} is quasi-balanced.
- Let us now prove, by induction on the balance $b(\mathbf{u})$, that L_{ε}^+ depends on $\tilde{L}_{\mathbf{u}}$ for each valid word \mathbf{u} of length at most $2k-1$. We have already seen this for $b(\mathbf{u}) = 1$. If $b(\mathbf{u}) = 0$, then $|\mathbf{u}| \leq 2k-2$, and for any letter a the word $\mathbf{w} := \mathbf{u}a$ is valid and quasi-balanced. The second term in (13) shows that $\tilde{L}_{\mathbf{w}}$ depends on $\tilde{L}_{\mathbf{u}}$. By transitivity, this implies that L_{ε}^+ depends on $\tilde{L}_{\mathbf{u}}$. We have thus set the initial cases of our induction, for balances 0 and 1. Now assume $b(\mathbf{u}) \geq 2$. There exists a letter a such that $\mathbf{w} := \mathbf{u}a$ is valid and has balance $b(\mathbf{u}) - 1$. If $b(\mathbf{w}) = 1$ (resp. $b(\mathbf{w}) > 1$), the second (resp. third) term in the equation (13) (resp. (6)) defining $\tilde{L}_{\mathbf{w}}$ shows that $\tilde{L}_{\mathbf{w}}$ depends on $\tilde{L}_{\mathbf{u}}$. By the induction hypothesis, L_{ε}^+ depends on $\tilde{L}_{\mathbf{w}}$, and thus by transitivity on $\tilde{L}_{\mathbf{u}}$.
- Finally, let \mathbf{u} be balanced of length between 2 and $2k$. Then $\mathbf{w} := \mathbf{u}_s$ is quasi-balanced. The first term of the equation (13) defining $\tilde{L}_{\mathbf{w}}$ involves $L_{\mathbf{w}} = K_{\mathbf{u}} + \tilde{L}_{\mathbf{w}}$, so that by transitivity, L_{ε}^+ depends on $K_{\mathbf{u}}$.

We have now checked all conditions of Theorem VII.6 of [30, p. 489]. Applying it gives our proposition. \square

3.4. BACK TO EXAMPLES

We now return to the cases $k = 1$ and $k = 2$ studied in Section 3.2. We refer to [30, Sec. VII.7] for generalities on the singularities of algebraic series, and on the asymptotic behaviour of their coefficients. When $k = 1$, we have obtained for L_{ε} the quadratic equation (10). Its dominant coefficient only vanishes at $t = 0$, and its discriminant is $\Delta_1(t) := t^4 + 4t^3 + 22t^2 - 12t + 1$. The radius ρ_1 must be one of the roots of Δ_1 . The only real positive roots are around 0.1032 and 0.3998. By solving (10) explicitly, we see that the smallest of these roots is indeed a singularity of L_{ε} . Hence $\rho_1 = 0.1032\dots$ and the corresponding growth rate is $\lambda_1 = 1/\rho_1 = 9.684\dots$, which improves on the lower bound 8 coming from Eulerian maps.

When $k = 2$, we have obtained for L_{ε} the quartic equation (12). Its dominant coefficient does not vanish away from 0, and its discriminant is

$$\begin{aligned} \Delta_2(t) := & 64t^{12}(t-1)(81t^{21} + 1863t^{20} + 11322t^{19} + 38592t^{18} + 101105t^{17} + 226631t^{16} + \\ & 393423t^{15} + 532907t^{14} + 665167t^{13} + 719797t^{12} + 454804t^{11} + 355710t^{10} + 360159t^9 - 262135t^8 - \\ & 239969t^7 + 723151t^6 - 1106764t^5 + 820832t^4 - 316644t^3 + 65424t^2 - 6780t + 268). \end{aligned}$$

The only roots in $(0,1)$ are 0.0984... and 0.2714.... The radius ρ_2 must be the first one (the other would give a growth rate smaller than 8). Hence the corresponding growth rate is $\lambda_2 = 1/\rho_2 = 10.16\dots$, which improves on the previous bound λ_1 .

We do not push our study to larger values of k , as we will obtain better bounds with the prime decomposition in the next section.

4. SUBSETS OF EULERIAN ORIENTATIONS, VIA THE PRIME DECOMPOSITION

In this section, we combine the restriction on allowed splits of the previous section with the prime decomposition of Section 2.4 to obtain a new family of subsets of Eulerian orientations. The results and proofs are similar to those of the previous section, and we give fewer details. The new subsets $\mathbb{L}^{(k)}$ satisfy $\mathcal{L}^{(k)} \subset \mathbb{L}^{(k)}$ (Proposition 8), hence they give better lower bounds on the growth rate μ than those obtained in the previous section. Moreover these bounds increase to μ as k increases (Proposition 7).

Recall from Section 2.4 that an Eulerian orientation is a sequence of prime Eulerian orientations, and that a prime (Eulerian) orientation can be obtained recursively from the atomic map by either:

- adding a loop, oriented in either way, around an orientation O_1 ,
- or performing a legal split on a prime orientation $O' \in \mathcal{O}$, followed by the concatenation of an arbitrary Eulerian orientation O'' at the new vertex created by the split (Figure 6).

Definition 6. Let $k \geq 1$. Let $\mathbb{L}^{(k)}$ be the set of planar orientations obtained recursively from the atomic map by either:

- concatenating a sequence of prime orientations of $\mathbb{L}^{(k)}$,
- or adding a loop, oriented in either way, around an orientation O_1 of $\mathbb{L}^{(k)}$,
- or performing a legal i -split on a prime orientation $O' \in \mathbb{L}^{(k)}$, with $i = \text{dv}(O')$ or $i \leq k$, followed by the concatenation of an arbitrary orientation of $\mathbb{L}^{(k)}$ at the new vertex created by the split.

Clearly, the sets $\mathbb{L}^{(k)}$ increase to the set \mathcal{O} of all Eulerian orientations as k increases, hence their growth rates $\bar{\lambda}_k$ form a non-decreasing sequence of lower bounds on μ . But we have in this case a stronger result.

Proposition 7. For $k \geq 1$, the sequence $(\bar{\ell}_n^{(k)})_{n \geq 0}$ that counts orientations of $\mathbb{L}^{(k)}$ by their size is super-multiplicative. Consequently, the associated growth rate

$$\bar{\lambda}_k := \lim_n (\bar{\ell}_n^{(k)})^{1/n} = \sup_n (\bar{\ell}_n^{(k)})^{1/n} \quad (14)$$

increases to μ as k tends to infinity.

Proof. By definition of $\mathbb{L}^{(k)}$, concatenating two orientations of $\mathbb{L}^{(k)}$ at their root vertex gives a new element of $\mathbb{L}^{(k)}$, which implies super-multiplicativity and the identity (14) (by Fekete's Lemma [59, p. 103]).

Now since $\mathbb{L}^{(k)}$ converges to \mathcal{O} , for any n , there exists k such that $o_n = \ell_n^{(k)}$ (one can take $k = n$, or even $k = n - 2$). Hence

$$o_n^{1/n} = (\bar{\ell}_n^{(k)})^{1/n} \leq \bar{\lambda}_k \leq \lim_k \bar{\lambda}_k,$$

and it follows now from (2) that $\mu \leq \lim_k \bar{\lambda}_k$. Since $\bar{\lambda}_k \leq \mu$, the proposition follows. \square

Proposition 8. For $k \geq 1$, the subset of orientations $\mathbb{L}^{(k)}$ includes the subset $\mathcal{L}^{(k)}$ defined in Section 3.

Proof. We prove this by induction on the number of edges. The inclusion is obvious for orientations with no edge. Now let $O \in \mathcal{L}^{(k)}$, having at least one edge.

If O is the merge of two orientations O_1 and O_2 , then the induction hypothesis implies that O_1 and O_2 are in $\mathbb{L}^{(k)}$. The structure of $\mathbb{L}^{(k)}$ implies that every prime sub-orientation of O_2 (attached at the root of O_2) also belongs to $\mathbb{L}^{(k)}$. Then O can be obtained as an orientation of $\mathbb{L}^{(k)}$ by first adding a loop around O_1 (this is the second construction in Definition 6), then concatenating one by one the prime sub-orientations of O_2 (first construction in Definition 6).

Otherwise, O is obtained by a legal split in an orientation O' formed of the prime sub-orientations P_1, \dots, P_ℓ . By the induction hypothesis, O' , and its prime sub-orientations P_1, \dots, P_ℓ ,

belong to $\mathbb{L}^{(k)}$. Let us say that the split occurs in P_i (this means that the sub-orientations P_1, \dots, P_{i-1} are attached to the new created vertex v' , while P_{i+1}, \dots, P_ℓ remain attached to the original vertex v , the root vertex of O). Then the orientation O_1 obtained by deleting from O the sub-orientations P_{i+1}, \dots, P_ℓ can be obtained by a legal split in the prime orientation P_i , followed by the concatenation of P_1, \dots, P_{i-1} at the new created vertex. This is the third construction in Definition 6, hence O_1 belongs to $\mathbb{L}^{(k)}$. It remains to concatenate P_{i+1}, \dots, P_ℓ at the root (first construction in Definition 6), and we recover O as an element of $\mathbb{L}^{(k)}$. \square

4.1. AN ALGEBRAIC SYSTEM FOR $\mathbb{L}^{(k)}$

We now fix $k \geq 1$. For \mathbf{w} a word on $\{0, 1\}$, let $L_{\mathbf{w}}^{(k)}(t) \equiv L_{\mathbf{w}}$ denote the generating function of orientations of $\mathbb{L}^{(k)}$ whose root word ends with \mathbf{w} . Let $K_{\mathbf{w}}$ be the generating function of those that have root word exactly \mathbf{w} . Let $L'_{\mathbf{w}}$ and $K'_{\mathbf{w}}$ be the corresponding series for *prime* orientations of $\mathbb{L}^{(k)}$. We are especially interested in the series L_ε that counts all orientations of $\mathbb{L}^{(k)}$.

If \mathbf{w} is not balanced, $K_{\mathbf{w}} = K'_{\mathbf{w}} = 0$, while if $\mathbf{w} = \varepsilon$, $K_{\mathbf{w}} = 1$ and $K'_{\mathbf{w}} = 0$. For \mathbf{w} non-empty and balanced, of length at most $2k$, we have

$$K_{\mathbf{w}} = \sum_{\mathbf{w}=\mathbf{u}\mathbf{v}} K_{\mathbf{u}}K'_{\mathbf{v}}, \quad (15)$$

since an orientation of $\mathbb{L}^{(k)}$ is a sequence of orientations of $\mathbb{L}^{(k)}$. Now the description of prime orientations of $\mathbb{L}^{(k)}$ (Definition 6) gives

$$K'_{\mathbf{w}} = tK_{\mathbf{w}_c} + tL_\varepsilon L'_{\mathbf{w}_s}. \quad (16)$$

The first term corresponds to adding a loop, and the second to a legal i -split, where $i \leq k$ is the half-length of \mathbf{w} . The factor L_ε accounts for the orientation O'' attached at the end of the root edge.

