STACK-SORTING WITH CONSECUTIVE-PATTERN-AVOIDING STACKS

COLIN DEFANT AND KAI ZHENG

ABSTRACT. We introduce consecutive-pattern-avoiding stack-sorting maps SC_{σ} , which are natural generalizations of West's stack-sorting map s and natural analogues of the classical-pattern-avoiding stack-sorting maps s_{σ} recently introduced by Cerbai, Claesson, and Ferrari. We characterize the patterns σ such that $Sort(SC_{\sigma})$, the set of permutations that are sortable via the map $s \circ SC_{\sigma}$, is a permutation class, and we enumerate the sets $Sort(SC_{\sigma})$ for $\sigma \in \{123, 132, 321\}$. We also study the maps SC_{σ} from a dynamical point of view, characterizing the periodic points of SC_{σ} for all $\sigma \in S_3$ and computing $\max_{\pi \in S_n} |SC_{\sigma}^{-1}(\pi)|$ for all $\sigma \in \{132, 213, 231, 312\}$. In addition, we characterize the periodic points of the classical-pattern-avoiding stack-sorting map s_{132} , and we show that the maximum number of iterations of s_{132} needed to send a permutation in S_n to a periodic point is n-1. The paper ends with numerous open problems and conjectures.

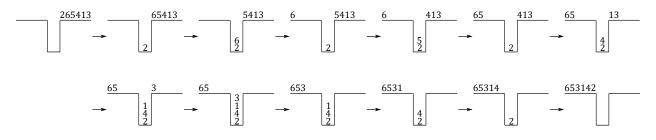
1. INTRODUCTION

1.1. Background. Let S_n denote the set of permutations of the set $[n] := \{1, \ldots, n\}$, which we write as words in one-line notation. The investigation of pattern avoidance in permutations began in 1968, when Knuth [28] introduced a stack-sorting machine and showed that a permutation can be sorted to the identity permutation using this machine if and only if it avoids the pattern 231 (we define pattern containment and avoidance formally below). In his 1990 Ph.D. dissertation, West [37] introduced a deterministic variant of Knuth's machine; this variant is a function $s: S_n \to S_n$. As in Knuth's case, a permutation $\pi \in S_n$ satisfies $s(\pi) = 123 \cdots n$ if and only if it avoids 231. More recently, Albert, Homberger, Pantone, Shar, and Vatter [2] introduced a vast generalization of Knuth's stack-sorting machine by considering \mathcal{C} -machines. Roughly speaking, a \mathcal{C} -machine is a sorting machine that uses a stack whose entries, when read from top to bottom, must have the same relative order as an element of the permutation class \mathcal{C} . Even more recently, Cerbai, Claesson, and Ferrari [11] generalized West's stack-sorting map in a similar manner. For each permutation pattern σ , they defined a map $s_{\sigma}: S_n \to S_n$ that sends each permutation through a stack in a right-greedy manner, insisting that the contents of the stack must avoid the pattern σ when read from top to bottom. In this more general setup, West's stack-sorting map is just s_{21} . Although the article [11] is quite recent, it has already spawned several subsequent papers [3, 5, 10, 12].

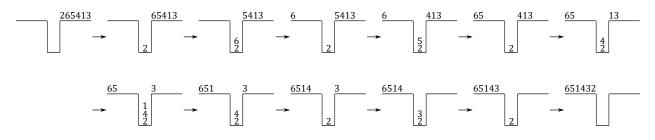
In this article, we introduce consecutive-pattern-avoiding stack-sorting maps SC_{σ} . In order to define the maps s_{σ} and SC_{σ} formally and simultaneously, it is helpful to define more general maps $T_{\mathcal{A}}$ dependent on sets \mathcal{A} of permutations. Assume we are given an input permutation $\pi = \pi_1 \cdots \pi_n$. Throughout this procedure, we consider the permutation obtained by reading the contents of a vertical stack from top to bottom. If moving the next entry in the input permutation to the top of the stack results in a stack whose contents have the same relative order as any element of \mathcal{A} , then the next entry in the input permutation is placed at the top of the stack (this operation is called a *push*). Otherwise, the entry at the top of the stack is annexed to the end of the growing output permutation (this operation is called a *pop*). This procedure stops when the output permutation has length n. We then define $T_{\mathcal{A}}(\pi)$ to be this output permutation. When \mathcal{A} is the set of permutations

that avoid σ , the map $T_{\mathcal{A}}$ is s_{σ} . Our focus in this article is on the case in which \mathcal{A} is the set of permutations that avoid σ as a consecutive pattern, in which case we write $T_{\mathcal{A}} = SC_{\sigma}$.

Example 1.1. The following diagram portrays the steps involved in sending the permutation 265413 through a consecutive-231-avoiding stack, showing that $SC_{231}(265413) = 653142$.



Notice that there is one step where the contents of the stack, when read from top to bottom, form the permutation 3142. The permutation 3142 contains the pattern 231 classically (i.e., nonconsecutively), but avoids 231 consecutively, which is why it is allowed to be in the stack. On the other hand, this would not be allowed if we were using a classical-231-avoiding stack. In that case, the steps would proceed as in the following diagram, which shows that $s_{231}(265413) = 651432$.



Our newly-defined maps SC_{σ} serve as perfectly legitimate generalizations of West's stack-sorting map because SC_{21} is the same as s, which is also the same as s_{21} (a permutation avoids the pattern 21 classically if and only if it avoids 21 consecutively). However, when the length of σ is at least 3, the maps SC_{σ} and s_{σ} exhibit different behaviors.

The decision to consider stacks that avoid consecutive patterns, rather than stacks that avoid classical patterns, is very natural. Indeed, when one actually works through the process of sending a permutation through a classical-21-avoiding stack, one checks at each step whether the next entry in the input permutation is greater than the entry currently sitting at the top of the stack. One does not bother checking if the next entry in the input is greater than all of the entries in the stack because the smallest number in the stack is necessarily the top entry (this is another manifestation of the simple observation that a permutation avoids 21 classically if and only if it avoids 21 consecutively). If σ has length $k \geq 3$, then when one sends a permutation through a classical- σ -avoiding stack, one must check all of the entries in the stack to see if the addition of the next entry produces an occurrence of the pattern σ . By contrast, when one sends a permutation through a consecutive- σ -avoiding stack, it is only necessary to compare the next entry in the input permutation with the top k-1 entries in the stack. In this way, the mechanics governing the maps SC_{σ} are in some sense "closer" to those governing West's stack-sorting map than those of the maps s_{σ} are.

Typically, researchers have approached West's stack-sorting map from a "sorting" point of view. The usual questions concern structural and enumerative aspects of t-stack-sortable permutations, which are permutations that get sorted to the identity using at most t iterations of the map

s [6–8, 13, 15, 17, 22, 37, 38]. On the other hand, there are now several articles that study the stacksorting map from a "dynamical" point of view (see [9, 14, 16, 18, 19] and the references therein). In this approach, we view the stack-sorting map as an interesting combinatorially-defined function, without caring so much about trying to sort permutations. The typical questions concern the number of preimages of a permutation under s, which is called the *fertility* of the permutation. The investigation of fertilities of permutations has revealed unexpected patterns and connections with several other aspects of combinatorics [9,14,16,19,31,33], including a very surprising and useful link with the combinatorics of cumulants in noncommutative probability theory [18]. This dynamical approach concerning fertilities has also shed new light on the *t*-stack-sortable permutations studied from the sorting point of view [15,17].

In this paper, we consider the maps SC_{σ} from both the sorting point of view and the dynamical point of view. All of the sorting questions we consider are analogues of questions asked about the maps s_{σ} in [11,12]. On the other hand, our dynamical approach is completely novel in the realm of pattern-avoiding stack-sorting maps (other than *s* itself). We will also see that there are several dynamical questions that one can ask about the maps SC_{σ} that are trivial when $\sigma = 21$ (i.e., when $SC_{\sigma} = s$), but which become much more interesting when the length of σ is at least 3.

1.2. Outline and Summary of Main Results. The articles [11, 12] focus primarily on the set Sort(s_{σ}) of permutations that get sorted to the identity using a σ -avoiding stack followed by a classical-21-avoiding stack. Since the set of permutations that s maps to the identity is the set Av(231) of 231-avoiding permutations, we have Sort(s_{σ}) = s_{σ}^{-1} (Av(231)). Note that Sort(s_{21}) = s^{-1} (Av(231)) consists of the 2-stack-sortable permutations, which have received an enormous amount of attention since their inception in West's dissertation [6–8, 15, 17, 22, 37, 38]. Let us also remark that the articles [3, 10] have studied other close variants of these sets; for example, the paper [3] explores the sets Sort($s_{\sigma,\tau}$) = $s_{\sigma,\tau}^{-1}$ (Av(231)), where $s_{\sigma,\tau}$ is the map that sends a permutation through a right-greedy stack whose contents avoid both σ and τ (using our earlier notation, $s_{\sigma,\tau} = T_{Av(\sigma,\tau)}$). In Section 2, we define Sort(SC $_{\sigma}$) = SC $_{\sigma}^{-1}$ (Av(231)), which can similarly be seen as the set of permutations that get sorted into the identity by a consecutive- σ -avoiding stack followed by a 21-avoiding stack. We will prove that if $\sigma \in S_k$ for some $k \ge 3$, then Sort(SC $_{\sigma}$) is a permutation class if and only if σ and $\hat{\sigma}$ both contain the pattern 231, where $\hat{\sigma}$ is the permutation obtained by swapping the first 2 entries in σ ; this is a direct analogue of one of the main results from [11].

The other main results from [11] concern the enumeration of the sets $Sort(s_{123})$ and $Sort(s_{321})$. Extending this work, the article [12] focuses on the enumeration of the set $Sort(s_{132})$. In Sections 4 and 6, we give analogues of these results by characterizing and enumerating the sets $Sort(SC_{123})$, $Sort(SC_{321})$, and $Sort(SC_{132})$. More precisely, we will see that $Sort(SC_{123})$ is enumerated by the first differences of Motzkin numbers, that $Sort(SC_{321})$ is enumerated by the Motzkin numbers, and that $Sort(SC_{132})$ is enumerated by the Catalan numbers.

In Section 3, we adopt the dynamical point of view and consider the periodic points of the maps SC_{σ} , which we view as functions from S_n to S_n for each $n \geq 1$. It is easy to show that the only periodic point of $s: S_n \to S_n$ is the identity permutation $123 \cdots n$. By contrast, if the length of σ is at least 3, then SC_{σ} has multiple periodic points. We will prove that if $\sigma \in S_3$, then the periodic points of $SC_{\sigma}: S_n \to S_n$ are precisely the permutations in S_n that avoid σ and its reverse as consecutive patterns.

It is known that the identity permutation $123 \cdots n$, which has $C_n := \frac{1}{n+1} \binom{2n}{n}$ preimages under s, has strictly more preimages under s than every other permutation in S_n (see [6, Chapter 8, Exercise

23]). The analogous results for the maps SC_{σ} are more difficult. In Section 4, we study preimages under the maps SC_{132} and SC_{312} , proving in particular that

$$\max_{\pi \in S_n} |\operatorname{SC}_{132}^{-1}(\pi)| = \max_{\pi \in S_n} |\operatorname{SC}_{312}^{-1}(\pi)| = \binom{n-1}{\lfloor \frac{n-1}{2} \rfloor}$$

(along the way, we obtain a simple formula for $|\operatorname{SC}_{132}^{-1}(\pi)|$ for all reverse-layered permutations π). In Section 5, we prove a similar result for the maps SC_{231} and SC_{213} , showing that

$$\max_{\pi \in S_n} |\operatorname{SC}_{231}^{-1}(\pi)| = \max_{\pi \in S_n} |\operatorname{SC}_{213}^{-1}(\pi)| = 2^{n-2}$$

when $n \geq 2$.

Section 7 is a little different from the rest of the paper because it focuses on classical-patternavoiding stack-sorting maps instead of consecutive-pattern-avoiding stack-sorting maps. The purpose of this section is to prove dynamical results about the maps s_{132} and s_{312} . Such questions have not previously been considered for these maps, and we believe our results open the gate for interesting further developments. Specifically, we prove that the periodic points of s_{132} (respectively, s_{312}) are precisely the permutations that avoid 132 and 231 (respectively, 312 and 213). We also prove the the maximum number of iterations needed for s_{132} (respectively, s_{312}) to send a permutation to one of its periodic points is n - 1 (these results are much more difficult than the analogous results for West's stack-sorting map).

Finally, Section 8 collects a large number of open problems, conjectures, and other suggestions for future work.

1.3. Terminology and Notation. Recall that in this article, a *permutation* is an ordering of the elements of the set [n] for some n. If π is a sequence of n distinct integers, then the *standardization* of π is the permutation in S_n obtained by replacing the *i*th-smallest entry in π with *i* for all *i*. For example, the standardization of 4829 is 2314. We say two sequences have the *same relative* order if their standardizations are equal. We say a permutation σ contains a permutation τ as a pattern if there is a (not necessarily consecutive) subsequence of σ that has the same relative order as τ ; otherwise, σ avoids τ . We say σ contains τ consecutively if σ has a consecutive subsequence with the same relative order as τ ; otherwise, σ avoids the consecutive, we will sometimes call it classical pattern containment (or avoidance) is nonconsecutive, we will sometimes call it classical information about classical and consecutive permutation patterns.

We denote consecutive patterns by underlining them (or referring to them as consecutive patterns). If $\tau^{(1)}, \tau^{(2)}, \ldots$ is a (finite or infinite) list of classical or consecutive patterns, then we write $\operatorname{Av}_n(\tau^{(1)}, \tau^{(2)}, \ldots)$ for the set of permutations in S_n that avoid $\tau^{(1)}, \tau^{(2)}, \ldots$. We also write $\operatorname{Av}(\tau^{(1)}, \tau^{(2)}, \ldots) = \bigcup_{n\geq 0} \operatorname{Av}_n(\tau^{(1)}, \tau^{(2)}, \ldots)$. For example, $\operatorname{Av}(231, 321)$ is the set of permutations that avoid 231 and 321, while $\operatorname{Av}(\underline{321})$ is the set of permutations that avoid the consecutive pattern 321. If τ is an arbitrary unspecified permutation, then we let $\operatorname{Av}(\underline{\tau})$ be the set of permutations that consecutively avoid τ .

