

# A Combinatorial Problem Solved by a Meta-Fibonacci Recurrence Relation

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## Abstract

We present a natural, combinatorial problem whose solution is given by the meta-Fibonacci recurrence relation  $a(n) = \sum_{i=1}^p a(n-i+1) - a(n-i)$ , where  $p$  is prime. This combinatorial problem is less general than those given in [3] and [4], but it has the advantage of having a simpler statement.

## 1 Introduction

Let  $M$  be a matrix with entries in  $\mathbb{Z}_2$ , such that every column contains at least one 1. We want to pick a subset of the rows such that when they are added together modulo 2, their sum  $\vec{s}$  has as many 1's as possible. If  $M$  has  $n$  columns, what is the largest number of 1's we can guarantee  $\vec{s}$  to have? For example, if  $n = 5$ , we can always find a set of rows whose sum  $\vec{s}$  contains at least four 1's. Let  $\lambda(n)$  denote the largest number of 1's  $\vec{s}$  can be guaranteed to have for any  $M$  with  $n$  nonzero columns. We will show that  $\lambda(n)$  satisfies the recurrence relation

$$\lambda(n) = \lambda(n - \lambda(n - 1)) + \lambda(n - 1 - \lambda(n - 2)). \quad (1)$$

More generally, for  $p$  prime, let  $\vec{v} = (v_1, \dots, v_n)$  satisfy  $v_i \in \mathbb{F}_p$  for  $1 \leq i \leq n$ . Let  $\text{supp}(\vec{v}) = \{i \in [n] : v_i \neq 0\}$  and let  $\|\vec{v}\| = |\text{supp}(\vec{v})|$ , i.e.,  $\|\vec{v}\|$  is the number of nonzero terms in  $\vec{v}$ . Let  $M$  be an  $m \times n$  matrix whose entries are in  $\mathbb{F}_p$ . Let  $\mathbf{row}(M)$  be the row space of  $M$ , i.e., the set of all linear combinations of the row vectors of  $M$  over the field  $\mathbb{F}_p$ . Let  $c(M)$  denote the *capacity* of  $M$ , which we define as follows,

$$c(M) = \max_{\vec{v} \in \mathbf{row}(M)} \|\vec{v}\|.$$

For each integer  $n \geq 1$ , let  $\lambda_p(n)$  be the minimum possible capacity of an  $\mathbb{F}_p$ -matrix consisting of  $n$  *nonzero* columns (i.e., no column equals  $\vec{0}$ ). Restated, let

$$\mathcal{M}_n^* = \{M \in \mathbb{F}_p^{m \times n} : 1 \leq m \leq p^n \text{ and no column of } M \text{ equals } \vec{0}\},$$

then

$$\lambda_p(n) = \min_{M \in \mathcal{M}_n^*} c(M).$$

We will see that  $\lambda_p$  satisfies the recurrence relation

$$\lambda_p(n) = \sum_{i=1}^p \lambda_p(n-i+1 - \lambda_p(n-i)). \quad (2)$$

This type of recurrence relation is called a meta-Fibonacci relation.

Meta-Fibonacci sequences have been studied by various authors, dating at least as far back as 1985, when Hofstadter [2] apparently coined the term “meta-Fibonacci.” These are integer sequences defined by “nested, Fibonacci-like” recurrence relations, such as relation (1), which was studied by Conolly [1], and (2). Generalizations of (2) were shown in [3] and [4] to be solutions to certain combinatorial problems involving  $k$ -ary infinite trees, and compositions of integers. The “matrix capacity” problem described above is a different combinatorial problem whose solution is also given by relation (2). This combinatorial problem is “natural” in the sense that it arose while the first named author was working on a problem in spatial graph theory. It was only later that we learned (through the OEIS [A046699](https://oeis.org/A046699)) that it can be characterized as a meta-Fibonacci sequence.

## 2 Main Result

We begin with a lemma which allows us to produce a lower bound on  $\lambda_p(n)$ . For the remainder of this paper, instead of writing  $\lambda_p$ , we will simply write  $\lambda$ . For a matrix  $M$ , let  $\mathbf{row}^*(M) = \mathbf{row}(M) - \{\vec{0}\}$ .