Now let \mathbf{w} be a valid word of length at most $2k - 1$, and let us write equations for the series $L_{\mathbf{w}}$ and $L'_{\mathbf{w}}$. For $L_{\mathbf{w}}$, the sequential structure of orientations of $\mathbb{L}^{(k)}$ gives

$$L_{\mathbf{w}} = \mathbb{1}_{\mathbf{w}=\varepsilon} + L_\varepsilon L'_{\mathbf{w}} + \sum_{\mathbf{w}=\mathbf{u}\mathbf{v}, \mathbf{v} \neq \mathbf{w}} L_{\mathbf{u}}K'_{\mathbf{v}}. \quad (17)$$

The second (resp. third) term counts orientations in which the root word of the last prime component ends with \mathbf{w} (resp. is shorter than \mathbf{w}). Finally, for the series $L'_{\mathbf{w}}$ we obtain the following counterpart of (6):

$$L'_{\mathbf{w}} = 2tL_\varepsilon \mathbb{1}_{\mathbf{w}=\varepsilon} + tL_{\mathbf{w}_p} \mathbb{1}_{\mathbf{w} \neq \varepsilon} + tK_{\mathbf{w}_c} \mathbb{1}_{\mathbf{w} \neq \varepsilon} \text{ balanced} + tL_\varepsilon \left(L'_{\mathbf{w}} + \sum_{\substack{\mathbf{u}=\mathbf{v}\mathbf{w} \\ 0 < |\mathbf{u}| \leq 2k \\ \mathbf{u} \text{ balanced}}} (L'_{\mathbf{u}_s} - K'_{\mathbf{u}}) \right). \quad (18)$$

The first three terms count orientations in which the root edge is a loop, and the last one those obtained by a split.

Proposition 9. *Consider the collection of equations consisting of:*

- Equation (15), written for all balanced words \mathbf{w} of length between 2 and $2k - 2$,
- Equation (16), written for all balanced words \mathbf{w} of length between 2 and $2k$,
- Equation (17), written for all valid words \mathbf{w} of length at most $2k - 2$,
- Equation (18), written for all valid words \mathbf{w} of length at most $2k - 1$.

In this collection, replace all trivial K - and K' -series by their value: $K_{\mathbf{w}} = K'_{\mathbf{w}} = 0$ when \mathbf{w} is not balanced, $K_\varepsilon = 1$, $K'_\varepsilon = 0$. Let S_0 denote the resulting system. The number of series it involves is $2f(k) - 2\binom{2k}{k}$, where $f(k)$ is given by (7). Moreover, S_0 defines uniquely all these series. Its size can be (roughly) divided by two upon exploiting the 0/1 symmetry.

Proof. To prove that all series are well defined by the system, we first check that every series occurring in the right-hand side of some equation is the left-hand side of another equation. Then we note that:

- the equations for prime orientations, namely (16) and (18), have a factor t in their right-hand sides,
- for the other two equations, (15) and (17), every non-trivial term in the right-hand side has a series of prime orientations as a factor.

Now the number of equations: every series that was occurring in the system S_0 of Proposition 3 now has two copies (one with a prime, one without), except for the series $K_{\mathbf{w}}$, for \mathbf{w} balanced of length $2k$, and $L_{\mathbf{w}}$, for \mathbf{w} quasi-balanced of length $2k - 1$, which have only one copy. Since there are $\binom{2k}{k}$ balanced words of length $2k$, and $2\binom{2k-1}{k} = \binom{2k}{k}$ quasi-balanced words of length $2k - 1$, the result follows. \square

Remark 10. As in Remark 4, if \mathbf{w} is such that $0\mathbf{w}$ and $1\mathbf{w}$ are both valid of length less than $2k - 2$ (resp. $2k - 1$), we can replace (17) (resp. (18)) by the simpler forward equation:

$$L_{\mathbf{w}} = K_{\mathbf{w}} + L_{0\mathbf{w}} + L_{1\mathbf{w}} \quad (\text{resp. } L'_{\mathbf{w}} = K'_{\mathbf{w}} + L'_{0\mathbf{w}} + L'_{1\mathbf{w}}).$$

This does not increase the size of the system.

4.2. EXAMPLES

4.2.1. **When $k = 1$,** the system S_0 contains $2(f(1) - 2) = 6$ equations, or 4 if we exploit the 0/1 symmetry:

$$\begin{cases} K'_{10} = t + tL_{\varepsilon}L'_0, \\ L_{\varepsilon} = 1 + L_{\varepsilon}L'_{\varepsilon}, \\ L'_{\varepsilon} = 2tL_{\varepsilon} + tL_{\varepsilon}(L'_{\varepsilon} + 2L'_0 - 2K'_{10}), \\ L'_0 = tL_{\varepsilon} + tL_{\varepsilon}(L'_0 + L'_0 - K'_{10}). \end{cases}$$

Eliminating all series but L_{ε} gives a cubic equation for the generating function $L_{\varepsilon} \equiv L_{\varepsilon}^{(1)}$ of Eulerian orientations in $\mathbb{L}^{(1)}$:

$$t^2L_{\varepsilon}^3 + t(t - 4)L_{\varepsilon}^2 + (2t + 1)L_{\varepsilon} - 1 = 0. \quad (19)$$

4.2.2. **When $k = 2$,** the system S_0 contains $2(f(2) - 6) = 30$ equations, or 16 if we exploit the 0/1 symmetry:

$$\left\{ \begin{array}{l} K_{01} = K_{10} = K'_{01}, \\ K'_{10} = K'_{01} = t + tL_\varepsilon L'_1, \\ K'_{1100} = tK_{10} + tL_\varepsilon L'_{100}, \\ K'_{1010} = tK_{01} + tL_\varepsilon L'_{010}, \\ K'_{0110} = tL_\varepsilon L'_{110}, \\ L_\varepsilon = 1 + L_\varepsilon L'_\varepsilon, \\ L_0 = L_1 = L_\varepsilon L'_0, \\ L_{00} = L_{11} = L_\varepsilon L'_{00}, \\ L_{01} = L_{10} = L_\varepsilon L'_{01}, \\ L'_\varepsilon = 2tL_\varepsilon + tL_\varepsilon(L'_\varepsilon + 2(L'_0 - K'_{10} + L'_{100} - K'_{1100} + L'_{010} - K'_{1010} + L'_{110} - K'_{0110})), \\ L'_0 = L'_1 = tL_\varepsilon + tL_\varepsilon(L'_0 + L'_0 - K'_{10} + L'_{100} - K'_{1100} + L'_{010} - K'_{1010} + L'_{110} - K'_{0110}), \\ L'_{00} = tL_0 + tL_\varepsilon(L'_{00} + L'_{100} - K'_{1100}), \\ L'_{10} = L'_{01} = tL_1 + t + tL_\varepsilon(L'_{10} + L'_1 - K'_{01} + L'_{010} - K'_{1010} + L'_{110} - K'_{0110}), \\ L'_{100} = tL_{10} + tL_\varepsilon(L'_{100} + L'_{100} - K'_{1100}), \\ L'_{010} = tL_{01} + tL_\varepsilon(L'_{010} + L'_{010} - K'_{1010}), \\ L'_{110} = tL_{11} + tL_\varepsilon(L'_{110} + L'_{110} - K'_{0110}). \end{array} \right. \quad (20)$$

Eliminating all series but L_ε gives an equation of degree 6 for the generating function $L_\varepsilon \equiv L_\varepsilon^{(2)}$ of Eulerian orientations in $\mathbb{L}^{(2)}$:

$$2t^5 L_\varepsilon^6 - t^4(t+8)L_\varepsilon^5 - t^3(3t^2 - 16)L_\varepsilon^4 + t^2(2t+3)(2t-5)L_\varepsilon^3 - t(2t^2 - 7t - 7)L_\varepsilon^2 - (5t+1)L_\varepsilon + 1 = 0. \quad (21)$$

4.3. ASYMPTOTIC ANALYSIS FOR SUBSETS OF EULERIAN ORIENTATIONS (PRIME DECOMPOSITION)

We now prove for the polynomial system of Proposition 9 an analogue of Proposition 5.

Proposition 11. *For $k \geq 1$, let $\bar{\rho}_k$ denote the radius of convergence of the series $L_\varepsilon^{(k)}$ that counts orientations of $\mathbb{L}^{(k)}$. Then $\bar{\rho}_k$ is the only singularity of $L_\varepsilon^{(k)}$ of minimal modulus, and it is of the square root type. Consequently, there exists a constant c such that the number $\bar{\ell}_n^{(k)}$ of orientations of size n in $\mathbb{L}^{(k)}$ satisfies, as n tends to infinity:*

$$\bar{\ell}_n^{(k)} \sim c \bar{\lambda}_k^n n^{-3/2},$$

with $\bar{\lambda}_k = 1/\bar{\rho}_k$.

Proof. Again, we apply the theory of positive irreducible polynomial systems [30, Sec. VII.6].

The system of Proposition 9 is not positive. To correct this, we replace the series $L_{\mathbf{w}}$ (for \mathbf{w} balanced) and $L'_{\mathbf{w}}$ (for \mathbf{w} balanced or quasi-balanced) by their “positive” versions:

$$L_{\mathbf{w}}^+ := L_{\mathbf{w}} - K_{\mathbf{w}}, \quad L'_{\mathbf{w}}^+ := L'_{\mathbf{w}} - K'_{\mathbf{w}} \quad (\mathbf{w} \text{ balanced}),$$

$$L'_{\mathbf{w}}^+ := L'_{\mathbf{w}} - K'_{\mathbf{w}} \quad (\mathbf{w} \text{ quasi-balanced}).$$

In particular, L_ε is replaced by $L_\varepsilon^+ := L_\varepsilon - 1$ and L'^+_ε coincides with L'_ε . We alter the original system S_0 as follows:

- (i) For \mathbf{w} balanced, we replace the equation (17) defining $L_{\mathbf{w}}$ by the difference between (17) and (15):

$$L_{\mathbf{w}}^+ = L_\varepsilon^+ L'_{\mathbf{w}} + L'^+_{\mathbf{w}} + \sum_{\mathbf{w}=\mathbf{uv}, \mathbf{v} \neq \mathbf{w}} L_{\mathbf{u}} K'_{\mathbf{v}}. \quad (22)$$

- (ii) For \mathbf{w} balanced or quasi-balanced, we replace the equation (18) defining $L'_\mathbf{w}$ by the difference between (18) and (16) (written for \mathbf{w} if \mathbf{w} is balanced, for $\overleftarrow{\mathbf{w}}$ otherwise):

$$L'^+_\mathbf{w} = 2tL_\varepsilon \mathbb{1}_{\mathbf{w}=\varepsilon} + t(L_{\mathbf{w}_p} - K_{\mathbf{w}_p}) \mathbb{1}_{\mathbf{w} \neq \varepsilon} + tL_\varepsilon \left(L'^+_\mathbf{w} + \sum_{\substack{\mathbf{u}=\mathbf{v}\mathbf{w}, \mathbf{u} \neq \mathbf{w} \\ |\mathbf{u}| \leq 2k-1 \\ \mathbf{u} \text{ quasi-balanced}}} L'^+_\mathbf{u} \right). \quad (23)$$

- (iii) In the new system thus obtained, we replace every series $K_\mathbf{w}$ such that \mathbf{w} is not balanced by 0, every series $L_\mathbf{w}$ (resp. $L'_\mathbf{w}$) such that \mathbf{w} is balanced by $K_\mathbf{w} + L^+_\mathbf{w}$ (resp. $K'_\mathbf{w} + L'^+_\mathbf{w}$), and every series $L'_\mathbf{w}$ such that \mathbf{w} is quasi-balanced by $K'_{\overleftarrow{\mathbf{w}}} + L'^+_\mathbf{w}$. In particular, the series $L_{\mathbf{w}_p} - K_{\mathbf{w}_p}$ occurring in (23) becomes $L^+_{\mathbf{w}_p}$ when \mathbf{w}_p is balanced, $L_{\mathbf{w}_p}$ otherwise. The series $L_\mathbf{w}$ (resp. $L'_\mathbf{w}$) that remain in the system are such that \mathbf{w} has balance at least 1 (resp. 2).

