

Counting Subwords Occurrences in Base- b Expansions

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Abstract

We count the number of distinct (scattered) subwords occurring in the base- b expansion of the non-negative integers. More precisely, we consider the sequence $(S_b(n))_{n \geq 0}$ counting the number of positive entries on each row of a generalization of the Pascal triangle to binomial coefficients of base- b expansions. By using a convenient tree structure, we provide recurrence relations for $(S_b(n))_{n \geq 0}$ leading to the b -regularity of the latter sequence. Then we deduce the asymptotics of the summatory function of the sequence $(S_b(n))_{n \geq 0}$.

1 Introduction

A *finite word* is a finite sequence of letters belonging to a finite set called the *alphabet*. The *binomial coefficient* $\binom{u}{v}$ of two finite words u and v is the number of times v occurs as a subsequence of u (meaning as a “scattered” subword). All along the paper, we let b denote an integer greater than 1. We let $\text{rep}_b(n)$ denote the (greedy) base- b expansion of $n \in \mathbb{N} \setminus \{0\}$ starting with a non-zero digit. We set $\text{rep}_b(0)$ to be the empty word denoted by ε . We let

$$L_b = \{1, \dots, b-1\} \{0, \dots, b-1\}^* \cup \{\varepsilon\}$$

be the set of base- b expansions of the non-negative integers. For all $w \in \{0, \dots, b-1\}^*$, we also define $\text{val}_b(w)$ to be the value of w in base b , i.e., if $w = w_n \dots w_0$ with $w_i \in \{0, \dots, b-1\}$ for all i , then $\text{val}_b(w) = \sum_{i=0}^n w_i b^i$.

Several generalizations and variations of the Pascal triangle exist and lead to interesting combinatorial, geometrical or dynamical properties [5, 6, 13, 14, 15]. Ordering the words of L_b by increasing genealogical order, we introduced Pascal-like triangles P_b [15] where the entry $P_b(m, n)$ is $\binom{\text{rep}_b(m)}{\text{rep}_b(n)}$. Clearly P_b contains $(b-1)$ copies of the usual Pascal triangle when only considering words of the form a^m with $a \in \{1, \dots, b-1\}$ and $m \geq 0$. In Figure 1, we depict the first few elements of P_3 [A284441](#) and its compressed version highlighting the number of positive elements on each line. The data provided by this compressed version is summed up in Definition 1.

Definition 1. For $n \geq 0$, we define the sequence $(S_b(n))_{n \geq 0}$ by setting

$$S_b(n) := \# \left\{ v \in L_b \mid \binom{\text{rep}_b(n)}{v} > 0 \right\}. \quad (1)$$

We also consider the summatory function $(A_b(n))_{n \geq 0}$ of the sequence $(S_b(n))_{n \geq 0}$ defined by $A_b(0) = 0$ and for all $n \geq 1$,

$$A_b(n) := \sum_{j=0}^{n-1} S_b(j).$$

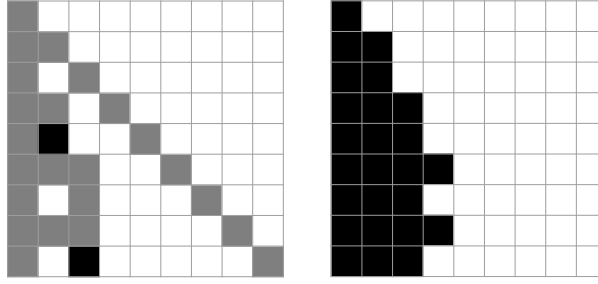


Figure 1: On the left, the first few rows of the generalized Pascal triangle P_3 (a white (resp., gray; resp., black) square corresponds to 0 (resp., 1; resp., 2)) and on the right, its compressed version.

The quantity $A_b(n)$ can be thought of as the total number of base- b expansions occurring as subwords in the base- b expansion of integers less than n (the same subword is counted k times if it occurs in the base- b expansion of k distinct integers).

In some sense, the sequences $(S_b(n))_{n \geq 0}$ and $(A_b(n))_{n \geq 0}$ measure the sparseness of P_b .

Example 2. If $b = 3$, then the first few terms of the sequence $(S_3(n))_{n \geq 0}$ [A282715](#) are

$$1, 2, 2, 3, 3, 4, 3, 4, 3, 4, 5, 6, 5, 4, 6, 7, 7, 6, 4, 6, 5, 7, 6, 7, 5, 6, 4, 5, 7, 8, 8, 7, 10, \dots$$

For instance, the subwords of the word 121 are $\varepsilon, 1, 2, 11, 12, 21, 121$. Thus, $S_3(\text{val}_3(121)) = S_3(16) = 7$. The first few terms of $(A_3(n))_{n \geq 0}$ [A284442](#) are

$$0, 1, 3, 5, 8, 11, 15, 18, 22, 25, 29, 34, 40, 45, 49, 55, \dots$$

We studied [16] the triangle P_2 [A282714](#) and the sequence $(S_2(n))_{n \geq 0}$ [A007306](#), which turns out to be the subsequence with odd indices of the Stern–Brocot sequence. The sequence $(S_2(n))_{n \geq 0}$ is 2-regular in the sense of Allouche and Shallit [1]. We studied [17] the behavior of $(A_2(n))_{n \geq 0}$ [A282720](#). To this aim, we exploited a particular decomposition of $A_2(2^\ell + r)$, for all $\ell \geq 1$ and all $0 \leq r < 2^\ell$, using powers of 3.

1.1 Our contribution

We conjectured six recurrence relations for $(S_3(n))_{n \geq 0}$ depending on the position of n between two consecutive powers of 3; see [16]. Using the heuristic from [3] suggesting recurrence relations, the sequence $(S_3(n))_{n \geq 0}$ was expected to be 3-regular. It was not obvious that we could derive general recurrence relations for $(S_b(n))_{n \geq 0}$ from the form of those satisfied by $(S_2(n))_{n \geq 0}$. We thought that $(b-1)b$ recurrence relations should be needed in the general case, leading to a cumbersome statement. Moreover it was computationally challenging to obtain many terms of $(S_b(n))_{n \geq 0}$ for large b because the number of words of length n in L_b grows like b^n . Therefore we lack data to conjecture the b -regularity of $(S_b(n))_{n \geq 0}$.

When studying $(A_2(n))_{n \geq 0}$, a possible extension seemed to emerge [17]. In particular, we prove that $A_2(2n) = 3A_2(n)$ and, sustained by computer experiments, we conjectured that $A_b(nb) = (2b-1)A_b(n)$.

Surprisingly, for all $b \geq 2$, we show in Section 2 that the recurrence relations satisfied by $(S_b(n))_{n \geq 0}$ reduce to three forms; see Proposition 3. In particular, this proves the conjecture stated in [16]. Therefore, in Section 3, we deduce the b -regularity of $(S_b(n))_{n \geq 0}$; see Theorem 16. Moreover we obtain a linear representation of the sequence with $b \times b$ matrices. We also show that $(S_b(n))_{n \geq 0}$ is palindromic over $[(b-1)b^\ell, b^{\ell+1}]$.

The key to study the asymptotics of $(A_b(n))_{n \geq 0}$ is to obtain specific recurrence relations for this sequence. In Proposition 26, we show that these relations involve powers of $(2b-1)$. Therefore, we prove the conjecture

$\text{rep}_b(n)$	ε	x	$x0$	xx	xy	$x00$	$x0x$	$x0y$	$xx0$	xxx	xyy	$xy0$	xyx	xyy	xyz
$S_b(n)$	1	2	3	3	4	4	5	6	5	4	6	7	7	6	8

Table 1: The first few values of $S_b(n)$ for $0 \leq n < b^3$, with pairwise distinct $x, y, z \in \{1, \dots, b-1\}$.

about $A_b(nb)$. In Section 4, using the so-called $(2b-1)$ -decompositions, we may apply the method introduced in [17].

We think that this paper motivates the quest for generalized Stern–Brocot sequences and analogues of the Farey tree [4, 7, 8, 11, 12, 18]. Namely can one reasonably define a tree structure, or some other combinatorial structure, in which the sequence $(S_b(n))_{n \geq 0}$ naturally appears?

Most of the results are proved by induction and the base case usually takes into account the values of $S_b(n)$ for $0 \leq n < b^2$. These values are easily obtained from Definition 1 and summarized in Table 1.

2 General recurrence relations in base b

The aim of this section is to prove the following result exhibiting recurrence relations satisfied by the sequence $(S_b(n))_{n \geq 0}$. This result is useful to prove that the summatory function of the latter sequence also satisfies recurrence relations; see Section 4.

Proposition 3. *The sequence $(S_b(n))_{n \geq 0}$ satisfies $S_b(0) = 1$, $S_b(1) = \dots = S_b(b-1) = 2$, and, for all $x, y \in \{1, \dots, b-1\}$ with $x \neq y$, all $\ell \geq 1$ and all $r \in \{0, \dots, b^{\ell-1} - 1\}$,*

$$S_b(xb^\ell + r) = S_b(xb^{\ell-1} + r) + S_b(r); \quad (2)$$

$$S_b(xb^\ell + xb^{\ell-1} + r) = 2S_b(xb^{\ell-1} + r) - S_b(r); \quad (3)$$

$$S_b(xb^\ell + yb^{\ell-1} + r) = S_b(xb^{\ell-1} + r) + 2S_b(yb^{\ell-1} + r) - 2S_b(r). \quad (4)$$

For the sake of completeness, we recall the definition of a particularly useful tool called *the trie of subwords* to prove Proposition 3. This tool is also useful to prove the b -regularity of the sequence $(S_b(n))_{n \geq 0}$; see Section 3.

Definition 4. Let w be a finite word over $\{0, \dots, b-1\}$. The language of its subwords is factorial, i.e., if xyz is a subword of w , then y is also a subword of w . Thus we may associate with w , the *trie¹ of its subwords*. The root is ε and if u and ua are two subwords of w with $a \in \{0, \dots, b-1\}$, then ua is a child of u . We let $\mathcal{T}(w)$ denote the subtree in which we only consider the children $1, \dots, b-1$ of the root ε and their successors, if they exist.

Remark 5. The number of nodes on level $\ell \geq 0$ in $\mathcal{T}(w)$ counts the number of subwords of length ℓ in L_b occurring in w . In particular, the number of nodes of the trie $\mathcal{T}(\text{rep}_b(n))$ is exactly $S_b(n)$ for all $n \geq 0$.

Definition 6. For each non-empty word $w \in L_b$, we consider a factorization of w into maximal blocks of consecutively distinct letters (i.e., $a_i \neq a_{i+1}$ for all i) of the form

$$w = a_1^{n_1} \cdots a_M^{n_M},$$

with $n_\ell \geq 1$ for all ℓ . For each $\ell \in \{0, \dots, M-1\}$, we consider the subtree T_ℓ of $\mathcal{T}(w)$ whose root is the node $a_1^{n_1} \cdots a_\ell^{n_\ell} a_{\ell+1}$. For convenience, we set T_M to be an empty tree with no node. Roughly speaking, we have a root of a new subtree T_ℓ for each new variation of digits in w . For each $\ell \in \{0, \dots, M-1\}$, we also let $\#T_\ell$ denote the number of nodes of the tree T_ℓ .