**Lemma 1.** *Let  $M$  be an  $\mathbb{F}_p$ -matrix with  $n$  nonzero columns, i.e.,  $M \in \mathcal{M}_n^*$ . Let  $\vec{v} \in \mathbf{row}^*(M)$ . If*

$$p\lambda(n - \|\vec{v}\|) > \|\vec{v}\|,$$

*then there is a vector  $\vec{z} \in \mathbf{row}^*(M)$  such that  $\|\vec{z}\| > \|\vec{v}\|$ .*

*Proof.* Let  $M$  be an  $\mathbb{F}_p$ -matrix with  $n$  nonzero columns. Let  $\vec{v} \in \mathbf{row}^*(M)$ , and let  $k = \|\vec{v}\|$ . Let  $\vec{v} = (v_1, \dots, v_n)$ . W.l.o.g., suppose  $v_i \neq 0$  for  $1 \leq i \leq k$  and  $v_i = 0$  for  $k+1 \leq i \leq n$ . Let  $\vec{w} \in \mathbf{row}^*(M)$  be such that  $w_i \neq 0$  for at least  $\lambda(n-k)$  coordinates  $i$ , where  $k+1 \leq i \leq n$ . In other words, if we let  $\vec{w}_L = (w_1, \dots, w_k)$  and  $\vec{w}_R = (w_{k+1}, \dots, w_n)$ , then  $\|\vec{w}_R\| \geq \lambda(n-k)$ . Since  $\|\vec{w}\| = \|\vec{w}_L\| + \|\vec{w}_R\|$ , if  $\|\vec{w}_L\| \geq (p-1)\lambda(n-k)$ , then  $\|\vec{w}\| \geq p\lambda(n-k) > \|\vec{v}\|$ , and we are done. So we may assume that  $\|\vec{w}_L\| < (p-1)\lambda(n-k)$ .

Our goal will be to prove that there exists a nonzero constant  $c$  such that  $\|c\vec{w}_L + \vec{v}_L\| > k - \lambda(n-k)$ , where  $\vec{v}_L = (v_1, \dots, v_k)$ . Once we establish that such a constant exists, then we will be done, because we will have  $\|c\vec{w} + \vec{v}\| = \|c\vec{w}_L + \vec{v}_L\| + \|\vec{w}_R\| > (k - \lambda(n-k)) + \lambda(n-k) = k$ .

For  $1 \leq a \leq p-1$ , let  $S_a = \{i \in [k] : aw_i + v_i = 0\}$ . Since  $v_i \neq 0$  for  $1 \leq i \leq k$ , then  $S_a \subseteq \text{supp}(\vec{w}_L)$ . Thus, if  $aw_i + v_i = 0 = bw_i + v_i$ , then  $w_i \neq 0$ , which allows us to conclude that  $a = b$ . Therefore, if  $a \neq b$ , then  $S_a \cap S_b = \emptyset$ . Since  $\bigcup_{a=1}^{p-1} S_a \subseteq \text{supp}(\vec{w}_L)$  and the  $S_a$  are pairwise disjoint, we have

$$\sum_{a=1}^{p-1} |S_a| \leq |\text{supp}(\vec{w}_L)| = \|\vec{w}_L\| < (p-1)\lambda(n-k).$$

Therefore, the average value of  $|S_a|$  is strictly less than  $\lambda(n-k)$ , and if we let  $c \in [p-1]$  be such that  $|S_c|$  is minimum, then  $|S_c| < \lambda(n-k)$ . Thus,  $\|c\vec{w}_L + \vec{v}_L\| = k - |S_c| > k - \lambda(n-k)$ , and as noted above, we are done. Specifically,  $\|c\vec{w} + \vec{v}\| > \|\vec{v}\|$ .  $\square$

It is easy to check that the following corollary holds.

**Corollary 1.** *If  $1 \leq n \leq p$ , then  $\lambda(n) = n$ .*

For an integer  $k \geq 0$ , let  $\sigma_k = \sum_{j=0}^k p^j$ .