We thus obtain a positive system, denoted S_1 , defining the following series:

- $K_\mathbf{w}$, for \mathbf{w} balanced of length between 2 and $2k - 2$,
- $K'_\mathbf{w}$, for \mathbf{w} balanced of length between 2 and $2k$,
- $L_\mathbf{w}$ for \mathbf{w} valid of balance at least 1 and length at most $2k - 2$,
- $L'_\mathbf{w}$ for \mathbf{w} valid of balance at least 2 and length at most $2k - 1$,
- $L^+_\mathbf{w}$, for \mathbf{w} balanced of length at most $2k - 2$,
- $L'^+_\mathbf{w}$, for \mathbf{w} balanced or quasi-balanced of length at most $2k - 1$.

For instance, when $k = 1$ we obtain the following system:

$$\begin{cases} K'_{10} = t + t(1 + L^+_\varepsilon)(K'_{10} + L'^+_{10}), \\ L^+_\varepsilon = L^+_\varepsilon L'^+_\varepsilon + L'^+_\varepsilon, \\ L'^+_\varepsilon = 2t(1 + L^+_\varepsilon) + t(1 + L^+_\varepsilon)(L'^+_\varepsilon + 2L'^+_{10}), \\ L'^+_{10} = tL^+_\varepsilon + t(1 + L^+_\varepsilon)L'^+_{10}. \end{cases}$$

Recall that the series we are interested in is L^+_ε . But then we can drop the first equation of the above system. *This size reduction occurs for any value of k* , and the positive system S_2 that we will study is finally obtained by performing one last change:

- (iv) Delete the equations defining the series $K'_\mathbf{w}$, for \mathbf{w} of length $2k$.

Observe that all the series involved in S_2 are well-defined by this system. This comes from the fact that all the series $K'_\mathbf{u}$, for \mathbf{u} of length $2k$, that occurred in S_0 came from the term $L'_{\mathbf{u}_s} - K'_\mathbf{u}$ of (18), which now reads $L'^+_{\mathbf{u}_s}$.

Here is for instance the system obtained for $k = 2$, which has three equations less than (20):

$$\left\{ \begin{array}{l} K_{01} = K_{10} = K'_{01}, \\ K'_{10} = K'_{01} = t + t(1 + L_\varepsilon^+)(K'_{01} + L'_{10}), \\ L_\varepsilon^+ = L_\varepsilon^+ L'_\varepsilon + L'_{10}, \\ L_0 = (1 + L_\varepsilon^+)(K'_{10} + L'_{01}), \\ L_{00} = L_{11} = (1 + L_\varepsilon^+)L'_{00}, \\ L'_{01} = L_\varepsilon^+(K'_{01} + L'_{01}) + L'_{01}, \\ L'_\varepsilon = 2t(1 + L_\varepsilon^+) + t(1 + L_\varepsilon^+)(L'_\varepsilon + 2(L'_{01} + L'_{100} + L'_{010} + L'_{110})), \\ L'_{01} = L'_{10} = tL_\varepsilon^+ + t(1 + L_\varepsilon^+)(L'_{01} + L'_{100} + L'_{010} + L'_{110}), \\ L'_{00} = tL_0 + t(1 + L_\varepsilon^+)(L'_{00} + L'_{100}), \\ L'_{10} = L'_{01} = tL_1 + t(1 + L_\varepsilon^+)(L'_{10} + L'_{010} + L'_{110}), \\ L'_{100} = tL'_{10} + t(1 + L_\varepsilon^+)L'_{100}, \\ L'_{010} = tL'_{01} + t(1 + L_\varepsilon^+)L'_{010}, \\ L'_{110} = tL'_{11} + t(1 + L_\varepsilon^+)L'_{110}. \end{array} \right.$$

Let us now discuss properness [30, p. 489]. The system S_2 that we have just obtained is not proper. However, the right-hand sides of the equations that define series with a prime (K' , L' and L'^+) are multiples of t (see (16), (18) and (23)). In the remaining equations, that is (15), (17) (for \mathbf{w} not balanced) and (22) (for \mathbf{w} balanced), each term on the right-hand side involves a series with a prime: hence after one iteration of S_2 , one obtains a new system S_3 which is positive and proper.

Aperiodicity holds as in the previous section, and we are left with irreducibility. Note that proving irreducibility for S_2 or its iterated version S_3 is equivalent, so we focus on S_2 . As in the previous section, we denote by $\tilde{L}_{\mathbf{w}}$ the series $L_{\mathbf{w}}^+$ or $L_{\mathbf{w}}$, depending on whether \mathbf{w} is balanced or not. Similarly, $\tilde{L}'_{\mathbf{w}}$ denotes $L'_{\mathbf{w}}^+$ if \mathbf{w} is balanced or quasi-balanced, and $L'_{\mathbf{w}}$ otherwise. Let us prove that all series depend on \tilde{L}_ε . We first observe that this holds for every K' - or \tilde{L} - or \tilde{L}' -series (see (16), (17), (18), (22), (23)). We are left with the series $K_{\mathbf{w}}$: but it depends on $K'_{\mathbf{w}}$ (see (15)), and hence on \tilde{L}_ε .

Conversely, let us prove that \tilde{L}_ε depends on every other series in the system. By (22), it depends on L'_ε . Then by (23) applied to $\mathbf{w} = \varepsilon$, it depends on every series $L'_{\mathbf{u}}$, where \mathbf{u} is quasi-balanced. Going back and forth between the equations defining the L -series and the L' -series (see (17), (18), (22), (23)), and using an induction on the balance, we then see that \tilde{L}_ε depends on all series $L_{\mathbf{u}}$ (for $|\mathbf{u}| \leq 2k - 2$) and all series $\tilde{L}'_{\mathbf{u}}$ (for $|\mathbf{u}| \leq 2k - 1$). Then the first term of (22), written as $L_\varepsilon^+(K'_{\mathbf{w}} + L'_{\mathbf{w}})$, shows that \tilde{L}_ε depends on all series $K'_{\mathbf{v}}$ with $|\mathbf{v}| \leq 2k - 2$. It remains to prove that \tilde{L}_ε depends on the K -series. Let $\mathbf{u} = a\mathbf{u}_s$ be balanced of length at most $2k - 2$, and define $\mathbf{w} = \mathbf{u}_s \bar{a}$. This word has balance 2. The second term of (18) involves $L_{\mathbf{w}_p} = L_{\mathbf{u}_s} = K_{\mathbf{u}} + L_{\mathbf{u}_s}^+$. Hence $L'_{\mathbf{w}}$ depends on $K_{\mathbf{u}}$, and by transitivity, \tilde{L}_ε depends on $K_{\mathbf{u}}$. This proves the irreducibility of the system and concludes the proof of the proposition. \square

4.4. BACK TO EXAMPLES

We first return to the cases $k = 1$ and $k = 2$ studied in Section 4.2. When $k = 1$, we obtained the cubic equation (19) for $L_\varepsilon^{(1)}$. The discriminant has three positive roots, which are 1, and (approximately) 0.094 and 15.9. The second one is the radius of convergence, and we obtain the lower bound $\bar{\lambda}_1 \simeq 10.603$ on the growth rate of Eulerian orientations. This improves significantly on the growth rate $\lambda_1 = 9.68 \dots$ obtained from the set $\mathcal{L}^{(1)}$.

For $k = 2$, we obtained the equation (21) satisfied by $L_\varepsilon^{(2)}$. The discriminant has two roots in $(0, 1)$, which are approximately 0.0911 and 0.414. The first one is the radius of convergence, and we obtain the lower bound $\bar{\lambda}_2 \simeq 10.9759$ on the growth rate of Eulerian orientations.

When $k = 3$, we find that $L_\varepsilon^{(3)}$ satisfies an equation of degree 20 (see the MAPLE sessions available on our web pages). The dominant coefficient only vanishes at $t = 8$, and the discriminant has only one relevant root, around 0.089. This gives the lower bound $\bar{\lambda}_3 \simeq 11.2289$ on the growth rate of Eulerian orientations.

When $k = 4$, we did not compute the equation satisfied by $L_\varepsilon^{(4)}$, but we estimated $\bar{\lambda}_4$ from the first 30 coefficients of $L_\varepsilon^{(4)}$ using quadratic approximants [19]. We predict $\bar{\lambda}_4 \simeq 11.41$. This value has then been confirmed by Bruno Salvy using the Maple package NewtonGF [48], with which he obtained 10 digits of $\bar{\lambda}_4$. This package also allows us to compute more coefficients in $L_\varepsilon^{(4)}$. Moreover, Jean-Charles Faugère [26] has finally been able to determine the equation for $L_\varepsilon^{(4)}$, which has degree 258 in $L_\varepsilon^{(4)}$.

Similarly, we predict

$$\bar{\lambda}_5 \simeq 11.56, \quad \bar{\lambda}_6 \simeq 11.68.$$

5. SUPERSETS OF EULERIAN ORIENTATIONS, VIA THE STANDARD DECOMPOSITION

We now want to define, and count, supersets of Eulerian orientations. Their generating functions will be described by functional equations involving divided differences (as in (3)). The proof of their algebraicity is non-trivial, relying on a deep result from Artin's approximation theory (Theorem 16).

Recall that Eulerian orientations can be obtained recursively from the atomic map by either:

- the merge of two orientations $O_1, O_2 \in \mathcal{O}$ (with the root loop oriented in either way),
- or a legal split on an orientation $O' \in \mathcal{O}$.

We now define the sets $\mathcal{U}^{(k)}$. The idea is that we allow illegal i -splits, provided i is larger than k .