Given a permutation $\sigma = \sigma_1 \cdots \sigma_k$ of length $k \ge 2$, we let $\hat{\sigma} = \sigma_2 \sigma_1 \sigma_3 \cdots \sigma_k$ denote the permutation obtained by swapping the first two entries in σ . We also let $\operatorname{rev}(\sigma) = \sigma_k \cdots \sigma_1$ denote the reverse of σ . If $\sigma \in S_k$, then we write $\operatorname{comp}(\sigma) = (k+1-\sigma_1)\cdots(k+1-\sigma_k)$ for the complement of σ . A descent (respectively, ascent) of σ is an index *i* such that $\sigma_i > \sigma_{i+1}$ (respectively, $\sigma_i < \sigma_{i+1}$). An ascending run of σ is a maximal consecutive increasing subsequence of π , while a descending run is a maximal consecutive decreasing subsequence. For example, the ascending runs of 7862351

are 78, 6, 235, and 1, while the descending runs of 7862351 are 7, 862, 3, and 51. We write $|\tau|$ for the length of a sequence τ of positive integers.

1.4. A Preliminary Lemma. Before we proceed, it will be helpful to record the following lemma, which will allow us to simplify many of our considerations. The proof of this lemma is immediate from the definitions of s_{σ} and SC_{σ}.

Lemma 1.2. If $\sigma \in S_k$ for some $k \ge 2$, then

 $s_{\operatorname{comp}(\sigma)} = \operatorname{comp} \circ s_{\sigma} \circ \operatorname{comp}$ and $\operatorname{SC}_{\operatorname{comp}(\sigma)} = \operatorname{comp} \circ \operatorname{SC}_{\sigma} \circ \operatorname{comp}$.

2. When is $Sort(SC_{\sigma})$ a Permutation Class?

A permutation class is a set C of permutations that is closed under classical pattern containment. Said differently, this means that if σ, τ are permutations such that σ contains τ classically and $\sigma \in C$, then $\tau \in C$. Equivalently, a set of permutations is a permutation class if and only if it is of the form $\operatorname{Av}(\tau^{(1)}, \tau^{(2)}, \ldots)$ for some classical patterns $\tau^{(1)}, \tau^{(2)}, \ldots$

Recall that $\hat{\sigma}$ is the permutation obtained by swapping the first two entries in σ . Furthermore, Sort $(s_{\sigma}) = s_{\sigma}^{-1}(\operatorname{Av}(231))$ is the set of permutations that get sorted into the identity by the map $s \circ s_{\sigma}$. The following theorem is one of the main results in [11].

Theorem 2.1 ([11]). Let $\sigma \in S_k$ for some $k \ge 2$. The set $Sort(s_{\sigma})$ is a permutation class if and only if $\sigma = 12$ or $\hat{\sigma}$ contains the pattern 231 classically.

Our goal in this section is to prove the following analogue of Theorem 2.1 for the consecutivepattern-avoiding stack-sorting maps SC_{σ} . Recall that $Sort(SC_{\sigma}) = SC_{\sigma}^{-1}(Av(231))$ is the set of permutations that get sorted into the identity by the map $s \circ SC_{\sigma}$.

Theorem 2.2. The set $Sort(SC_{12})$ is equal to the permutation class Av(213), while $Sort(SC_{21})$ is not a permutation class. If $\sigma \in S_k$ for some $k \ge 3$, then the set $Sort(SC_{\sigma})$ is a permutation class if and only if σ and $\hat{\sigma}$ both contain the pattern 231 classically. In this case, $Sort(SC_{\sigma}) = Av(132)$.

We will prove Theorem 2.2 via the following two lemmas. Recall that $rev(\sigma)$ denotes the reverse of the permutation σ .

Lemma 2.3. Let $\sigma \in S_k$ for some $k \geq 3$. If $\widehat{\sigma}$ contains 231 classically, then $Sort(SC_{\sigma})$ is the set $Av(132, rev(\sigma))$ of permutations that avoid 132 classically and avoid $rev(\sigma)$ consecutively.

Proof. In general, it is immediate from the definition of SC_{σ} that $SC_{\sigma}(\pi) = rev(\pi)$ whenever π avoids $rev(\sigma)$ consecutively. Indeed, in this case, all of the entries in π enter the stack and then proceed to exit the stack in the reverse of the order in which they entered. It follows that if π avoids $rev(\sigma)$ consecutively and avoids 132 classically, then $SC_{\sigma}(\pi) = rev(\pi) \in Av(231)$. Furthermore, if π contains 132 classically and avoids $rev(\sigma)$ consecutively, then $SC_{\sigma}(\pi) = rev(\pi)$ contains 231 classically. This implies that $\pi \notin Sort(SC_{\sigma})$.

Suppose now that π contains rev (σ) consecutively, and let *i* be the minimal index such that $\pi_i \cdots \pi_{i+k-1}$ is a consecutive occurrence of the pattern rev (σ) . Consider sending π through a consecutive- σ -avoiding stack. When π_{i+k-1} is about to enter the stack, the top k-1 entries in the stack, read from top to bottom, are $\pi_{i+k-2}, \ldots, \pi_i$. The entry π_{i+k-2} pops out of the stack, and the

entry π_{i+k-1} enters the stack immediately afterward. Once all of the entries have finally exited the stack, the entries $\pi_{i+k-2}, \pi_{i+k-1}, \pi_{i+k-3}, \ldots, \pi_i$ will appear in this order in the output permutation $\mathrm{SC}_{\sigma}(\pi)$. However, $\pi_{i+k-2}\pi_{i+k-1}\pi_{i+k-3}\cdots\pi_i$ has the same relative order as $\hat{\sigma}$, which in turn contains 231 classically by hypothesis. This shows that $\mathrm{SC}_{\sigma}(\pi) \notin \mathrm{Av}(231)$, so $\pi \notin \mathrm{Sort}(\mathrm{SC}_{\sigma})$.

Lemma 2.4. If $\hat{\sigma} \in \operatorname{Av}_k(231)$ for some $k \geq 3$, then $\operatorname{Sort}(\operatorname{SC}_{\sigma})$ is not a permutation class.

Proof. We first dispose of the case in which $\sigma = 231$. A simple computation yields $SC_{231}(2413) = 3142 \notin Av(231)$ and $SC_{231}(25314) = 54132 \in Av(231)$. This shows that 25314 is in Sort(SC_{231}) while 2413 is not. Since 25314 classically contains 2413, it follows that $Sort(SC_{231})$ is not a permutation class. We may now assume $\sigma \neq 231$. Because the length of σ is at least 3, we have $SC_{\sigma}(132) = 231$. This shows that $132 \notin Sort(SC_{\sigma})$.

Let us now assume that σ classically contains 231 but is not equal to 231. One can easily check that $SC_{\sigma}(rev(\sigma)) = \hat{\sigma}$, and we assumed that $\hat{\sigma} \in Av(231)$. This shows that $rev(\sigma) \in Sort(SC_{\sigma})$. Now note that $rev(\sigma)$ contains 132 classically and that $132 \notin Sort(SC_{\sigma})$. It follows that $Sort(SC_{\sigma})$ is not a permutation class in this case.

Next, assume $\sigma \in \operatorname{Av}(231)$, $\sigma_1 < \sigma_2$, and $\sigma_2 > 2$. These assumptions force $\sigma_1 = 1$ since, otherwise, the entries $\sigma_1, \sigma_2, 1$ would form a classical 231 pattern in σ . Let $\pi = \sigma'_k \cdots \sigma'_2 \sigma'_1 \sigma_2$, where $\sigma'_i = \sigma_i$ if $\sigma_i < \sigma_2$ and $\sigma'_i = \sigma_i + 1$ if $\sigma_i \ge \sigma_2$. We will prove that $\pi \in \operatorname{Sort}(\operatorname{SC}_{\sigma})$. Since the entries $2, \sigma'_2, \sigma_2$ form an occurrence of the classical pattern 132 in π and $132 \notin \operatorname{Sort}(\operatorname{SC}_{\sigma})$, this will imply that $\operatorname{Sort}(\operatorname{SC}_{\sigma})$ is not a permutation class. One can easily compute that $\operatorname{SC}_{\sigma}(\pi) = \sigma'_2 \sigma_2 \sigma'_1 \sigma'_3 \cdots \sigma'_k$. Notice that the permutation $\sigma_2 \sigma'_1 \sigma'_3 \cdots \sigma'_k$ has the same relative order as $\hat{\sigma}$, which avoids 231 classically by hypothesis. Therefore, if $\operatorname{SC}_{\sigma}(\pi)$ contains 231, then the first entry in the occurrence of the 231 pattern must be σ'_2 . Since $\sigma'_2 = \sigma_2 + 1$, we can replace σ'_2 with σ_2 in this subsequence in order to obtain an occurrence of the pattern 231 in $\operatorname{SC}_{\sigma}(\pi)$ that starts with the entry σ_2 . However, this contradicts our earlier observation that $\sigma_2 \sigma'_1 \sigma'_3 \cdots \sigma'_k$ avoids 231.

For our next case, we assume $\sigma \in \operatorname{Av}(231)$, $\sigma_1 = 1$, and $\sigma_2 = 2$. Let $\pi = \sigma_k \cdots \sigma_3 21(k+1)$ be the concatenation of $\operatorname{rev}(\sigma)$ with the entry k+1. Let $\pi' = \sigma_k \cdots \sigma_3 2(k+1)1(k+2)$. The first 3 entries of the permutation $\operatorname{SC}_{\sigma}(\pi) = 2(k+1)1\sigma_3 \cdots \sigma_k$ form a 231 pattern, so $\pi \notin \operatorname{Sort}(\operatorname{SC}_{\sigma})$. On the other hand, since π' certainly avoids the pattern $\operatorname{rev}(\sigma)$ consecutively, we have $\operatorname{SC}_{\sigma}(\pi') = \operatorname{rev}(\pi') = (k+2)1(k+1)2\sigma_3 \cdots \sigma_k$. The assumption $\sigma \in \operatorname{Av}(231)$ guarantees that $\operatorname{SC}_{\sigma}(\pi') \in \operatorname{Av}(231)$, so $\pi' \in \operatorname{Sort}(\operatorname{SC}_{\sigma})$. The permutation π' contains π classically, so $\operatorname{Sort}(\operatorname{SC}_{\sigma})$ is not a permutation class in this case.

The final case we need to consider is that in which $\sigma \in \operatorname{Av}(231)$ and $\sigma_1 > \sigma_2$. Notice that the hypothesis $\hat{\sigma} \in \operatorname{Av}(231)$ forces $\sigma_2 = 1$ in this case. Let $\pi = \sigma'_k \cdots \sigma'_2 \sigma'_1 2$, where $\sigma'_2 = 1$ and $\sigma'_i = \sigma_i + 1$ for all $i \neq 2$. We have $\operatorname{SC}_{\sigma}(\pi) = 12\sigma'_1\sigma'_3\cdots\sigma'_k$; note that this permutation avoids 231 because σ avoids 231. Therefore, $\pi \in \operatorname{Sort}(\operatorname{SC}_{\sigma})$. The entries $\sigma'_2, \sigma'_1, 2$ form an occurrence of the pattern 132 in π , and we have seen that $132 \notin \operatorname{Sort}(\operatorname{SC}_{\sigma})$. Therefore, $\operatorname{Sort}(\operatorname{SC}_{\sigma})$ is not a permutation class in this final case.

Proof of Theorem 2.2. It is straightforward to check that if $\pi \in S_n$ for some $n \geq 1$, then the last entry in $SC_{12}(\pi)$ is 1. It follows that $SC_{12}(\pi)$ avoids 231 if and only if it avoids 12. Using Lemma 1.2, we find that

$$Sort(SC_{12}) = SC_{12}^{-1}(Av(231)) = SC_{12}^{-1}(Av(12)) = comp(SC_{21}^{-1}(comp(Av(12))))$$
$$= comp(s^{-1}(Av(21))) = comp(Av(231)) = Av(213).$$

On the other hand, the set $Sort(SC_{21})$ is not a permutation class because it includes the permutation 35241 but not the permutation 3241.

Now assume $\sigma \in S_k$ for some $k \geq 3$. If σ and $\hat{\sigma}$ both contain 231 classically, then it follows from Lemma 2.3 that

$$\operatorname{Sort}(\operatorname{SC}_{\sigma}) = \operatorname{Av}\left(132, \underline{\operatorname{rev}(\sigma)}\right) = \operatorname{Av}(132).$$

If $\hat{\sigma} \in \operatorname{Av}(231)$, then Lemma 2.4 tells us that $\operatorname{Sort}(\operatorname{SC}_{\sigma})$ is not a permutation class. Now assume $\sigma \in \operatorname{Av}(231)$ and $\hat{\sigma} \notin \operatorname{Av}(231)$. Note that this implies that $\sigma_1 > \sigma_2 > 1$. Since $\operatorname{SC}_{\sigma}(\operatorname{rev}(\sigma)) = \hat{\sigma}$, the permutation $\operatorname{rev}(\sigma)$ is not in $\operatorname{Sort}(\operatorname{SC}_{\sigma})$. However, the permutation $\pi = (\sigma_k + 1) \cdots (\sigma_2 + 1)1(\sigma_1 + 1)$, which contains $\operatorname{rev}(\sigma)$ classically, is in $\operatorname{Sort}(\operatorname{SC}_{\sigma})$. To see this, note that the condition $\sigma_1 > \sigma_2 > 1$ guarantees that π avoids $\operatorname{rev}(\sigma)$ consecutively. Therefore, $\operatorname{SC}_{\sigma}(\pi) = 1(\sigma_1 + 1)(\sigma_2 + 1) \cdots (\sigma_k + 1)$, and this permutation avoids 231 because σ does. Hence, $\operatorname{Sort}(\operatorname{SC}_{\sigma})$ is not a permutation class. \Box

3. Periodic Points

In this section, we prove the following theorem, which completely classifies the periodic points of the consecutive- σ -avoiding stack-sorting map for $\sigma \in S_3$.

Theorem 3.1. Let $\sigma \in S_3$. The periodic points of the map $SC_{\sigma} : S_n \to S_n$ are precisely the permutations in $Av_n\left(\underline{\sigma}, \underline{rev}(\sigma)\right)$. When $n \ge 2$, each of these periodic points has period 2.