**Proposition 1.** *Suppose*

$$n = \sum_{j=\ell}^k b_j \sigma_j,$$

where  $b_k \geq 1$ , and  $0 \leq b_j \leq p-1$  for  $j \neq \ell$ , and  $1 \leq b_\ell \leq p$ . Then

$$\lambda(n) \geq \sum_{j=\ell}^k b_j p^j.$$

*Proof of Proposition 1.* We proceed by induction on  $k$ . When  $k = 0$ , then  $n = b_0 \sigma_0 = b_0$ . Since  $1 \leq b_0 \leq p$ , then  $\lambda(n) = b_0$  by Corollary 1, thus,  $\lambda(n) = b_0 p^0$  and the result holds. Now suppose  $k \geq 1$ . Our inductive hypothesis will be if

$$n = \sum_{j=\ell}^m b_j \sigma_j,$$

where  $b_m \geq 1$ , and  $0 \leq b_j \leq p-1$  for  $j \neq \ell$ , and  $1 \leq b_\ell \leq p$ , and  $m < k$ , then

$$\lambda(n) \geq \sum_{j=\ell}^m b_j p^j.$$

Let  $M$  be an  $\mathbb{F}_p$ -matrix with  $n$  nonzero columns. Suppose  $\vec{v} \in \mathbf{row}^*(M)$  with

$$\|\vec{v}\| < \sum_{j=\ell}^k b_j p^j.$$

Then

$$\begin{aligned} n - \|\vec{v}\| &> n - \sum_{j=\ell}^k b_j p^j = \sum_{j=\ell}^k b_j \sigma_j - \sum_{j=\ell}^k b_j p^j \\ &= \sum_{j=\ell}^k b_j (\sigma_j - p^j) \\ &= \sum_{j=\ell}^k b_j \sigma_{j-1}, \end{aligned}$$

where we define  $\sigma_{-1} = 0$  to handle the case  $j = 0$ , since  $\sigma_0 - p^0 = 0$ . Thus,

$$n - \|\vec{v}\| \geq \sum_{\ell-1 \leq j \leq k-1} b_{j+1}\sigma_j + 1.$$

We want to determine a lower bound on  $p\lambda\left(\sum_{j=\ell-1}^{k-1} b_{j+1}\sigma_j + 1\right)$  that allows us to conclude that  $p\lambda(n - \|\vec{v}\|) > \|\vec{v}\|$  so that we may use Lemma 1. We consider the case where  $b_\ell = p$  and the case where  $1 \leq b_\ell \leq p - 1$  separately.

Suppose  $b_\ell = p$ . Then

$$\begin{aligned} \sum_{\ell-1 \leq j \leq k-1} b_{j+1}\sigma_j + 1 &= \sum_{\ell \leq j \leq k-1} b_{j+1}\sigma_j + b_\ell\sigma_{\ell-1} + 1 \\ &= \sum_{\ell \leq j \leq k-1} b_{j+1}\sigma_j + (p\sigma_{\ell-1} + 1) \\ &= \sum_{\ell \leq j \leq k-1} b_{j+1}\sigma_j + \sigma_\ell \\ &= \sum_{\ell+1 \leq j \leq k-1} b_{j+1}\sigma_j + b_{\ell+1}\sigma_\ell + \sigma_\ell \\ &= \sum_{\ell+1 \leq j \leq k-1} b_{j+1}\sigma_j + (b_{\ell+1} + 1)\sigma_\ell. \end{aligned}$$

Notice that our sum satisfies all of the criteria for the inductive hypothesis. Specifically, the coefficient of its lowest sigma-term  $\sigma_\ell$  is  $b_{\ell+1} + 1$ , which satisfies  $1 \leq b_{\ell+1} + 1 \leq p$ ; the coefficient of  $\sigma_j$  is  $b_{j+1}$  and  $0 \leq b_{j+1} \leq p - 1$  for  $j \neq \ell$ ; the coefficient of the largest sigma-term  $\sigma_{k-1}$  is  $b_k$ , which satisfies  $b_k \geq 1$ ; and finally, the index of its largest sigma term is  $k - 1$  which is strictly less than  $k$ . Therefore, by the inductive hypothesis,

$$\begin{aligned} p \cdot \lambda\left(\sum_{\ell+1 \leq j \leq k-1} b_{j+1}\sigma_j + (b_{\ell+1} + 1)\sigma_\ell\right) &\geq p \left(\sum_{\ell+1 \leq j \leq k-1} b_{j+1}p^j + (b_{\ell+1} + 1)p^\ell\right) \\ &= \sum_{\ell+1 \leq j \leq k-1} b_{j+1}p^{j+1} + (b_{\ell+1} + 1)p^{\ell+1} \\ &= \sum_{\ell \leq j \leq k-1} b_{j+1}p^{j+1} + p \cdot p^\ell \\ &= \sum_{\ell+1 \leq j \leq k} b_j p^j + p \cdot p^\ell \\ &= \sum_{\ell \leq j \leq k} b_j p^j, \end{aligned}$$

where the last equality holds because  $b_\ell = p$ . Since  $\lambda$  is a nondecreasing function, our previous work implies

$$\begin{aligned} p\lambda(n - \|\vec{v}\|) &\geq p \cdot \lambda\left(\sum_{\ell+1 \leq j \leq k-1} b_{j+1}\sigma_j + (b_{\ell+1} + 1)\sigma_\ell\right) \\ &\geq \sum_{\ell \leq j \leq k} b_j p^j > \|\vec{v}\|. \end{aligned}$$