Definition 12. *Let $k \geq 1$. Let $\mathcal{U}^{(k)}$ be the set of planar orientations obtained recursively from the atomic map by either:*

- *the merge of two orientations $O_1, O_2 \in \mathcal{U}^{(k)}$ (with the root loop oriented in either way),*
- *or a legal i -split on a map $O' \in \mathcal{U}^{(k)}$ with $i \leq k$ (small split),*
- *or an arbitrary split on a map $O' \in \mathcal{U}^{(k)}$ with $i > k$ (large split). If the split is legal, the root edge is oriented in the only way that makes the new orientation Eulerian. Otherwise, it is oriented away from the root vertex.*

Observe that all Eulerian orientations belong to $\mathcal{U}^{(k)}$. Moreover, the sets $\mathcal{U}^{(k)}$ form a decreasing sequence, as fewer illegal splits are performed as k grows. Finally, for $k \geq n$ (and even for $k \geq n - 2$), all orientations of size n in $\mathcal{U}^{(k)}$ are Eulerian. Hence the limit of the sets $\mathcal{U}^{(k)}$ is the set \mathcal{O} of all Eulerian orientations.

Another important observation is that, if the root vertex of an orientation of $\mathcal{U}^{(k)}$ has degree at most $2k$, then the root word of this orientation is balanced.

5.1. FUNCTIONAL EQUATIONS FOR $\mathcal{U}^{(k)}$

We now fix an integer k . For a word \mathbf{w} on $\{0, 1\}$, let $U_{\mathbf{w}}^{(k)}(t; x) \equiv U_{\mathbf{w}}(x)$ denote the generating function of orientations of $\mathcal{U}^{(k)}$ whose root word ends with \mathbf{w} , counted by the edge number (variable t) and the half-degree of the root vertex (variable x). Let $T_{\mathbf{w}}^{(k)}(t) \equiv T_{\mathbf{w}}$ denote the generating function of orientations of $\mathcal{U}^{(k)}$ having root word exactly \mathbf{w} . We do not record in this series the root degree (which is the length of \mathbf{w}). To lighten notation, we often denote simply by $U_{\mathbf{w}}$ the edge generating function $U_{\mathbf{w}}(1) \equiv U_{\mathbf{w}}(t, 1)$, and by $U_{\mathbf{w}}^x$ the refined generating function $U_{\mathbf{w}}(x) \equiv U_{\mathbf{w}}(t; x)$.

Note that $T_{\mathbf{w}} = 0$ if \mathbf{w} is not balanced and that $T_\varepsilon = 1$. Now, for \mathbf{w} balanced of length between 2 and $2k$, we have

$$T_{\mathbf{w}} = t \sum_{\mathbf{u}\mathbf{a}\mathbf{v}=\mathbf{w}} T_{\mathbf{u}}T_{\mathbf{v}} + tU_{\mathbf{w}_s}. \quad (24)$$

The first term counts orientations obtained by a merge. The second one counts those obtained by a split, which is necessarily small since we have assumed $|\mathbf{w}| \leq 2k$. Note the analogy with (5).

For \mathbf{w} valid of length at most $2k - 1$, let us now prove the following identity:

$$U_{\mathbf{w}}^x = \mathbb{1}_{\mathbf{w}=\varepsilon} + 2txU_{\varepsilon}^xU_{\mathbf{w}}^x + tx \sum_{\mathbf{w}=\mathbf{u}\mathbf{a}\mathbf{v}} U_{\mathbf{u}}^x x^{|\mathbf{v}|/2} T_{\mathbf{v}} + tx^{|\mathbf{w}|/2} \sum_{\mathbf{w}=\mathbf{a}\mathbf{u}\mathbf{a}\mathbf{v}} T_{\mathbf{u}} T_{\mathbf{v}} \\ + t \sum_{\substack{\mathbf{u}=\mathbf{v}\mathbf{w} \\ 2 \leq |\mathbf{u}| \leq 2k \\ \mathbf{u} \text{ balanced}}} x^{|\mathbf{u}|/2} U_{\mathbf{u}} + \frac{tx}{x-1} (U_{\mathbf{w}}^x - x^k U_{\mathbf{w}}) - \frac{tx}{x-1} \sum_{\substack{\mathbf{u}=\mathbf{v}\mathbf{w} \\ |\mathbf{u}| \leq 2k-2}} T_{\mathbf{u}} (x^{|\mathbf{u}|/2} - x^k). \quad (25)$$

The first line is similar to the first line of (6): it counts the atomic map and orientations obtained from a merge. The only difference is that we now record the root degree. On the second line, the first sum counts orientations obtained by a small split (with root word \mathbf{u}). Let us explain the remaining terms, which count orientations obtained by a large split, legal or not, of an orientation O' whose root word ends (necessarily) with \mathbf{w} . Given an orientation O' with root vertex degree $2d$, with $d > k$, the generating function of orientations obtained from O' by a large split is

$$t^{1+e(O')} (x^{k+1} + x^{k+2} + \dots + x^d) = t^{1+e(O')} \frac{x^{d+1} - x^{k+1}}{x-1}.$$

Let us underline that we cannot apply a large split to an orientation O' whose root word \mathbf{u} satisfies $|\mathbf{u}| \leq 2k$. Hence the generating function of orientations obtained by a large split is

$$\frac{tx}{x-1} \left(\left(U_{\mathbf{w}}^x - \sum_{\mathbf{u}=\mathbf{v}\mathbf{w}, |\mathbf{u}| \leq 2k} x^{|\mathbf{u}|/2} T_{\mathbf{u}} \right) - x^k \left(U_{\mathbf{w}} - \sum_{\mathbf{u}=\mathbf{v}\mathbf{w}, |\mathbf{u}| \leq 2k} T_{\mathbf{u}} \right) \right),$$

which gives the last two terms of (25) (the terms $T_{\mathbf{u}}$ with $|\mathbf{u}| = 2k$ do not contribute).

Remark 13. In the proof of (25), we have tried to follow the same steps as in the proof of (6). However, comparing (24) and (25) suggests to replace (25) by a lighter equation:

$$U_{\mathbf{w}}^x = x^{|\mathbf{w}|/2} T_{\mathbf{w}} + 2txU_{\varepsilon}^xU_{\mathbf{w}}^x + tx \sum_{\mathbf{w}=\mathbf{u}\mathbf{a}\mathbf{v}} U_{\mathbf{u}}^x x^{|\mathbf{v}|/2} T_{\mathbf{v}} \\ + t \sum_{\substack{\mathbf{u}=\mathbf{v}\mathbf{w} \\ |\mathbf{u}| \leq 2k-1 \\ \mathbf{u} \text{ quasi-balanced}}} x^{(1+|\mathbf{u}|)/2} U_{\mathbf{u}} + \frac{tx}{x-1} (U_{\mathbf{w}}^x - x^k U_{\mathbf{w}}) - \frac{tx}{x-1} \sum_{\substack{\mathbf{u}=\mathbf{v}\mathbf{w} \\ |\mathbf{u}| \leq 2k-2}} T_{\mathbf{u}} (x^{|\mathbf{u}|/2} - x^k). \quad (26)$$

Proposition 14. *Consider the collection of equations consisting of:*

- Equation (24), written for all balanced words \mathbf{w} of length between 2 and $2k - 2$,
- Equation (26), written for all valid words \mathbf{w} of length at most $2k - 1$.

In this collection, replace all trivial T -series by their value: $T_{\mathbf{w}} = 0$ when \mathbf{w} is not balanced, $T_{\varepsilon} = 1$. Let R_0 denote the resulting system. The number of series it involves is $f(k) - \binom{2k}{k}$, where $f(k)$ is given by (7). Moreover, R_0 defines uniquely these series. Its size can be (roughly) divided by two upon exploiting the 0/1 symmetry.

The proof is similar to the proof of Proposition 3.

Remark 15. As in Remark 4, if \mathbf{w} is such that $0\mathbf{w}$ and $1\mathbf{w}$ are both valid of length less than $2k$, we can replace (26) by the simpler forward equation:

$$U_{\mathbf{w}}^x = x^{|\mathbf{w}|/2} T_{\mathbf{w}} + U_{0\mathbf{w}}^x + U_{1\mathbf{w}}^x.$$

This does not increase the size of the system.

5.2. EXAMPLES

5.2.1. **When $k = 1$,** the system of Proposition 14 contains $f(1) - 2 = 3$ equations. Upon exploiting the 0/1 symmetry, it reads:

$$\begin{cases} U_\varepsilon^x = 1 + 2tx(U_\varepsilon^x)^2 + 2txU_0 + \frac{tx}{x-1}(U_\varepsilon^x - xU_\varepsilon) + tx, \\ U_0^x = 2txU_\varepsilon^xU_0^x + txU_\varepsilon^x + txU_0 + \frac{tx}{x-1}(U_0^x - xU_0). \end{cases} \quad (27)$$

The first equation can be replaced by the forward equation $U_\varepsilon^x = 1 + 2U_0^x$. We explain in Section 5.4 how to solve this system.

5.2.2. **When $k = 2$,** the system of Proposition 14 contains $f(2) - 6 = 15$ equations. Upon exploiting the 0/1 symmetry and (some) forward equations, it reads:

$$\begin{cases} T_{10} &= t + tU_0, \\ U_\varepsilon^x &= 1 + 2U_0^x, \\ U_0^x = U_1^x &= 2txU_\varepsilon^xU_0^x + txU_\varepsilon^x + txU_0 + tx^2(U_{100} + U_{010} + U_{110}) + \frac{tx}{x-1}(U_0^x - x^2U_0) + tx^2T_{10}, \\ U_{10}^x = U_{01}^x &= xT_{10} + U_{110}^x + U_{010}^x, \\ &= xT_{10} + 2txU_\varepsilon^xU_{10}^x + txU_{11}^x + tx^2(U_{010} + U_{110}) + \frac{tx}{x-1}(U_{10}^x - x^2U_{10}) + tx^2T_{10}, \\ U_{00}^x = U_{11}^x &= 2txU_\varepsilon^xU_{00}^x + txU_0^x + tx^2U_{100} + \frac{tx}{x-1}(U_{00}^x - x^2U_{00}), \\ U_{100}^x &= 2txU_\varepsilon^xU_{100}^x + txU_{10}^x + tx^2U_{100} + \frac{tx}{x-1}(U_{100}^x - x^2U_{100}), \\ U_{010}^x &= 2txU_\varepsilon^xU_{010}^x + tx(U_{01}^x + U_\varepsilon^x xT_{10}) + tx^2U_{010} + \frac{tx}{x-1}(U_{010}^x - x^2U_{010}), \\ U_{110}^x &= 2txU_\varepsilon^xU_{110}^x + tx(U_{11}^x + U_\varepsilon^x xT_{10}) + tx^2U_{110} + \frac{tx}{x-1}(U_{110}^x - x^2U_{110}). \end{cases} \quad (28)$$

We explain in Section 5.4 below how to solve this system.

5.3. ALGEBRAICITY

Since the early work of Brown in the sixties on the *quadratic method* [20], a lot has been known about equations involving divided differences of the form $(F(t; x) - F(t; 1))/(x - 1)$. However, most of the literature deals with a single equation, not with a system [14, 34]. In order to prove that the series $U_\varepsilon^{(k)}(t; x)$ that counts orientations of $\mathcal{U}^{(k)}$ is algebraic, we use a deep theorem from Artin's approximation theory, due to Popescu [43, 50]. The form we will need is given below. We recall that $\mathbb{C}[[z_1, \dots, z_\ell]]$ is the ring of formal power series in the variables z_1, \dots, z_ℓ , with complex coefficients, and that a series Z in this ring is *algebraic* if it satisfies a non-trivial polynomial equation $\text{Pol}(z_1, \dots, z_\ell, Z) = 0$.