Note that for $k-i \geq 2$, one could possibly have $a_k = a_i$. For each $\ell \in \{0, \dots, M-1\}$, we let $\text{Alph}(\ell)$ denote the set of letters occurring in $a_{\ell+1} \cdots a_M$. Then for each letter $a \in \text{Alph}(\ell)$, we let $j(a, \ell)$ denote the smallest index in $\{\ell+1, \dots, M\}$ such that $a_{j(a, \ell)} = a$.

¹This tree is also called prefix tree or radix tree. All successors of a node have a common prefix and the root is the empty word.

Example 7. In this example, we set $b = 3$ and $w = 22000112 \in L_3$. Using the previous notation, we have $M = 4$, $a_1 = 2$, $a_2 = 0$, $a_3 = 1$ and $a_4 = 2$. For instance, $\text{Alph}(0) = \{0, 1, 2\}$, $\text{Alph}(2) = \{1, 2\}$ and $j(0, 0) = 2$, $j(1, 0) = 3$, $j(2, 0) = 1$ and $j(2, 1) = 4$.

The following result describes the structure of the tree $\mathcal{T}(w)$. It directly follows from the definition.

Proposition 8 ([16, Proposition 27]). *Let w be a finite word in L_b . With the above notation about M and the subtrees T_ℓ , the tree $\mathcal{T}(w)$ has the following properties.*

1. *The node of label ε has $\#(\text{Alph}(0) \setminus \{0\})$ children that are a for $a \in \text{Alph}(0) \setminus \{0\}$. Each child a is the root of a tree isomorphic $T_{j(a,0)-1}$.*
2. *For each $\ell \in \{0, \dots, M-1\}$ and each $i \in \{0, \dots, n_{\ell+1}-1\}$ with $(\ell, i) \neq (0, 0)$, the node of label $x = a_1^{n_1} \dots a_\ell^{n_\ell} a_{\ell+1}^i$ has $\#(\text{Alph}(\ell))$ children that are xa for $a \in \text{Alph}(\ell)$. Each child xa with $a \neq a_{\ell+1}$ is the root of a tree isomorphic to $T_{j(a,\ell)-1}$.*

Example 9. Let us continue Example 7. The tree $\mathcal{T}(22000112)$ is depicted in Figure 2. We use three different colors to represent the letters 0, 1, 2. The tree T_0 (resp., T_1 ; resp., T_2 ; resp., T_3) is the subtree of

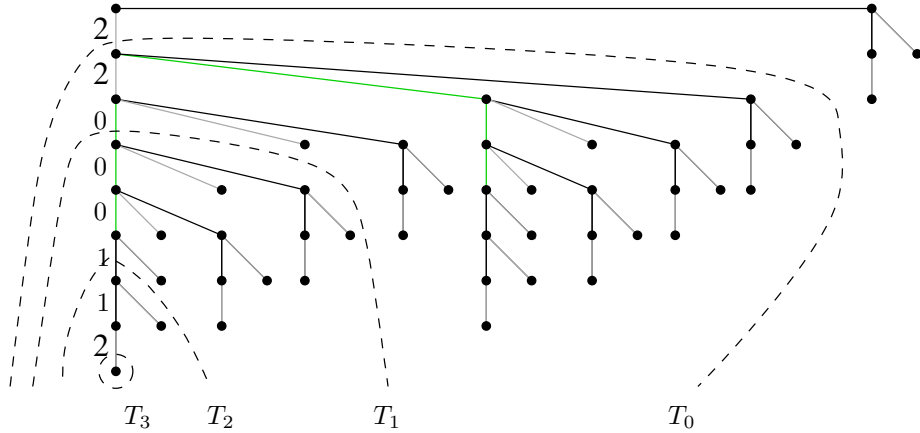


Figure 2: The trie $\mathcal{T}(22000112)$.

$\mathcal{T}(w)$ with root 2 (resp., 2^20 ; resp., 2^20^31 ; resp., $2^20^31^22$). These subtrees are represented in Figure 2 using dashed lines. The tree T_3 is limited to a single node since the number of nodes of T_{M-1} is n_M , which is equal to 1 in this example.

Using tries of subwords, we prove the following five lemmas. Their proofs are essentially the same, so we only prove two of them.

Lemma 10. *For each letter $x \in \{1, \dots, b-1\}$ and each word $u \in \{0, \dots, b-1\}^*$, we have*

$$\# \left\{ v \in L_b \mid \binom{x00u}{v} > 0 \right\} = 2 \cdot \# \left\{ v \in L_b \mid \binom{x0u}{v} > 0 \right\} - \# \left\{ v \in L_b \mid \binom{xu}{v} > 0 \right\}.$$

Proof. Recall that from Remark 5, we need to prove that $\#\mathcal{T}(x00u) = 2\#\mathcal{T}(x0u) - \#\mathcal{T}(xu)$.

Assume first that u is of the form $u = 0^n$, $n \geq 0$. The tree $\mathcal{T}(xu)$ is linear and has $n+2$ nodes, $\mathcal{T}(x0u)$ has $n+3$ nodes and $\mathcal{T}(x00u)$ has $n+4$ nodes. The formula holds.

Now suppose that u contains other letters than 0. We let a_1, \dots, a_m denote all the pairwise distinct letters of u different from 0. They are implicitly ordered with respect to their first appearance in u . If

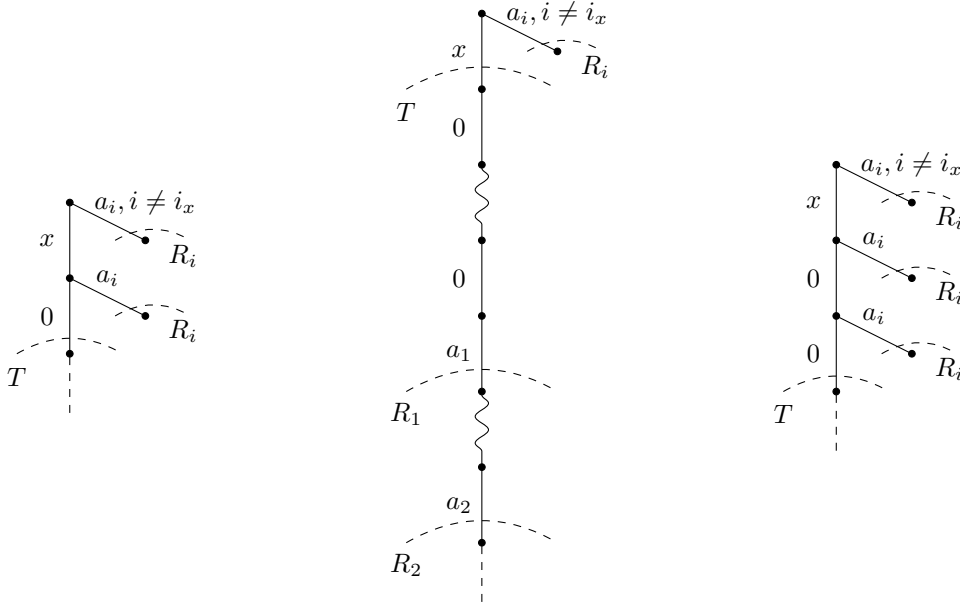


Figure 3: Schematic structure of the trees $\mathcal{T}(x0u)$, $\mathcal{T}(xu)$ and $\mathcal{T}(x00u)$.

$x \in \{a_1, \dots, a_m\}$, we let $i_x \in \{1, \dots, m\}$ denote the index such that $a_{i_x} = x$. For all $i \in \{1, \dots, m\}$, we let $u_i a_i$ denote the prefix of u that ends with the first occurrence of the letter a_i in u , and we let R_i denote the subtree of $\mathcal{T}(xu)$ with root $xu_i a_i$.

First, observe that the subtree T of $\mathcal{T}(xu)$ with root x is equal to the subtree of $\mathcal{T}(x0u)$ with root $x0$ and also to the subtree of $\mathcal{T}(x00u)$ with root $x00$.

Secondly, for all $i \in \{1, \dots, m\}$, the subtree of $\mathcal{T}(x0u)$ with root xa_i is R_i . Similarly, $\mathcal{T}(x00u)$ contains two copies of R_i : the subtrees of root xa_i and $x0a_i$.

Finally, for all $i \in \{1, \dots, m\}$ with $i \neq i_x$, the subtree of $\mathcal{T}(x0u)$ with root a_i is R_i and the subtree of $\mathcal{T}(x00u)$ with root a_i is R_i .

The situation is depicted in Figure 3 where we put a unique edge for several indices when necessary, e.g., the edge labeled by a_i stands for m edges labeled by a_1, \dots, a_m . The claimed formula holds since

$$2 \cdot (2 + \#T + 2 \sum_{\substack{1 \leq i \leq m \\ i \neq i_x}} \#R_i + \#R_{i_x}) - (1 + \#T + \sum_{\substack{1 \leq i \leq m \\ i \neq i_x}} \#R_i) = 3 + \#T + 3 \sum_{\substack{1 \leq i \leq m \\ i \neq i_x}} \#R_i + 2\#R_{i_x}.$$

□

Lemma 11. For each letter $x \in \{1, \dots, b-1\}$ and each word $u \in \{0, \dots, b-1\}^*$, we have

$$\# \left\{ v \in L_b \mid \binom{xx0u}{v} > 0 \right\} = \# \left\{ v \in L_b \mid \binom{x0u}{v} > 0 \right\} + \# \left\{ v \in L_b \mid \binom{xu}{v} > 0 \right\}.$$

Proof. The proof is similar to the proof of Lemma 10. □

Lemma 12. For all letters $x, y \in \{1, \dots, b-1\}$ and each word $u \in \{0, \dots, b-1\}^*$, we have

$$\# \left\{ v \in L_b \mid \binom{x0yu}{v} > 0 \right\} = \# \left\{ v \in L_b \mid \binom{xyu}{v} > 0 \right\} + \# \left\{ v \in L_b \mid \binom{yu}{v} > 0 \right\}.$$

Proof. The proof is similar to the proof of Lemma 10. Observe that one needs to divide the proof into two cases according to whether x is equal to y or not. As a first case, also consider $u = y^n$ with $n \geq 0$ instead of $u = 0^n$ with $n \geq 0$. \square

Lemma 13. For all letters $x, y \in \{1, \dots, b-1\}$ and each word $u \in \{0, \dots, b-1\}^*$, we have

$$\# \left\{ v \in L_b \mid \binom{xyxu}{v} > 0 \right\} = 2 \cdot \# \left\{ v \in L_b \mid \binom{xyu}{v} > 0 \right\} - \# \left\{ v \in L_b \mid \binom{yu}{v} > 0 \right\}.$$

Proof. The proof is similar to the proof of Lemma 12. \square

The next lemma having a slightly more technical proof, we present it.

Lemma 14. For all letters $x, y \in \{1, \dots, b-1\}$ with $x \neq y$, $z \in \{0, \dots, b-1\}$ and each word $u \in \{0, \dots, b-1\}^*$, we have

$$\begin{aligned} \# \left\{ v \in L_b \mid \binom{xyzxu}{v} > 0 \right\} &= \# \left\{ v \in L_b \mid \binom{xzcu}{v} > 0 \right\} + 2 \cdot \# \left\{ v \in L_b \mid \binom{yzcu}{v} > 0 \right\} \\ &\quad - 2 \cdot \# \left\{ v \in L_b \mid \binom{zcu}{v} > 0 \right\}. \end{aligned}$$

Proof. Let $x, y \in \{1, \dots, b-1\}$ with $x \neq y$, $z \in \{0, \dots, b-1\}$, and let $u \in \{0, \dots, b-1\}^*$. Our reasoning is again based on the structure of the associated trees. The proof is divided into two cases depending on the fact that $z = 0$ or not.