Thus, by Lemma 1, there is a vector  $\vec{z} \in \mathbf{row}^*(M)$  such that  $\|\vec{z}\| > \|\vec{v}\|$ .

Now suppose  $1 \leq b_\ell \leq p - 1$ . Recall that our sum is

$$\sum_{\ell-1 \leq j \leq k-1} b_{j+1} \sigma_j + 1 = \sum_{\ell-1 \leq j \leq k-1} b_{j+1} \sigma_j + 1 \cdot \sigma_0.$$

In this case, the smallest sigma-term is  $\sigma_0$ , and its coefficient is  $b_1 + 1$ , where  $b_1 = 0$  if  $\ell \geq 2$ . We note that our sum satisfies all of the criteria for the inductive hypothesis. Since each  $b_j$  satisfies  $0 \leq b_j \leq p - 1$ , then  $1 \leq b_1 + 1 \leq p$ ; when  $j \geq 1$ , the coefficient of each  $\sigma_j$  is  $b_{j+1}$  and  $0 \leq b_{j+1} \leq p - 1$ ; the coefficient of the largest sigma-term  $\sigma_{k-1}$  is  $b_k$ , which satisfies  $b_k \geq 1$ ; and finally, the index of its largest sigma term is  $k - 1$  which is strictly less than  $k$ .

When  $\ell \geq 2$ , the coefficient of  $\sigma_0$  is 1, and we apply the inductive hypothesis to obtain

$$\begin{aligned} p \cdot \lambda \left( \sum_{\ell-1 \leq j \leq k-1} b_{j+1} \sigma_j + 1 \right) &\geq p \left( \sum_{\ell-1 \leq j \leq k-1} b_{j+1} p^j + 1 \right) \\ &= \sum_{\ell-1 \leq j \leq k-1} b_{j+1} p^{j+1} + p \\ &= \sum_{\ell \leq j \leq k} b_j p^j + p. \end{aligned}$$

Thus,

$$\begin{aligned} p \lambda(n - \|\vec{v}\|) &\geq p \cdot \lambda \left( \sum_{\ell-1 \leq j \leq k-1} b_{j+1} \sigma_j + 1 \right) \\ &\geq \sum_{\ell \leq j \leq k} b_j p^j + p > \|\vec{v}\|. \end{aligned}$$

When  $\ell \in \{0, 1\}$ , our sum is  $\sum_{j=0}^{k-1} b_{j+1} \sigma_j + 1$ , and we apply the inductive hypothesis to obtain

$$\begin{aligned} p \cdot \lambda \left( \sum_{0 \leq j \leq k-1} b_{j+1} \sigma_j + 1 \right) &= p \cdot \lambda \left( \sum_{1 \leq j \leq k-1} b_{j+1} \sigma_j + (b_1 + 1) \right) \\ &\geq p \left( \sum_{1 \leq j \leq k-1} b_{j+1} p^j + b_1 + 1 \right) \\ &= \sum_{1 \leq j \leq k-1} b_{j+1} p^{j+1} + b_1 p + p \\ &= \sum_{1 \leq j \leq k} b_j p^j + p. \end{aligned}$$

Thus,

$$\begin{aligned}
p\lambda(n - \|\vec{v}\|) &\geq p \cdot \lambda\left(\sum_{0 \leq j \leq k-1} b_{j+1}\sigma_j + 1\right) \\
&\geq \sum_{1 \leq j \leq k} b_j p^j + p \\
&> \sum_{\ell \leq j \leq k} b_j p^j > \|\vec{v}\|.
\end{aligned}$$

Thus, by Lemma 1, there is a vector  $\vec{z} \in \mathbf{row}^*(M)$  such that  $\|\vec{z}\| > \|\vec{v}\|$ . Therefore  $\lambda(n) \geq \sum_{j=\ell}^k b_j p^j$ .  $\square$

Now we show that every  $n \geq 1$  can be written in the form described in Proposition 1.