Theorem 16 ([50], Thm. 1.4). *Consider a polynomial system of n equations in $\ell + n$ variables over \mathbb{C} , written as $P_i(z_1, \dots, z_\ell, y_1, \dots, y_n) = 0$, for $1 \leq i \leq n$. Let (d_1, \dots, d_n) be a sequence of integers in $\{0, 1, \dots, \ell\}$. Assume that there exists an n -tuple $\mathcal{Y} = (Y_1, \dots, Y_n)$ of series in $\mathbb{C}[[z_1, \dots, z_\ell]]$ that satisfies the following conditions:*

- the n -tuple \mathcal{Y} solves this system, that is,

$$P_i(z_1, \dots, z_\ell, Y_1, \dots, Y_n) = 0 \quad \text{for } 1 \leq i \leq n,$$

- for $1 \leq i \leq n$, the series Y_i does not depend on the variables z_j such that $j > d_i$ (if $d_i = \ell$, then there is no condition on the series Y_i).

Then there exists an n -tuple (Z_1, \dots, Z_n) of algebraic series in $\mathbb{C}[[z_1, \dots, z_\ell]]$ that solves the system and satisfies the same dependence conditions as \mathcal{Y} .

In particular, if the system has a unique solution satisfying the dependence conditions, then this solution is algebraic.

An application. To our knowledge, this theorem has not been applied yet in a combinatorial context. So, before we use it to prove the algebraicity of $U_\varepsilon^{(k)}(t; x)$, let us examine its application to a simple equation, namely (3).

First, observe that the algebraicity of $M(t; x)$ is not obvious. Clearly, if we could prove that $M(t; 1)$ is algebraic, we would be done with $M(t; x)$ as well, but why should $M(t; 1)$ be algebraic? We can apply the above theorem as follows. Let us denote $t = z_1$ and $x = 1 + z_2$ (we shall explain later why we need to translate the variable x). We consider the system in z_1, z_2, y_1 and y_2 consisting of the following (single) equation:

$$z_2 y_2 = z_2 + z_1 z_2 (1 + z_2) y_2^2 + z_1 (1 + z_2) (y_2 - y_1).$$

Take $d_1 = 1$ and $d_2 = 2$. Then (3) shows that the pair $(Y_1, Y_2) := (M(t; 1), M(t; x))$ solves the above equation. Moreover $Y_1 = M(t; 1)$ is independent of $z_2 = x - 1$, while $Y_2 = M(t; x)$ depends on both variables z_1 and z_2 , in accordance with $d_1 = 1$ and $d_2 = 2$.

Let us now prove that there cannot be another solution of this system in the ring $\mathbb{C}[[z_1, z_2]]$ such that Y_1 is independent of z_2 . First, setting $z_2 = 0$ in the equation shows that Y_1 must be the specialization of Y_2 at $z_2 = 0$. This, combined with the factor z_1 occurring in every non-initial term in the right-hand side, implies that the coefficient of z_1^n in Y_2 can be computed by induction of n , starting from the constant coefficient 1. Hence the uniqueness of (Y_1, Y_2) . The algebraicity of $M(t; x)$ now follows from the above theorem.

Note that, if we had used $z_2 = x$ instead of $z_2 = x - 1$, we could not apply the last part of Theorem 16. The equation would read

$$(z_2 - 1)y_2 = (z_2 - 1) + z_1(z_2 - 1)z_2 y_2^2 + z_1 z_2 (y_2 - y_1),$$

but this equation has many solutions in the ring $\mathbb{C}[[z_1, z_2]]$ of formal power series in $z_1 = t$ and $z_2 = x$. For instance, one can take $Y_1 = 0$ and

$$Y_2 = \frac{1 - x + tx - \sqrt{(1 - x + tx)^2 - 4tx(1 - x)^2}}{2tx(1 - x)}.$$

Theorem 16 tells us that at least one of these solutions is algebraic, but we need uniqueness to conclude that *our* solution is algebraic. The key point is that a series in $\mathbb{C}[[z_1, z_2]]$ can always be specialized at $z_2 = 0$, but not at $z_2 = 1$.

We now apply Theorem 16 to the larger example of orientations of $\mathcal{U}^{(k)}$.

Proposition 17. *For any $k \geq 1$, the generating function $U_\varepsilon^{(k)}(t; x)$ that counts orientations of $\mathcal{U}^{(k)}$ is algebraic.*

Proof. Again, we take as variables $z_1 = t$ and $z_2 = x - 1$. For short, we denote z_2 by z . We consider the polynomial system consisting of the following equations, which mimic (24) and (26). For \mathbf{w} balanced of length between 2 and $2k - 2$,

$$A_{\mathbf{w}} = t \sum_{a\mathbf{u}\bar{a}\mathbf{v}=\mathbf{w}} A_{\mathbf{u}} A_{\mathbf{v}} + t C_{\mathbf{w}_s},$$

and for \mathbf{w} valid of length at most $2k - 1$,

$$\begin{aligned} z B_{\mathbf{w}} &= z(1+z)^{|\mathbf{w}|/2} A_{\mathbf{w}} + 2tz(1+z) B_\varepsilon B_{\mathbf{w}} + tz(1+z) \sum_{\mathbf{w}=\mathbf{u}\mathbf{a}\mathbf{v}} B_{\mathbf{u}} (1+z)^{|\mathbf{v}|/2} A_{\mathbf{v}} \\ &\quad + tz \sum_{\substack{\mathbf{u}=\mathbf{v}\mathbf{w} \\ |\mathbf{u}| \leq 2k-1 \\ \mathbf{u} \text{ quasi-balanced}}} (1+z)^{(1+|\mathbf{u}|)/2} C_{\mathbf{u}} \\ &\quad + t(1+z) (B_{\mathbf{w}} - (1+z)^k C_{\mathbf{w}}) - t(1+z) \sum_{\substack{\mathbf{u}=\mathbf{v}\mathbf{w} \\ |\mathbf{u}| \leq 2k-2}} A_{\mathbf{u}} ((1+z)^{|\mathbf{u}|/2} - (1+z)^k), \end{aligned} \quad (29)$$

where $A_\varepsilon = 1$. The variables $A_{\mathbf{w}}$, $B_{\mathbf{w}}$ and $C_{\mathbf{w}}$ play the role of the y_i in Theorem 16. By construction, the series

$$A_{\mathbf{w}} := T_{\mathbf{w}}(t), \quad B_{\mathbf{w}} := U_{\mathbf{w}}(t; 1+z), \quad C_{\mathbf{w}} := U_{\mathbf{w}}(t; 1)$$

solve the system. Moreover, $A_{\mathbf{w}}$ and $C_{\mathbf{w}}$ do not depend on $z_2 = z$.

By Theorem 16, it suffices to prove that the system in A, B, C has a unique solution in $\mathbb{C}[[t, z]]$ satisfying these dependence relations to conclude that all our series T and U counting orientations are algebraic.

So assume that $A_{\mathbf{w}}, B_{\mathbf{w}}$ and $C_{\mathbf{w}}$ solve the system and satisfy the required dependences. Then by setting $z = 0$ in (29), we see that $C_{\mathbf{w}}$ must be the specialization of $B_{\mathbf{w}}$ at $z = 0$, for all \mathbf{w} valid of length at most $2k - 1$. Then the form of the system implies that the coefficient of t^n in all series can be computed by induction on n , the initial values being $B_{\varepsilon} = C_{\varepsilon} = 1 + O(t)$ and $A_{\mathbf{w}} = B_{\mathbf{w}} = C_{\mathbf{w}} = O(t)$ for \mathbf{w} non-empty (recall that we have set $A_{\varepsilon} = 1$). This proves the uniqueness of the solution (with the required dependences) and concludes the proof. \square

5.4. BACK TO EXAMPLES

5.4.1. The case $k = 1$. Let us go back to System (27). In the first equation, replace U_0 by $(U_{\varepsilon} - 1)/2$ to obtain a single equation involving only $U_{\varepsilon}^x = U_{\varepsilon}(x)$ and $U_{\varepsilon} = U_{\varepsilon}(1)$. For simplicity, we now drop the index ε . This equation reads:

$$\text{Pol}(U(x), U(1), t, x) = 0,$$

with

$$\text{Pol}(x_0, x_1, t, x) = (x - 1)(-x_0 + 1 + 2txx_0^2 + tx(x_1 - 1)) + tx(x_0 - xx_1) + tx(x - 1).$$

We apply Brown's *quadratic method*. Its principle is the following: if there exists a formal power series $X \equiv X(t)$ such that

$$\text{Pol}_{x_0}(U(X), U(1), t, X) = 0, \quad (30)$$

then this series X must be a double root of the discriminant $\Delta(U(1), t, x)$ of $\text{Pol}(x_0, U(1), t, x)$ with respect to x_0 (the notation Pol_{x_0} stands for the derivative of Pol with respect to its first variable). The proof is elementary (see [34, Sec. 2.9] or [14]). Equation (30) reads

$$X = 1 + tX + 4tX(X - 1)U(X),$$

and has a unique power series solution $X(t)$, whose coefficients can be computed by induction from those of $U(x)$ (we do not need to determine X , just to know that it exists). Thus X is a double root of $\Delta(U(1), t, x)$, and hence the discriminant in x of Δ must vanish. This gives the following cubic equation for $U(1)$ (see our MAPLE sessions):

$$64t^3U(1)^3 + 2t(24t^2 - 36t + 1)U(1)^2 + (-15t^3 + 9t^2 + 19t - 1)U(1) + t^3 + 27t^2 - 19t + 1 = 0. \quad (31)$$

The series $U(1)$ has a unique positive singularity τ_1 , around 0.0765, which is a root of $216t^3 - 81t^2 + 18t - 1$. This gives the upper bound $\mu_1 = 1/\tau_1 = 13.0659 \dots$ on the growth rate of Eulerian orientations. Expanding the series near τ_1 (using for instance the MAPLE function `algeqtoser` [53]) shows that it has a singularity in $(1 - \mu_1 t)^{3/2}$, as the generating function of many families of planar maps.

5.4.2. The case $k = 2$. We now return to the system (28). Observe that we can reduce it to a system of three equations defining the series $U_{\varepsilon}^x, U_{10}^x$ and U_{100}^x :

$$\begin{cases} U_{\varepsilon}^x &= 1 + 4txU_{\varepsilon}^xU_0^x + 2txU_{\varepsilon}^x + 2txU_0 + 2tx^2(U_{100} + U_{10}) + \frac{2tx}{x-1}(U_0^x - x^2U_0), \\ U_{10}^x &= xt(1 + U_0) + 2txU_{\varepsilon}^xU_{10}^x + txU_0^x + tx^2U_{10} + \frac{tx}{x-1}(U_{10}^x - x^2U_{10}), \\ U_{100}^x &= 2txU_{\varepsilon}^xU_{100}^x + txU_{10}^x + tx^2U_{100} + \frac{tx}{x-1}(U_{100}^x - x^2U_{100}), \end{cases} \quad (32)$$

in which we still need to plug

$$U_0 = \frac{U_{\varepsilon} - 1}{2} \quad \text{and} \quad U_0^x = \frac{U_{\varepsilon}^x - 1}{2}. \quad (33)$$

To solve this system, we could develop a matrix analogue of the quadratic method, where (30) would be replaced by the cancellation of the Jacobian of the system. However, we prefer a step by step approach here, among other reasons because our system is not generic (its Jacobian has a multiple root).

