

Prime and Prejudice: Primality Testing Under Adversarial Conditions

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ABSTRACT

This work provides a systematic analysis of primality testing under adversarial conditions, where the numbers being tested for primality are not generated randomly, but instead provided by a possibly malicious party. Such a situation can arise in secure messaging protocols where a server supplies Diffie-Hellman parameters to the peers, or in a secure communications protocol like TLS where a developer can insert such a number to be able to later passively spy on client-server data. We study a broad range of cryptographic libraries and assess their performance in this adversarial setting. As examples of our findings, we are able to construct 2048-bit composites that are declared prime with probability $1/16$ by OpenSSL's primality testing in its default configuration; the advertised performance is 2^{-80} . We can also construct 1024-bit composites that *always* pass the primality testing routine in GNU GMP when configured with the recommended minimum number of rounds. And, for a number of libraries (Cryptlib, LibTomCrypt, JavaScript Big Number, WolfSSL), we can construct composites that *always* pass the supplied primality tests. We explore the implications of these security failures in applications, focusing on the construction of malicious Diffie-Hellman parameters. We show that, unless careful primality testing is performed, an adversary can supply parameters (p, q, g) which on the surface look secure, but where the discrete logarithm problem in the subgroup of order q generated by g is easy. We close by making recommendations for users and developers. In particular, we promote the Baillie-PSW primality test which is both efficient and conjectured to be robust even in the adversarial setting for numbers up to a few thousand bits.

KEYWORDS

Primality testing; Miller-Rabin test; Lucas test; Baillie-PSW test; Diffie-Hellman; TLS

ACM Reference Format:

Martin R. Albrecht, Jake Massimo, Kenneth G. Paterson, and Juraj Somorovsky. 2018. Prime and Prejudice: Primality Testing Under Adversarial

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CCS '18, October 15–19, 2018, Toronto, ON, Canada

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ACM ISBN 978-1-4503-5693-0/18/10...\$15.00

<https://doi.org/10.1145/3243734.3243787>

Conditions. In *CCS '18: 2018 ACM SIGSAC Conference on Computer & Communications Security, Oct. 15–19, 2018, Toronto, ON, Canada*. ACM, New York, NY, USA, 18 pages. <https://doi.org/10.1145/3243734.3243787>

1 INTRODUCTION

Many cryptographic primitives rely on prime numbers, with RSA being the most famous example. However, even in constructions that do not rely on the difficulty of factoring integers into prime factors, primality is often relied upon to prevent an adversary from applying a divide-and-conquer approach (e.g. in the Pohlig-Hellman algorithm or in a Lim-Lee small subgroup attack [60]) or to prevent the existence of degenerate cases such as zero divisors (which may complicate security proofs or reduce output entropy).

One approach to obtaining prime numbers in instantiations of these cryptographic primitives is to produce such numbers as they are needed on whatever device requires them. This is accomplished by sampling random integers and checking for primality. This process can be computationally intensive to the point of being prohibitively so. The high cost of producing prime numbers led implementations to seek ways to reduce this cost and, as demonstrated in [44], these performance improvements may then lead to devastating attacks.

If the required prime numbers are public, an alternative approach is possible: (low-power) devices are provisioned with prime numbers from a server or a standard. For example, the popular Telegram messenger [35] uses Diffie-Hellman (DH) parameters provided by the server to establish end-to-end encryption *between peers*. If the peers do not validate the correctness of the supplied DH parameters,¹ the Telegram server can provide malicious DH parameters with composite group orders and thereby passively obtain the established secrets.

Another example is the Transport Layer Security protocol [18] which can use Diffie-Hellman key exchange to establish master secrets in the handshake protocol. The DH parameters are generated by the TLS server and sent to the client during each TLS handshake.² It is clear that the TLS server provider does not gain any advantage by sending malicious DH parameters to the client since it knows the established master key. However, we can consider an adversarial developer who implements a malicious sever with backdoored DH parameter generation, cf. [19, 63]. If such parameters are accepted by TLS clients and used in the DH key exchange, a passive adversary can observe the traffic and obtain the master key.