Lemma 1.2 tells us that for each permutation pattern σ , the maps SC_{σ} and $SC_{comp(\sigma)}$ are topologically conjugate (with the discrete topology), with the complementation map comp serving as the topological conjugacy. It follows that if Theorem 3.1 holds when σ is some specific pattern τ , then it also holds when $\sigma = comp(\tau)$. Thus, we really only need to prove the theorem when $\sigma \in \{123, 132, 231\}$. We first handle the case in which $\sigma \in \{132, 231\}$.

Proposition 3.2. For $\sigma \in \{132, 231\}$, the set of periodic points of the map $SC_{\sigma} : S_n \to S_n$ is the set $Av_n(\underline{132}, \underline{231})$ of permutations in S_n that avoid 132 and 231 consecutively (these permutations also avoid 132 and 231 classically). When $n \geq 2$, these points have period 2.

Proof. Note that $\{132, 231\} = \{\sigma, \operatorname{rev}(\sigma)\}$. It is clear that if $\pi \in \operatorname{Av}(\underline{132}, \underline{231})$, then $\operatorname{SC}_{\sigma}(\pi) = \operatorname{rev}(\pi) \in \operatorname{Av}(\underline{132}, \underline{231})$, so $\operatorname{SC}_{\sigma}^2(\pi) = \pi$. Hence, π is a periodic point of $\operatorname{SC}_{\sigma}$, and the period is 2 if $n \geq 2$.

We now prove that every periodic point of $\mathrm{SC}_{\sigma} : S_n \to S_n$ is in $\mathrm{Av}(\underline{132},\underline{231})$. This is trivial if $n \leq 2$, so we may assume $n \geq 3$ and induct on n. Given $\tau \in S_n$, let τ^* be the permutation in S_{n-1} obtained by deleting the entry 1 from τ and then decreasing all remaining entries by 1. Observe that if the entries 1 and 2 appear consecutively in τ , then they also appear consecutively in $\mathrm{SC}_{\sigma}(\tau)$ and that $(\mathrm{SC}_{\sigma}(\tau))^* = \mathrm{SC}_{\sigma}(\tau^*)$. Now fix $\pi \in S_n$. By induction, we have $\mathrm{SC}_{\sigma}^t(\pi^*) \in \mathrm{Av}(\underline{132},\underline{231})$ for all sufficiently large t. We will prove that if t is sufficiently large, then the entries 1 and 2 appear consecutively in $\mathrm{SC}_{\sigma}^t(\pi)$. It will then follow that $(\mathrm{SC}_{\sigma}^t(\pi))^* = \mathrm{SC}_{\sigma}^t(\pi^*) \in \mathrm{Av}(\underline{132},\underline{231})$ for all sufficiently large t, which will imply that $\mathrm{SC}_{\sigma}^t(\pi) \in \mathrm{Av}(\underline{132},\underline{231})$ for all sufficiently large t.

Let $D(\tau)$ denote the number of entries between 1 and 2 in τ (regardless of whether 1 appears before or after 2 in τ). First, note that neither 1 nor 2 will get popped from the stack until after all of the entries of τ have entered the stack because neither 1 nor 2 can trigger the consecutive- σ avoiding restriction when at the top of the stack. This immediately implies that $D(\tau) \leq D(SC_{\sigma}(\tau))$. We further show that if $D(\tau) > 0$, then $D(SC_{\sigma}^2(\tau)) < D(\tau)$, implying that 1 and 2 are consecutive entries in $\operatorname{SC}_{\sigma}^{t}(\pi)$ for all sufficiently large t. Let $\{1,2\} = \{\tau_{i}, \tau_{i+c}\}$ so that $D(\pi) = c-1$, and assume $c \geq 2$. The consecutive subsequence $\tau_{i}\tau_{i+1}\cdots\tau_{i+c}$ of τ contains a consecutive occurrence of either σ or $\operatorname{rev}(\sigma)$. If it contains a consecutive occurrence of $\operatorname{rev}(\sigma)$, then when τ_{i+c} enters the stack, the entries between τ_{i} and τ_{i+c} form a strict subset of $\{\tau_{i+1},\ldots,\tau_{i+c-1}\}$. Moreover, these entries are precisely the entries between 1 and 2 in $\operatorname{SC}_{\sigma}(\tau)$ since neither 1 nor 2 is popped until all entries of τ have entered the stack. This implies that $D(\operatorname{SC}_{\sigma}(\tau)) < D(\tau)$. If there is no consecutive occurrence of $\operatorname{rev}(\sigma)$ in the consecutive subsequence $\tau_{i}\tau_{i+1}\cdots\tau_{i+c}$ of τ , then $\tau_{i+c}\tau_{i+c-1}\cdots\tau_{i}$ is a consecutive subsequence that $D(\operatorname{SC}_{\sigma}(\tau)) < D(\tau)$.

Next, we consider the periodic points of SC_{123} and SC_{321} . By complementing, it suffices to just consider SC_{321} . To understand this map, it helps to decompose a permutation π by writing $\pi = a_1 \cdots a_k$, where the a_i 's are the ascending runs (maximal consecutive increasing subsequences). In addition, let a_i^m be the subsequence of a_i consisting of the entries that are not the first or last entries in a_i , and let a_i^e be the sequence obtained from a_i by deleting a_i^m . If $|a_i|$ is 1 or 2, then $a_i^e = a_i$ and a_i^m is an empty sequence. For example, the ascending runs of 4572136 are $a_1 = 457$, $a_2 = 2$, and $a_3 = 136$, with $a_1^e = 47$, $a_2^e = 2$, $a_3^e = 16$, $a_1^m = 5$, a_2^m empty, and $a_3^m = 3$.

Lemma 3.3. Keeping the notation above, let $\pi = a_1 \cdots a_k$, where the a_i 's are ascending runs. Then $SC_{321}(\pi) = a_1^m \cdots a_k^m \operatorname{rev}(a_1^e \cdots a_k^e) = a_1^m \cdots a_k^m \operatorname{rev}(a_k^e) \cdots \operatorname{rev}(a_1^e)$.

Proof. When the ascending run a_i enters the consecutive-321-avoiding stack, each entry in a^m is popped immediately after it enters while a_i^e remains in the stack. After all entries have been pushed into the stack, the entries remaining in the stack are popped out in the reverse of the order in which they entered and thus the output is $a_1^m \cdots a_k^m \operatorname{rev}(a_1^e \cdots a_k^e) = a_1^m \cdots a_k^m \operatorname{rev}(a_k^e) \cdots \operatorname{rev}(a_1^e)$.

To identify the periodic points of SC₃₂₁, we analyze how the number of peaks and valleys of a permutation π changes as SC₃₂₁ is iteratively applied to π . A *peak* (respectively, *valley*) of $\pi = \pi_1 \cdots \pi_n$ is an index $i \in \{2, \ldots, n-1\}$ such that $\pi_{i-1} < \pi_i > \pi_{i+1}$ (respectively, $\pi_{i-1} > \pi_i < \pi_{i+1}$). If i is a peak (respectively, valley) of π , then the entry π_i is called a *peak top* (respectively, *valley valley bottom*) of π . We let $q(\pi)$ denote the total number of indices that are either peaks or valleys of π .

Lemma 3.4. Let $\pi = a_1 \cdots a_k \in S_n$, where a_1, \ldots, a_k are the ascending runs of π . We have $g(\pi) \leq g(\mathrm{SC}_{321}(\pi))$. Moreover, if $|a_k| \geq 2$, then this inequality is an equality if and only if $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e)$ is a decreasing sequence. If $|a_k| = 1$, then $g(\pi) = g(\mathrm{SC}_{321}(\pi))$ if and only if $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e)$ is an increasing sequence.

Proof. We specifically show that the peak tops and valley bottoms of π are still peak tops and valley bottoms in $SC_{321}(\pi)$. The key observation is that each peak top (valley bottom) of π is the second (first) entry of a_j^e for some ascending run a_j of length at least 2, and it is not the first or last entry of π .

Suppose the entry π_i is a peak top of π . Then π_i is the last entry of some ascending run a_j of length at least 2. We have $a_j^e = \pi_h \pi_i$ for some $\pi_h < \pi_i$. Similarly, π_{i+1} is the first entry of the ascending run a_{j+1} and is less than π_i . By Lemma 3.3, the sequence rev (a_{j+1}^e) rev (a_j^e) , which itself has $\pi_{i+1}\pi_i\pi_h$ as a consecutive subsequence, is a consecutive subsequence of $SC_{321}(\pi)$. Consequently, π_i is still a peak top of $SC_{321}(\pi)$. A completely analogous argument shows that every valley bottom of π is also a valley bottom of $SC_{321}(\pi)$, so $g(\pi) \leq g(SC_{321}(\pi))$. For the equality cases, suppose first that a_k has length at least 2. Note that $g(\pi) = g(\mathrm{SC}_{321}(\pi))$ if and only if no entry in $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e)$ except for the last entry is a peak top or valley bottom in $\mathrm{SC}_{321}(\pi)$. This happens if and only if $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e)$ is an increasing or decreasing sequence. Since $\operatorname{rev}(a_k^e)$ is a decreasing sequence, it follows that $g(\pi) = g(\mathrm{SC}_{321}(\pi))$ if and only if $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e)$ is a decreasing sequence.

Now suppose a_k has length 1. We again have $g(\pi) = g(\operatorname{SC}_{321}(\pi))$ if and only if no element in $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e)$ is a peak top or valley bottom in $\operatorname{SC}_{321}(\pi)$. Note that $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e)$ is part of the consecutive subsequence $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e) \operatorname{rev}(a_{k-1}^e)$ in $\operatorname{SC}_{321}(\pi)$. Since a_k has length 1, it is a single entry that is less than the entry before it in π ; this entry before it is also the entry after it in $a_1^m \cdots a_k^m \operatorname{rev}(a_{k-1}^e)$. Thus, no entry in $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e)$ is a peak top or valley bottom if and only if $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e)$ is an increasing sequence.

By the previous lemma, it is clear that if $g(\pi) < g(SC_{321}(\pi))$, then π cannot be a periodic point. We make use of this fact when proving the next few results.

Corollary 3.5. Let $\pi = a_1 \cdots a_k$, where the a_i 's denote ascending runs. If $|a_k| \ge 3$ or $|a_1| = |a_2| = 1$, then π is not a periodic point of SC₃₂₁.

Proof. If $|a_k| \geq 3$, then $a_k^m \operatorname{rev}(a_k^e)$ is not a decreasing sequence, so $g(\pi) < g(\operatorname{SC}_{321}(\pi))$ by Lemma 3.4. Hence, π is not a periodic point of SC_{321} . If $|a_1| = |a_2| = 1$, then $\operatorname{SC}_{321}(\pi)$ is a permutation ending in an ascending run of length at least 3, so it follows from the preceding argument that $\operatorname{SC}_{321}(\pi)$ is not a periodic point of SC_{321} . Hence, π is not a periodic point in this case either.

Corollary 3.6. Let $\pi = a_1 \cdots a_k \in S_n$ for some $n \ge 3$, where the a_i 's denote ascending runs. If $|a_1| \ge 3$ or $|a_{k-1}| = |a_k| = 1$, then π is not a periodic point of SC₃₂₁.

Proof. Suppose π is a periodic point of SC_{321} and $|a_1| \geq 2$. Let ℓ be the period of π so that $\mathrm{SC}_{321}^{\ell}(\pi) = \pi$. Let π_1 denote the first entry of π . By Lemma 3.4, we must have $g(\pi) = g(\mathrm{SC}_{321}^t(\pi))$ for all $t \geq 1$. By Lemma 3.3, the last two entries of $\mathrm{SC}_{321}(\pi)$ form the sequence $\mathrm{rev}(a_1^e)$; in particular, the last ascending run of $\mathrm{SC}_{321}(\pi)$ has length 1 and is simply π_1 . Since $\mathrm{SC}_{321}(\pi)$ ends in an ascending run of length 1 and $g(\pi) = g(\mathrm{SC}_{321}(\pi))$, Lemma 3.4 guarantees that $\mathrm{SC}_{321}^2(\pi)$ begins with an ascending run that contains π_1 , so $\mathrm{SC}_{321}^2(\pi)_1 \leq \pi_1$. In fact, the first ascending run of $\mathrm{SC}_{321}^2(\pi)$ must actually contain a_1^e , so $\mathrm{SC}_{321}^2(\pi)$ has an ascending run of length at least 2. Moreover, if $\mathrm{SC}_{321}(\pi)$ has an ascending run of length at least 2. Moreover, if $\mathrm{SC}_{321}(\pi)$ has an ascending run of length at least 3, then the first ascending run of $\mathrm{SC}_{321}^2(\pi)$ actually begins with an entry that is smaller than π_1 , implying that $\mathrm{SC}_{321}^2(\pi)_1 < \pi_1$. Repeating this argument shows that ($\mathrm{SC}_{321}^{2t}(\pi)_{1}_{1,t \geq 1}$ is a weakly decreasing sequence and is nonconstant if $\mathrm{SC}_{321}^{2t-1}(\pi)$ contains an ascending run of length at least 3 for some $t \geq 1$. Since $\mathrm{SC}_{321}^{k\ell}(\pi) = \pi$ for all integers $k \geq 1$, it follows that for every $t \geq 1$, the permutation $\mathrm{SC}_{321}^{2t-1}(\pi)$ does not have an ascending run of length at least 3. This implies that $\mathrm{SC}_{321}^{2t-1}(\pi)$ for all $t \geq 1$.

Now consider π_2 , the second entry of π . As previously stated, $SC_{321}(\pi)$ ends in $rev(a_1^e)$, so the second-to-last entry of $SC_{321}(\pi)$ is greater than or equal to π_2 . This implies that the second entry of $SC_{321}^2(\pi)$ is greater than or equal to π_2 , as $SC_{321}^2(\pi) = rev(SC_{321}(\pi))$. Repeating this argument shows that the sequence $(SC_{321}^{2t}(\pi)_2)_{t\geq 1}$ is weakly increasing. As $SC_{321}^{2k\ell}(\pi) = \pi$ for all $k \geq 1$, this sequence must actually be constant. If $|a_1| \geq 3$, then the second-to-last entry of $SC_{321}(\pi)$, which is also the second entry of $SC_{321}^2(\pi)$, is actually strictly greater than π_2 . This is a contradiction.