- As a first case, suppose that $z \neq 0$. Now assume that u is of the form $u = z^n$, $n \geq 0$. If $x \neq z$ and $y \neq z$, the tree $\mathcal{T}(zu)$ is linear and has $n+2$ nodes, $\mathcal{T}(xzu)$ and $\mathcal{T}(yzu)$ have $2(n+2)$ nodes and $\mathcal{T}(xyzzu)$ has $4(n+2)$ nodes and the claimed formula holds. If $x \neq z$ and $y = z$, the tree $\mathcal{T}(zu)$ is linear and has $n+2$ nodes, $\mathcal{T}(xzu)$ has $2(n+2)$ nodes, $\mathcal{T}(yzu)$ has $n+3$ nodes and $\mathcal{T}(xyzzu)$ has $2(n+3)$ nodes and the claimed formula holds. If $x = z$ and $y \neq z$, the tree $\mathcal{T}(zu)$ is linear and has $n+2$ nodes, $\mathcal{T}(xzu)$ has $n+3$ nodes, $\mathcal{T}(yzu)$ has $2(n+2)$ nodes and $\mathcal{T}(xyzzu)$ has $3(n+2)+1$ nodes and the claimed formula holds.

Now suppose that u contains other letters than z . We let a_1, \dots, a_m denote all the pairwise distinct letters of u different from z . They are implicitly ordered with respect to their first appearance in u . If $x, y, 0 \in \{a_1, \dots, a_m\}$, we let $i_x, i_y, i_0 \in \{1, \dots, m\}$ respectively denote the indices such that $a_{i_x} = x$, $a_{i_y} = y$ and $a_{i_0} = 0$. For all $i \in \{1, \dots, m\}$, we let $u_i a_i$ denote the prefix of u that ends with the first occurrence of the letter a_i in u , and we let R_i denote the subtree of $\mathcal{T}(zu)$ with root $zu_i a_i$.

First, observe that the subtree T of $\mathcal{T}(zu)$ with root z is equal to the subtree of $\mathcal{T}(xzu)$ with root xz , to the subtree of $\mathcal{T}(yzu)$ with root yz and also to the subtree of $\mathcal{T}(xyzzu)$ with root xyz .

Suppose that $x \neq z$ and $y \neq z$. Using the same reasoning as in the proof of Lemma 10, the situation is depicted in Figure 4. The claimed formula holds since

$$\begin{aligned} &(2 + 2\#T + 2 \sum_{\substack{1 \leq i \leq m \\ i \neq i_x, i_y, i_0}} \#R_i + \#R_{i_x} + 2\#R_{i_y} + \#R_{i_0}) \\ &+ 2 \cdot (2 + 2\#T + 2 \sum_{\substack{1 \leq i \leq m \\ i \neq i_x, i_y, i_0}} \#R_i + 2\#R_{i_x} + \#R_{i_y} + \#R_{i_0}) \\ &- 2 \cdot (1 + \#T + \sum_{\substack{1 \leq i \leq m \\ i \neq i_x, i_y, i_0}} \#R_i + \#R_{i_x} + \#R_{i_y}) \\ &= 4 + 4\#T + 4 \sum_{\substack{1 \leq i \leq m \\ i \neq i_x, i_y, i_0}} \#R_i + 3\#R_{i_x} + 2\#R_{i_y} + 3\#R_{i_0}. \end{aligned}$$

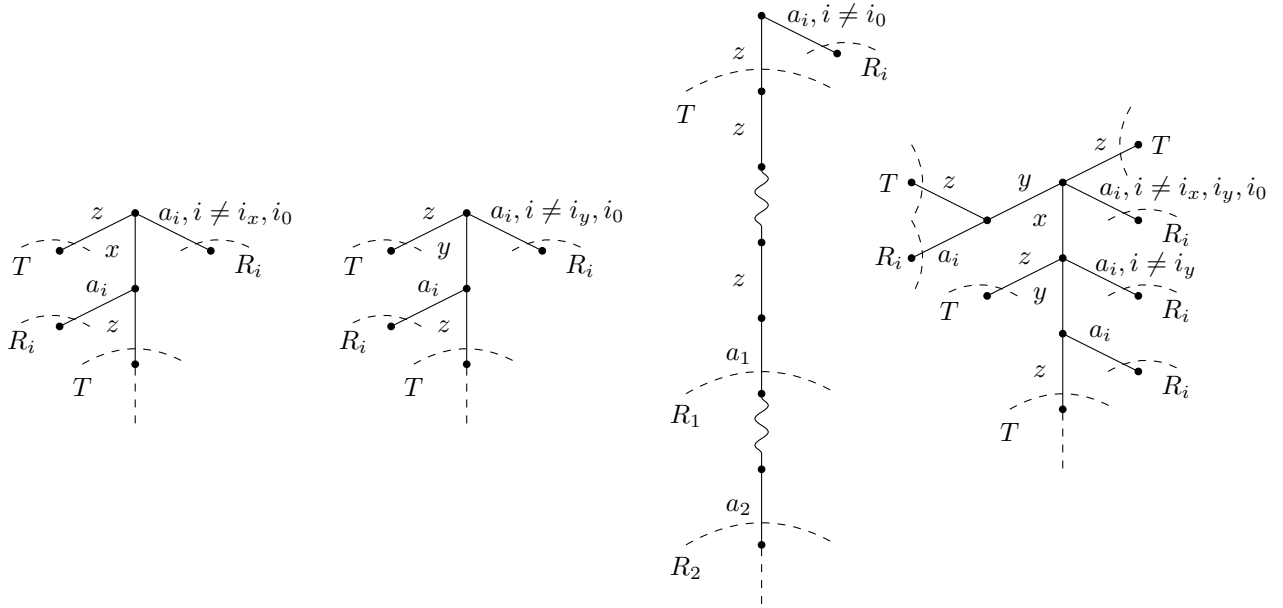


Figure 4: Schematic structure of the trees $\mathcal{T}(xzu)$, $\mathcal{T}(yzu)$, $\mathcal{T}(zu)$ and $\mathcal{T}(xyzu)$ when $x \neq z$, $y \neq z$ and $z \neq 0$.

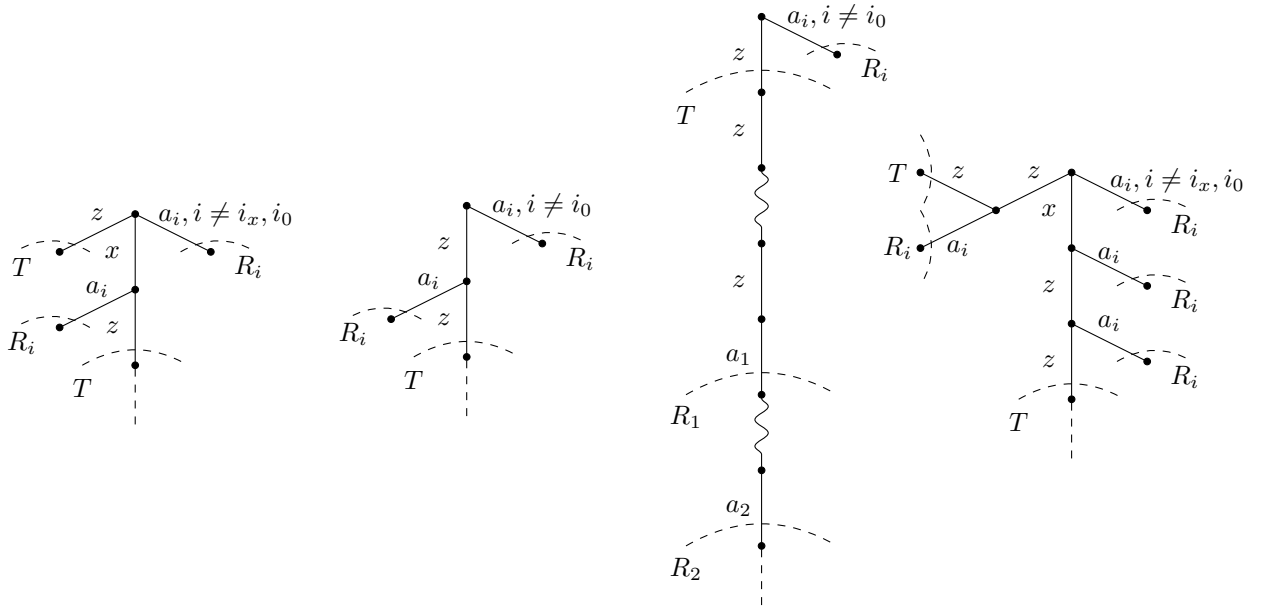


Figure 5: Schematic structure of the trees $\mathcal{T}(xzu)$, $\mathcal{T}(yzu)$, $\mathcal{T}(zu)$ and $\mathcal{T}(xyzu)$ when $x \neq z$, $y = z$ and $z \neq 0$.

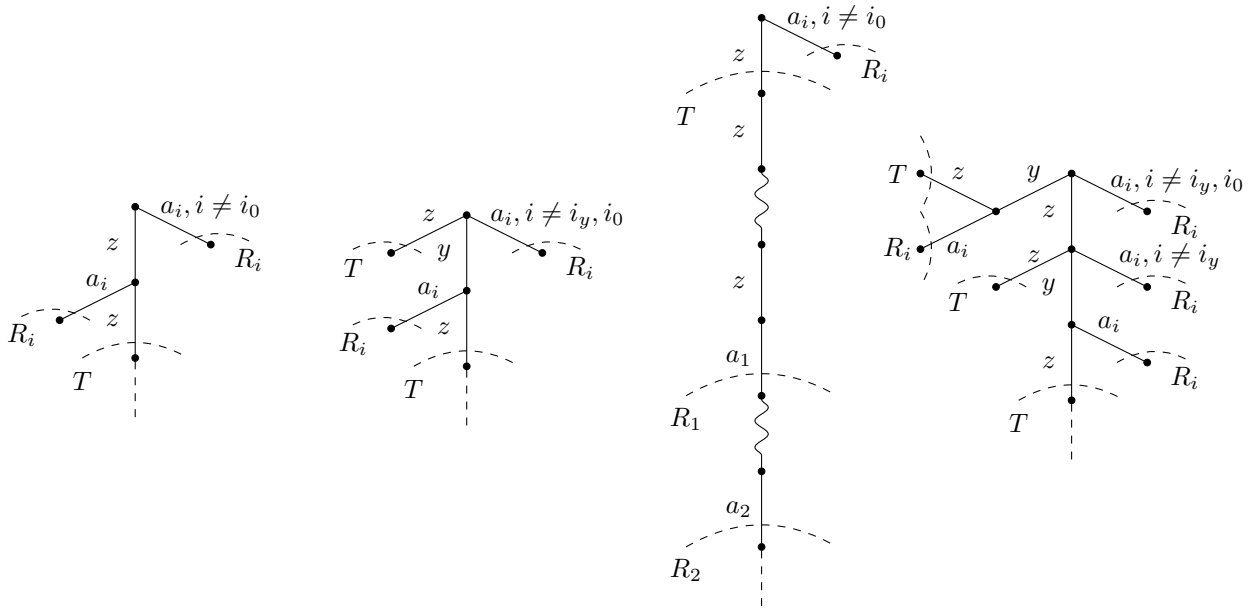


Figure 6: Schematic structure of the trees $\mathcal{T}(xzu)$, $\mathcal{T}(yzu)$, $\mathcal{T}(zu)$ and $\mathcal{T}(xyzu)$ when $x = z$, $y \neq z$ and $z \neq 0$.