¹We stress that they *do* perform validation in the default implementation.

²Up to version 1.2 (inclusive) of the protocol.

Here, weak DH parameters that still pass tests by trusted tools offer a sense of plausible deniability. Moreover, if an application simply silently rejects bad parameters then any countermeasures could be overcome by repeatedly sending malicious parameter sets having a reasonable probability of fooling those countermeasures, until the target client accepts them.

In recent years we have seen several backdoors in cryptographic implementations. For example, NIST standardised the Dual EC pseudorandom number generator (PRNG) which allows an adversary to predict generated random values if it can select a generator point Q and collect enough PRNG output [10]. In 2016 it was shown that Juniper implemented this PRNG which enabled an adversary to passively decrypt VPN sessions [9].

A notable example of a potential backdoor involving a composite number is the security advisory [53] pushed by command-line data transfer utility socat, which is popular with security professionals such as penetration testers. There, the DH prime p parameter was replaced with a new 2048 bit value because “*the hard coded 1024 bit DH p parameter was not prime*”. The advisory goes on to state “*since there is no indication of how these parameters were chosen, the existence of a trapdoor that makes possible for an eavesdropper to recover the shared secret from a key exchange that uses them cannot be ruled out*”, which highlights a real world application of this attack model. Similarly, the prime group parameter p given by Group 23 of RFC5114 [32] for use in DH key exchanges has been found to be partially vulnerable to small subgroup attacks [60]. It might seem that code reviews and the availability of rigorous primality testing (in, say, mathematical software packages, cf. Appendix J) impose high rates of detectability for malicious parameter sets in code or standards, but as these examples highlight, such sets still occur in practice.

Given these incidents we can assume a motivated adversary who is able to implement software serving maliciously generated primes and/or DH parameters. Thus, there is a need for cryptographic applications that rely on third-party primes to perform primality testing. Indeed, many cryptographic libraries incorporate primality testing facilities and thus it appears this requirement is easy to satisfy. However, the primary application of these tests is to check primality (or, more precisely, compositeness) for locally-generated, random inputs during prime generation. Thus, it is a natural question to ask whether these libraries are robust against malicious inputs, i.e. inputs designed to fool the library into accepting a composite number as prime. We refer to this setting as *primality testing under adversarial conditions*.

1.1 Overview of Primality Testing

One of the most widely used primality tests is the Miller-Rabin [41, 51] test. Based upon modular exponentiation by repeated squaring, Miller-Rabin is an efficient polynomial-time algorithm with complexity $O(t \log^3 n)$ where t is the number of trials performed. Yet due to its probabilistic nature, it is well known that a t -trial Miller-Rabin test is only accurate in declaring a given composite number to be composite with probability at least $1 - (1/4)^t$. Arnault [2], Pomerance [50] and Narayanan [43] all explore methods of producing Miller-Rabin pseudoprimes, that is, composite numbers

that when tested by Miller-Rabin, achieve the highest probability of $(1/4)^t$ of being wrongly classified as “probably prime”.

Another common choice is the Lucas test [7], and its more stringent variant the strong Lucas probable prime test. Similarly to the Miller-Rabin test, t trials of a strong Lucas test will declare a given composite number as being composite with probability at least $1 - (4/15)^t$ and as being prime with probability at most $(4/15)^t$ [3]. As with the Miller-Rabin test, there are known methods for constructing strong Lucas pseudoprimes [2].

The Lucas test (strong or standard) can be combined with a single Miller-Rabin test (on base 2) to form what is known as the Baillie-PSW test [49]. Due to slightly longer running times, this test is often only adopted for use in mathematical software packages and seen less in cryptographic libraries. Unlike the Miller-Rabin and Lucas tests when performed alone, there are no known pseudoprimes for the Baillie-PSW test (yet there is no proof that they cannot exist).