**Claim 1.** *Let  $n \in \mathbb{Z}^+$ . Suppose  $n < \sigma_{k+1}$ . Let  $n_{k+1} = n$ , and for  $0 \leq j \leq k$ , assuming  $n_{j+1}$  is defined, let  $b_j$  be the largest integer such that  $b_j \sigma_j \leq n_{j+1}$ , and let  $n_j = n_{j+1} - b_j \sigma_j$ . Then for  $0 \leq j \leq k$ , we have  $0 \leq n_{j+1} \leq p\sigma_j$  and  $0 \leq b_j \leq p$ . Moreover,*

$$n = \sum_{j=0}^k b_j \sigma_j,$$

and if  $b_j = p$ , then  $b_i = 0$  for  $i < j$ .

*Proof of Claim 1.* Suppose  $n \in \mathbb{Z}^+$  and  $n < \sigma_{k+1}$ . Then  $n \leq \sigma_{k+1} - 1 = p\sigma_k$ . Let  $n_{k+1} = n$ , and for  $0 \leq j \leq k$ , assuming  $n_{j+1}$  is defined, let  $b_j$  be the largest integer such that  $b_j \sigma_j \leq n_{j+1}$ , and let  $n_j = n_{j+1} - b_j \sigma_j$ . We proceed by induction on  $k - j$ . Assume  $0 \leq n_{j+1} \leq p\sigma_j$  and let  $b_j$  and  $n_j$  be defined as above. Since  $0 \leq n_{j+1}$ , then  $b_j \geq 0$ . Since  $n_{j+1} \leq p\sigma_j$  and  $b_j \sigma_j \leq n_{j+1}$ , then  $b_j \sigma_j \leq p\sigma_j$ . Thus, since  $\sigma_j \geq 1$ , we have  $b_j \leq p$ . Since  $b_j \sigma_j \leq n_{j+1}$  and  $n_j = n_{j+1} - b_j \sigma_j$ , then  $n_j \geq 0$ . Since  $n_{j+1} < (b_j + 1)\sigma_j$ , then  $n_{j+1} - b_j \sigma_j < \sigma_j$ , i.e.,  $n_j < \sigma_j - 1 = p\sigma_{j-1}$ . Therefore, by induction,  $0 \leq n_{j+1} \leq p\sigma_j$  and  $0 \leq b_j \leq p$  for  $0 \leq j \leq k$ .

Now suppose  $b_j = p$ . Since  $b_j \sigma_j \leq n_{j+1} \leq p\sigma_j$ , then  $n_{j+1} = p\sigma_j$  and  $n_j = n_{j+1} - b_j \sigma_j = 0$ . Moreover,  $b_i = 0$  and  $n_i = 0$  for all  $i < j$ .

To see that  $n = \sum_{j=0}^k b_j \sigma_j$ , observe that  $b_j \sigma_j = n_{j+1} - n_j$  for  $0 \leq j \leq k$ , because of the definition of  $n_j$ . Thus,

$$\sum_{j=0}^k b_j \sigma_j = \sum_{j=0}^k (n_{j+1} - n_j) = n_{k+1} - n_0 = n - n_0.$$

Since  $0 \leq n_1 \leq p\sigma_0 = p$ , then, by definition,  $b_0 = n_1$  and  $n_0 = n_1 - b_0 \sigma_0 = n_1 - n_1(1) = 0$ . Thus,  $\sum_{j=0}^k b_j \sigma_j = n$   $\square$

With Proposition 1 and Claim 1, we have established a lower bound on  $\lambda(n)$  for all  $n \geq 1$ . We need to prove the corresponding upper bound. We will do so by constructing a matrix with  $n$  columns whose capacity equals the lower bound given in Proposition 1. We begin by constructing such a matrix for certain values of  $n$ , namely, when  $n = \sigma_k$  for some  $k \geq 0$ .