Suppose that $x \neq z$ and $y = z$. The situation is depicted in Figure 5. The claimed formula holds since

$$\begin{aligned}
& (2 + 2\#T + 2 \sum_{\substack{1 \leq i \leq m \\ i \neq i_x, i_0}} \#R_i + \#R_{i_x} + \#R_{i_0}) \\
& + 2 \cdot (2 + \#T + 2 \sum_{\substack{1 \leq i \leq m \\ i \neq i_x, i_0}} \#R_i + 2\#R_{i_x} + \#R_{i_0}) \\
& - 2 \cdot (1 + \#T + \sum_{\substack{1 \leq i \leq m \\ i \neq i_x, i_0}} \#R_i + \#R_{i_x}) \\
& = 4 + 2\#T + 4 \sum_{\substack{1 \leq i \leq m \\ i \neq i_x, i_0}} \#R_i + 3\#R_{i_x} + 3\#R_{i_0}.
\end{aligned}$$

Suppose that $x = z$ and $y \neq z$. The situation is depicted in Figure 6. The claimed formula holds since

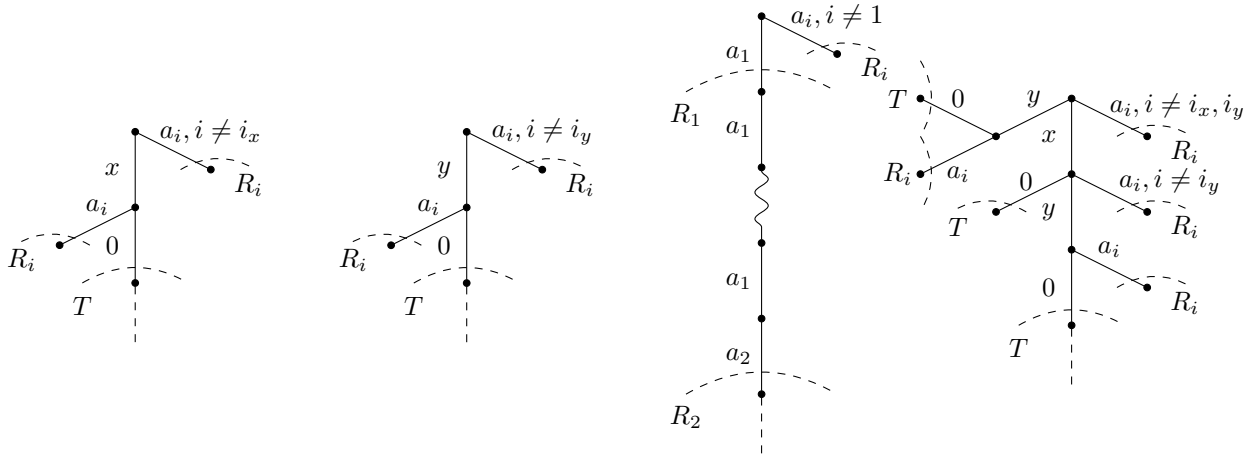


Figure 7: Schematic structure of the trees $\mathcal{T}(x0u)$, $\mathcal{T}(y0u)$, $\mathcal{T}(\text{rep}_b(\text{val}_b(u)))$ and $\mathcal{T}(xy0u)$.

$$\begin{aligned}
& (2 + \#T + 2 \sum_{\substack{1 \leq i \leq m \\ i \neq i_y, i_0}} \#R_i + 2\#R_{i_y} + \#R_{i_0}) \\
& + 2 \cdot (2 + 2\#T + 2 \sum_{\substack{1 \leq i \leq m \\ i \neq i_y, i_0}} \#R_i + \#R_{i_y} + \#R_{i_0}) \\
& - 2 \cdot (1 + \#T + \sum_{\substack{1 \leq i \leq m \\ i \neq i_y, i_0}} \#R_i + \#R_{i_y}) \\
& = 4 + 3\#T + 4 \sum_{\substack{1 \leq i \leq m \\ i \neq i_y, i_0}} \#R_i + 2\#R_{i_y} + 3\#R_{i_0}.
\end{aligned}$$

• As a second case, suppose that $z = 0$. Then, by convention, leading zeroes are not allowed in base- b expansions and we must prove that the following formula holds

$$\begin{aligned}
\# \left\{ v \in L_b \mid \binom{xy0u}{v} > 0 \right\} &= \# \left\{ v \in L_b \mid \binom{x0u}{v} > 0 \right\} + 2 \cdot \# \left\{ v \in L_b \mid \binom{y0u}{v} > 0 \right\} \\
&\quad - 2 \cdot \# \left\{ v \in L_b \mid \binom{\text{rep}_b(\text{val}_b(u))}{v} > 0 \right\}.
\end{aligned}$$

It is useful to note that $\text{rep}_b(\text{val}_b(\cdot)) : \{0, \dots, b-1\}^* \mapsto L_b$ plays a normalization role. It removes leading zeroes.

If $u = 0^n$, with $n \geq 0$, then $\text{rep}_b(\text{val}_b(u)) = \varepsilon$ and the tree $\mathcal{T}(\text{rep}_b(\text{val}_b(u)))$ has only one node. The trees $\mathcal{T}(x0u)$ and $\mathcal{T}(y0u)$ both have $n + 3$ nodes and the tree $\mathcal{T}(xy0u)$ has $3(n + 2) + 1$ nodes and the claimed formula holds.

Now suppose that u contains other letters than 0. We let a_1, \dots, a_m denote all the pairwise distinct letters of u different from 0. They are implicitly ordered with respect to their first appearance in u . If $x, y \in \{a_1, \dots, a_m\}$, we let $i_x, i_y \in \{1, \dots, m\}$ respectively denote the indices such that $a_{i_x} = x$ and $a_{i_y} = y$. For all $i \in \{1, \dots, m\}$, we let $u'_i a_i$ denote the prefix of $\text{rep}_b(\text{val}_b(u))$ that ends with the first occurrence of the letter a_i in $\text{rep}_b(\text{val}_b(u))$, and we let R_i denote the subtree of $\mathcal{T}(\text{rep}_b(\text{val}_b(u)))$ with root $u'_i a_i$.

The situation is depicted in Figure 7. Observe that the subtree T of $\mathcal{T}(y0u)$ with root $y0$ is equal to the subtree of $\mathcal{T}(x0u)$ with root $x0$ and to the subtree of $\mathcal{T}(xy0u)$ with root $xy0$. The claimed formula holds since

$$\begin{aligned}
& (2 + \#T + 2 \sum_{\substack{1 \leq i \leq m \\ i \neq i_x, i_y}} \#R_i + \#R_{i_x} + 2\#R_{i_y}) \\
& + 2 \cdot (2 + \#T + 2 \sum_{\substack{1 \leq i \leq m \\ i \neq i_x, i_y}} \#R_i + 2\#R_{i_x} + \#R_{i_y}) \\
& - 2 \cdot (1 + \sum_{\substack{1 \leq i \leq m \\ i \neq i_x, i_y}} \#R_i + \#R_{i_x} + \#R_{i_y}) \\
& = 4 + 3\#T + 4 \sum_{\substack{1 \leq i \leq m \\ i \neq i_x, i_y}} \#R_i + 3\#R_{i_x} + 2\#R_{i_y}.
\end{aligned}$$

□

Those five lemmas can be translated into recurrence relations satisfied by the sequence $(S_b(n))_{n \geq 0}$ using Definition 1.

Proof of Proposition 3. The first part is clear using Table 1. Let $x, y \in \{1, \dots, b-1\}$ with $x \neq y$. Proceed by induction on $\ell \geq 1$.

Let us first prove (2). If $\ell = 1$, then $r = 0$ and (2) follows from Table 1. Now suppose that $\ell \geq 2$ and assume that (2) holds for all $\ell' < \ell$. Let $r \in \{0, \dots, b^{\ell-1} - 1\}$, and let u be a word in $\{0, \dots, b-1\}^*$ such that $|u| \geq 1$ and $\text{rep}_b(xb^\ell + r) = x0u$. The proof is divided into two parts according to the first letter of u . If $u = 0u'$ with $u' \in \{0, \dots, b-1\}^*$, then

$$\begin{aligned}
S_b(xb^\ell + r) &= 2S_b(xb^{\ell-1} + r) - S_b(xb^{\ell-2} + r) && \text{(by Lemma 10)} \\
&= 2(S_b(xb^{\ell-2} + r) + S_b(r)) - S_b(xb^{\ell-2} + r) && \text{(by induction hypothesis)} \\
&= S_b(xb^{\ell-2} + r) + S_b(r) + S_b(r) \\
&= S_b(xb^{\ell-1} + r) + S_b(r), && \text{(by induction hypothesis)}
\end{aligned}$$

which proves (2). Now if $u = zu'$ with $z \in \{1, \dots, b-1\}$ and $u' \in \{0, \dots, b-1\}^*$, then (2) directly follows from Definition 1 and Lemma 12.

Let us prove (3). If $\ell = 1$, then $r = 0$ and (2) follows from Table 1. Now suppose that $\ell \geq 2$ and assume that (3) holds for all $\ell' < \ell$. Let $r \in \{0, \dots, b^{\ell-1} - 1\}$, and let u be a word in $\{0, \dots, b-1\}^*$ such that $|u| \geq 1$ and $\text{rep}_b(xb^\ell + xb^{\ell-1} + r) = xxu$. The proof is divided into two parts according to the first letter of u . If $u = 0u'$ with $u' \in \{0, \dots, b-1\}^*$, then

$$\begin{aligned}
S_b(xb^\ell + xb^{\ell-1} + r) &= S_b(xb^{\ell-1} + r) + S_b(xb^{\ell-2} + r) && \text{(by Lemma 11)} \\
&= S_b(xb^{\ell-2} + r) + S_b(r) + S_b(xb^{\ell-2} + r) && \text{(using (2))} \\
&= 2(S_b(xb^{\ell-2} + r) + S_b(r)) - S_b(r) \\
&= 2S_b(xb^{\ell-1} + r) - S_b(r), && \text{(using (2))}
\end{aligned}$$

which proves (3). Now if $u = zu'$ with $z \in \{1, \dots, b-1\}$ and $u' \in \{0, \dots, b-1\}^*$, then (3) directly follows from Definition 1 and Lemma 13.