Clearly, when conducting a Miller-Rabin or Lucas test, the choice of the parameter t (the number of trials) is critical. Many cryptographic libraries, for example OpenSSL [48], use test parameters originating from [14] as popularised in the *Handbook of Applied Cryptography* [40]. These give the number of iterations of Miller-Rabin needed for an error rate less than 2^{-80} , when testing a random input n . A main result of [14] is that if n is a randomly selected b -bit odd integer, then t independent rounds of Miller-Rabin testing to give an error probability:

$$P(X|Y_t) < b^{3/2} 2^t t^{-1/2} 4^{2-\sqrt{t}b} \quad \text{for } 3 \leq t \leq b/9 \text{ and } b \geq 21,$$

where X denotes the event that n is composite, and Y_t the event that t rounds of Miller-Rabin declares n to be prime. This bound enables the computation of the minimum value t needed to obtain $P(X|Y_t) \leq 2^{-80}$ for a range of bit-sizes b ; see Table 2.

However, these error estimates are for primality testing with Miller-Rabin on randomly generated n . In the *adversarial* setting, we are actually concerned with the probability that t trials of Miller-Rabin (or some other test) declare a *given* n to be prime, given that it is composite. This probability is independent of bit-size, and is at most $(1/4)^t$ if random bases are used in Miller-Rabin tests. Similar remarks apply for both variants of the Lucas test.

1.2 Contributions & Outline

We investigate the implementation landscape of primality testing in both cryptographic libraries and mathematical software packages, and measure the security impact of the widespread failure of implementations to achieve robust primality testing in the adversarial setting.

We review primality testing in Section 2. In Section 3, we then review known techniques for constructing pseudoprimes and extend them with our target applications in mind. In Section 4, we then survey primality testing in cryptographic libraries and mathematical software, evaluating their performance in the adversarial setting. We propose techniques to defeat their tests where we can. Overall, our finding is that most libraries are not robust in the adversarial setting. Our main results in this direction are summarised in Table 1.

As one highlight of our results, we find that OpenSSL with its default primality testing routine will declare certain composites

Table 1: Results of our analysis of cryptographic libraries. This shows how the number of rounds of Miller-Rabin used is determined, whether a Baillie-PSW test is implemented, the documented failure rate of the primality test (that is, the probability that it wrongly declares a composite to be prime), and our highest achieved failure rate for composite input.

Library	Rounds of MR testing	Baillie-PSW?	Documented Failure Rate	Our Highest Failure Rate
OpenSSL 1.1.1-pre6	Default bit-size based	No	$< 2^{-80}$	1/16
GNU GMP 6.1.2	User-defined t	No	$(1/4)^t$	100% for $t \leq 15$
GNU Mini-GMP 6.1.2	User-defined t	No	$(1/4)^t$	100% for $t \leq 101$
Java 10	User-defined t	Yes (≥ 100 bits)	$< (1/2)^t$	0% for ≥ 100 bits
JSBN 1.4	User-defined t	No	$< (1/2)^t$	100%
Libcrypt 1.8.2	User-defined t	No	Not given	1/1024 [†]
Cryptlib 3.4.4	User-defined $t \leq 100$	No	Not given	100%
LibTomMath 1.0.1	User-defined $t \leq 256$	No	$(1/4)^t$	100%
LibTomCrypt 1.18.1	User-defined $t \leq 256$	No	$(1/4)^t$	100%
WolfSSL 3.13.0	User-defined $t \leq 256$	No	$(1/4)^t$	100%
Bouncy Castle C# 1.8.2	User-defined t	No	$(1/4)^t$	$(1/4)^t$
Botan 2.6.0	User-defined t	No	$\leq (1/2)^t$	$(1/4)^t$
Crypto++ 7.0	2 or 12	Yes	Not given	0%
GoLang 1.10.3	User-defined t	Yes	$< (1/4)^t$	0%
GoLang pre-1.8	User-defined t	No	$< (1/4)^t$	100% for $t \leq 13$

[†] When calling the `check_prime` function as opposed to `gcry_prime_check` (or calling `gcry_prime_check` in versions prior to 1.3.0).

n of cryptographic size to be prime with probability 1/16, while the documented failure rate is 2^{-80} . This arises from OpenSSL's reliance on Table 2 to compute the number of rounds of Miller-Rabin testing required, and this number decreases as the size of n increases. As another highlight, we construct a 1024-bit composite that is guaranteed to be declared prime by the GNU GMP library [23] for anything up to and including 15 rounds of testing (the recommended minimum by GMP). This is as a result of GNU GMP initialising its PRNG to a static state and consequently using bases in its Miller-Rabin testing that depend only on n , the number being tested. We also show how base selection by randomly sampling from a fixed list of primes, as in Cryptlib, LibTomCrypt, JavaScript Big Number (JSBN) and WolfSSL, can be subverted: we construct composites n of cryptographic size that are guaranteed to be declared prime by these libraries regardless of how many rounds of testing are performed.