Now suppose π is a periodic point and $|a_{k-1}| = |a_k| = 1$. This means we must have $g(\pi) = g(\operatorname{SC}_{321}(\pi))$, so it follows from Lemma 3.4 that $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e)$ is an increasing sequence. Notice also that the first 3 entries of the sequence $\operatorname{rev}(a_k^e) \operatorname{rev}(a_{k-1}^e) \operatorname{rev}(a_{k-2}^e)$ appear in increasing order (we must have $k \geq 3$ since $n \geq 3$ and $|a_{k-1}| = |a_k| = 1$). This shows that $\operatorname{SC}_{321}(\pi)$ is a periodic point beginning with an ascending run of length at least 3, which contradicts what we found in the previous two paragraphs.

The following proposition proves Theorem 3.1 for $\sigma = 321$, thereby completing the proof of the entire theorem. Indeed, we already proved the theorem for $\sigma \in \{132, 213, 231, 312\}$, and we can use Lemma 1.2 to see that the next proposition implies Theorem 3.1 for $\sigma = 123$ as well.

Proposition 3.7. The set of periodic points of $SC_{321} : S_n \to S_n$ is $Av_n(\underline{123}, \underline{321})$, the set of permutations in S_n that avoid 123 and 321 consecutively. When $n \ge 2$, these points have period 2.

Proof. It is clear that if $\pi \in Av(\underline{123}, \underline{321})$, then $SC_{321}(\pi) = rev(\pi) \in Av(\underline{123}, \underline{321})$, so $SC_{321}^2(\pi) = \pi$. Hence, π is a periodic point of SC_{321} , and the period is 2 if $n \ge 2$.

Suppose that $\pi \notin \operatorname{Av}(\underline{123}, \underline{321})$; we will show that π is not a periodic point of SC_{321} . If π does not contain a consecutive 321 pattern but does contain a consecutive 123 pattern, then $\operatorname{SC}_{321}(\pi) = \operatorname{rev}(\pi)$ contains a consecutive 321, so we may simply assume that $\pi \notin \operatorname{Av}(\underline{321})$ (otherwise, the same analysis applies to $\operatorname{SC}_{321}(\pi)$). Suppose instead that π is a periodic point, and let a_1, \ldots, a_k be its ascending runs. Note that we have $g(\pi) = g(\operatorname{SC}_{321}(\pi))$. Furthermore, the sequence $a_1^m \cdots a_k^m$ is nonempty because π contains a consecutive 321 pattern. It then follows from Lemma 3.3 that $\operatorname{SC}_{321}(\pi)$ is a periodic point of SC_{321} that begins with the sequence $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e)$. If $|a_k| \geq 2$, then $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e)$ has length at least 3, and Lemma 3.4 tells us that this sequence is decreasing. This implies that the periodic point $\operatorname{SC}_{321}(\pi)$ starts with a decreasing sequence of length at least 3, contradicting Corollary 3.5. Hence, we must have $|a_k| = 1$. Consider π_{n-1} , which is last entry in a_{k-1} . We must have $\pi_{n-1} > \pi_n$, so it follows from Lemma 3.4 that $a_1^m \cdots a_k^m \operatorname{rev}(a_k^e)\pi_{n-1}$ is an increasing sequence of length at least 3. This shows that $\operatorname{SC}_{321}(\pi)$ is a periodic point whose first ascending run has length at least 3, which contradicts Corollary 3.6.

4. The Maps SC_{132} and SC_{312}

In this section, we explore the maps SC_{132} and SC_{312} from both a sorting point of view and a dynamical point of view. We do not have much to say about SC_{312} from a sorting point of view because we were not able to enumerate the set $Sort(SC_{312})$. The initial values of the sequence $(|Sort_n(SC_{312})|)_{n\geq 0}$ are

1, 1, 2, 5, 15, 50, 179, 675, 2649, 10734;

this sequence appears to be new. However, we *will* be able to characterize and enumerate the permutations in $Sort(SC_{132})$. Specifically, we show that they are in bijection with Dyck paths and are thus enumerated by the Catalan numbers.

From a dynamical point of view, we will be interested in the number of preimages of permutations under SC₁₃₂ and SC₃₁₂. In this setup, it suffices to focus our attention on SC₁₃₂. Indeed, since 312 = comp(132), Lemma 1.2 implies that $|\text{SC}_{312}^{-1}(\pi)| = |\text{SC}_{132}^{-1}(\text{comp}(\pi))|$ for all $\pi \in S_n$. Consequently, we will not need to discuss the map SC₃₁₂ directly in this section. Let us begin with the set Sort(SC₁₃₂). Recall that a *left-to-right minimum* of a permutation $\pi = \pi_1 \cdots \pi_n$ is an entry π_i such that $\pi_i < \pi_j$ for all $1 \le j < i$. For example, the left-to-right minima of 4572163 are 4, 2, and 1.

Proposition 4.1. Let $\pi = a_1 \cdots a_k \in S_n$, where the subsequences a_i are the ascending runs of π . For each *i*, let us write $a_i = m_i t_i$, where m_i is the first entry of a_i and t_i is the (possibly empty) string obtained by removing m_i from a_i . Then $SC_{132}(\pi) \in Av(231)$ if and only if each m_i is a left-to-right minimum of π and the concatenation $rev(t_1) \cdots rev(t_k)$ is a decreasing sequence.

Proof. Let us first assume $\mathrm{SC}_{132}(\pi) \in \mathrm{Av}(231)$. We will prove by induction that the entries m_1, \ldots, m_k are all left-to-right minima of π . Since m_1 is the first entry of π , it is trivially a left-to-right minimum. Now consider $1 \leq i < k$, and suppose that m_1, \ldots, m_i are all left-to-right minimum. This implies that if we write $t_i = b_1 \cdots b_j$, where $j = |t_i|$, then $j \geq 1$ and $m_i < m_{i+1} < b_j$. Let ℓ be the smallest index such that $b_\ell > m_{i+1}$. When we send π through the consecutive-132-avoiding stack, the consecutive subsequence a_i will enter the stack consecutively, so $\mathrm{rev}(a_i) = b_j \cdots b_1 m_i$ will be at the top of the stack when m_{i+1} is next in line to be pushed into the stack. The entries $b_j, \ldots, b_{\ell+1}$ will be popped, and then m_{i+1} will be pushed into the stack on top of b_ℓ . Since $b_\ell > m_{i+1}$, the entry m_{i+1} will not get popped out of the stack until after all of the entries of π have entered the stack. It follows that m_{i+1}, b_ℓ, m_i form an occurrence of the classical pattern 231 in $\mathrm{SC}_{132}(\pi)$; this is our desired contradiction. It is straightforward to see that $\mathrm{SC}_{132}(a_1 \cdots a_k) = \mathrm{rev}(t_1) \cdots \mathrm{rev}(t_k) m_k \cdots m_1$, where $m_k = 1$. Since we are assuming $\mathrm{SC}_{132}(a_1 \cdots a_k)$ avoids 231, the sequence $\mathrm{rev}(t_1) \cdots \mathrm{rev}(t_k)$ must be decreasing.

For the reverse direction, suppose that the entries m_1, \ldots, m_k are left-to-right minima and that the sequence $\operatorname{rev}(t_1) \cdots \operatorname{rev}(t_k)$ is decreasing. As above, $\operatorname{SC}_{132}(\pi) = \operatorname{rev}(t_1) \cdots \operatorname{rev}(t_k) m_k \cdots m_1$, which is a decreasing sequence followed by an increasing sequence. Such a permutation necessarily avoids 231.

Remark 4.2. The proof of Proposition 4.1 yields the somewhat-unexpected fact that if a permutation is in the image of SC_{132} and avoids the pattern 231, then it also avoids 132.

Proposition 4.1 tells us that each $\pi \in \text{Sort}(\text{SC}_{132})$ is determined by the first entries m_i of its ascending runs as well as the lengths $|a_1|, \ldots, |a_k|$ of its ascending runs. Furthermore, these quantities must satisfy $n + 1 - m_i \ge |a_1| + \cdots + |a_i|$ for all *i*. Indeed, the entries in t_1 must be the $|a_1| - 1$ largest elements of $[n] \setminus \{m_1, \ldots, m_k\}$, the entries in t_2 must be the next $|a_2| - 1$ largest elements, and so on. The condition that $n + 1 - m_i \ge |a_1| + \cdots + |a_i|$ for all *i* guarantees that the m_i 's are left-to-right minima.

A Dyck path of semilength n is a word over the alphabet $\{U, D\}$ that contains n copies of each letter and has the additional property that each of its prefixes contains at least as many occurrences of U as occurrences of D. Let \mathbf{D}_n denote the set of Dyck paths of semilength n. It is well known that $|\mathbf{D}_n| = C_n = \frac{1}{n+1} {\binom{2n}{n}}$ is the nth Catalan number. Given $\pi \in S_n$ such that $\mathrm{SC}_{132}(\pi) \in \mathrm{Av}(231)$, write $\pi = a_1 \cdots a_k$ and $a_i = m_i t_i$ as in Proposition 4.1. Let $\Phi(\pi)$ be the Dyck path $U^{m_0-m_1}D^{|a_1|}U^{m_1-m_2}D^{|a_2|}\cdots U^{m_{k-1}-m_k}D^{|a_k|}$, where we make the convention $m_0 = n + 1$.

Example 4.3. Let $\pi = 589436712$, and note that $SC_{132}(\pi) = 987621345 \in Av(231)$. We have $a_1 = 589, a_2 = 4, a_3 = 367$, and $a_4 = 12$. The first entries of these ascending runs are $m_1 = 5$, $m_2 = 4, m_3 = 3$, and $m_4 = 1$. Therefore, $\Phi(\pi) = UUUUUDDDUUDDDUUDDDUUDD$.

Theorem 4.4. The map Φ : Sort_n(SC₁₃₂) \rightarrow **D**_n is a bijection. Consequently,

$$|\operatorname{Sort}_n(\operatorname{SC}_{132})| = C_n$$

Proof. Preserve the notation from above. The total number of occurrences of the letter U in $\Phi(\pi)$ is $(m_0 - m_1) + \cdots + (m_{k-1} - m_k) = m_0 - m_k = (n+1) - 1 = n$, and the total number of occurrences of D is $|a_1| + \cdots + |a_k| = n$. As mentioned above, we have $n + 1 - m_i \ge |a_1| + \cdots + |a_i|$ for all i; this guarantees that $\Phi(\pi)$ is indeed a Dyck path of semilength n. Moreover, by the previous discussion, Φ is injective.

To prove surjectivity, suppose we are given a Dyck path $\Lambda = U^{\gamma_1} D^{\delta_1} U^{\gamma_2} D^{\delta_2} \cdots U^{\gamma_k} D^{\delta_k}$ for some positive integers $\gamma_1, \ldots, \gamma_k, \delta_1, \ldots, \delta_k$. Define $m_i = n + 1 - \gamma_1 - \cdots - \gamma_i$. Using the fact that Λ is a Dyck path, one can verify that there exists a permutation $\pi \in S_n$ with ascending runs a_1, \ldots, a_k satisfying the conditions in Proposition 4.1 with $a_i = m_i t_i$ and $|a_i| = \delta_i$ for all i. This permutation is in Sort_n(SC₁₃₂) and satisfies $\Phi(\pi) = \Lambda$.

The next series of results determines the maximum number of preimages of permutations under SC_{132} and SC_{312} . As mentioned above, it is only necessary to consider SC_{132} .

Lemma 4.5. Let $\pi \in S_n$. Suppose there exists an entry $i \in [n-1]$ that occurs before i+1 but not immediately before i+1 in π . Let π' be the permutation obtained by swapping i and i+1 in π . Then,

$$|\operatorname{SC}_{132}^{-1}(\pi)| \le |\operatorname{SC}_{132}^{-1}(\pi')|.$$

Proof. Suppose $SC_{132}(\tau) = \pi$, and let τ' be the permutation obtained by swapping *i* and *i* + 1 in τ . We will prove that $SC_{132}(\tau') = \pi'$. To do so, it suffices to show that the sequence of "push" and "pop" operations that is used to send τ through the consecutive-132-avoiding stack is exactly the same as the sequence used to send τ' through the stack.

When sending τ through the consecutive-132-avoiding stack, the entries *i* and *i* + 1 are never consecutive entries in the stack. Indeed, if they were, then *i* would sit on top of *i* + 1 because *i* appears before *i* + 1 in the output permutation π . However, if *i* sits on top of *i* + 1, then these two entries will appear consecutively in the output π , contrary to our hypothesis. Now consider the sequences of "push" and "pop" operations that send τ and τ' through the stack. Note that in each sequence, the first two operations are necessarily pushes. Now suppose the first two sequences agree in their first *k* operations. Suppose further that after the first *k* operations, the stack-sorting procedure that is being applied to τ (respectively, τ') has *b* (respectively, *b'*) as the top entry in the stack, has *c* (respectively, *c'*) as the second-to-top entry in the stack, and has *a* (respectively, *a'*) as the entry next in line to enter the stack. We show that the next operation is the same as well.

If none of a, b, and c are i or i + 1, then it must be that a = a', b = b', and c = c', so the claim is trivially true. Next, if exactly one of a, b, and c is i or i + 1, then abc and a'b'c' have the same relative order, so the claim is true is in this case as well. Finally, suppose exactly two of a, b, c are i or i + 1. Because i and i + 1 cannot appear consecutively in the stack, we cannot have $\{b, c\} = \{i, i + 1\}$. If $\{a, b\} = \{i, i + 1\}$, then one of the two sorting procedures (sorting τ or τ') would have i and i + 1 appear consecutively in the stack, which is again impossible. Therefore, it must be that (a, c) = (i, i + 1) or (a', c') = (i, i + 1). The two cases are symmetric, so we may assume the former. Note that abc cannot have the same relative order as 132 as, otherwise, i and i + 1 would end up being consecutive in the stack once b gets popped out and c = i + 1 gets pushed in. Thus, the next operation for the τ case is to push a into the stack. Likewise a' = i + 1 and c' = i, so a'b'c' cannot have the same relative order as 132. The next operation in the τ' case is therefore to push a' into the stack. By induction, the sequences of operations when sending τ and τ' through the consecutive-132-avoiding stack are identical.