Let us finally prove (4). If $\ell = 1$, then $r = 0$ and (2) follows from Table 1. Now suppose that $\ell \geq 2$ and assume that (4) holds for all $\ell' < \ell$. Let $r \in \{0, \dots, b^{\ell-1} - 1\}$, let z be a letter of $\{1, \dots, b-1\}$ and let u be a word in $\{0, \dots, b-1\}^*$ such that $\text{rep}_b(xb^\ell + yb^{\ell-1} + r) = xyz u$. Using Definition 1 and Lemma 14, we directly have that

$$S_b(xb^\ell + yb^{\ell-1} + r) = S_b(xb^{\ell-1} + r) + S_b(yb^{\ell-1} + r) - 2S_b(r),$$

which proves (4). □

3 Regularity of the sequence $(S_b(n))_{n \geq 0}$

The sequence $(S_2(n))_{n \geq 0}$ is shown to be 2-regular; see [16]. We recall that the b -kernel of a sequence $s = (s(n))_{n \geq 0}$ is the set

$$\mathcal{K}_b(s) = \{(s(b^i n + j))_{n \geq 0} \mid i \geq 0 \text{ and } 0 \leq j < b^i\}.$$

A sequence $s = (s(n))_{n \geq 0} \in \mathbb{Z}^{\mathbb{N}}$ is b -regular if there exists a finite number of sequences $(t_1(n))_{n \geq 0}, \dots, (t_\ell(n))_{n \geq 0}$ such that every sequence in the \mathbb{Z} -module $\langle \mathcal{K}_b(s) \rangle$ generated by the b -kernel $\mathcal{K}_b(s)$ is a \mathbb{Z} -linear combination of the t_r 's. In this section, we prove that the sequence $(S_b(n))_{n \geq 0}$ is b -regular. As a consequence, one can get matrices to compute $S_b(n)$ in a number of matrix multiplications proportional to $\log_b(n)$. To prove the b -regularity of the sequence $(S_b(n))_{n \geq 0}$ for any base b , we first need a lemma involving some matrix manipulations.

Lemma 15. *Let I and 0 respectively be the identity matrix of size $b^2 \times b^2$ and the zero matrix of size $b^2 \times b^2$. Let M_b be the block-matrix of size $b^3 \times b^3$*

$$M_b := \begin{pmatrix} I & I & 2I & \cdots & \cdots & \cdots & 2I \\ 2I & 3I & 3I & 4I & \cdots & \cdots & 4I \\ \vdots & \vdots & 4I & \ddots & \ddots & & \vdots \\ \vdots & \vdots & \vdots & \ddots & \ddots & \ddots & \vdots \\ \vdots & \vdots & \vdots & & \ddots & \ddots & 4I \\ \vdots & \vdots & \vdots & & & \ddots & 3I \\ 2I & 3I & 4I & \cdots & \cdots & \cdots & 4I \end{pmatrix}.$$

This matrix is invertible and its inverse is given by

$$M_b^{-1} := \begin{pmatrix} 3I & 2I & \cdots & \cdots & 2I & -(2b-3)I \\ -2I & 0 & \cdots & \cdots & 0 & I \\ 0 & -I & \ddots & & \vdots & \vdots \\ \vdots & \ddots & \ddots & \ddots & \vdots & \vdots \\ \vdots & & \ddots & \ddots & 0 & \vdots \\ 0 & \cdots & \cdots & 0 & -I & I \end{pmatrix}.$$

For the proof of the previous lemma, simply proceed to the multiplication of the two matrices. Using this lemma, we prove that the sequence $(S_b(n))_{n \geq 0}$ is b -regular.

Theorem 16. *For all $r \in \{0, \dots, b^2 - 1\}$, we have*

$$S_b(nb^2 + r) = a_r S_b(n) + \sum_{s=0}^{b-2} c_{r,s} S_b(nb + s) \quad \forall n \geq 0, \quad (5)$$

where the coefficients a_r and $c_{r,s}$ are unambiguously determined by the first few values $S_b(0), S_b(1), \dots, S_b(b^3 - 1)$ and given in Table 2, Table 3 and Table 4. In particular, the sequence $(S_b(n))_{n \geq 0}$ is b -regular. Moreover, a choice of generators for $\langle \mathcal{K}_b(s) \rangle$ is given by the b sequences $(S_b(n))_{n \geq 0}, (S_b(bn))_{n \geq 0}, (S_b(bn + 1))_{n \geq 0}, \dots, (S_b(bn + b - 2))_{n \geq 0}$.

Proof. We proceed by induction on $n \geq 0$. For the base case $n \in \{0, 1, \dots, b^2 - 1\}$, we first compute the coefficients a_r and $c_{r,s}$ using the values of $S_b(nb^2 + r)$ for $n \in \{0, \dots, b - 1\}$ and $r \in \{0, \dots, b^2 - 1\}$. Then we show that (5) also holds with these coefficients for $n \in \{b, \dots, b^2 - 1\}$.

$\text{rep}_b(r)$	ε	x	$b-1$	$x0$	$(b-1)0$	xx	$(b-1)(b-1)$	xy	$(b-1)x$	$x(b-1)$
a_r	-1	-2	$2b-3$	-2	$4b-4$	-1	$4b-3$	-2	$4b-4$	$2b-3$

Table 2: Values of a_r for $0 \leq r < b^2$ with $x, y \in \{1, \dots, b-2\}$ and $x \neq y$.

$\text{rep}_b(r)$	ε	x	$b-1$	$x0$	$(b-1)0$	xx	$(b-1)(b-1)$	xy	$(b-1)x$	$x(b-1)$
$c_{r,0}$	2	2	1	1	-1	0	-2	0	-2	-1

Table 3: Values of $c_{r,0}$ for $0 \leq r < b^2$ with $x, y \in \{1, \dots, b-2\}$ and $x \neq y$.

Base case. Let I denote the identity matrix of size $b^2 \times b^2$. The system of b^3 equations (5) when $n \in \{0, \dots, b-1\}$ and $r \in \{0, \dots, b^2-1\}$ can be written as $MX = V$ where the matrix $M \in \mathbb{Z}_{b^3}^{b^3}$ is equal to

$$\begin{pmatrix} S_b(0)I & S_b(0)I & S_b(1)I & S_b(2)I & \cdots & S_b(b-2)I \\ S_b(1)I & S_b(b)I & S_b(b+1)I & S_b(b+2)I & \cdots & S_b(2b-2)I \\ \vdots & \vdots & \vdots & \vdots & \ddots & \vdots \\ S_b(b-1)I & S_b(b(b-1))I & S_b(b(b-1)+1)I & S_b(b(b-1)+2)I & \cdots & S_b(b(b-1)+b-2)I \end{pmatrix}$$

and the vectors $X, V \in \mathbb{Z}^{b^3}$ are respectively given by

$$\begin{aligned} X^T &= (a_0 \quad \cdots \quad a_{b^2-1} \quad c_{0,0} \quad c_{1,0} \quad \cdots \quad c_{b^2-1,0} \quad \cdots \quad c_{0,b-2} \quad c_{1,b-2} \quad \cdots \quad c_{b^2-1,b-2}), \\ V^T &= (S_b(0) \quad S_b(1) \quad \cdots \quad S_b(b^3-1)). \end{aligned}$$

Observe that in the vector X , the coefficients $c_{r,s}$ are first sorted by s then by r . Using Table 1, the matrix M is equal to the matrix M_b of Lemma 15. By this lemma, the previous system has a unique solution given by $X = M_b^{-1}V$. Consequently, using Lemma 15, we have, for all $r \in \{0, \dots, b^2-1\}$ and all $s \in \{1, \dots, b-2\}$,

$$\begin{aligned} a_r &= 3S_b(r) + 2 \sum_{j=1}^{b-2} S_b(jb^2+r) - (2b-3)S_b((b-1)b^2+r), \\ c_{r,0} &= -2S_b(r) + S_b((b-1)b^2+r), \\ c_{r,s} &= -S_b(sb^2+r) + S_b((b-1)b^2+r). \end{aligned}$$

The values of the coefficients can then be computed using Table 1 and are stored in Table 2, Table 3 and Table 4.

For $n \in \{b, \dots, b^2-1\}$, the values of $S_b(nb^2+r)$ are given in Table 5, Table 6 and Table 7 according to whether $\text{rep}_b(n)$ is of the form $x0, xx$ or xy with $x \neq y$. The proof that (5) holds for each $n \in \{b, \dots, b^2-1\}$ only requires easy computations that are left to the reader.

$\text{rep}_b(r)$	ε	x		$b-1$	$x0$		$(b-1)0$	xx	
s	z	x	z	z	x	z	z	x	z
$c_{r,s}$	0	1	0	-1	2	0	-2	2	0

$\text{rep}_b(r)$	$(b-1)(b-1)$	xy		$x(b-1)$		$(b-1)x$	
s	z	x	y	x	z	x	z
$c_{r,s}$	-2	2	1	0	1	-1	-2

Table 4: Values of $c_{r,s}$ for $0 \leq r < b^2$ and $1 \leq s \leq b-2$ with $x, y, z \in \{1, \dots, b-2\}$ pairwise distinct.

$\text{rep}_b(r)$	ε	x	y	$x0$	$y0$	xx	yy	xy	yx	yz
$S_b(nb^2 + r)$	5	7	8	8	10	7	9	10	11	12

Table 5: Values of $S_b(nb^2 + r)$ for $b \leq n < b^2$ with $\text{rep}_b(n) = x0$ and $x, y, z \in \{1, \dots, b-1\}$ pairwise distinct.

$\text{rep}_b(r)$	ε	x	y	$x0$	$y0$	xx	yy	xy	yx	yz
$S_b(nb^2 + r)$	7	8	10	7	11	5	9	8	10	12

Table 6: Values of $S_b(nb^2 + r)$ for $b \leq n < b^2$ with $\text{rep}_b(n) = xx$ and $x, y, z \in \{1, \dots, b-1\}$ pairwise distinct.

Inductive step. Consider $n \geq b^2$ and suppose that the relation (5) holds for all $m < n$. Then $|\text{rep}_b(n)| \geq 3$. Like for the base case, we need to consider several cases according to the form of the base- b expansion of n . More precisely, we need to consider the following five forms, where $u \in \{0, \dots, b-1\}^*$, $x, y, z \in \{1, \dots, b-1\}$, $x \neq z$, and $t \in \{0, \dots, b-1\}$:

$$x00u \text{ or } xx0u \text{ or } x0yu \text{ or } xxyu \text{ or } xztu.$$

Let us focus on the first form of $\text{rep}_b(n)$ since the same reasoning can be applied for the other ones. Assume that $\text{rep}_b(n) = x00u$ where $x \in \{1, \dots, b-1\}$ and $u \in \{0, \dots, b-1\}^*$. For all $r \in \{0, \dots, b^2-1\}$, there exist $r_1, r_2 \in \{0, \dots, b-1\}$ such that $\text{val}_b(r_1r_2) = r$. We have

$$\begin{aligned}
S_b(nb^2 + r) &= S_b(\text{val}_b(x00ur_1r_2)) \\
&= 2S_b(\text{val}_b(x0ur_1r_2)) - S_b(\text{val}_b(xur_1r_2)) && \text{(by Lemma 10)} \\
&= a_r 2S_b(\text{val}_b(x0u)) + \sum_{s=0}^{b-2} c_{r,s} 2S_b(\text{val}_b(x0us)) \\
&\quad - a_r S_b(\text{val}_b(xu)) - \sum_{s=0}^{b-2} c_{r,s} S_b(\text{val}_b(xus)) && \text{(by induction hypothesis)} \\
&= a_r S_b(\text{val}_b(x00u)) + \sum_{s=0}^{b-2} c_{r,s} S_b(\text{val}_b(x00us)) && \text{(by Lemma 10)} \\
&= a_r S_b(n) + \sum_{s=0}^{b-2} c_{r,s} S_b(nb + s), && \text{(by Lemma 10)}
\end{aligned}$$

which proves (5).