We go on to examine the implications of our findings for applications, focussing on DH parameter testing. The good news is that OpenSSL is not impacted because of its insistence on safe primes for use in DH; that is, it requires DH parameters (p, q, g) for which $q = (p-1)/2$ and both p, q are tested for primality. Our current techniques cannot produce malicious parameters in this case. On the other hand, when more liberal choices of parameter are permitted, as is the case in Bouncy Castle and Botan, we are able to construct malicious DH parameter sets which pass the libraries' testing but for which the discrete logarithm problem in the subgroup generated by g is easy.

We close by discussing avenues for improving the robustness of primality testing in the adversarial setting in Section 6.

1.3 Disclosure and Mitigations

We reported our findings and suggested suitable mitigations based on the outcome of our analysis to OpenSSL, GMP, JSBN, Cryptlib, LibTomMath, LibTomCrypt, WolfSSL, Bouncy Castle and Botan. We give a short review of the outcome of these discussions.

When we reached out to the OpenSSL developers, they were in the process of amending their primality testing code to make it FIPS-compliant [47]. However, these changes do not consider the adversarial scenario on which our paper focuses, and the default settings in OpenSSL remain weak in that scenario. Thus, it is left to the user to choose parameters suitable for this scenario. LibTomMath and LibTomCrypt developers are also in the process of adjusting the primality testing functions within their library. They plan to remove the fixed base Miller-Rabin testing and replace the function with a Baillie-PSW test in accordance with our recommendations [33]. WolfSSL have made several adaptations in an upcoming release [26] to their primality testing in response to our findings. This includes now performing Miller-Rabin with pseudo-random bases, not overriding the user's choice of iterations and increasing the number of rounds performed on prime parameters in DH and DSA check functions. Bouncy Castle have also made changes based upon our findings, by removing the DH verification function and replacing it with a whitelisting approach in upcoming release 1.8.3. They are also looking into performing Baillie-PSW in future versions as per our suggestion. Botan version 2.7.0 [37] has increased the number of rounds of Miller-Rabin performed in DH verification and includes the addition of the Lucas test to perform Baillie-PSW as per our suggestions. GNU GMP, Mini-GMP and Cryptlib all remain unchanged, but the authors of Cryptlib pointed out a code comment that indicates the limitations of their primality test. We received no correspondence from JSBN.

2 BACKGROUND ON PRIMALITY TESTING

A primality test is an algorithm used to determine whether or not a given number is prime. These primality tests come in two different varieties; deterministic and probabilistic. Deterministic primality testing algorithms prove conclusively that a number is prime, but they tend to be slow and are not widely used in practice. A famous example is the AKS test [1]. We do not discuss such tests further in this paper, except where they arise in certain mathematical software.

Probabilistic primality tests make use of arithmetic conditions that all primes must satisfy, and test these conditions for the number n of interest. If the condition *does not* hold, we learn that n must be composite. However, if it does hold we may only infer that n is *probably* prime, since some composite numbers may also pass the test. By making repeated tests, the probability that n is composite conditioned on it having passed some number t of tests can be made sufficiently small for cryptographic applications. A typical target probability is 2^{-80} , cf. [40, 4.49]. A critical consideration here is whether n was generated adversarially or not, since the bounds that can be inferred on probability may be radically different in the two cases; more on this below.

We now discuss three widely-used tests: the Fermat, Miller-Rabin, and Lucas tests.

2.1 Fermat Test

The Fermat primality test is based upon the following theorem.