From now on, we lighten notation by denoting $A = U_\varepsilon^x$, $A_1 = U_\varepsilon$, $B = U_{10}^x$, $B_1 = U_{10}$, $C = U_{100}^x$ and $C_1 = U_{100}$. We will determine three polynomial equations relating the one-variable series A_1, B_1 and C_1 , and then eliminate B_1 and C_1 to obtain a polynomial equation satisfied by $A_1 = U_\varepsilon$.

We now describe the various steps of our calculation, without giving the intermediate equations: we refer to our web pages for a Maple session where the calculations are performed.

The first equation of (32), after injecting (33), involves only one x -dependent series, namely $A = U_\varepsilon^x = U_\varepsilon(x)$. Once the denominators are cleared out, the degree in A is 2, and we can apply the quadratic method of Section 5.4.1: the discriminant (in x) of a certain discriminant (in x_0) vanishes, and this gives a first equation between A_1, B_1 and C_1 .

We then move to the second equation of (32), which (after injecting (33)) involves two x -dependent series, namely A and B . It is linear in the latter series, with coefficient:

$$1 - x + tx + 2tx(x - 1)A. \quad (34)$$

This coefficient vanishes for a (unique) series in t , denoted X , satisfying

$$X = 1 + tX + 2tX(X - 1)A_2, \quad \text{with } A_2 := U_\varepsilon(X).$$

Replacing x by X in the second equation of (32) gives another equation between X and A_2 , from which we compute

$$\begin{aligned} 2t(A_1 - 2B_1)X^2 + (1 - t - 2tA_1)X &= 1, \\ A_2 &= \frac{2XB_1}{X - 1} - A_1. \end{aligned} \quad (35)$$

We now eliminate X and A_2 between the last two identities and the first equation of (32), specialized at $x = X$. This gives a second equation between our three main unknown series A_1, B_1 and C_1 .

We finally consider the third equation of (32) (after injecting (33)), which now involves all three x -dependent series. It is linear in C , again with coefficient (34). Setting $x = X$ in this equation gives an expression of $U_{10}(X)$:

$$B_2 := U_{10}(X) = XC_1/(X - 1).$$

We now get back to the second equation of (32), differentiate it with respect to x and set $x = X$. Replacing B_2 and A_2 by the above expressions gives:

$$A_2' := \frac{\partial U_\varepsilon}{\partial x}(X) = 2 \frac{2C_1t(2X - 1)(X - 1)A_1 - 4Xt(2X - 1)B_1C_1 - B_1t(X - 1) - (t - 1)(X - 1)C_1}{t(X - 1)^2(4C_1X + X - 1)}.$$

It remains to differentiate the first equation of (32) with respect to x , specialize it at $x = X$, and plug the above values of A_2', B_2 and A_2 to obtain one more equation between A_1, B_1, C_1 and X . Eliminating X thanks to (35) gives our third and last equation between A_1, B_1 and C_1 .

From this system, we eliminate B_1 and C_1 , and obtain an equation of degree 27 for $A_1 = U_\varepsilon$. Its dominant coefficient does not vanish away from 0, and its discriminant has three roots in $[1/10, 1/16]$ (where we know that the radius must be found), respectively located around 0.07509, 0.07658 and 0.07727. Following numerically the branches that start from 1 at $t = 0$ shows that the radius of U_ε is the second one, giving the upper bound $\mu_2 = 13.057\dots$ on the growth rate μ of Eulerian orientations. From numerical estimates of the singular exponent, we predict that the series has again a ‘‘planar map’’ singularity in $(1 - \mu_2 t)^{3/2}$. This is known to hold for many series satisfying an equation with divided differences [24]. This leads us to complete Proposition 17 as follows.

Conjecture 18. *For every k , the algebraic series $U_\varepsilon^{(k)}(t; 1)$ that counts orientations of $\mathcal{U}^{(k)}$ has a unique dominant singularity $\tau_k = 1/\mu_k$ which is of the planar map type: as t approaches τ_k from below,*

$$U_\varepsilon^{(k)}(t; 1) = c_0 + c_1(1 - \mu_k t) + c_2(1 - \mu_k t)^{3/2}(1 + o(1))$$

with $c_2 \neq 0$.

6. SUPERSETS OF EULERIAN ORIENTATIONS, VIA THE PRIME DECOMPOSITION

In this section, we combine the illegal large splits of the previous section with the prime decomposition of Section 2.4 to obtain a new family of supersets of Eulerian orientations. These new supersets $\mathbb{U}^{(k)}$ satisfy $\mathbb{U}^{(k)} \subset \mathcal{U}^{(k)}$ (Proposition 20), hence they give better bounds on the growth rate μ than those obtained from the standard decomposition. Many arguments are similar to those of the previous section, and we give fewer details.

Recall from Section 2.4 that an Eulerian orientation is a sequence of prime Eulerian orientations, and that a prime (Eulerian) orientation can be obtained recursively from the atomic map by either:

- adding a loop, oriented in either way, around an orientation O_1 ,
- or a legal split on a prime orientation $O' \in \mathcal{O}$, followed by the concatenation of an arbitrary Eulerian orientation O'' at the new vertex created by the split (Figure 6).

Definition 19. *Let $k \geq 1$. Let $\mathbb{U}^{(k)}$ be the set of planar orientations obtained recursively from the atomic map by either:*

- *concatenating a sequence of prime orientations of $\mathbb{U}^{(k)}$,*
- *or adding a loop, oriented in either way, around an orientation O_1 of $\mathbb{U}^{(k)}$,*
- *or performing a legal i -split on a prime orientation $O' \in \mathbb{U}^{(k)}$, with $i \leq k$, followed by the concatenation of an arbitrary orientation O'' of $\mathbb{U}^{(k)}$ at the new vertex created by the split (small split),*
- *or performing an arbitrary i -split on a prime orientation $O' \in \mathbb{U}^{(k)}$, with $i > k$, followed by the concatenation of an arbitrary orientation O'' of $\mathbb{U}^{(k)}$ at the new vertex created by the split (large split). If the split is legal, then the new edge is given the only orientation that makes the root word balanced, otherwise the root edge is oriented away from the root vertex.*

Again, the sets $\mathbb{U}^{(k)}$ decrease to the set \mathcal{O} of all Eulerian orientations as k increases, hence their growth rates $\bar{\mu}_k$ form a non-increasing sequence of upper bounds on μ . We do not know if this sequence converges to μ . At any rate, the convergence appears to be rather slow, as shown by the estimates of $\bar{\mu}_k$ in Table 1.

Proposition 20. *For $k \geq 1$, the superset of orientations $\mathbb{U}^{(k)}$ is contained in the superset $\mathcal{U}^{(k)}$ defined in Section 5.*

Proof. We prove this by induction on the number of edges. The inclusion is obvious for orientations with no edges. Now let $O \in \mathbb{U}^{(k)}$, having at least one edge.

If O is prime and is obtained by adding a loop around a smaller orientation O_1 of $\mathbb{U}^{(k)}$ (second construction in Definition 19), then O_1 belongs to $\mathcal{U}^{(k)}$ by the induction hypothesis, and so does O (first construction in Definition 12).

Assume now that O is prime and is obtained by an i -split in a prime orientation O' of $\mathbb{U}^{(k)}$, followed by the concatenation of an orientation O'' of $\mathbb{U}^{(k)}$ at the new vertex (third or fourth construction in Definition 19). Then the orientation \tilde{O} obtained by concatenating O' and O'' at their root belongs to $\mathbb{U}^{(k)}$ (first construction in $\mathbb{U}^{(k)}$) and hence to $\mathcal{U}^{(k)}$ by the induction hypothesis. But then one can recover O by performing an i -split in \tilde{O} , which is allowed in \tilde{O} as it was allowed in O' . This is the second construction in Definition 12, hence O is in $\mathcal{U}^{(k)}$.

Assume finally that O is obtained by concatenating a prime orientation P of $\mathbb{U}^{(k)}$ and another orientation O_2 of $\mathbb{U}^{(k)}$ (first construction in Definition 19). By the induction hypothesis, both P and O_2 are in $\mathcal{U}^{(k)}$. If the root edge of P is a loop, deleting it from P leaves an orientation O_1 which is in $\mathcal{U}^{(k)}$. Then we can reconstruct O by a merge of O_1 and O_2 as in the first construction of Definition 12. If the root edge of P is not a loop (Figure 8), then P was obtained by the third or fourth construction in Definition 19: allowed split in a prime orientation P' of $\mathbb{U}^{(k)}$, followed by the concatenation of some $O'' \in \mathbb{U}^{(k)}$ at the new vertex. Let \tilde{O} be obtained by concatenating O'' , P' and O_2 (in counterclockwise order) at their roots. Then \tilde{O} is in $\mathbb{U}^{(k)}$, but also in $\mathcal{U}^{(k)}$ by

the induction hypothesis. Then O can be recovered by a split in \tilde{O} , which is allowed in \tilde{O} as it was allowed in P' (the split may have been small in P' and become large in \tilde{O} , because of the orientation O_2 , but the converse is not possible). This is the second construction in Definition 12, hence O is in $\mathcal{U}^{(k)}$. \square

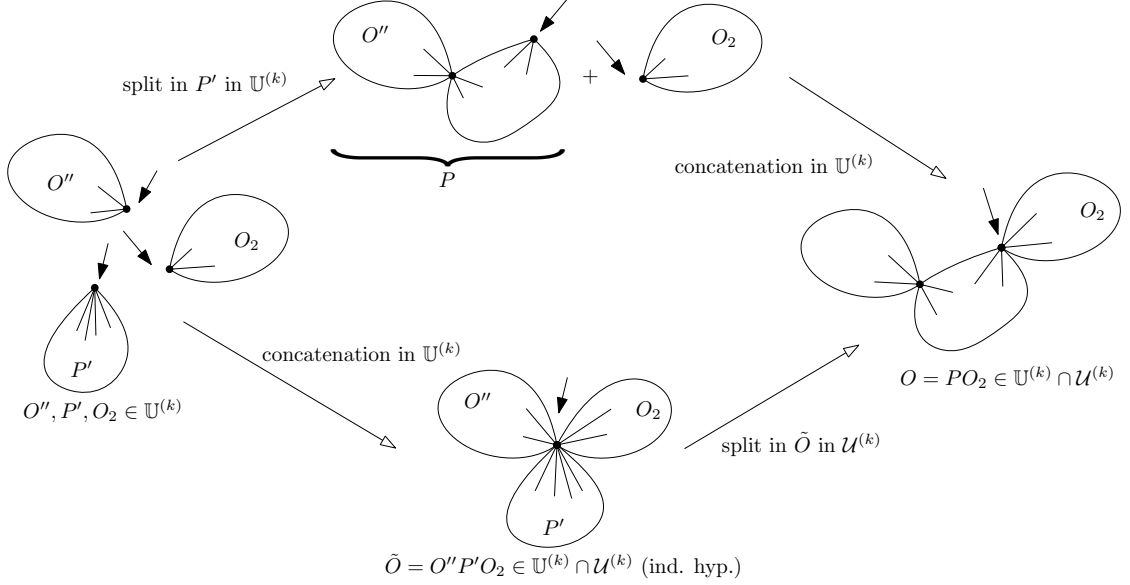


FIGURE 8. Two constructions of the orientation O : (top) in $\mathbb{U}^{(k)}$, via an i -split, (bottom) in $\mathcal{U}^{(k)}$, via a j -split, with $j = i + \text{dv}(O_2)$.