For each integer  $k \geq 0$ , we define a  $(k+1) \times \sigma_k$  matrix  $B_k$ , recursively, as follows. The matrix  $B_0$  is the  $1 \times 1$  matrix whose sole entry is 1. For  $k \geq 1$ ,  $B_k$  can be defined as a block matrix with a “row” consisting of  $p$  copies of  $B_{k-1}$  followed by a  $k \times 1$  column of 0’s, then one more row of dimensions  $1 \times \sigma_k$  with its first  $\sigma_{k-1}$  entries equal to 0 (below the first  $B_{k-1}$ ), then  $\sigma_{k-1}$  entries equal to 1 (below the next  $B_{k-1}$ ),  $\dots$ , then  $\sigma_{k-1}$  entries equal to  $p-1$  (below the last  $B_{k-1}$ ), and one last entry equal to 1, i.e.,

$$B_k = \left[ \begin{array}{c|c|c|c|c} B_{k-1} & B_{k-1} & \cdots & B_{k-1} & \begin{array}{c} 0 \\ \vdots \\ 0 \end{array} \\ \hline 0 \dots 0 & 1 \dots 1 & \cdots & (p-1) \dots (p-1) & 1 \end{array} \right].$$

For  $k \geq 1$ , let  $B'_k$  be the  $k \times \sigma_k$  matrix obtained from  $B_k$  by removing its last row, i.e.,

$$B'_k = \left[ \begin{array}{c|c|c|c|c} B_{k-1} & B_{k-1} & \cdots & B_{k-1} & \begin{array}{c} 0 \\ \vdots \\ 0 \end{array} \end{array} \right].$$

**Lemma 2.** For each  $\vec{v} \in \mathbf{row}^*(B_k)$ ,  $\|\vec{v}\| = p^k$ .

*Proof.* We proceed by induction on  $k$ . When  $k = 0$ , the result is trivial. Let  $k \geq 1$ . Assume the result for  $j < k$ . Let  $\vec{v} \in \mathbf{row}^*(B_k)$ . We first consider the case where  $\vec{v} \in \mathbf{row}^*(B'_k)$ . Then we can write

$$\vec{v} = (v_1^{(0)}, \dots, v_{\sigma_{k-1}}^{(0)}, v_1^{(1)}, \dots, v_{\sigma_{k-1}}^{(1)}, \dots, v_1^{(p-1)}, \dots, v_{\sigma_{k-1}}^{(p-1)}, 0).$$

To shorten notation, we will write

$$\vec{v} = (\vec{v}_0, \vec{v}_1, \dots, \vec{v}_{p-1}, 0), \tag{3}$$

where  $\vec{v}_i = (v_1^{(i)}, \dots, v_{\sigma_{k-1}}^{(i)})$  for  $0 \leq i \leq p-1$ . Technically, in equation (3),  $\vec{v}_i$  simply represents the coordinates  $v_1^{(i)}, \dots, v_{\sigma_{k-1}}^{(i)}$ . We observe that  $\vec{v}_0 = \vec{v}_1 = \dots = \vec{v}_{p-1}$  based on how  $B'_k$  and  $\vec{v}$  are defined. We also observe that  $\vec{v}_i \in \mathbf{row}^*(B_{k-1})$ . By the inductive hypothesis,  $\|\vec{v}_i\| = p^{k-1}$ , therefore,  $\|\vec{v}\| = p^k$ .

We now show the result holds for  $\vec{w} \in \mathbf{row}^*(B_k) - \mathbf{row}^*(B'_k)$ . Let  $\vec{u}$  be the last row in  $B_k$ , i.e.,  $\vec{u} = (0, \dots, 0, 1, \dots, 1, \dots, p-1, \dots, p-1, 1)$ . We observe that  $\|\vec{u}\| = \sigma_k - \sigma_{k-1} = p^k$ , thus, the result holds when  $\vec{w} = \vec{u}$ . To illustrate our argument, we next consider the special case where  $\vec{w} = \vec{v} + \vec{u}$  for some  $\vec{v} \in \mathbf{row}^*(B'_k)$ . Again, we slightly abuse notation and write  $\vec{u} = (\vec{0}, \vec{1}, \dots, (p-1)\vec{1}, 1)$ , where  $\vec{c}$  (or  $c\vec{1}$ ) represents the  $\sigma_{k-1}$ -dimensional vector  $(c, \dots, c)$ . Then we can write  $\vec{v} + \vec{u} = (\vec{v}_0 + \vec{0}, \vec{v}_1 + \vec{1}, \dots, \vec{v}_{p-1} + (p-1)\vec{1}, 1)$ . Since we are working modulo  $p$ , a coordinate of  $\vec{v}_j + j\vec{1}$  is congruent to 0 if and only if the corresponding coordinate of  $\vec{v}_j$  is congruent to  $p-j$ . Thus, we can count the total number of coordinates that are congruent to 0 in  $\vec{v} + \vec{u}$  as follows

$$\left( \begin{array}{c} \text{Total \# of 0-coordinates} \\ \text{in } \vec{v} + \vec{u} \end{array} \right) = \sum_{j=0}^{p-1} (\# \text{ of } (p-j)\text{-coordinates in } \vec{v}_j). \tag{4}$$