THEOREM 2.1 (FERMAT'S LITTLE THEOREM). *If p is prime and a is not divisible by p , then*

$$a^{p-1} \equiv 1 \pmod{p}.$$

To test n for primality, one simply chooses a base a and computes $a^{n-1} \pmod{n}$. If $a^{n-1} \not\equiv 1 \pmod{n}$, then we can be certain that n is composite. If after testing a variety of bases a_i , we find that they all satisfy $a_i^{n-1} \equiv 1 \pmod{n}$, we may conclude that n is probably prime.

It is well known that there exists composite numbers that satisfy $a^{n-1} \equiv 1 \pmod{n}$ for all integers a that are not divisible by n . These numbers completely thwart the Fermat test, and are known as Carmichael numbers. These will be of relevance in the sequel. The following result is fundamental in the construction of Carmichael numbers.

THEOREM 2.2 (KORSELT'S CRITERION). *A positive composite integer n is a Carmichael number if and only if n is square-free, and $p-1 \mid n-1$ for all prime divisors p of n .*

2.2 Miller-Rabin Test

The Miller-Rabin [41, 51] primality test is based upon the fact that there are no non-trivial roots of unity modulo a prime. Let $n > 1$ be an odd integer to be tested and write $n = 2^e d + 1$ where d is odd. If n is prime, then for any integer a with $1 \leq a < n$, we have:

$$a^d \equiv 1 \pmod{n} \text{ or } a^{2^i d} \equiv -1 \pmod{n} \text{ for some } 0 \leq i < e.$$

The Miller-Rabin test then consists of checking the above conditions, declaring a number to be (probably) prime if one of the two conditions holds, and to be composite if both fail. If one condition

holds, then we say n is a *pseudoprime to base a* , or that a is a *non-witness to the compositeness of n* (since n may be composite, but a does not demonstrate this fact).

For a composite n , let $S(n)$ denote the number of non-witnesses $a \in [1, n-1]$. An upper-bound on $S(n)$ is given by results of [42, 51]:

THEOREM 2.3 (MONIER-RABIN BOUND). *Let $n \neq 9$ be odd and composite. Then*

$$S(n) \leq \frac{\varphi(n)}{4}$$

where φ denotes the Euler totient function.

This bound will be critical in determining the probability that an adversarially generated n passes the Miller-Rabin test. Since for large n , we have $\varphi(n) \approx n$, it indicates that no composite n can pass the Miller-Rabin test for t random bases with probability greater than $(1/4)^t$. Hence achieving a target probability of 2^{-80} requires $t \geq 40$. The test is commonly implemented using either (a) a set of fixed bases (e.g. JSBN) or (b) randomly chosen bases (e.g. OpenSSL). Of course, the $(1/4)^t$ bound only holds in the case of randomly chosen bases.

2.3 Lucas Test

The Lucas primality test [7] makes use of Lucas sequences, defined as follows:

Definition 2.4 (Lucas sequence [3]). Let P and Q be integers and $D = P^2 - 4Q$. Then the Lucas sequences (U_k) and (V_k) (with $k \geq 0$) are defined recursively by:

$$\begin{aligned} U_{k+2} &= PU_{k+1} - QU_k & \text{where,} & & U_0 &= 0, U_1 = 1, \\ V_{k+2} &= PV_{k+1} - QV_k & & & V_0 &= 2, V_1 = P. \end{aligned}$$

The Lucas probable prime test then relies on the following theorem (in which $\left(\frac{x}{p}\right)$ denotes the Legendre symbol, with value 1 if x is a square modulo p and value -1 otherwise):

THEOREM 2.5 ([12]). *Let P, Q and D and the Lucas sequences $(U_k), (V_k)$ be defined as above. If p is a prime with $\gcd(p, 2QD) = 1$, then*

$$U_{p-\left(\frac{x}{p}\right)} \equiv 0 \pmod{p}. \tag{1}$$

The Lucas probable prime test repeatedly tests property (1) for different pairs (P, Q) . This leads to the notion of a Lucas pseudoprime with respect to such a pair.