Since the map $\tau \mapsto \tau'$ is one-to-one, this shows that $|\operatorname{SC}_{132}^{-1}(\pi)| \leq |\operatorname{SC}_{132}^{-1}(\pi')|$, as desired. \Box

Corollary 4.6. For every positive integer n, we have

$$\max_{\pi \in \operatorname{Av}_n(132,213)} |\operatorname{SC}_{132}^{-1}(\pi)| = \max_{\pi \in S_n} |\operatorname{SC}_{132}^{-1}(\pi)|.$$

Proof. Suppose $\pi \in S_n$. If π classically contains an occurrence of either 132 or 213, then there exists some $i \in [n-1]$ such that i and i+1 are not consecutive in π and i occurs before i+1. Lemma 4.5 tells us that swapping these entries yields a permutation with at least as many preimages under SC₁₃₂. By repeatedly performing such swaps, we must eventually reach a permutation in Av_n(132, 213) (the sequence of swaps must terminate because each swap decreases the number of inversions in the permutation by 1), and this permutation has at least as many preimages under SC₁₃₂ as π .

A permutation in S_n is called *reverse-layered* if the set of entries in each of its ascending runs forms an interval of consecutive integers. For example, 5674231 is reverse-layered because the sets of entries in its ascending runs are $\{5, 6, 7\}$, $\{4\}$, $\{2, 3\}$, and $\{1\}$. It is well known [30] that $Av_n(132, 213)$ is precisely the set of reverse-layered permutations in S_n . Note that each reverselayered permutation is uniquely determined by its set of descents, which we can encode via a word of length n over the alphabet $\{A, D\}$ that starts with A. More precisely, the reverse-layered permutation $\pi = \pi_1 \cdots \pi_n$ corresponds to the word $b_1 \cdots b_n$, where $b_1 = A$ and for $2 \le j \le n$, we have $b_j = A$ if $\pi_{j-1} < \pi_j$ and $b_j = D$ if $\pi_{j-1} > \pi_j$. Let $d_{\pi}(k)$ denote the number of occurrences of Din the subword $b_1 \cdots b_k$. Let $a_{\pi}(k)$ denote the number of occurrences of A in the word $b_{k+1} \cdots b_n$ that appear to the right of an occurrence of the letter D in $b_{k+1} \cdots b_n$. For example, the reverse-layered permutation $\pi = 78634512$ corresponds to the word AADDAADA. We have

 $(d_{\pi}(0), \dots, d_{\pi}(8)) = (0, 0, 0, 1, 2, 2, 2, 3, 3)$ and $(a_{\pi}(0), \dots, a_{\pi}(8)) = (3, 3, 3, 3, 1, 1, 1, 0, 0).$

Using this notation, we can obtain a nice formula for $|SC_{132}^{-1}(\pi)|$ when π is reverse-layered. The reader may find it helpful to refer to Example 4.8 while reading the following proof.

Theorem 4.7. Let $\pi \in Av_n(132, 213)$ be a reverse-layered permutation in S_n , and preserve the definitions of $d_{\pi}(k)$ and $a_{\pi}(k)$ from above. We have

$$|\operatorname{SC}_{132}^{-1}(\pi)| = \sum_{k=0}^{n} \binom{d_{\pi}(k) + a_{\pi}(k)}{k}.$$

Proof. Let $SC_{132,k}^{-1}(\pi)$ denote the set of permutations $\tau \in SC_{132}^{-1}(\pi)$ such that when τ is sent through the consecutive-132-avoiding stack, exactly k entries get popped out of the stack before all entries have entered the stack. We will show that

$$|\operatorname{SC}_{132,k}^{-1}(\pi)| = \binom{d_{\pi}(k) + a_{\pi}(k)}{k},$$

which will complete the proof. Note that $|\operatorname{SC}_{132,0}^{-1}(\pi)| = 1$ because the only element of $\operatorname{SC}_{132,0}^{-1}(\pi)$ is $\operatorname{rev}(\pi)$. Indeed, SC_{132} does actually send $\operatorname{rev}(\pi)$ to π because π is reverse-layered (hence, $\operatorname{rev}(\pi)$ avoids 231). Therefore, we may assume that $k \geq 1$ in what follows.

Suppose $\tau \in \mathrm{SC}_{132,k}^{-1}(\pi)$. When we send τ through the consecutive-132-avoiding stack, the k entries popped before all entries have entered the stack must be exactly π_1, \ldots, π_k , and they must be popped in that order. The other entries π_{k+1}, \ldots, π_n get popped out of the stack later in this order, so they must appear in τ in the order π_n, \ldots, π_{k+1} . Thus, when constructing such a preimage τ , we first insert each entry π_i with $1 \leq i \leq k$ between a pair of entries $\pi_{j(i)}, \pi_{j(i)+1}$ with

 $k+1 \leq j(i) \leq n$ and then reverse the entire permutation. When we do so, the following conditions must be satisfied:

- the string $\pi_{i(i)}\pi_i\pi_{i(i)+1}$ has the same relative order as 132;
- $j(1) \geq \cdots \geq j(k);$
- if j(i) = j(i+1), then $\pi_i > \pi_{i+1}$ and $\pi_{i+1}\pi_i$ is a consecutive subsequence of τ .

The first and third conditions ensure that the entries π_1, \ldots, π_k actually get popped before all entries have entered the stack; the second and third conditions ensure that they are popped in the correct order. Moreover, each placement of the entries π_1, \ldots, π_k between consecutive entries in $\pi_{k+1} \cdots \pi_n$ such that the three conditions are satisfied yields a unique preimage $\tau \in \mathrm{SC}_{132,k}^{-1}(\pi)$.

To determine the number of valid placements, let π_{ℓ} be the last entry in the ascending run of π that includes π_k . When we insert the entry π_k between the entries $\pi_{j(k)}, \pi_{j(k)+1}$, the first condition above guarantees that $j(k) \geq \ell + 1$, so it follows from the second condition that $j(1) \geq \cdots \geq j(k) \geq \ell + 1$. Thus, for each $i \in [k]$, the possible choices for the index j(i) are precisely the ascents counted by $a_{\pi}(k)$. To satisfy the third condition, note that we can have j(i) = j(i+1) only if i is a descent of π . In summary, this proves that the number of ways to choose a preimages τ is equal to the number of ways to choose the indices $j(1), \ldots, j(k)$ from among the ascents counted by $a_{\pi}(k)$ such that $j(1) \geq \cdots \geq j(k)$, where the inequality $j(i) \geq j(i+1)$ is actually strict whenever i is an ascent of π . If we decide that there are precisely t descents i such that j(i) = j(i+1), then there are $\binom{d_{\pi}(k)}{t}$ choices for these descents. There are then $\binom{k}{k-t}$ choices for the set $\{j(1), \ldots, j(k)\}$ of values taken by the indices $j(1), \ldots, j(k)$. These choices uniquely determine $j(1), \ldots, j(k)$. Together, this shows that

$$|\operatorname{SC}_{132,k}^{-1}(\pi)| = \sum_{t=0}^{d_{\pi}(k)} {\binom{d_{\pi}(k)}{t} \binom{a_{\pi}(k)}{k-t}} = {\binom{d_{\pi}(k) + a_{\pi}(k)}{k}},$$

where the second equality is due to Vandermonde's Identity.

Example 4.8. Let us take π to be the permutation 12131189106751234. Suppose we want to construct a preimage $\tau \in \mathrm{SC}_{132,5}^{-1}(\pi)$. The ascents counted by $a_{\pi}(5)$ are 7, 10, 11, 12, which correspond to the following blanks: $6_751_23_4$. Therefore, $a_{\pi}(5) = 4$. We need to insert the entries 12, 13, 11, 8, 9 into some of these blanks. As an example illustrating the notation from the proof of Theorem 4.7, placing the entry $\pi_2 = 13$ in the blank between the entries $\pi_{10} = 1$ and $\pi_{11} = 2$ would correspond to setting j(2) = 10. The indices in $\{1, \ldots, 5\}$ that are descents of π are 2 and 3 (meaning $d_{\pi}(5) = 2$), so the proof of the theorem tells us that we must choose $j(1), \ldots, j(5) \in \{7, 10, 11, 12\}$ such that $j(1) > j(2) \ge j(3) \ge j(4) > j(5)$. The number of ways to make this choice is

$$\sum_{t=0}^{2} \binom{2}{t} \binom{4}{5-t} = \binom{2+4}{5}.$$

For one specific choice, we can take j(1) = 12, j(2) = j(3) = j(4) = 10, and j(5) = 7. This corresponds to inserting the entries into the blanks to produce the permutation 697511311823124. Reversing this permutation produces $\tau = 412328111315796$, which is indeed in SC⁻¹₁₃₂₅(π).

Lemma 4.9. Let $\pi = \pi_1 \cdots \pi_n \in S_n \setminus \{123 \cdots n\}$, and let $b = b_1 \cdots b_n$ be its corresponding word over the alphabet $\{A, D\}$ as above. Let *i* be the smallest positive integer such that $b_{i+1} = A$, and let $a = a_{\pi}(i-1)$. Finally, let π' be the permutation whose corresponding word over $\{A, D\}$ is the

$$b' = \begin{cases} \underbrace{AD \cdots DD}_{i+1} b_{i+2} \cdots b_n & \text{if } a \ge i; \\ \underbrace{AD \cdots D}_{i} b_{i+2} \cdots b_n A & \text{if } a < i. \end{cases}$$

Then $|\operatorname{SC}_{132}^{-1}(\pi)| \le |\operatorname{SC}_{132}^{-1}(\pi')|.$

Proof. Suppose $a \ge i$. Then $b = \underbrace{AD \cdots DA}_{i+1} b_{i+2} \cdots b_n$ and $b' = \underbrace{AD \cdots DD}_{i+1} b_{i+2} \cdots b_n$, with permutations π and π' corresponding to b and b' respectively. Note that $a = a_{\pi}(i-1) = \cdots = a_{\pi}(1)$. Furthermore, $d_{\pi}(r) = r - 1$ for all $1 \le r \le i$. By Theorem 4.7 and the Hockey Stick Identity,

$$|\operatorname{SC}_{132}^{-1}(\pi)| = 1 + \binom{a}{1} + \dots + \binom{a+i-2}{i-1} + \sum_{r=i}^{n} \binom{d_{\pi}(r) + a_{\pi}(r)}{r}$$
$$= 1 + \binom{a+i-1}{i-1} + \sum_{r=i}^{n} \binom{d_{\pi}(r) + a_{\pi}(r)}{r}.$$

It is straightforward to see that $d_{\pi}(i) = d_{\pi}(i+1)$ and $a_{\pi}(i) = a_{\pi}(i+1)$, so by Pascal's Identity,

$$|\operatorname{SC}_{132}^{-1}(\pi)| = 1 + \binom{a+i-1}{i-1} + \binom{d_{\pi}(i+1) + a_{\pi}(i+1) + 1}{i+1} + \sum_{r=i+2}^{n} \binom{d_{\pi}(r) + a_{\pi}(r)}{r}.$$

As for the preimages of π' , notice that for $r \ge i+1$, we have $d_{\pi'}(r) = d_{\pi}(r)+1$ and $a_{\pi'}(r) = a_{\pi}(r)$. Furthermore, we have $d_{\pi'}(r) = d_{\pi}(r) = r-1$ and $a_{\pi'}(r) \ge a_{\pi}(r)-1$ for all $1 \le r \le i$. Thus,

$$|\operatorname{SC}_{132}^{-1}(\pi')| \ge 1 + \binom{a-1}{1} + \dots + \binom{a+i-3}{i-1} + \binom{a+i-2}{i} + \binom{d_{\pi}(i+1) + a_{\pi}(i+1) + 1}{i+1} + \sum_{r=i+2}^{n} \binom{d_{\pi}(r) + a_{\pi}(r)}{r} = 1 + \binom{a+i-1}{i} + \binom{d_{\pi}(i+1) + a_{\pi}(i+1) + 1}{i+1} + \sum_{r=i+2}^{n} \binom{d_{\pi}(r) + a_{\pi}(r)}{r}.$$

This proves that if $a = a_{\pi}(i-1) \ge i$, then $|\operatorname{SC}_{132}^{-1}(\pi)| \le |\operatorname{SC}_{132}^{-1}(\pi')|$ as desired.

Now suppose a < i so that $b = \underbrace{AD \cdots DA}_{i+1} b_{i+2} \cdots b_n$ and $b' = \underbrace{AD \cdots D}_i b_{i+2} \cdots b_n A$, with permutations π and π' corresponding to b and b' respectively. We may suppose that $b_j = D$ for some $j \ge i+2$ as, otherwise, b = b' and the inequality is trivially true. Thus, $a_{\pi'}(i) = a_{\pi}(i) + 1$ and $d_{\pi'}(i) = d_{\pi}(i)$. It is straightforward to see that $d_{\pi'}(r) = d_{\pi}(r)$ and $a_{\pi'}(r) = a_{\pi}(r)$ for $1 \le r \le i-1$. Furthermore, $d_{\pi}(r) = d_{\pi'}(r-1)$ and $a_{\pi}(r) \le a_{\pi'}(r-1)$ for $r \ge i+2$. Observe also that $\binom{d_{\pi'}(n)+a_{\pi'}(n)}{n} = \binom{d_{\pi'}(n)}{n} = 0$. It follows that,

$$|\operatorname{SC}_{132}^{-1}(\pi')| - |\operatorname{SC}_{132}^{-1}(\pi)| = \binom{d_{\pi}(i) + a_{\pi}(i) + 1}{i} - \binom{d_{\pi}(i) + a_{\pi}(i)}{i} - \binom{d_{\pi}(i+1) + a_{\pi}(i+1)}{i+1} + \sum_{r=i+1}^{n-1} \binom{d_{\pi'}(r) + a_{\pi'}(r)}{r} - \sum_{r=i+2}^{n} \binom{d_{\pi}(r) + a_{\pi}(r)}{r}$$

16

$$\geq \binom{d_{\pi}(i) + a_{\pi}(i) + 1}{i} - \binom{d_{\pi}(i) + a_{\pi}(i)}{i} - \binom{d_{\pi}(i+1) + a_{\pi}(i+1)}{i+1} + \sum_{r=i+2}^{n} \left[\binom{d_{\pi}(r) + a_{\pi}(r)}{r-1} - \binom{d_{\pi}(r) + a_{\pi}(r)}{r} \right].$$

Since a < i by assumption, $a_{\pi}(r) < a < i < r$ for all $r \ge i + 2$. Moreover, $d_{\pi}(r) < r$ trivially, so it follows that $2r > d_{\pi}(r) + a_{\pi}(r)$ and

$$\binom{d_{\pi}(r) + a_{\pi}(r)}{r-1} \ge \binom{d_{\pi}(r) + a_{\pi}(r)}{r}$$

for all $r \ge i + 2$. This shows that the summation in the above inequality is nonnegative.