b -regularity. From the first part of the proof, we directly deduce that the \mathbb{Z} -module $\langle \mathcal{K}_b(S_b) \rangle$ is generated by the $(b+1)$ sequences

$$(S_b(n))_{n \geq 0}, (S_b(bn))_{n \geq 0}, (S_b(bn+1))_{n \geq 0}, \dots, (S_b(bn+b-1))_{n \geq 0}.$$

We now show that we can reduce the number of generators. To that aim, we prove that

$$S_b(nb + b - 1) = (2b - 1)S_b(n) - \sum_{s=0}^{b-2} S_b(nb + s) \quad \forall n \geq 0. \quad (6)$$

We proceed by induction on $n \geq 0$. As a base case, the proof that (6) holds for each $n \in \{b, \dots, b^2-1\}$ only requires easy computations that are left to the reader (using Table 1). Now consider $n \geq b^2$ and suppose that the relation (6) holds for all $m < n$. Then $|\text{rep}_b(n)| \geq 3$. Mimicking the first induction step of this proof, we need to consider several cases according to the form of the base- b expansion of n . More precisely, we need to consider the following five forms, where $u \in \{0, \dots, b-1\}^*$, $x, y, z \in \{1, \dots, b-1\}$, $x \neq z$, and $t \in \{0, \dots, b-1\}$:

$$x00u \text{ or } xx0u \text{ or } x0yu \text{ or } xxyu \text{ or } xztu.$$

$\text{rep}_b(r)$	ε	x	y	z	$x0$	$y0$	$z0$	xx	yy	zz	xy	xz	yx	yz	zx	zy	zt
$S_b(nb^2 + r)$	10	13	12	14	13	11	15	10	8	12	12	14	11	12	15	14	16

Table 7: Values of $S_b(nb^2 + r)$ for $b \leq n < b^2$ with $\text{rep}_b(n) = xy$ and $x, y, z, t \in \{1, \dots, b-1\}$ pairwise distinct.

Let us focus on the first form of $\text{rep}_b(n)$ since the same reasoning can be applied for the other ones. Assume that $\text{rep}_b(n) = x00u$ where $x \in \{1, \dots, b-1\}$ and $u \in \{0, \dots, b-1\}^*$. We have

$$\begin{aligned}
S_b(nb + b - 1) &= S_b(\text{val}_b(x00u(b-1))) \\
&= 2S_b(\text{val}_b(x0u(b-1))) - S_b(\text{val}_b(xu(b-1))) && \text{(by Lemma 10)} \\
&= (2b-1)2S_b(\text{val}_b(x0u)) - \sum_{s=0}^{b-2} 2S_b(\text{val}_b(x0us)) \\
&\quad - (2b-1)S_b(\text{val}_b(xu)) + \sum_{s=0}^{b-2} S_b(\text{val}_b(xus)) && \text{(by induction hypothesis)} \\
&= (2b-1)S_b(\text{val}_b(x00u)) - \sum_{s=0}^{b-2} S_b(\text{val}_b(x00us)) && \text{(by Lemma 10)} \\
&= (2b-1)S_b(n) - \sum_{s=0}^{b-2} S_b(nb + s), && \text{(by Lemma 10)}
\end{aligned}$$

which proves (5).

The \mathbb{Z} -module $\langle \mathcal{K}_b(S_b) \rangle$ is thus generated by the b sequences

$$(S_b(n))_{n \geq 0}, (S_b(bn))_{n \geq 0}, (S_b(bn+1))_{n \geq 0}, \dots, (S_b(bn+b-2))_{n \geq 0}.$$

□

Example 17. Let $b = 2$. Using Table 2, Table 3 and Table 4, we find that $a_0 = -1$, $a_1 = 1$, $a_2 = 4$, $a_3 = 5$, $c_{0,0} = 2$, $c_{1,0} = 1$, $c_{2,0} = -1$ and $c_{3,0} = -2$. In this case, there are no $c_{r,s}$ with $s > 0$. Applying Theorem 16 and from (6), we get

$$\begin{aligned}
S_2(2n+1) &= 3S_2(n) - S_2(2n), \\
S_2(4n) &= -S_2(n) + 2S_2(2n), \\
S_2(4n+1) &= S_2(n) + S_2(2n), \\
S_2(4n+2) &= 4S_2(n) - S_2(2n), \\
S_2(4n+3) &= 5S_2(n) - 2S_2(2n)
\end{aligned}$$

for all $n \geq 0$. This result is a rewriting of [16, Theorem 21]. Observe that the third and the fifth identities are redundant: they follow from the other ones.

Example 18. Let $b = 3$. Using Table 2, Table 3 and Table 4, the values of the coefficients a_r , $c_{r,0}$ and $c_{r,1}$ can be found in Table 8. Applying Theorem 16 and from (6), we get

r	0	1	2	3	4	5	6	7	8
a_r	-1	-2	3	-2	-1	3	8	8	9
$c_{r,0}$	2	2	1	1	0	-1	-1	-2	-2
$c_{r,1}$	0	1	-1	2	2	1	-2	-1	-2

Table 8: The values of $a_r, c_{r,0}, c_{r,1}$ when $b = 3$ and $r \in \{0, \dots, 8\}$.

$$\begin{aligned}
S_3(3n+2) &= 5S_3(n) - S_3(3n) - S_3(3n+1), \\
S_3(9n) &= -S_3(n) + 2S_3(3n), \\
S_3(9n+1) &= -2S_3(n) + 2S_3(3n) + S_3(3n+1), \\
S_3(9n+2) &= 3S_3(n) + S_3(3n) - S_3(3n+1), \\
S_3(9n+3) &= -2S_3(n) + S_3(3n) + 2S_3(3n+1), \\
S_3(9n+4) &= -S_3(n) + 2S_3(3n+1), \\
S_3(9n+5) &= 3S_3(n) - S_3(3n) + S_3(3n+1), \\
S_3(9n+6) &= 8S_3(n) - S_3(3n) - 2S_3(3n+1), \\
S_3(9n+7) &= 8S_3(n) - 2S_3(3n) - S_3(3n+1), \\
S_3(9n+8) &= 9S_3(n) - 2S_3(3n) - 2S_3(3n+1)
\end{aligned}$$

for all $n \geq 0$. This result is a proof of [16, Conjecture 26]. Observe that the fourth, the seventh and the tenth identities are redundant.

Remark 19. Combining (5) and (6) yield $b^2 + 1$ identities to generate the \mathbb{Z} -module $\langle \mathcal{K}_b(S_b) \rangle$. However, as illustrated in Example 17 and Example 18, only $b^2 - b + 1$ identities are useful: the relations established for the sequences $(S_b(b^2n + br + b - 1))_{n \geq 0}$, with $r \in \{0, \dots, b - 1\}$, can be deduced from the other identities.

Remark 20. Using Theorem 16 and (6) and the set of b generators of the \mathbb{Z} -module $\langle \mathcal{K}_b(S_b) \rangle$ being

$$\{(S_b(n))_{n \geq 0}, (S_b(bn))_{n \geq 0}, (S_b(bn + 1))_{n \geq 0}, \dots, (S_b(bn + b - 2))_{n \geq 0}\},$$

we get matrices to compute $S_b(n)$ in a number of steps proportional to $\log_b(n)$. For all $n \geq 0$, let

$$V_b(n) = \begin{pmatrix} S_b(n) \\ S_b(bn) \\ S_b(bn + 1) \\ \vdots \\ S_b(bn + b - 2) \end{pmatrix} \in \mathbb{Z}^b.$$

Consider the matrix-valued morphism $\mu_b : \{0, 1, \dots, b - 1\}^* \rightarrow \mathbb{Z}_b^b$ defined, for all $s \in \{0, \dots, b - 2\}$, by

$$\mu_b(s) = \begin{pmatrix} 0 & 0 & \cdots & 0 & 1 & 0 & \cdots & 0 \\ a_{bs} & c_{bs,0} & \cdots & c_{bs,s-1} & c_{bs,s} & c_{bs,s+1} & \cdots & c_{bs,b-2} \\ a_{bs+1} & c_{bs+1,0} & \cdots & c_{bs+1,s-1} & c_{bs+1,s} & c_{bs+1,s+1} & \cdots & c_{bs+1,b-2} \\ \vdots & \vdots & \cdots & \vdots & \vdots & \vdots & \cdots & \vdots \\ a_{bs+b-2} & c_{bs+b-2,0} & \cdots & c_{bs+b-2,s-1} & c_{bs+b-2,s} & c_{bs+b-2,s+1} & \cdots & c_{bs+b-2,b-2} \end{pmatrix}$$

and

$$\mu_b(b - 1) = \begin{pmatrix} (2b - 1) & -1 & -1 & \cdots & -1 \\ a_{b(b-1)} & c_{b(b-1),0} & c_{b(b-1),1} & \cdots & c_{b(b-1),b-2} \\ a_{b(b-1)+1} & c_{b(b-1)+1,0} & c_{b(b-1)+1,1} & \cdots & c_{b(b-1)+1,b-2} \\ \vdots & \vdots & \vdots & \cdots & \vdots \\ a_{b(b-1)+b-2} & c_{b(b-1)+b-2,0} & c_{b(b-1)+b-2,1} & \cdots & c_{b(b-1)+b-2,b-2} \end{pmatrix}.$$

Observe that the number of generators explains the size of the matrices above. For each $s \in \{0, \dots, b - 2\}$, exactly $b - 1$ identities from Theorem 16 are used to define the matrix $\mu_b(s)$. If $s, s' \in \{0, \dots, b - 2\}$ are such that $s \neq s'$, then the relations used to define the matrices $\mu_b(s)$ and $\mu_b(s')$ are pairwise distinct. Finally, the first row of the matrix $\mu_b(b - 1)$ is (6) and the other rows are $b - 1$ identities from Theorem 16, which are distinct from the previous relations. Consequently, $(b - 1)(b - 1) + b$ identities are used, which corroborates Remark 19.

Using the definition of the morphism μ , we can show that $V_b(bn + s) = \mu_b(s)V_b(n)$ for all $s \in \{0, \dots, b - 1\}$ and $n \geq 0$. Consequently, if $\text{rep}_b(n) = n_k \cdots n_0$, then

$$S_b(n) = (1 \ 0 \ \cdots \ 0) \mu_b(n_0) \cdots \mu_b(n_k) V_b(0).$$

For example, when $b = 2$, the matrices $\mu_2(0)$ and $\mu_2(1)$ are those given in [16, Corollary 22]. When $b = 3$, we get

$$\mu_3(0) = \begin{pmatrix} 0 & 1 & 0 \\ -1 & 2 & 0 \\ -2 & 2 & 1 \end{pmatrix}, \quad \mu_3(1) = \begin{pmatrix} 0 & 0 & 1 \\ -2 & 1 & 2 \\ -1 & 0 & 2 \end{pmatrix}, \quad \mu_3(2) = \begin{pmatrix} 5 & -1 & -1 \\ 8 & -1 & -2 \\ 8 & -2 & -1 \end{pmatrix}.$$

The class of b -synchronized sequences is intermediate between the classes of b -automatic sequences and b -regular sequences. These sequences were first introduced in [9].