Definition 2.6 (Lucas pseudoprime). Let n be a composite number such that $\gcd(n, 2QD) = 1$. If $U_{n-\left(\frac{x}{n}\right)} \equiv 0 \pmod{n}$, then n is called a *Lucas pseudoprime* with respect to parameters (P, Q) .

We can now introduce the notion of a strong Lucas probable prime and strong Lucas pseudoprime with respect to parameters (P, Q) by the following theorem.

THEOREM 2.7 ([3]). *Let p be a prime number not dividing $2QD$. Set $p - \left(\frac{D}{p}\right) = 2^k q$ with q odd. Then one of the following conditions is satisfied:*

$$p \mid U_q \text{ or } \exists i \text{ such that } 0 \leq i < k \text{ and } p \mid V_{2^i q}. \tag{2}$$

The strong Lucas probable prime test repeatedly tests property (2) for different pairs (P, Q) . This leads to the definition of a strong Lucas pseudoprime with respect to parameters (P, Q) as follows.

Definition 2.8 (strong Lucas pseudoprime). Let n be a composite number such that $\gcd(n, 2QD) = 1$. Set $n - \left(\frac{D}{n}\right) = 2^k q$ with q odd. Suppose that:

$$n \mid U_q \quad \text{or} \quad \exists i \text{ such that } 0 \leq i < k \text{ and } n \mid V_{2^i q}.$$

Then n is called a *strong Lucas pseudoprime* with respect to parameters (P, Q) .

A strong Lucas pseudoprime is also a Lucas pseudoprime (for the same (P, Q) pair), but the converse is not necessarily true. The strong version of the test is therefore seen as the more stringent option.

REMARK. *The Lucas pseudoprime and strong Lucas pseudoprime tests are also known as a Lucas-Selfridge test and a strong Lucas-Selfridge test respectively, specifically when used with Selfridge's parameters $P = 1, Q = -1$.*

Analogously to the Monier-Rabin theorem for pseudoprimes for the Miller-Rabin primality test, Arnault [3] showed that for an integer D and n a composite with $\gcd(D, n) = 1$ and $n \neq 9$, the number of pairs (P, Q) with $0 \leq P, Q < n, \gcd(Q, n) = 1, P^2 - 4Q \equiv D \pmod{n}$ such that n is strong Lucas pseudoprime with respect to (P, Q) is at most $4n/15$. There is an exception to this result for certain forms of twin primes (we omit the details here), but Arnault goes on to prove that even these particular forms of twin prime n have at most $n/2$ pairs (P, Q) such that n is a strong Lucas pseudoprime with respect to (P, Q) . From this, we can infer that t applications of the strong Lucas test would declare a composite n to be probably prime with a probability at most $(4/15)^t$.

2.4 Baillie-PSW

The Baillie-PSW [49] test is a probabilistic primality test formed by combining a single Miller-Rabin test with base 2 with either a Lucas or strong Lucas pseudoprime test. The idea of this test is that the two components are "orthogonal" and so it is very unlikely that a number n will pass both parts. Indeed, there are no known composite n that pass the Baillie-PSW test. Gilchrist [20] confirmed that there are no Baillie-PSW pseudoprimes less than 2^{64} . PRIMO [38] is an elliptic curve based primality proving program that uses the Baillie-PSW test to check all intermediate probable primes. If any of these values were indeed composite, the final certification would necessarily have failed. Since this has never occurred during its use, PRIMO's author Martin estimates [62] that there are no Baillie-PSW pseudoprimes with less than about 10000 digits. This empirical evidence suggests that numbers of cryptographic size for use in Diffie-Hellman and RSA are unlikely to be Baillie-PSW pseudoprimes. However, Pomerance gives a heuristic argument in [49] that there are in-fact infinitely many Baillie-PSW pseudoprimes. The construction of a single example is a significant open problem in number theory.

3 CONSTRUCTING PSEUDOPRIMES

In this section, we review known methods of constructing pseudoprimes for the Miller-Rabin and Lucas tests. We also provide variations on these methods. We will use the results of this section in the next one, where we study the robustness of cryptographic libraries for primality testing in the adversarial setting.