Finally, we also have that $d_{\pi}(i) = d_{\pi}(i+1) = i-1$ and $a_{\pi}(i) = a_{\pi}(i+1)$, so by Pascal's Identity,

$$|\operatorname{SC}_{132}^{-1}(\pi')| - |\operatorname{SC}_{132}^{-1}(\pi)| \ge \binom{d_{\pi}(i) + a_{\pi}(i) + 1}{i} - \binom{d_{\pi}(i) + a_{\pi}(i)}{i} - \binom{d_{\pi}(i+1) + a_{\pi}(i+1)}{i+1} = \binom{i + a_{\pi}(i)}{i} - \binom{i + a_{\pi}(i)}{i+1} \ge 0,$$

as desired. The last inequality follows from the fact that $a_{\pi}(i) \leq a_{\pi}(i-1) = a < i$.

If we start with a permutation $\pi \in Av(132, 213)$ whose corresponding word over $\{A, D\}$ is not of the form $AD \cdots DA \cdots A$ (i.e., the permutation is not simply a decreasing sequence followed by an increasing sequence), then the previous lemma allows us to alter the permutation and obtain a new permutation π' with at least as many preimages under SC₁₃₂ as π . This is stated in the following corollary. It is well known that a permutation can be written as a decreasing sequence followed by an increasing sequence if and only if it avoids 132 and 231 classically. Therefore, the set of permutations in Av(132, 213) that can be written as a decreasing sequence followed by an increasing sequence is precisely Av(132, 213, 231).

Corollary 4.10. For every positive integer n, we have

$$\max_{\pi \in \operatorname{Av}_n(132,213,231)} |\operatorname{SC}_{132}^{-1}(\pi)| = \max_{\pi \in S_n} |\operatorname{SC}_{132}^{-1}(\pi)|.$$

Proof. Let $M = \max_{\pi \in S_n} |\operatorname{SC}_{132}^{-1}(\pi)|$. By Corollary 4.6, there is a permutation $\pi \in \operatorname{Av}_n(132, 213)$ such that $|\operatorname{SC}_{132}^{-1}(\pi)| = M$. Suppose π contains the pattern 231. We can alter π using Lemma 4.9 in order to obtain a new permutation π' with $|\operatorname{SC}_{132}^{-1}(\pi)| = M$. Let b and b' be the words over $\{A, D\}$ corresponding to π and π' respectively. Given a word w over $\{A, D\}$, let $f_1(w)$ denote the length of the longest consecutive string of D's in w starting in the second position of w, and let $f_2(w)$ be the length of the longest suffix of w that uses only the letter A. Because π contains 231, it is straightforward to check (using the definition of b' given in the proof of Lemma 4.9) that $f_1(b) \leq f_1(b')$ and $f_2(b) \leq f_2(b')$, where at least one of these inequalities must be strict. This shows that if we repeatedly use Lemma 4.9 to alter permutations, we must eventually reach a permutation $\tau \in \operatorname{Av}_n(132, 213, 231)$ with $|\operatorname{SC}_{132}^{-1}(\tau)| = M$.

Theorem 4.11. The maximum number of preimages under SC_{132} or SC_{312} that a permutation in S_n can have is given by

$$\max_{\pi \in S_n} |\operatorname{SC}_{132}^{-1}(\pi)| = \max_{\pi \in S_n} |\operatorname{SC}_{312}^{-1}(\pi)| = \binom{n-1}{\lfloor \frac{n-1}{2} \rfloor}$$

Proof. By complementing, it suffices to consider only the map SC_{132} . Let $M = \max_{\pi \in S_n} |SC_{132}^{-1}(\pi)|$. By Corollary 4.10, there exists $\pi \in \operatorname{Av}_n(132, 213, 231)$ with $|\operatorname{SC}_{132}^{-1}(\pi)| = M$. Since π avoids 132, 213, and 231, it must be of the form

$$n(n-1)\cdots(n-r+1)123\cdots(n-r)$$

The word over $\{A, D\}$ corresponding to π is $A \underbrace{D \cdots D}_{r} \underbrace{A \cdots A}_{n-r-1}$. Observe that $d_{\pi}(k) = \min(k-1, r)$ for all $1 \leq k \leq n$. Furthermore, we have $a_{\pi}(k) = n-r-1$ if $1 \leq k \leq r$ and $a_{\pi}(k) = 0$ if $r+1 \leq k \leq n$. By Theorem 4.7 and the Hockey Stick Identity, we have

$$M = |\operatorname{SC}_{132}^{-1}(\pi)| = 1 + \binom{n-r-1}{1} + \binom{n-r}{2} + \dots + \binom{n-2}{r} = \binom{n-1}{r} \leq \binom{n-1}{\lfloor \frac{n-1}{2} \rfloor}.$$

eover, equality is achieved when $r = \lfloor \frac{n-1}{2} \rfloor$; this completes the proof.

Moreover, equality is achieved when $r = \lfloor \frac{n-1}{2} \rfloor$; this completes the proof.

5. The Maps SC_{213} and SC_{231}

We begin this section with a brief discussion of the sets $Sort(SC_{213})$ and $Sort(SC_{231})$. The initial terms of the sequence $(|\operatorname{Sort}_n(\operatorname{SC}_{213})|)_{n>0}$ are

1, 1, 2, 5, 15, 50, 180, 686, 2731, 11254.

This sequence appears to be new. On the other hand, the first 10 terms of $(|Sort_n(SC_{213})|)_{n>0}$ are

1, 1, 2, 6, 21, 79, 311, 1265, 5275, 22431.

These numbers match the initial terms in the OEIS sequence A033321, which is the binomial transform of Fine's sequence. Fine's sequence $(F_k)_{k>0}$ can be defined via its generating function

$$\sum_{k \ge 0} F_k x^k = \frac{1 - \sqrt{1 - 4x}}{3 - \sqrt{1 - 4x}}.$$

Conjecture 5.1. For each positive integer n, we have

$$|\operatorname{Sort}_{n}(\operatorname{SC}_{231})| = \sum_{k=0}^{n} \binom{n}{k} F_{k+1}.$$

Our main focus in this section will be on the maximum number of preimages a permutation in S_n can have under either SC₂₁₃ or SC₂₃₁.

Theorem 5.2. For every $n \ge 2$, we have

$$\max_{\pi \in S_n} |\operatorname{SC}_{213}^{-1}(\pi)| = \max_{\pi \in S_n} |\operatorname{SC}_{231}^{-1}(\pi)| = 2^{n-2}.$$

Proof. The equality $\max_{\pi \in S_n} |\operatorname{SC}_{213}^{-1}(\pi)| = \max_{\pi \in S_n} |\operatorname{SC}_{231}^{-1}(\pi)|$ follows directly from Lemma 1.2, so we can focus on the map $\operatorname{SC}_{231}^{n \in \mathbb{S}_n}$. Now consider the process of sending a permutation $\tau = \tau_1 \cdots \tau_n \in S_n$ through a consecutive-231-avoiding stack. Say that the entry τ_i is premature if it is popped out of the stack before all of the entries of τ enter the stack. The effect of sending τ through the stack is to shift all of the premature entries to the left while preserving their relative order and reverse the remaining (non-premature) entries. More formally, if $\{\tau_{j_1}, \cdots, \tau_{j_k}\}$ are the premature entries of τ , then $\operatorname{SC}_{231}(\tau) = \tau_{j_1} \cdots \tau_{j_k} \operatorname{rev}(\tau \setminus \{\tau_{j_1}, \cdots, \tau_{j_k}\})$, where $\tau \setminus \{\tau_{j_1}, \cdots, \tau_{j_k}\}$ is simply the sequence τ with the terms in $\{\tau_{j_1}, \cdots, \tau_{j_k}\}$ removed.

Now fix $\pi = \pi_1 \cdots \pi_n \in S_n$, and consider a preimage τ of π under SC₂₃₁. Suppose that $\tau_{j_1}, \ldots, \tau_{j_k}$ are the premature entries of τ (with $j_1 < \cdots < j_k$). By the previous paragraph, τ is obtained by inserting the entries π_1, \ldots, π_k into rev $(\pi_{k+1} \cdots \pi_n)$ so that the entry in position j_i of τ is π_i for all $i \in [k]$. This shows that τ is uniquely determined by specifying π and the set of indices $\{j_1, \ldots, j_k\}$. Finally, notice that the first and last entries of a permutation cannot be premature. Therefore, the positions of the premature entries of τ must form a subset of $\{2, \ldots, n-1\}$. It follows that $|\operatorname{SC}_{231}^{-1}(\pi)| \leq 2^{n-2}$

To complete the proof, it suffices to show that $|\operatorname{SC}_{231}^{-1}(n(n-1)\cdots 1)| \geq 2^{n-2}$. For each set $J = \{j_1, \ldots, j_k\} \subseteq \{2, \ldots, n-1\}$ (with $j_1 < \cdots < j_k$), let τ^J be the permutation obtained by inserting the entries $n, n-1, \ldots, n-k+1$ into the sequence $123 \cdots (n-k)$ so that for each $i \in [k]$, the entry n+1-i is in position j_i . For example, if n=8 and $J = \{2,3,5\}$, then $\tau^J = 18726345$. It is straightforward to check that $\operatorname{SC}_{231}(\tau^J) = n(n-1)\cdots 1$ and that $\tau^J \neq \tau^{J'}$ whenever $J \neq J'$. Consequently, $|\operatorname{SC}_{231}^{-1}(n(n-1)\cdots 1)| \geq 2^{n-2}$.

6. The MAPS SC_{123} and SC_{321}

Our goal in this section is to enumerate the sets $Sort(SC_{123})$ and $Sort(SC_{321})$, thereby proving analogues of two more of the main theorems from [11]. Our enumeration of $Sort(SC_{321})$ generalizes to $Sort(SC_{k(k-1)\cdots 1})$, so we focus on $Sort(SC_{k(k-1)\cdots 1})$ first instead.

It follows immediately from Lemma 2.3 that $\operatorname{Sort}(\operatorname{SC}_{k(k-1)\dots 1}) = \operatorname{Av}(132, \underline{12\dots k})$ for $k \geq 3$. In Proposition 3.2 of [24] it was shown that this set is enumerated by the so-called *generalized Motzkin* numbers $M_{k-1,n}$, which can be defined via the formula

$$M_{k-1,n} = \frac{1}{n+1} \sum_{j=0}^{\lfloor n/k \rfloor} (-1)^j \binom{n+1}{j} \binom{2n-jk}{n}.$$

For k = 3, 4, 5, and 6, these numbers are given by OEIS sequences A001006, A036765, A036766, and A036767 respectively [32].

Theorem 6.1 ([24]). For $k \ge 3$, the sets $\operatorname{Sort}_n(\operatorname{SC}_{k(k-1)\cdots 1}) = \operatorname{Av}_n(132, \underline{12\cdots k})$ are enumerated by the generalized Motzkin numbers. More precisely,

$$|\operatorname{Sort}_n(\operatorname{SC}_{k(k-1)\cdots 1})| = M_{k-1,n}.$$

The above theorem completes our enumeration of $\text{Sort}(\text{SC}_{k(k-1)\cdots 1})$. We remark that when k = 3, the generalized Motzkin numbers $M_{2,n}$ are the same as the classical Motzkin numbers M_n . These numbers will also be related to our enumeration of $\text{Sort}(\text{SC}_{123})$.

Corollary 6.2. The set $Sort(SC_{321})$ is enumerated by the Motzkin numbers. That is,

$$|\operatorname{Sort}_n(\operatorname{SC}_{321})| = M_n.$$

The following is a well-known recurrence for the Motzkin numbers.

Lemma 6.3. We have $M_0 = M_1 = 1$ and, for $n \ge 2$,

$$M_n = M_{n-1} + \sum_{i=0}^{n-2} M_i M_{n-2-i}.$$

To begin enumerating Sort(SC₁₂₃), we first classify the set. In order to do this, it is helpful to use the notion of a *vincular pattern*. We refer the reader to [34] for a formal treatment of vincular patterns; we will only need to consider the set Av(132, 3214, 4213, 4312). This is the set of permutations that avoid 132 and also do not contain any occurrences of the patterns 3214, 4213, or 4312 in which the first 3 entries in the occurrence of the pattern appear consecutively in the permutation. For example, 51324 belongs to this set of permutations because, although 5, 3, 2, 4 form an occurrence of the pattern 4213, the entries 5, 3, 2 do not appear consecutively in 51324.

Theorem 6.4. The set $Sort(SC_{123})$ consists of the permutations $\pi = \pi_1 \cdots \pi_n$, for some positive integer n, that satisfy the following properties:

- $\pi \in Av(132)$
- if $\pi_i \pi_{i+1} \cdots \pi_j$ is a descending run of π of length at least 3, then there are no entries bigger than π_{j-1} to the right of π_{j-1} in π , and the entries $\pi_{i+1}, ..., \pi_{j-1}$ are consecutive integers.

Alternatively, $Sort(SC_{123}) = Av(132, \underline{3214}, \underline{4213}, \underline{4312}).$

Proof. It is straightforward to check that a permutation satisfies the two listed conditions if and only if it is in Av(132, <u>3214</u>, <u>4213</u>, <u>4312</u>). The proof of the first statement is straightforward once we find an expression for SC₁₂₃(π). Fortunately, this was done in Lemma 3.3 for the map SC₃₂₁, so we may use Lemma 1.2 to find a similar expression for SC₁₂₃(π) by complementing. Namely, suppose $\pi = d_1 \cdots d_k$, where each d_i is a descending run. We have SC₁₂₃(π) = $d_1^m \cdots d_k^m \operatorname{rev}(d_k^e) \cdots \operatorname{rev}(d_1^e)$, where d_i^m is the subsequence of d_i consisting of the entries that are not the first or last entries in d_i and d_i^e is the sequence obtained from d_i by deleting d_i^m .