Proposition 21. *The sequence $(S_b(n))_{n \geq 0}$ is not b -synchronized.*

Proof. The proof is exactly the same as [16, Proposition 24]. \square

To conclude this section, the following result proves that the sequence $(S_b(n))_{n \geq 0}$ has a partial palindromic structure as the sequence $(S_2(n))_{n \geq 0}$; see [16]. For instance, the sequence $(S_3(n))_{n \geq 0}$ is depicted in Figure 8 inside the interval $[2 \cdot 3^4, 3^5]$.

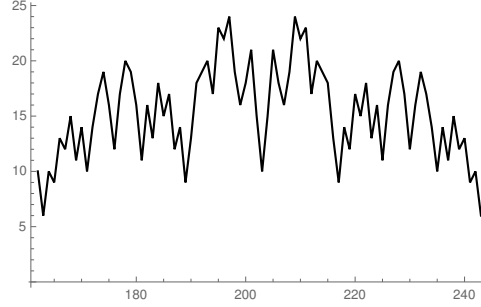


Figure 8: The sequence $(S_3(n))_{n \geq 0}$ inside the interval $[2 \cdot 3^4, 3^5]$.

Proposition 22. *Let u be a word in $\{0, 1, \dots, b-1\}^*$. Define \bar{u} by replacing in u every letter $a \in \{0, 1, \dots, b-1\}$ by the letter $(b-1) - a \in \{0, 1, \dots, b-1\}$. Then*

$$\# \left\{ v \in L_b \mid \binom{(b-1)u}{v} > 0 \right\} = \# \left\{ v \in L_b \mid \binom{(b-1)\bar{u}}{v} > 0 \right\}.$$

In particular, there exists a palindromic substructure inside of the sequence $(S_b(n))_{n \geq 0}$, i.e., for all $\ell \geq 1$ and $0 \leq r < b^\ell$,

$$S_b((b-1) \cdot b^\ell + r) = S_b((b-1) \cdot b^\ell + b^\ell - r - 1).$$

Proof. The trees $\mathcal{T}((b-1)u)$ and $\mathcal{T}((b-1)\bar{u})$ are isomorphic. Indeed, on the one hand, each node of the form $(b-1)x$ in the first tree corresponds to the node $(b-1)\bar{x}$ in the second one and conversely. On the other hand, if there exist letters $a \in \{1, \dots, b-2\}$ in the word $(b-1)u$, the position of the first letter a in the word $(b-1)u$ is equal to the position of the first letter $(b-1) - a$ in the word $(b-1)\bar{u}$ and conversely. Consequently, the node of the form ax in the first tree corresponds to the node of the form $((b-1) - a)\bar{x}$ in the second tree and conversely.

For the special case, note that for every word z of length ℓ , there exists $r \in \{0, \dots, b^\ell - 1\}$ such that $\text{rep}_b((b-1) \cdot b^\ell + r) = (b-1)z$ and

$$\text{val}_b(\bar{z}) = b^\ell - 1 - r \in \{0, \dots, b^\ell - 1\}.$$

Hence, $(b-1)\bar{z} = \text{rep}_b((b-1) \cdot b^\ell + b^\ell - 1 - r)$. Using (1), we obtain the desired result. \square

4 Asymptotics of the summatory function $(A_b(n))_{n \geq 0}$

In this section, we consider the summatory function $(A_b(n))_{n \geq 0}$ of the sequence $(S_b(n))_{n \geq 0}$; see Definition 1. The aim of this section is to apply the method introduced in [17] to obtain the asymptotic behavior of $(A_b(n))_{n \geq 0}$. As an easy consequence of the b -regularity of $(S_b(n))_{n \geq 0}$, we have the following result.

Proposition 23. *For all $b \geq 2$, the sequence $(A_b(n))_{n \geq 0}$ is b -regular.*

Proof. This is a direct consequence of Theorem 16 and of the fact that the summatory function of a b -regular sequence is also b -regular; see [2, Theorem 16.4.1]. \square

From a linear representation with matrices of size $d \times d$ associated with a b -regular sequence, one can derive a linear representation with matrices of size $2d \times 2d$ associated with its summatory function; see [10, Lemma 1]. Consequently, using Remark 20, one can obtain a linear representation with matrices of size $2b \times 2b$ for the summatory function $(A_b(n))_{n \geq 0}$. The goal is to decompose $(A_b(n))_{n \geq 0}$ into linear combinations of powers of $(2b - 1)$. We need the following two lemmas.

Lemma 24. *For all $\ell \geq 0$ and all $x \in \{1, \dots, b - 1\}$, we have*

$$A_b(xb^\ell) = (2x - 1) \cdot (2b - 1)^\ell.$$

Proof. We proceed by induction on $\ell \geq 0$. If $\ell = 0$ and $x \in \{1, \dots, b - 1\}$, then using Table 1, we have

$$A_b(x) = S_b(0) + \sum_{j=1}^{x-1} S_b(j) = 2x - 1.$$

If $\ell = 1$ and $x \in \{1, \dots, b - 1\}$, then we have

$$A_b(xb) = A_b(b) + \sum_{y=1}^{x-1} \sum_{j=0}^{b-1} S_b(yb + j).$$

Using Table 1, we get $A_b(xb) = (2x - 1)(2b - 1)$.

Now suppose that $\ell \geq 1$ and assume that the result holds for all $\ell' \leq \ell$. To prove the result, we again proceed by induction on $x \in \{1, \dots, b - 1\}$. When $x = 1$, we must show that $A_b(b^{\ell+1}) = (2b - 1)^{\ell+1}$. We have

$$A_b(b^{\ell+1}) = A_b(b^\ell) + \sum_{y=1}^{b-1} \sum_{j=0}^{b^\ell-1} S_b(yb^\ell + j).$$

By decomposing the sum into three parts accordingly to Proposition 3, we get

$$\begin{aligned} A_b(b^{\ell+1}) &= A_b(b^\ell) + \sum_{y=1}^{b-1} \sum_{j=0}^{b^\ell-1} S_b(yb^\ell + j) + \sum_{y=1}^{b-1} \sum_{j=0}^{b^{\ell-1}-1} S_b(yb^\ell + yb^{\ell-1} + j) \\ &+ \sum_{y=1}^{b-1} \sum_{\substack{1 \leq z \leq b-1 \\ z \neq y}} \sum_{j=0}^{b^{\ell-1}-1} S_b(yb^\ell + zb^{\ell-1} + j), \end{aligned}$$

and, using Proposition 3,

$$\begin{aligned} A_b(b^{\ell+1}) &= A_b(b^\ell) \\ &+ \sum_{y=1}^{b-1} \sum_{j=0}^{b^{\ell-1}-1} (S_b(yb^{\ell-1} + j) + S_b(j)) \end{aligned} \tag{7}$$

$$+ \sum_{y=1}^{b-1} \sum_{j=0}^{b^{\ell-1}-1} (2S_b(yb^{\ell-1} + j) - S_b(j)) \tag{8}$$

$$+ \sum_{y=1}^{b-1} \sum_{\substack{1 \leq z \leq b-1 \\ z \neq y}} \sum_{j=0}^{b^{\ell-1}-1} (S_b(yb^{\ell-1} + j) + 2S_b(zb^{\ell-1} + j) - 2S_b(j)). \tag{9}$$

By observing that for all y ,

$$\sum_{j=0}^{b^{\ell-1}-1} S_b(yb^{\ell-1} + j) = A_b((y+1)b^{\ell-1}) - A_b(yb^{\ell-1}) \quad \text{and} \quad \sum_{j=0}^{b^{\ell-1}-1} S_b(j) = A_b(b^{\ell-1}), \quad (10)$$

and that

$$\sum_{y=1}^{b-1} (A_b((y+1)b^{\ell-1}) - A_b(yb^{\ell-1})) = A_b(b^\ell) - A_b(b^{\ell-1}), \quad (11)$$

we obtain

$$(7) = A_b(b^\ell) + (b-2)A_b(b^{\ell-1}),$$

$$(8) = 2A_b(b^\ell) - (b+1)A_b(b^{\ell-1}),$$

$$(9) = 3(b-2)(A_b(b^\ell) - A_b(b^{\ell-1})) - 2(b-1)(b-2)A_b(b^{\ell-1}) = 3(b-2)A_b(b^\ell) - (b-2)(2b+1)A_b(b^{\ell-1}),$$

and finally

$$A_b(b^{\ell+1}) = (3b-2)A_b(b^\ell) - (2b^2-3b+1)A_b(b^{\ell-1}).$$

Using the induction hypothesis, we obtain

$$A_b(b^{\ell+1}) = (3b-2)(2b-1)^\ell - (2b^2-3b+1)(2b-1)^{\ell-1} = (2b-1)^{\ell+1},$$

which ends the case where $x = 1$.

Now suppose that $x \in \{2, \dots, b-1\}$ and assume that the result holds for all $x' < x$. The proof follows the same lines as in the case $x = 1$ with the difference that we decompose the sum into

$$\begin{aligned} A_b(xb^{\ell+1}) &= A_b((x-1)b^{\ell+1}) + \sum_{j=0}^{b^{\ell+1}-1} S_b((x-1)b^{\ell+1} + j) \\ &= A_b((x-1)b^{\ell+1}) + \sum_{j=0}^{b^\ell-1} S_b((x-1)b^{\ell+1} + j) + \sum_{j=0}^{b^\ell-1} S_b((x-1)b^{\ell+1} + (x-1)b^\ell + j) \\ &\quad + \sum_{\substack{1 \leq y \leq b-1 \\ y \neq x-1}} \sum_{j=0}^{b^\ell-1} S_b((x-1)b^{\ell+1} + yb^\ell + j). \end{aligned}$$

Applying Proposition 3 and using (10) and (11) leads to the equality

$$A_b(xb^{\ell+1}) = A_b((x-1)b^{\ell+1}) + (b-1)A_b(xb^\ell) - (b-1)A_b((x-1)b^\ell) + 2A_b(b^{\ell+1}) - 2(b-1)A_b(b^\ell).$$

The induction hypothesis ends the computation. \square

Lemma 25. For all $\ell \geq 1$ and all $x, y \in \{1, \dots, b-1\}$, we have

$$A_b(xb^\ell + yb^{\ell-1}) = \begin{cases} (4xb - 2x + 4y - 2b) \cdot (2b-1)^{\ell-1}, & \text{if } y \leq x; \\ (4xb - 2x + 4y - 2b - 1) \cdot (2b-1)^{\ell-1}, & \text{if } y > x. \end{cases}$$

Proof. The proof of this lemma is similar to the proof of Lemma 24 so we only proof the formula for $A_b(xb^\ell + yb^{\ell-1})$, the other being similarly handled. We proceed by induction on $\ell \geq 1$. If $\ell = 1$, the result follows from Table 1. Assume that $\ell \geq 2$ and that the formulas hold for all $\ell' < \ell$. We have

$$A_b(xb^\ell + yb^{\ell-1}) = A_b(xb^\ell) + \sum_{j=0}^{b^{\ell-1}-1} S_b(xb^\ell + j) + \sum_{y=1}^{x-1} \sum_{j=0}^{b^{\ell-1}-1} S_b(xb^\ell + yb^{\ell-1} + j).$$

Applying Proposition 3 and using (10) and (11) leads to the equality

$$A_b(xb^\ell + xb^{\ell-1}) = A_b(xb^\ell) + xA_b((x+1)b^{\ell-1}) + (2-x)A_b(xb^{\ell-1}) + (1-2x)A_b(b^{\ell-1}).$$

Using Lemma 24 completes the computation. \square

Lemma 24 and Lemma 25 give rise to recurrence relations satisfied by the summatory function $(A_b(n))_{n \geq 0}$ as stated below. This is a key result that permits us to introduce $(2b-1)$ -decompositions (Definition 28 below) of the summatory function $(A_b(n))_{n \geq 0}$ and allows us to easily deduce Theorem 30; see [17] for similar results in base 2.