Since  $\vec{v}_0 = \vec{v}_1 = \dots = \vec{v}_{p-1}$ , equation (4) reduces to

$$\binom{\text{Total \# of 0-coordinates}}{\text{in } \vec{v} + \vec{u}} = \binom{\text{Total \# of coordinates}}{\text{in } \vec{v}_0} = \sigma_{k-1}.$$

Thus,  $\|\vec{v} + \vec{u}\| = \sigma_k - \sigma_{k-1} = p^k$ . In general,  $\vec{w} \in \mathbf{row}^*(B_k) - \mathbf{row}^*(B'_k)$  satisfies  $\vec{w} = \vec{v} + c\vec{u}$  for some  $\vec{v} \in \mathbf{row}^*(B'_k)$  and  $c \not\equiv 0 \pmod{p}$ . In this case,  $\vec{w} = (\vec{v}_0 + c\vec{0}, \vec{v}_1 + c\vec{1}, \dots, \vec{v}_{p-1} + c(p-1)\vec{1}, 1)$ , and equation (4) becomes

$$\binom{\text{Total \# of 0-coordinates}}{\text{in } \vec{w}} = \sum_{j=0}^{p-1} (\# \text{ of } (p-cj)\text{-coordinates in } \vec{v}_j), \quad (5)$$

where arithmetic is modulo  $p$ . Since  $\vec{v}_0 = \vec{v}_1 = \dots = \vec{v}_{p-1}$ , we obtain

$$\binom{\text{Total \# of 0-coordinates}}{\text{in } \vec{w}} = \sum_{j=0}^{p-1} (\# \text{ of } (p-cj)\text{-coordinates in } \vec{v}_0).$$

Since  $p$  is prime and  $c \not\equiv 0 \pmod{p}$ , then  $\{p, p-c, p-2c, \dots, p-(p-1)c\}$  is a equivalent to  $\{0, 1, \dots, p-1\}$  modulo  $p$ , thus,

$$\binom{\text{Total \# of 0-coordinates}}{\text{in } \vec{w}} = \binom{\text{Total \# of coordinates}}{\text{in } \vec{v}_0} = \sigma_{k-1}.$$

Therefore,  $\|\vec{w}\| = \sigma_k - \sigma_{k-1} = p^k$ , and we can conclude that for each  $\vec{v} \in \mathbf{row}^*(B_k)$ ,  $\|\vec{v}\| = p^k$ .  $\square$

Since  $B_k$  has  $\sigma_k$  columns, Lemma 2 implies that  $\lambda(n) \leq p^k$  when  $n = \sigma_k$  for some nonnegative integer  $k$ . We would like a similar upper bound on  $\lambda(n)$  for all positive integers  $n$ . Thus, we provide the following proposition.

**Proposition 2.** *If  $n = \sum_{j=0}^k b_j \sigma_j$ , then*

$$\lambda(n) \leq \sum_{j=0}^k b_j p^j.$$

*Proof of Proposition 2.* We will construct a matrix  $M$  with  $n$  columns such that  $c(M) = \sum_{j=0}^k b_j p^j$ . The matrix  $M$  will essentially be a block matrix with  $b_j$  copies of  $B_j$  for  $0 \leq j \leq k$ . However, the number of rows of  $B_j$  does not equal the number of rows of  $B_\ell$  when  $j \neq \ell$ . Thus, for  $0 \leq j \leq k$ , we define the  $(k+1) \times \sigma_j$  matrix  $B_j^{(k)}$  where the first  $j$  rows of  $B_j^{(k)}$  match the first  $j$  rows of  $B_j$  and the last  $k+1-j$  rows of  $B_j^{(k)}$  all equal the last row of  $B_j$ . Thus,  $B_0^{(k)}$  is a  $(k+1) \times 1$  column of 1's, and for  $1 \leq j \leq k$ ,