3.1 Miller-Rabin Pseudoprimes

The exact number of non-witnesses $S(n)$ for any composite number n can be computed given the factorisation of n [12]. Generating composites n that have large numbers of non-witnesses is not so straightforward. In empirical work, Pomerance *et al.* [50] showed that many composite numbers that pass a Miller-Rabin primality test have the form $n = (k + 1)(rk + 1)$ where r is small and $k + 1$ is prime. More recently, Höglund [25] and Nicely [45] used the Miller-Rabin primality test as implemented in GNU GMP to test randomly generated numbers of this form for various values of r and for various different sizes of k . Their results support the claims made by [50].

We now consider existing methods for producing composites which have many non-witnesses, for two forms of the Miller-Rabin test: firstly where the bases are chosen randomly and secondly where a fixed set of bases is used.

3.1.1 Random Bases. For random bases, we are interested in constructing composite n that have large numbers of non-witnesses, i.e. for which $S(n)$ is large. Such numbers will pass the Miller-Rabin test with probability $S(n)/n$ per trial; of course, this probability is bounded by $\varphi(n)/4n \approx 1/4$ by the Monier-Rabin theorem, but we are interested in how close to this bound we can get. We rely on the following:

THEOREM 3.1 ([43]). *Consider an odd composite integer n with m distinct prime factors p_1, \dots, p_m . Suppose that $n = 2^e \cdot d + 1$ where d is odd. Also suppose that $n = \prod_{i=1}^m p_i^{t_i}$ where each p_i can be expressed as $2^{e_i} \cdot d_i + 1$ with each d_i odd. Then*

$$S(n) = \prod_{i=1}^m \gcd(d, d_i) \cdot \left(\frac{2^{\min(e_i) \cdot m} - 1}{2^m - 1} + 1 \right). \quad (3)$$

Note how the bound in this theorem does not depend on the exponents t_i , indicating that square-free numbers will have relatively large $S(n)$. Also note the dependence on the terms $\gcd(d, d_i)$, indicating that ensuring that the odd part of each prime factor p_i has a large gcd with the odd part of n is necessary for large n . As an easy corollary of this theorem, we obtain:

COROLLARY 3.2 ([43]). *Let x be an odd integer such that $2x + 1$ and $4x + 1$ are both prime. Then $n = (2x + 1)(4x + 1)$ has $\varphi(n) = 8x^2$ and achieves the Monier-Rabin bound, i.e. it satisfies $S(n) = \varphi(n)/4$.*

The proof of this corollary follows easily on observing that we may take $m = 2$ and $d = d_1 = d_2 = x$ in the preceding theorem. Narayanan [43] also showed that if n is a Carmichael number of the form $p_1 p_2 p_3$, where each p_i is a distinct prime with $p_i \equiv 3 \pmod{4}$, then $S(n)$ achieves the Monier-Rabin bound. He also gave further results showing that these two forms for n are the only ones achieving the Monier-Rabin bound, with all other n satisfying $S(n) \leq \varphi(n)/6$.

3.1.2 Fixed Bases. Some implementations of the Miller-Rabin primality test select bases from a fixed list (often of primes), rather than choosing them at random. For example, until 2010, the Py-Crypto 2.1.0 (2009) [34] primality test `isPrime()` performed 7 rounds of Miller-Rabin using the first 7 primes as bases, while LibTomMath chooses the first t entries from a hard-coded list of primes as bases.

Arnault [2] presented a method for producing composite numbers $n = p_1 p_2 \cdots p_h$ that are guaranteed to be declared prime by Miller-Rabin for any fixed set of prime bases $A = \{a_1, a_2, \dots, a_t\}$. We give an overview and examples of Arnault's method in Appendix A.

Since fixed base Miller-Rabin tests are relatively uncommon in implementations, it might seem that Arnault's method would not be very useful. We shall however see that this method is particularly helpful when an implementation chooses bases randomly from a large fixed list of possibilities. For example, an implementation might select prime bases randomly from a list of primes below 1000; since Arnault's method scales well (we simply need to solve more congruences simultaneously with the CRT) we can use this method to produce a composite n such that all primes below 1000 are non-witnesses for n . We shall see applications of this approach for different libraries in Sections 4.3, 4.5, 4.7, 4.8, 4.9 and 4.10.