To see that all permutations in Sort(SC₁₂₃) must satisfy the two conditions, note first that if *abc* forms an occurrence of the classical pattern 132 in π and a,b, and c are in the descending runs d_i, d_j , and d_ℓ respectively, then c, the second entry of $\operatorname{rev}(d_j^e)$, and the first entry of $\operatorname{rev}(d_i^e)$ form an occurrence of the classical pattern 231 in SC₁₂₃(π). This shows that Sort(SC₁₂₃) \subseteq Av(132). For the first part of the second condition, suppose $\pi_i \pi_{i+1} \cdots \pi_j$ is a descending run of length at least 3 and that π_ℓ is an entry to the right of π_j that is also the first entry of a different descending run. Then $\pi_{j-1}\pi_\ell\pi_j$ is a subsequence of SC₁₂₃(π), and we certainly have $\pi_j < \pi_{j-1}$. It follows that if SC₁₂₃(π) \in Av(231), then $\pi_\ell < \pi_{j-1}$. Hence, no entry to the right of π_{j-1} in π is bigger than π_{j-1} . Finally, the previous two points imply the second part of the second condition. If the entries $\pi_{i+1}, \ldots, \pi_{j-1}$ are not consecutive integers, then there is some entry $\pi_{i'}$ that is not part of the descending run containing these entries and that satisfies $\pi_{i+1} > \pi_{i'} > \pi_{j-1}$. Then either $\pi_{i'}\pi_i\pi_{i+1}$ is a 132-subsequence or there is an entry to the right of π_{j-1} that is bigger than π_{j-1} .

To see that every permutation satisfying these conditions is indeed in Sort(SC₁₂₃), again decompose π into its descending runs and write SC₁₂₃(π) = $d_1^m \cdots d_k^m \operatorname{rev}(d_k^e) \cdots \operatorname{rev}(d_1^e)$. The first condition ensures that $\operatorname{rev}(d_k^e) \cdots \operatorname{rev}(d_1^e)$ does not contain an occurrence of the classical pattern 231. The first part of the second condition ensures that $d_1^m \cdots d_k^m$ is a decreasing sequence, so it also does not contain an occurrence of 231. Finally, the first condition and first part of the second condition ensure that there is no occurrence of the pattern 231 containing entries from both $d_1^m \cdots d_k^m$ and $\operatorname{rev}(d_k^e) \cdots \operatorname{rev}(d_1^e)$. Indeed, suppose *abc* is an occurrence such a pattern. As $d_1^m \cdots d_k^m$ is a decreasing sequence, the entry *a* must be in $d_1^m \cdots d_k^m$ while *b* and *c* are in $\operatorname{rev}(d_k^e) \cdots \operatorname{rev}(d_1^e)$. Let *i*, *j*, ℓ be the indices such that *a* is in d_i^m , *b* is in d_j^e , and *c* is in d_ℓ^e . Since *c* appears to the right of *b* in SC₁₂₃(π), we must have $\ell < j$. The first part of the second condition guarantees that *j* $\leq i$. However, this implies that *cba* is an occurrence of the classical pattern 132 in π , which is a contradiction. Before enumerating the permutations in $Sort(SC_{123})$, we need one more lemma that counts permutations avoiding 132 classically and 123 consecutively. We obtain this result via a bijection between Av(132) and Av(231). For our purposes, we only need the existence of this bijection and one of its properties. The bijection is described in detail in Lemma 4.1 of [14].

Lemma 6.5 ([14]). For each $n \ge 1$, there exists a bijection swd : $Av_n(132) \rightarrow Av_n(231)$ that preserves the set of descents. That is, an index i is a descent of π if and only if it is a descent of $swd(\pi)$.

Corollary 6.6. For each $n \ge 1$, the sets of permutations $\operatorname{Av}_n(132, \underline{321})$, $\operatorname{Av}_n(231, \underline{321})$, and $\operatorname{Av}_n(132, \underline{123})$ are in bijection. Consequently, $|\operatorname{Av}_n(132, \underline{321})| = M_n$.

Proof. Using the previous lemma, we obtain a chain of bijections

 $\begin{array}{c} \operatorname{Av}_n(132,\underline{321}) \xrightarrow{\operatorname{swd}} \operatorname{Av}_n(231,\underline{321}) \xrightarrow{\operatorname{rev}} \operatorname{Av}_n(132,\underline{123}). \\ \end{array}$ By Theorem 6.2, $|\operatorname{Av}_n(132,\underline{321})| = |\operatorname{Av}_n(132,\underline{123})| = M_n.$

We need one final lemma before completing the enumeration of $Sort(SC_{123})$.

Lemma 6.7. The number of permutations in $Sort_n(SC_{123})$ beginning with the entry n is M_{n-1} .

Proof. We proceed by induction on n, noting first that the lemma is trivial if $n \leq 2$. Assume $n \geq 3$. Suppose $\pi \in \operatorname{Sort}_n(\operatorname{SC}_{123})$ starts with the entry n, and let π' be the permutation obtained by deleting the leading entry n from π . It follows from Theorem 6.4 that $\pi' \in \operatorname{Sort}_{n-1}(\operatorname{SC}_{123})$ and that π' either starts with the entry n-1 or starts with an ascent. On the other hand, if $\tau \in \operatorname{Sort}_{n-1}(\operatorname{SC}_{123})$ starts with the entry n-1 or starts with an ascent, then the permutation $n\tau$ is in $\operatorname{Sort}_n(\operatorname{SC}_{123})$. By the induction hypothesis, there are M_{n-2} permutations in $\operatorname{Sort}_{n-1}(\operatorname{SC}_{123})$ that start with n-1.

We now need to determine the number of permutations $\tau \in \text{Sort}_{n-1}(\text{SC}_{123})$ that start with an ascent. Let $i \in \{2, \ldots, n-1\}$ be such that $\tau_i = n-1$. We may write τ as A(n-1)B, where A does not start with a descent (meaning it either starts with an ascent or consists of a single entry). Since τ must avoid 132 by Theorem 6.4, every entry in A is greater than every entry in B, so the entries in A are $n-i+2,\ldots,n-2$ and the entries in B are $1,\ldots,n-i+1$. Since (n-1)B has length n-i+1 and has the same relative order as (n-i+1)B, the number of possibilities for (n-1)Bis the same as the number of possibilities for (n-i+1)B. Now (n-i+1)B can be any element of $\operatorname{Sort}_{n-i+1}(\operatorname{SC}_{123})$ that starts with n-i+1, so it follows from the induction hypothesis that the number of possibilities for (n-1)B is M_{n-i} . Next, for each possible choice of A, we can subtract n-i+1 from each entry of A. Theorem 6.4 implies that this operation yields a bijection between the set of possible choices for A and the set of permutations in $Av_{i-1}(132, 321)$ that do not start with a descent. If $i \geq 3$, then since A avoids 132 and starts with an ascent, the first two entries of A must be consecutive integers. As a result, removing the first entry and standardizing gives a bijection between permutations in $Av_{i-1}(132, 321)$ that do not start with a descent and permutations in $Av_{i-2}(132, 321)$. This implies that there are M_{i-2} possibilities for A by Lemma 6.6. If i = 2, then it is also certainly true that the number of choices for A is $M_{i-2} = M_0 = 1$.

Putting this all together and invoking the Motzkin number recurrence in Lemma 6.3, we find that the number of permutations in $\text{Sort}_n(\text{SC}_{123})$ beginning with n is

$$M_{n-2} + \sum_{i=2}^{n-1} M_{i-2}M_{n-i} = M_{n-2} + \sum_{i=0}^{n-3} M_i M_{n-3-i} = M_{n-1}.$$

We can now prove that $Sort(SC_{123})$ is enumerated by the first differences of Motzkin numbers, which form the OEIS sequence A002026 [32].

Theorem 6.8. The permutations in $Sort(SC_{123})$ are counted by the first differences of Motzkin numbers. That is,

$$|\operatorname{Sort}_{n}(\operatorname{SC}_{123})| = M_{n+1} - M_{n}$$

Proof. Let X_i^n denote the set of permutations in $\operatorname{Sort}_n(\operatorname{SC}_{123})$ in which the *i*th entry is *n*. If $\pi \in X_i^n$, then we may write $\pi = AnB$, where *A* has length i - 1 and *B* has length n - i. By Theorem 6.4, AnB avoids the pattern 132, so every entry in *A* is bigger than every entry in *B*, meaning that *A* is a permutation of the numbers $n - i + 1, \ldots, n - 1$ and *B* is a permutation of $1, \ldots, n - i$. Moreover, the standardization of *A* can be an arbitrary element of $\operatorname{Av}(132, \underline{321})$, so the number of possibilities for *A* is M_{i-1} by Corollary 6.6. Meanwhile, the standardization of nB can be an arbitrary permutation in $\operatorname{Sort}_{n-i+1}(\operatorname{SC}_{123})$ that starts with n - i + 1, so by Lemma 6.7, there are M_{n-i} possibilities for nB. Thus, $|X_i^n| = M_{i-1}M_{n-i}$, and

$$|\operatorname{Sort}_{n}(\operatorname{SC}_{123})| = \sum_{i=1}^{n} M_{i-1} M_{n-i} = \sum_{i=0}^{n-1} M_{i} M_{n-1-i} = M_{n+1} - M_{n}.$$

7. Dynamics of the Maps s_{132} and s_{312}

The set $\text{Sort}(s_{132})$ was the primary focus of the recent article [12], where it was shown that $\text{Sort}(s_{132})$ is equal to the set of permutations avoiding the classical pattern 2314 as well as a certain mesh pattern. The same article proved that $|\text{Sort}_n(s_{132})| = \sum_{k=0}^{n-1} {n-1 \choose k} C_k$, where C_k denotes the kth Catalan number. In this section, we briefly discuss some dynamical properties of the map s_{132} and, by complementation, the map s_{312} .

Theorem 7.1. The periodic points of the map $s_{132} : S_n \to S_n$ are precisely the permutations in $Av_n(132, 231)$. The periodic points of the map $s_{312} : S_n \to S_n$ are precisely the permutations in $Av_n(213, 312)$. When $n \ge 2$, each of these periodic points has period 2. Furthermore, for every permutation $\pi \in S_n$, we have $s_{132}^{n-1}(\pi) \in Av_n(132, 231)$ and $s_{312}^{n-1}(\pi) \in Av_n(213, 312)$.

Proof. By Lemma 1.2, it suffices to prove the desired properties of s_{132} . Note that if we have $\pi \in Av(132, 231)$, then $s_{132}(\pi) = rev(\pi) \in Av(132, 231)$, so $s_{132}^2(\pi) = \pi$. Hence, π is a periodic point of s_{132} , and the period is 2 if $n \ge 2$.

It remains to prove that $s_{132}^{n-1}(\pi) \in \operatorname{Av}(132, 231)$ for every $\pi \in S_n$. This is trivial if n = 1, so we may assume $n \geq 2$ and induct on n. Consider sending a permutation $\tau \in S_n$ through a classical-132-avoiding stack. When the entry 1 enters the stack, the entries below it in the stack must appear in increasing order from top to bottom. Furthermore, the entry 1 will not leave the stack until after all entries have entered the stack. This implies that the last ascending run of $s_{132}(\tau)$ begins with the entry 1. Therefore, it suffices to prove that if $\pi \in S_n$ is a permutation whose last ascending run begins with 1, then $s_{132}^{n-2}(\pi) \in \operatorname{Av}(132, 231)$. We prove this by induction on n, noting first that it is trivial if n = 2. Now suppose $n \geq 3$.

For each $\tau \in S_n$, let τ^* be the permutation in S_{n-1} obtained by deleting the entry 1 from τ and then decreasing all remaining entries by 1. Observe that if the entries 1 and 2 appear consecutively in τ , then they also appear consecutively in $s_{132}(\tau)$ and that $(s_{132}(\tau))^* = s_{132}(\tau^*)$. By induction, we have $s_{132}^{n-2}(\tau^*) \in \operatorname{Av}(132, 231)$. If the last ascending run of τ begins with 1, then the above analysis of the classical-132-avoiding stack shows that the entries 1 and 2 appear consecutively in $s_{132}(\tau)$. It then follows that the entries 1 and 2 appear consecutively in $s_{132}^{n-2}(\tau)$ and that $(s_{132}^{n-2}(\tau))^* = s_{132}^{n-2}(\tau^*) \in \operatorname{Av}(132, 231)$. This implies that $s_{132}^{n-2}(\tau) \in \operatorname{Av}(132, 231)$, as desired. \Box

Whenever one is confronted with a noninvertible finite dynamical system, it is natural to ask for the maximum number of iterations of the map needed to send every point to a periodic point. For example, Knuth [29, pages 106–110] studied this problem for the classical bubble sort map. For West's stack-sorting map s, permutations requiring close to the maximum number of iterations to get sorted into the identity were studied by West and Claesson–Dukes–Steingrímsson [13,37]. Ungar gave a rather nontrivial argument proving that the maximum number of iterations needed to sort a permutation in S_n using the pop-stack-sorting map is n-1 [36], and the permutations requiring n-1 iterations were investigated further by Asinowski, Banderier, and Hackl [1]. Some other papers that have studied similar questions for other finite dynamical systems include [4, 21, 23, 25, 26, 35].

Our second result in this section will determine the maximum number of iterations of s_{132} (equivalently, s_{312}) needed to send a permutation to a periodic point. In what follows, we let $sd_{132}(\pi)$ denote the smallest nonnegative integer t such that $s_{132}^t(\pi) \in Av(132, 231)$. Note that $sd_{132}(\pi)$ is also the smallest nonnegative integer t such that $s_{312}^t(comp(\pi)) \in Av(213, 312)$.

Theorem 7.2. For $n \ge 3$, we have

$$\max_{\pi \in S_n} \operatorname{sd}_{132}(\pi) = n - 1.$$

Proof. Theorem 7.1 tells us that $\max_{\pi \in S_n} \operatorname{sd}_{132}(\pi) \leq n-1$, so it suffices to find a permutation $\pi \in S_n$ such that $\operatorname{sd}_{132}(\pi) \geq n-1$. Let n = 3m + r, where $r \in \{0, 1, 2\}$. Let λ_m denote the permutation in S_{n-r} obtained by taking the skew sum of m copies of the permutation 132. In other words, $\lambda_m = \xi_m \xi_{m-1} \cdots \xi_1$, where ξ_j denotes the sequence (3j-2)(3j)(3j-1). Let

$$\pi = \begin{cases} \lambda_m & \text{if } r = 0; \\ n\lambda_m & \text{if } r = 1; \\ (n-1)n\lambda_m & \text{if } r = 2. \end{cases}$$

We claim that $s_{132}^{n-2}(\pi) \notin \operatorname{Av}(132, 231)$; that is, $\operatorname{sd}_{132}(\pi) \ge n-1$. For the sake of simplicity, we will prove this in the case r = 0; the proofs in the cases r = 1 and r = 2 are similar. Thus, n = 3m and $\pi = \lambda_m$. If n = 3, then we are done because $s_{132}(\pi) = s_{132}(132) = 231 \notin \operatorname{Av}(132, 231)$. Therefore, we may assume $n = 3m \ge 6$.