$$B_j^{(k)} = \left[ \begin{array}{c|c|c|c|c} & & & & 0 \\ & B_{j-1} & B_{j-1} & \cdots & B_{j-1} \\ & & & & \vdots \\ & & & & 0 \\ \hline 0 \dots 0 & 1 \dots 1 & \cdots & (p-1) \dots (p-1) & 1 \\ \hline 0 \dots 0 & 1 \dots 1 & \cdots & (p-1) \dots (p-1) & 1 \\ \hline \vdots & \vdots & \cdots & \vdots & \vdots \\ \hline 0 \dots 0 & 1 \dots 1 & \cdots & (p-1) \dots (p-1) & 1 \end{array} \right]$$



where the last row is repeated  $(k+1) - j$  times. After comparing  $B_j^{(k)}$  with  $B_j$ , it is easy to see that  $\mathbf{row}^*(B_j^{(k)}) = \mathbf{row}^*(B_j)$ .

Let  $n$  be a positive integer such that  $n = \sum_{j=0}^k b_j \sigma_j$ . Let  $M$  be the  $(k+1) \times n$  matrix defined as a block matrix with  $b_j$  copies of  $B_j^{(k)}$  for  $0 \leq j \leq k$ , where the blocks appear in a single row in nondecreasing order according to their lower index, i.e.,

$$M = \left[ \underbrace{B_0^{(k)} \cdots B_0^{(k)}}_{b_0} \mid \underbrace{B_1^{(k)} \cdots B_1^{(k)}}_{b_1} \mid \cdots \mid \underbrace{B_k^{(k)} \cdots B_k^{(k)}}_{b_k} \right].$$

Let  $\vec{v} \in \mathbf{row}^*(M)$ . Then we can (essentially) write

$$\vec{v} = (\vec{v}_1^{(0)}, \dots, \vec{v}_{b_0}^{(0)}, \vec{v}_1^{(1)}, \dots, \vec{v}_{b_1}^{(1)}, \dots, \vec{v}_1^{(k)}, \dots, \vec{v}_{b_k}^{(k)})$$

where  $\vec{v}_i^{(j)} \in \mathbf{row}^*(B_j)$  for  $0 \leq j \leq k$  and  $1 \leq i \leq b_j$ . Moreover, for  $1 \leq i \leq b_j$ , we have  $\vec{v}_i^{(j)} = \vec{v}_{b_j}^{(j)}$ . Thus,

$$\|\vec{v}\| = \sum_{j=0}^k b_j \|\vec{v}_{b_j}^{(j)}\|.$$

Because  $\vec{v}_{b_j}^{(j)} \in \mathbf{row}^*(B_j)$ , Lemma 2 implies  $\|\vec{v}_{b_j}^{(j)}\| = p^j$ , therefore,

$$\|\vec{v}\| = \sum_{j=0}^k b_j p^j.$$

Thus,  $c(M) = \sum_{j=0}^k b_j p^j$ , and  $\lambda(n) \leq \sum_{j=0}^k b_j p^j$ .  $\square$

Thus, we can combine Propositions 1 and 2 with Claim 1 to obtain the following corollary.

**Corollary 2.** *Let  $n \in \mathbb{Z}^+$ . Suppose  $n < \sigma_{k+1}$ . Then*

$$n = \sum_{j=0}^k b_j \sigma_j,$$

where  $0 \leq b_j \leq p$  for  $0 \leq j \leq k$ , and if  $b_j = p$ , then  $b_i = 0$  for  $i < j$ . Moreover,

$$\lambda(n) = \sum_{j=0}^k b_j p^j.$$

**Corollary 3.** *The sequence  $\lambda(n)$  satisfies the meta-Fibonacci recurrence relation*

$$\lambda(n) = \sum_{i=1}^p \lambda(n - i + 1 - \lambda(n - i)).$$

*Proof of Corollary 3.* We refer to Corollary 32 in [4], which implies that a sequence which is defined by the meta-Fibonacci recurrence relation (2) is also defined by the recurrence relation

$$\lambda(n) = p^k + \lambda(n - \sigma_k), \tag{6}$$

for  $\sigma_k \leq n < \sigma_{k+1}$ . Based on Corollary 2, it is clear that  $\lambda(n)$  satisfies recurrence (6). Therefore,  $\lambda(n)$  satisfies the meta-Fibonacci recurrence (2).  $\square$

## References

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