6.1. FUNCTIONAL EQUATIONS FOR $\mathbb{U}^{(k)}$

We now fix an integer k . For a word \mathbf{w} on $\{0, 1\}$, let $\mathbf{U}_{\mathbf{w}}^{(k)}(t; x) \equiv \mathbf{U}_{\mathbf{w}}(x)$ denote the generating function of orientations of $\mathbb{U}^{(k)}$ whose root word ends with \mathbf{w} , counted by the edge number (variable t) and the half-degree of the root vertex (variable x). Let $\mathbf{T}_{\mathbf{w}}^{(k)}(t) \equiv \mathbf{T}_{\mathbf{w}}$ denote the generating function of orientations of $\mathbb{U}^{(k)}$ having root word exactly \mathbf{w} . We define analogous generating functions $\mathbf{U}'_{\mathbf{w}}(x)$ and $\mathbf{T}'_{\mathbf{w}}$ for prime orientations. As in the previous section, we often denote simply by $\mathbf{U}_{\mathbf{w}}$ (resp. $\mathbf{U}'_{\mathbf{w}}$) the edge generating function $\mathbf{U}_{\mathbf{w}}(t; 1)$ (resp. $\mathbf{U}'_{\mathbf{w}}(t; 1)$), and by $\mathbf{U}_{\mathbf{w}}^x$ (resp. $\mathbf{U}'_{\mathbf{w}}^x$) the refined generating function $\mathbf{U}_{\mathbf{w}}(t; x)$ (resp. $\mathbf{U}'_{\mathbf{w}}(t; x)$).

Note that $\mathbf{T}_{\mathbf{w}} = \mathbf{T}'_{\mathbf{w}} = 0$ if \mathbf{w} is not balanced and that $\mathbf{T}_{\varepsilon} = 1$, $\mathbf{T}'_{\varepsilon} = 0$. For \mathbf{w} balanced of length between 2 and $2k$, we have both a sequential equation

$$\mathbf{T}_{\mathbf{w}} = \sum_{\mathbf{w}=\mathbf{uv}} \mathbf{T}_{\mathbf{u}} \mathbf{T}'_{\mathbf{v}} \quad (36)$$

analogous to (15), and an equation for prime orientations:

$$\mathbf{T}'_{\mathbf{w}} = t \mathbf{T}_{\mathbf{w}_c} + t \mathbf{U}_{\varepsilon} \mathbf{U}'_{\mathbf{w}_s}, \quad (37)$$

analogous to (16). The factor \mathbf{U}_{ε} accounts for the orientation concatenated after a split.

For \mathbf{w} valid of length at most $2k - 1$, we have a sequential equation, analogous to (17) but taking care of the root degree:

$$\mathbf{U}_{\mathbf{w}}^x = \mathbf{1}_{\mathbf{w}=\varepsilon} + \mathbf{U}_{\varepsilon}^x \mathbf{U}'_{\mathbf{w}}^x + \sum_{\mathbf{w}=\mathbf{uv}, \mathbf{v} \neq \mathbf{w}} \mathbf{U}_{\mathbf{u}}^x x^{|\mathbf{v}|/2} \mathbf{T}'_{\mathbf{v}}. \quad (38)$$

Finally, we have the following equation for prime orientations, which is the counterpart of (18) and involves ingredients of (25) for orientations obtained by a split:

$$\begin{aligned} U_{\mathbf{w}}^x &= 2txU_{\varepsilon}^x \mathbf{1}_{\mathbf{w}=\varepsilon} + txU_{\mathbf{w}_p}^x \mathbf{1}_{\mathbf{w} \neq \varepsilon} + tx^{|\mathbf{w}|/2} \mathbb{T}_{\mathbf{w}\varepsilon} \mathbf{1}_{\mathbf{w} \neq \varepsilon} \text{ balanced} \\ &+ tU_{\varepsilon} \left(\sum_{\substack{\mathbf{u}=\mathbf{v}\mathbf{w} \\ 2 \leq |\mathbf{u}| \leq 2k \\ \mathbf{u} \text{ balanced}}} x^{|\mathbf{u}|/2} U'_{\mathbf{u}_s} + \frac{x}{x-1} (U'_{\mathbf{w}} - x^k U'_{\mathbf{w}}) - \frac{x}{x-1} \sum_{\substack{\mathbf{u}=\mathbf{v}\mathbf{w} \\ |\mathbf{u}| \leq 2k-2}} \mathbb{T}'_{\mathbf{u}} (x^{|\mathbf{u}|/2} - x^k) \right). \end{aligned} \quad (39)$$

The first line counts orientations obtained by adding a loop, and the second those obtained by a split.

Remark 21. As in the previous section, we can use (37) to replace (39) by a slightly lighter equation:

$$\begin{aligned} U_{\mathbf{w}}^x &= x^{|\mathbf{w}|/2} \mathbb{T}'_{\mathbf{w}} + 2txU_{\varepsilon}^x \mathbf{1}_{\mathbf{w}=\varepsilon} + txU_{\mathbf{w}_p}^x \mathbf{1}_{\mathbf{w} \neq \varepsilon} + tU_{\varepsilon} \left(\sum_{\substack{\mathbf{u}=\mathbf{v}\mathbf{w} \\ |\mathbf{u}| \leq 2k-1 \\ \mathbf{u} \text{ quasi-balanced}}} x^{(1+|\mathbf{u}|)/2} U'_{\mathbf{u}} \right. \\ &\quad \left. + \frac{x}{x-1} (U'_{\mathbf{w}} - x^k U'_{\mathbf{w}}) - \frac{x}{x-1} \sum_{\substack{\mathbf{u}=\mathbf{v}\mathbf{w} \\ |\mathbf{u}| \leq 2k-2}} \mathbb{T}'_{\mathbf{u}} (x^{|\mathbf{u}|/2} - x^k) \right). \end{aligned} \quad (40)$$

Proposition 22. Consider the collection of equations consisting of:

- Equation (36), written for all balanced words \mathbf{w} of length between 2 and $2k-4$,
- Equation (37), written for all balanced words \mathbf{w} of length between 2 and $2k-2$,
- Equation (38), written for all valid words \mathbf{w} of length at most $2k-2$,
- Equation (40), written for all valid words \mathbf{w} of length at most $2k-1$.

In this collection, replace all trivial \mathbb{T} - and \mathbb{T}' -series by their value: $\mathbb{T}_{\mathbf{w}} = \mathbb{T}'_{\mathbf{w}} = 0$ when \mathbf{w} is not balanced, $\mathbb{T}_{\varepsilon} = 1$, $\mathbb{T}'_{\varepsilon} = 0$. Let R_0 denote the resulting system. The number of series it involves is $2f(k) - 3\binom{2k}{k} - \binom{2k-2}{k-1} \mathbf{1}_{k>1}$, where $f(k)$ is given by (7). Moreover, R_0 defines uniquely all these series. Its size can be (roughly) divided by two upon exploiting the 0/1 symmetry.

The proof is similar to the proofs of Propositions 3 and 9.

Remark 23. As always, we can alternatively write forward equations:

$$U_{\mathbf{w}}^x = x^{|\mathbf{w}|/2} \mathbb{T}_{\mathbf{w}} + U_{0\mathbf{w}}^x + U_{1\mathbf{w}}^x, \quad U_{\mathbf{w}}^x = x^{|\mathbf{w}|/2} \mathbb{T}'_{\mathbf{w}} + U_{0\mathbf{w}}^x + U_{1\mathbf{w}}^x.$$

6.2. EXAMPLES

6.2.1. **When $k=1$,** the system of Proposition 22 contains $2f(1) - 3 \cdot 2 = 4$ equations. Upon exploiting the 0/1 symmetry, it reads:

$$\begin{cases} U_{\varepsilon}^x = 1 + U_{\varepsilon}^x U'_{\varepsilon}, \\ U'_{\varepsilon} = 2txU_{\varepsilon}^x + tU_{\varepsilon} \left(2xU'_0 + \frac{x}{x-1} (U'_{\varepsilon} - xU'_{\varepsilon}) \right), \\ U'_{0} = txU_{\varepsilon}^x + tU_{\varepsilon} \left(xU'_0 + \frac{x}{x-1} (U'_{0} - xU'_0) \right). \end{cases} \quad (41)$$

The second equation can be replaced by the forward equation $U'_{\varepsilon} = 2U'_{0}$.

We solve this system in Section 6.4.

6.2.2. **When $k = 2$,** the system of Proposition 22 contains $2f(2) - 3 \cdot 6 - 2 = 22$ equations. Upon exploiting the 0/1 symmetry and the forward equations, it reads:

$$\left\{ \begin{array}{l} \mathsf{T}'_{01} = t + t\mathsf{U}_\varepsilon \mathsf{U}'_1, \\ \mathsf{U}_\varepsilon^x = 1 + \mathsf{U}_\varepsilon^x \mathsf{U}'_\varepsilon, \\ \mathsf{U}_0^x = \mathsf{U}_\varepsilon^x \mathsf{U}'_0, \\ \mathsf{U}_{10}^x = \mathsf{U}_{01}^x = \mathsf{U}_\varepsilon^x \mathsf{U}'_{10}, \\ \mathsf{U}_{00}^x = \mathsf{U}_{11}^x = \mathsf{U}_\varepsilon^x \mathsf{U}'_{00}, \\ \mathsf{U}'_\varepsilon = 2\mathsf{U}'_0, \\ \mathsf{U}'_0 = \mathsf{U}'_1 = t\mathsf{U}_\varepsilon^x + t\mathsf{U}_\varepsilon \left(x\mathsf{U}'_0 + x^2(\mathsf{U}'_{100} + \mathsf{U}'_{010} + \mathsf{U}'_{110}) + \frac{x}{x-1}(\mathsf{U}'_0 - x^2\mathsf{U}'_0) + x^2\mathsf{T}'_{10} \right), \\ \mathsf{U}'_{10} = \mathsf{U}'_{01} = x\mathsf{T}'_{10} + \mathsf{U}'_{110} + \mathsf{U}'_{010}, \\ \phantom{\mathsf{U}'_{10} = \mathsf{U}'_{01}} = x\mathsf{T}'_{10} + t\mathsf{U}_\varepsilon^x + t\mathsf{U}_\varepsilon \left(x^2\mathsf{U}'_{010} + x^2\mathsf{U}'_{110} + \frac{x}{x-1}(\mathsf{U}'_{10} - x^2\mathsf{U}'_{10}) + x^2\mathsf{T}'_{10} \right), \\ \mathsf{U}'_{00} = t\mathsf{U}_\varepsilon^x + t\mathsf{U}_\varepsilon \left(x^2\mathsf{U}'_{100} + \frac{x}{x-1}(\mathsf{U}'_{00} - x^2\mathsf{U}'_{00}) \right), \\ \mathsf{U}'_{100} = t\mathsf{U}_{10}^x + t\mathsf{U}_\varepsilon \left(x^2\mathsf{U}'_{100} + \frac{x}{x-1}(\mathsf{U}'_{100} - x^2\mathsf{U}'_{100}) \right), \\ \mathsf{U}'_{010} = t\mathsf{U}_{01}^x + t\mathsf{U}_\varepsilon \left(x^2\mathsf{U}'_{010} + \frac{x}{x-1}(\mathsf{U}'_{010} - x^2\mathsf{U}'_{010}) \right), \\ \mathsf{U}'_{110} = t\mathsf{U}_{11}^x + t\mathsf{U}_\varepsilon \left(x^2\mathsf{U}'_{110} + \frac{x}{x-1}(\mathsf{U}'_{110} - x^2\mathsf{U}'_{110}) \right). \end{array} \right. \quad (42)$$

We explain in Section 6.4 below how to solve this system.