It will be helpful to introduce a little more notation. If $k \ge 1$ is odd (respectively, even), we let V_k denote the permutation in S_{3k} obtained by listing the odd (respectively, even) numbers in [3k + 1] in decreasing order and then listing the even (respectively, odd) numbers in [3k + 1] in increasing order. For example, we have $V_1 = 3124$, $V_2 = 6421357$, and $V_3 = 97531246810$. For $i \in \{0, 1, 2\}$, let $D_{k,m}^{(i)}$ be the sequence obtained by writing the elements of $\{3k + 2, 3k + 3, \ldots, 3m\}$ that are congruent to *i* modulo 3 in decreasing order. For example, $D_{2,5}^{(0)} = 15129$, $D_{2,5}^{(1)} = 1310$, and $D_{2,5}^{(2)} = 14118$. The following two claims can be verified in a straightforward manner using nothing more than the definition of the map s_{132} ; we leave this verification to the reader.

Claim 1: We have $s_{132}^4(\lambda_m) = D_{1,m}^{(2)} D_{1,m}^{(0)} V_1 \operatorname{rev} \left(D_{1,m}^{(1)} \right).$

Claim 2: For $1 \le k \le m - 2$, we have

$$s_{132}^3 \left(D_{k,m}^{(2)} D_{k,m}^{(0)} V_k \operatorname{rev} \left(D_{k,m}^{(1)} \right) \right) = D_{k+1,m}^{(2)} D_{k+1,m}^{(0)} V_{k+1} \operatorname{rev} \left(D_{k+1,m}^{(1)} \right)$$

These two claims imply that $s_{132}^{n-2}(\pi) = s_{132}^{3m-2}(\lambda_m) = D_{m-1,m}^{(2)} D_{m-1,m}^{(0)} V_{m-1} \operatorname{rev} \left(D_{m-1,m}^{(1)} \right)$. Note that $D_{m-1,m}^{(2)}$ consists of the single entry 3m-1. Similarly, $D_{m-1,m}^{(0)}$ consists of the single entry 3m. Finally, $D_{m-1,m}^{(1)}$ is empty. Thus, $s_{132}^{n-2}(\pi) = (3m-1)(3m)V_{m-1} \notin \operatorname{Av}(132,231)$.

8. FUTURE DIRECTIONS

We have initiated the study of the consecutive-pattern-avoiding stack-sorting maps SC_{σ} and analyzed them as both dynamical systems and sorting procedures. These maps are generalizations of West's stack-sorting map [37] and are variants of the classical-pattern-avoiding stack-sorting maps studied by Cerbai, Claesson, and Ferrari in [11].

Our main results on SC_{σ} as a sorting procedure were the characterization of when $Sort(SC_{\sigma})$ is a permutation class and the enumeration of $Sort(SC_{\sigma})$ summarized in the table below.

$\sigma \setminus n$	0	1	2	3	4	5	6	7	8	9	OEIS
123	1	1	2	5	12	30	76	196	512	1353	A002006
132	1	1	2	5	14	42	132	429	1430	4862	A000108
213	1	1	2	5	15	50	180	686	2731	11254	unknown
231	1	1	2	6	21	79	311	1265	5275	22431	unknown
312	1	1	2	5	15	50	179	675	2649	10734	unknown
321	1	1	2	4	9	21	51	127	323	835	A001006

It would be interesting to have nontrivial information about the asymptotic behavior of the unknown sequences in this table. Recall that we also have Conjecture 5.1, which states that $Sort(SC_{231})$ is enumerated by the binomial transform of Fine's sequence (OEIS sequence A033321).

From the dynamical point of view, we first proved Theorem 3.1, which characterized the periodic points of SC_{σ} for each $\sigma \in S_3$. This theorem has motivated us to formulate the following conjecture, which we have additionally confirmed when k = 4 and $n \leq 8$.

Conjecture 8.1. Let $\sigma \in S_k$ for some $k \geq 3$. The periodic points of the map $SC_{\sigma} : S_n \to S_n$ are precisely the permutations in $Av_n\left(\underline{\sigma}, \underline{rev}(\sigma)\right)$.

We then asked for the maximum value of $|SC_{\sigma}^{-1}(\pi)|$ for $\pi \in S_n$. The following table summarizes what we know about these values; note that by Lemma 1.2, we only need to consider one pattern from each of the pairs $\{123, 321\}, \{132, 312\}, \{213, 231\}$.

$\sigma \setminus n$	1	2	3	4	5	6	7	8	9	OEIS
123	1	1	2	3	4	7	11	16	26	unknown
132	1	1	2	3	6	10	20	35	70	A001405
231	1	1	2	4	8	16	32	64	128	A011782

It would be interesting to have nontrivial estimates for the numbers $\max_{\pi \in S_n} |\operatorname{SC}_{123}^{-1}(\pi)|$ appearing in the first row of the above table.

In Section 7, we studied the maximum number of iterations of s_{132} (equivalently, s_{312}) needed to send a permutation to a periodic point. This line of investigation is completely open for the

consecutive-pattern-avoiding stack-sorting maps. However, we do have the following conjecture, which we have verified for $n \leq 9$. Recall that the set of periodic points of $SC_{231} : S_n \to S_n$ is $Av_n(\underline{132}, \underline{231}) = Av_n(\underline{132}, \underline{231})$.

Conjecture 8.2. Let $n \geq 3$. We have $\operatorname{SC}_{231}^{2n-4}(\pi) \in \operatorname{Av}_n(132, 231)$ for every $\pi \in S_n$. Furthermore, there exists a permutation $\tau \in S_n$ such that $\operatorname{SC}_{231}^{2n-5}(\tau) \notin \operatorname{Av}_n(132, 231)$.

In [16], the first author showed that the numbers 3, 7, 11, 15, 19, 23 are *infertility numbers*, meaning that there does not exist a permutation π such that $|s^{-1}(\pi)| \in \{3, 7, 11, 15, 19, 23\}$. By contrast, we have the following conjecture for the maps SC_{σ} with $\sigma \in S_3$.

Conjecture 8.3. For every $\sigma \in S_3$ and every positive integer f, there exists a permutation π such that $|\operatorname{SC}_{\sigma}^{-1}(\pi)| = f$.

Section 7 contains the first theorems of a dynamical nature concerning maps of the form s_{σ} . It would be interesting to have other dynamical results for these maps, including analogues of Theorems 7.1 and 7.2 for the maps s_{σ} with $\sigma \in \{123, 213, 231, 321\}$. In Theorems 7.1 and 7.2, we proved that every permutation $\pi \in S_n$ satisfies $s_{132}^{n-1}(\pi) \in \operatorname{Av}(132, 231)$ and that there exists at least one permutation $\tau \in S_n$ such that $s_{132}^{n-2}(\tau) \notin \operatorname{Av}(132, 231)$. It would be interesting to obtain nontrivial information about these permutations τ . We also have the following conjecture. If n is odd (respectively, even), let \mathbf{V}_n denote the permutation in S_n obtained by listing the odd (respectively, even) elements of [n] in decreasing order and then listing the even (respectively, odd) elements of [n] in increasing order. For example, $\mathbf{V}_6 = 642135$ and $\mathbf{V}_7 = 7531246$.

Conjecture 8.4. If $\tau \in S_n$ is such that $s_{132}^{n-2}(\tau) \notin Av(132, 231)$, then $s_{132}^{n-1}(\tau) = \mathbf{V}_n$.

For example, one can check that if $\tau \in S_6$ is such that $s_{132}^4(\tau) \notin Av(132, 231)$, then $s_{132}^5(\tau) = \mathbf{V}_6 = 642135$.

Of course, it would also be very natural to consider the maps of the form $SC_{\sigma^{(1)},\ldots,\sigma^{(r)}}$, where $SC_{\sigma^{(1)},\ldots,\sigma^{(r)}}$ is defined using a stack whose contents must consecutively avoid all of the patterns $\sigma^{(1)},\ldots,\sigma^{(r)}$. It would also be exciting to have results about the maps SC_{σ} for patterns σ of length at least 4 (such as progress toward Conjecture 8.1).

9. Acknowledgments

This research was conducted as part of the University of Minnesota Duluth Mathematics REU and was supported by NSF-DMS Grant 1949884 and NSA Grant H98230-20-1-0009. The authors thank Joe Gallian for running the REU program and providing encouragement. We also thank Giulio Cerbai for very helpful discussions. The first author was supported by a Fannie and John Hertz Foundation Fellowship and an NSF Graduate Research Fellowship.

References

- A. Asinowski, C. Banderier, and B. Hackl, Flip-sort and combinatorial aspects of pop-stack sorting. arXiv:2003.04912.
- [2] M. Albert, C. Homberger, J. Pantone, N. Shar, and V. Vatter, Generating permutations with restricted containers. J. Combin. Theory Ser. A, 157 (2018).
- [3] J.-L. Baril, C. Khalil, V. Vajnovszki, Catalan and Schröder permutations sortable by two restricted stacks. arXiv:2004.01812.

- 25
- [4] H.-J. Bentz, Proof of the Bulgarian Solitaire conjectures. Ars Combin., 23 (1987), 151–170.
- [5] K. Berlow, Restricted stacks as functions. arXiv:2008.01164.
- [6] M. Bóna, Combinatorics of permutations. CRC Press, 2012.
- [7] M. Bóna, A survey of stack-sorting disciplines. Electron. J. Combin., 9 (2003).
- [8] M. Bousquet-Mélou, Multi-statistic enumeration of two-stack sortable permutations. *Electron. J. Combin.*, 5 (1998).
- [9] M. Bousquet-Mélou, Sorted and/or sortable permutations. Discrete Math., 225 (2000), 25–50.
- [10] G. Cerbai, Sorting Cayley permutations with pattern-avoiding machines. arXiv:2006.05692.
- [11] G. Cerbai, A. Claesson, and L. Ferrari, Stack sorting with restricted stacks. J. Combin. Theory Ser. A, 173 (2020).
- [12] G. Cerbai, A. Claesson, L. Ferrari, and E. Steingrímsson, Sorting with pattern-avoiding stacks: the 132-machine. *Electron. J. Combin.*, 27 (2020).
- [13] A. Claesson, M. Dukes, and E. Steingrímsson, Permutations sortable by n 4 passes through a stack. Ann. Combin., 14 (2010), 45–51.
- [14] C. Defant, Catalan intervals and uniquely sorted permutations. J. Combin. Theory Ser. A, 174 (2020).
- [15] C. Defant, Counting 3-stack-sortable permutations. J. Combin. Theory Ser. A, 172 (2020).
- [16] C. Defant, Fertility numbers. J. Comb., **11** (2020), 511–526.
- [17] C. Defant, Preimages under the stack-sorting algorithm. Graphs Combin., 33 (2017), 103–122.
- [18] C. Defant, Troupes, cumulants, and stack-sorting. arXiv:2004.11367.
- [19] C. Defant, M. Engen, and J. A. Miller, Stack-sorting, set partitions, and Lassalle's sequence. J. Combin. Theory Ser. A, 175 (2020).
- [20] S. Elizalde. A survey of consecutive patterns in permutations. In: Recent trends in combinatorics, 601–618. Springer, 2016.
- [21] G. Etienne, Tableux de Young et Solitaire Bulgare, J. Combin. Theory Ser. A, 58 (1991), 181–197.
- [22] I. Goulden and J. West, Raney paths and a combinatorial relationship between rooted nonseparable planar maps and two-stack-sortable permutations, J. Combin. Theory Ser. A, 75 (1996), 220–242.
- [23] J. R. Griggs and C.-C. Ho, The cycling of partitions and composition under repeated shifts. Adv. Appl. Math., 21 (1998), 205–227.
- [24] N. Hein and J. Huang, Modular Catalan numbers, European J. Combin., 61 (2017), 197–218.
- [25] J. D. Hobby, D. Knuth, Problem 1: Bulgarian Solitaire, In: A Programming and Problem-Solving Seminar, Department of Computer Science, Stanford University, 1983 (December), 6–13.
- [26] K. Igusa, Solution of the Bulgarian solitaire conjecture, Math. Mag., 58 (1985), 259–271.
- [27] S. Kitaev, Patterns in Permutations and Words. Monographs in Theoretical Computer Science. Springer, Heidelberg, 2011.
- [28] D. E. Knuth, The Art of Computer Programming, Volume I: Fundamental Algorithms. Addison-Wesley, 1973.
- [29] D. E. Knuth, The Art of Computer Programming, Volume 3: Sorting and Searching. Addison-Wesley, Reading, 2nd ed., 1998.
- [30] S. Linton, N. Ruškuc, V. Vatter, Permutation Patterns, London Mathematical Society Lecture Note Series, Volume 376. Cambridge University Press, 2010.
- [31] H. Mularczyk, Lattice paths and pattern-avoiding uniquely sorted permutations. arXiv:1908.04025.
- [32] The On-Line Encyclopedia of Integer Sequences, published electronically at http://oeis.org, 2020.
- [33] M. Sankar, Further bijections to pattern-avoiding valid hook configurations. To appear in Adv. Appl. Math.
- [34] E. Steingrímsson, Generalized permutation patterns a short survey. In: Permutation Patterns (2010), S. Linton, N. Ruškuc, and V. Vatter (eds.), vol. 376 of London Math. Soc. Lecture Note Ser., 137–152.
- [35] A. L. Toom, Problem Solution M655. Kvant, 12 (1981), 28–30. (In Russian.)
- [36] P. Ungar, 2N noncollinear points determine at least 2N directions. J. Combin. Theory Ser. A, 33 (1982), 343–347.
- [37] J. West, Permutations with restricted subsequences and stack-sortable permutations, Ph.D. Thesis, MIT, 1990.
- [38] D. Zeilberger, A proof of Julian West's conjecture that the number of two-stack-sortable permutations of length n is 2(3n)!/((n+1)!(2n+1)!). Discrete Math., **102** (1992), 85–93.

FINE HALL, 304 WASHINGTON RD., PRINCETON, NJ 08544 E-mail address: cdefant@princeton.edu

Fine Hall, 304 Washington Rd., Princeton, NJ 08544

 $E\text{-}mail \ address: \texttt{kzheng@princeton.edu}$