Proposition 26. *For all $x, y \in \{1, \dots, b-1\}$ with $x \neq y$, all $\ell \geq 1$ and all $r \in \{0, \dots, b^{\ell-1}\}$,*

$$A_b(xb^\ell + r) = (2b-2) \cdot (2x-1) \cdot (2b-1)^{\ell-1} + A_b(xb^{\ell-1} + r) + A_b(r); \quad (12)$$

$$A_b(xb^\ell + xb^{\ell-1} + r) = (4xb-2x-2b+2) \cdot (2b-1)^{\ell-1} + 2A_b(xb^{\ell-1} + r) - A_b(r); \quad (13)$$

$$A_b(xb^\ell + yb^{\ell-1} + r) = \begin{cases} (4xb-4x-2b+3) \cdot (2b-1)^{\ell-1} + A_b(xb^{\ell-1} + r) \\ \quad + 2A_b(yb^{\ell-1} + r) - 2A_b(r), & \text{if } y < x; \\ (4xb-4x-2b+2) \cdot (2b-1)^{\ell-1} + A_b(xb^{\ell-1} + r) \\ \quad + 2A_b(yb^{\ell-1} + r) - 2A_b(r), & \text{if } y > x. \end{cases} \quad (14)$$

Proof. We first prove (12). Let $x \in \{1, \dots, b-1\}$, $\ell \geq 1$ and $r \in \{0, \dots, b^{\ell-1}\}$. If $r = 0$, then (12) holds using Lemma 24. Now suppose that $r \in \{1, \dots, b^{\ell-1}\}$. Applying successively Proposition 3 and Lemma 24, we have

$$\begin{aligned} A_b(xb^\ell + r) &= A_b(xb^\ell) + \sum_{j=0}^{r-1} S_b(xb^\ell + j) \\ &= A_b(xb^\ell) + \sum_{j=0}^{r-1} (S_b(xb^{\ell-1} + j) + S_b(j)) \\ &= A_b(xb^\ell) + (A_b(xb^{\ell-1} + r) - A_b(xb^{\ell-1})) + A_b(r) \\ &= (2b-2)(2x-1)(2b-1)^{\ell-1} + A_b(xb^{\ell-1} + r) + A_b(r), \end{aligned}$$

which proves (12).

The proof of (13) and (14) are similar, thus we only prove (13). Let $x \in \{1, \dots, b-1\}$, $\ell \geq 1$ and $r \in \{0, \dots, b^{\ell-1}\}$. If $r = 0$, then (13) holds using Lemma 25. Now suppose that $r \in \{1, \dots, b^{\ell-1}\}$. Applying Proposition 3, we have

$$\begin{aligned} A_b(xb^\ell + xb^{\ell-1} + r) &= A_b(xb^\ell + xb^{\ell-1}) + \sum_{j=0}^{r-1} S_b(xb^\ell + xb^{\ell-1} + j) \\ &= A_b(xb^\ell + xb^{\ell-1}) + \sum_{j=0}^{r-1} (2S_b(xb^{\ell-1} + j) - S_b(j)) \\ &= A_b(xb^\ell + xb^{\ell-1}) + 2(A_b(xb^{\ell-1} + r) - A_b(xb^{\ell-1})) - A_b(r). \end{aligned}$$

Using Lemma 24 and Lemma 25, we get

$$\begin{aligned} A_b(xb^\ell + xb^{\ell-1} + r) &= (4xb+2x-2b)(2b-1)^{\ell-1} - 2(2x-1)(2b-1)^{\ell-1} + 2A_b(xb^{\ell-1} + r) - A_b(r) \\ &= (4xb-2x-2b+2)(2b-1)^{\ell-1} + 2A_b(xb^{\ell-1} + r) - A_b(r), \end{aligned}$$

which proves (13). \square

The following corollary was conjectured in [17].

Corollary 27. *For all $n \geq 0$, we have $A_b(nb) = (2b - 1)A_b(n)$.*

Proof. Let us proceed by induction on $n \geq 0$. It is easy to check by hand that the result holds for $n \in \{0, \dots, b-1\}$. Thus consider $n \geq b$ and suppose that the result holds for all $n' < n$. The reasoning is divided into three cases according to the form of the base- b expansion of n . As a first case, we write $n = xb^\ell + r$ with $x \in \{1, \dots, b-1\}$, $\ell \geq 1$ and $0 \leq r < b^{\ell-1}$. By Proposition 26, we have

$$A_b(nb) - (2b - 1)A_b(n) = (2b - 2) \cdot (2x - 1) \cdot (2b - 1)^\ell + A_b(xb^\ell + br) + A_b(br) - (2b - 2) \cdot (2x - 1) \cdot (2b - 1)^\ell - (2b - 1)A_b(xb^{\ell-1} + r) - (2b - 1)A_b(r)$$

We conclude this case by using the induction hypothesis. The other cases can be handled using the same technique. \square

Using Proposition 26, we can define $(2b - 1)$ -decompositions as follows.

Definition 28. Let $n \geq b$. Applying iteratively Proposition 26 provides a unique decomposition of the form

$$A_b(n) = \sum_{i=0}^{\ell_b(n)} d_i(n) (2b - 1)^{\ell_b(n)-i}$$

where $d_i(n)$ are integers, $d_0(n) \neq 0$ and $\ell_b(n)$ stands for $\lfloor \log_b n \rfloor - 1$. We say that the word

$$d_0(n) \cdots d_{\ell_b(n)}(n)$$

is the $(2b - 1)$ -decomposition of $A_b(n)$. For the sake of clarity, we also write $(d_0(n), \dots, d_{\ell_b(n)}(n))$. Also notice that the notion of $(2b - 1)$ -decomposition is only valid for integers in the sequence $(A_b(n))_{n \geq 0}$.

Example 29. Let $b = 3$. Let us compute the 5-decomposition of $A_3(150) = 1665$. We have $\text{rep}_3(150) = 12120$ and $\ell_3(150) = 3$. Applying once Proposition 26 leads to

$$A_3(150) = A_3(3^4 + 2 \cdot 3^3 + 15) = 4 \cdot 5^3 + A_3(3^3 + 15) + 2A_3(2 \cdot 3^3 + 15) - 2A_3(15). \quad (15)$$

Applying again Proposition 26, we get

$$\begin{aligned} A_3(3^3 + 15) &= A_3(3^3 + 3^2 + 6) = 6 \cdot 3^2 + 2A_3(3^2 + 6) - A_3(6), \\ A_3(2 \cdot 3^3 + 15) &= A_3(2 \cdot 3^3 + 3^2 + 6) = 13 \cdot 3^2 + A_3(2 \cdot 3^2 + 6) + 2A_3(3^2 + 6) - 2A_3(6), \\ A_3(15) &= A_3(3^2 + 2 \cdot 3^1) = 4 \cdot 5^1 + A_3(3^1) + 2A_3(2 \cdot 3^1) - 2A_3(0). \end{aligned}$$

Using Proposition 26, we find

$$\begin{aligned} A_3(3^2 + 6) &= A_3(3^2 + 2 \cdot 3^1) = 4 \cdot 5^1 + A_3(3^1) + 2A_3(2 \cdot 3^1) - 2A_3(0), \\ A_3(2 \cdot 3^2 + 6) &= A_3(2 \cdot 3^2 + 2 \cdot 3^1) = 16 \cdot 5^1 + 2A_3(2 \cdot 3^1) - A_3(0), \\ A_3(6) &= A_3(2 \cdot 3^1) = 12 \cdot 5^0 + A_3(2 \cdot 3^0) + A_3(0) = 15 \cdot 5^0. \end{aligned}$$

Using Lemma 24, we have $A_3(3^1) = 5^1$ and $A_3(2 \cdot 3^1) = 3 \cdot 5^1$. Plugging all those values together in (15), we finally have

$$A_3(150) = 4 \cdot 5^3 + 32 \cdot 5^2 + 82 \cdot 5^1 - 45 \cdot 5^0.$$

The 5-decomposition of $A_3(150)$ is thus $(4, 32, 82, -45)$.

The proof of the next result follows the same lines as the proof of [17, Theorem 1]. Therefore we only sketch it.

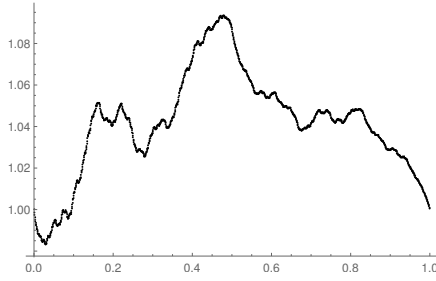


Figure 9: The function \mathcal{H}_3 over one period.

Theorem 30. *There exists a continuous and periodic function \mathcal{H}_b of period 1 such that, for all large enough n ,*

$$A_b(n) = (2b - 1)^{\log_b n} \mathcal{H}_b(\log_b n).$$

As an example, when $b = 3$, the function \mathcal{H}_3 is depicted in Figure 9 over one period.

Sketch of the proof of Theorem 30. Let us start by defining the function \mathcal{H}_b . Given any integer $n \geq 1$, we let ϕ_n denote the function

$$\phi_n(\alpha) = \frac{A_b(e_n(\alpha))}{(2b - 1)^{\log_b(e_n(\alpha))}}, \quad \alpha \in [0, 1)$$

where $e_n(\alpha) = b^{n+1} + b \lfloor b^n \alpha \rfloor + 1$. With a proof analogous to the one of [17, Proposition 20], the sequence of functions $(\phi_n)_{n \geq 1}$ uniformly converges to a function Φ_b . As in [17, Theorem 5], this function is continuous on $[0, 1]$ and such that $\Phi_b(0) = \Phi_b(1) = 1$. Furthermore, it satisfies

$$A_b(b^k + r) = (2b - 1)^{\log_b(b^k + r)} \Phi_b\left(\frac{r}{b^k}\right) \quad k \geq 1, 0 \leq r < b^k;$$

see [17, Lemma 24]. Using Corollary 27, we get that, for all $n = b^j(b^k + r)$, $j, k \geq 0$ and $r \in \{0, \dots, b^k - 1\}$,

$$A_b(n) = (2b - 1)^{\log_b(n)} \Phi_b\left(\frac{r}{b^k}\right).$$

The function \mathcal{H}_b is defined by $\mathcal{H}_b(x) = \Phi_b(b^{\{x\}} - 1)$ for all real x ($\{\cdot\}$ stands for the fractional part). \square

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