6.3. ALGEBRAICITY

The analogue of Proposition 17 holds for the supersets obtained via the prime decomposition.

Proposition 24. *For any $k \geq 1$, the generating function $\mathsf{U}_\varepsilon^{(k)}(t; x)$ that counts orientations of $\mathbb{U}^{(k)}$ is algebraic.*

Proof. Again, the idea is to apply Theorem 16 to the system of Proposition 22, after writing $x = 1 + z$. The proof is roughly the same as that of Proposition 17: we define a polynomial system involving two variables, t and z , and six families of unknowns $A_{\mathbf{w}}, A'_{\mathbf{w}}, B_{\mathbf{w}}, B'_{\mathbf{w}}, C_{\mathbf{w}}, C'_{\mathbf{w}}$. The equations they satisfy are those of Proposition 22, rewritten with

$$\mathsf{T}_{\mathbf{w}} \rightarrow A_{\mathbf{w}}, \quad \mathsf{T}'_{\mathbf{w}} \rightarrow A'_{\mathbf{w}}, \quad \mathsf{U}_{\mathbf{w}}^x \rightarrow B_{\mathbf{w}}, \quad \mathsf{U}'_{\mathbf{w}}^x \rightarrow B'_{\mathbf{w}}, \quad \mathsf{U}_{\mathbf{w}} \rightarrow C_{\mathbf{w}}, \quad \mathsf{U}'_{\mathbf{w}} \rightarrow C'_{\mathbf{w}}.$$

In fact, the only series $\mathsf{U}_{\mathbf{w}}$ occurring in our system is U_ε , so that the polynomial system we construct involves C_ε , but no other C -series. The prescribed dependences are that the A, A', C and C' series are independent of z . For instance, when $k = 1$ we convert (41) into:

$$\begin{cases} B_\varepsilon = 1 + B_\varepsilon B'_\varepsilon, \\ B'_\varepsilon = 2t(1+z)B_\varepsilon + tC_\varepsilon \left(2(1+z)C'_0 + \frac{1+z}{z}(B'_\varepsilon - (1+z)C'_\varepsilon) \right), \\ B'_0 = t(1+z)B_\varepsilon + tC_\varepsilon \left((1+z)C'_0 + \frac{1+z}{z}(B'_0 - (1+z)C'_0) \right). \end{cases}$$

However, this is not sufficient, because this system does *not* imply that C_ε is B_ε at $z = 0$ (it *does* however imply that C'_ε is B'_ε at $z = 0$, and similarly for C'_0 and B'_0). Hence, in order to apply Popescu's theorem, we need to add to our collection of equations the case $x = 1$, $\mathbf{w} = \varepsilon$ of (38), namely $C_\varepsilon = 1 + C_\varepsilon C'_\varepsilon$. The rest of the argument mimics the proof of Proposition 17. \square

6.4. BACK TO EXAMPLES

We first return to the system (41) obtained for $k = 1$. In the second equation, replace U_ε^x by $1/(1 - \mathsf{U}'_\varepsilon)$, U_ε by $1/(1 - \mathsf{U}'_\varepsilon)$ and U'_0 by $\mathsf{U}'_\varepsilon/2$. This gives a polynomial equation involving only U'_ε and U'_0 , which can be solved by the quadratic method already used in Section 5.4. This gives for U'_ε a cubic equation. Getting back to $\mathsf{U}_\varepsilon = 1/(1 - \mathsf{U}'_\varepsilon)$, we obtain for the generating

function $U_\varepsilon^{(1)}$ of orientations in $\mathbb{U}^{(1)}$ the same cubic equation (31) as for orientations of $\mathcal{U}^{(1)}$. In fact, one can check that $\mathbb{U}^{(1)} = \mathcal{U}^{(1)}$. Of course, the upper bound on μ is $\bar{\mu}_1 = \mu_1 = 13.0659\dots$

Let us now solve the system (42) obtained for $k = 2$. It can be reduced to a system of three equations defining the series U'_ε, U'_{10} and U'_{100} :

$$\begin{cases} U'_\varepsilon &= 2txU_\varepsilon^x + tU_\varepsilon \left(xU'_\varepsilon + 2x^2(U'_{100} + U'_{10}) + \frac{x}{x-1}(U'^x_\varepsilon - x^2U'_\varepsilon) \right), \\ U'^x_{10} &= xt(1 + U_\varepsilon U'_\varepsilon/2) + tx(U_\varepsilon^x - 1)/2 + tU_\varepsilon \left(x^2U'_{10} + \frac{x}{x-1}(U'^x_{10} - x^2U'_{10}) \right), \\ U'^x_{100} &= txU'_{10} + tU_\varepsilon \left(x^2U'_{100} + \frac{x}{x-1}(U'^x_{100} - x^2U'_{100}) \right), \end{cases} \quad (43)$$

in which we inject

$$U_\varepsilon = \frac{1}{1 - U'_\varepsilon}, \quad U_\varepsilon^x = \frac{1}{1 - U'^x_\varepsilon} \quad \text{and} \quad U^x_{10} = \frac{U'^x_{10}}{1 - U'^x_\varepsilon}. \quad (44)$$

We lighten notation by denoting $A = U'^x_\varepsilon$, $A_1 = U'_\varepsilon$, $B = U'^x_{10}$, $B_1 = U'_{10}$, $C = U'^x_{100}$ and $C_1 = U'_{100}$, and we follow the steps used in Section 5.4.2 to solve System (28). Again, we refer to our web pages for the corresponding Maple session. The intermediate steps are as follows. We first apply the quadratic method to the first equation. We then turn to the second one. The equation satisfied by X is

$$X = 1 + \frac{t}{1 - t - A_1},$$

(Note that it gives X explicitly in terms of A_1 , whereas we had a quadratic equation (35) in the previous case.) We then derive

$$A_2 := U'_\varepsilon(X) = 1 + t \frac{1 - A_1}{t - 2B_1(1 - A_1)}.$$

We finally consider the third equation of (43) (after injecting (44)), and derive:

$$B_2 := U'_{10}(X) = (1 - A_2)C_1/t.$$

Then it follows from the second equation that:

$$A'_2 := \frac{\partial U'_\varepsilon}{\partial x}(X) = \frac{2(1 - A_1)(1 - t - A_1)^2((1 - t - A_1)C_1 + 2(-1 + A_1)B_1^2 + tB_1)}{(t - 2B_1(1 - A_1))^3}.$$

At the end, we obtain an equation of degree 28 for $A_1 = U'_\varepsilon$, and then for U_ε . Its dominant coefficient does not vanish away from 0, and its discriminant has only one root in $[1/10, 1/16]$ (where we know that the radius must be found), around 0.0766. This gives the upper bound $\bar{\mu}_2 = 13.047\dots$ on the growth rate μ of Eulerian orientations. From numerical estimates of the singular exponent, we predict that the series has again a ‘‘planar map’’ singularity in $(1 - \bar{\mu}_2 t)^{3/2}$.

For $k = 3, 4$ and 5 , we have generated our systems of equations and computed the first 100 coefficients of $U_\varepsilon^{(k)}(t; x)$. From this we get the estimates of the growth rates $\bar{\mu}_k$ shown in Table 2. The singularity still appears to be in $(1 - \bar{\mu}_k t)^{3/2}$. We conjecture that this holds for any k .

Conjecture 25. *For every k , the algebraic series $U_\varepsilon^{(k)}(t; 1)$ that counts orientations of $\mathbb{U}^{(k)}$ has a unique dominant singularity $\bar{\tau}_k = 1/\bar{\mu}_k$ which is of the planar map type: as t approaches $\bar{\tau}_k$ from below,*

$$U_\varepsilon^{(k)}(t; 1) = c_0 + c_1(1 - \mu_k t) + c_2(1 - \bar{\mu}_k t)^{3/2}(1 + o(1))$$

for $c_2 \neq 0$.

7. FINAL COMMENTS

As mentioned at the end of the introduction, it might be easier to count Eulerian orientation of 4-valent maps. In such orientations, each vertex has two in-going and two outgoing edges, so that counting Eulerian orientations means solving the so-called *ice-model* on (random) 4-valent maps [3, Chap. 8]. The number of Eulerian orientations of a 4-valent planar map is known to be a third of the number of proper 3-colourings of its dual [60]. Thus counting these orientations is equivalent to counting 3-coloured planar quadrangulations. A number of enumerative results are already known about coloured maps. In particular, 3-coloured planar maps, and 3-coloured planar triangulations, have algebraic generating functions [6, 17, 54, 58]. More generally, q -coloured planar maps and triangulations have differentially algebraic generating functions. So this could be true for 3-coloured quadrangulations as well, and hence for Eulerian orientations of 4-valent maps. Several results on this problem appear in the physics literature, but there does not seem to be an explicit exact solution at the moment [42, 61].

To finish, let us recall two questions left open by this paper (beyond the enumeration of Eulerian orientations!).

- Do the growth rates $\bar{\mu}_k$ of orientations of $\mathbb{U}^{(k)}$ decrease to the growth rate μ of Eulerian orientations, or to a larger value?
- In Sections 3.3 and 4.3, we have used a general result about positive irreducible systems of polynomial equations to prove that the generating functions of our subsets of Eulerian orientations have a square root singularity. Could one define a notion of positive irreducible system of polynomial equations *with divided differences* whose solutions would systematically exhibit a singularity in $(1 - \mu t)^{3/2}$? Hopefully this would apply to our supersets of orientations and prove Conjectures 18 and 25. A first step in this direction, applicable to a single equation, is achieved in [24].

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