3.1.3 Hybrid Technique. The method above produces composites that are in fact always Carmichael numbers. We know from Section 3.1.1 that if n is a Carmichael number with 3 distinct prime factors all congruent to 3 (mod 4), then n has the maximum number of non-witnesses, $\varphi(n)/4$. We can set $h = 3$ in Arnault's method and tweak it slightly to ensure that, as well as producing n with a specified set A of non-witnesses, it produces an n meeting the Monier-Rabin bound, so that random base Miller-Rabin tests will also pass with the maximum probability. The tweak is very simple: we ensure that $2 \in A$; this forces $p_1 \equiv 3$ or $5 \pmod{8}$; we then select $p_1 \equiv 3 \pmod{8}$ so that $p_1 \equiv 3 \pmod{4}$. Arnault's method sets $p_i = k_i(p_1 - 1) + 1$ where the k_i are co-prime to all the elements of A . Since $2 \in A$, the k_i must all be odd; it is easy to see that this forces $p_i \equiv 3 \pmod{4}$ too.

We will give an application of this technique in Section 4.6.

3.1.4 Extension For Composite Fixed Bases. The method of Arnault [2] works (as presented) only for prime bases, and not for composite bases. Although less common, some implementations use both prime and composite bases in their Miller-Rabin testing. By setting $n \equiv 3 \pmod{4}$, we know that $e = 1$ when writing $n = 2^e \cdot d + 1$ for d odd. In this case, the conditions to pass the Miller-Rabin test simply become $a^{(n-1)/2} \equiv \pm 1 \pmod{n}$. Hence, if $n \equiv 3 \pmod{4}$ is pseudoprime to some set of bases $\{a_1, a_2, \dots, a_t\}$, then n is also pseudoprime for any base b arising as a product $b = a_1^{e_1} \cdot a_2^{e_2} \cdots a_t^{e_t} \pmod{n}$ (for any set of indices $e_i \in \mathbb{Z}$). Therefore we can construct a composite n that is pseudoprime with respect to any list of bases $\{b_1, \dots, b_t\}$ (of which any number can be composite) by using the hybrid method described in Section 3.1.3, but with set A in that method being the complete set of prime factors arising in the b_i . Note that in this method, n is of the form $n = p_1 p_2 p_3$ where each $p_i \equiv 3 \pmod{4}$, so we have $n \equiv 3 \pmod{4}$ as needed. Moreover,

because of the form of n , the composites generated in this manner will also meet the Monier-Rabin bound.

We will give an application of this technique in Section 4.3, where we study Mini-GMP [23] which uses Euler's polynomial to generate Miller-Rabin bases.

3.2 Lucas Pseudoprimes

Like Miller-Rabin pseudoprimes, Lucas pseudoprimes are with respect to some choice of test parameters. Throughout this work we follow Selfridge's Method A [7] of parameter selection, which is summarised as follows:

Definition 3.3 (Selfridge's Method A [7]). Let D be the first element of the sequence $5, -7, 9, -11, 13, \dots$ for which $\left(\frac{D}{n}\right) = -1$. Then set $P = 1$ and $Q = (1 - D)/4$.

There are two reasons for studying this particular method for setting parameters. The first is that it is the parameter choice used when performing the Lucas part of the Baillie-PSW primality test [7, 50]. The second is that this is the method that both Java [11] and Crypto++ [13] libraries that we study use in their implementation of the Lucas test.

The Lucas and strong Lucas-probable prime tests with this parameter choice are commonly referred to in the literature as Lucas and strong Lucas-Selfridge probable prime tests. Pseudoprimes for this parameter choice are well-documented. The OEIS sequence A217120 [5] presents a small list of them, referring to a table of all Lucas pseudoprimes below $10^{14} \approx 2^{47}$ compiled by Jacobsen [28]. There is an equivalent sequence A217255 [6] for strong Lucas pseudoprimes. Any pseudoprime for the strong Lucas probable prime test with respect to some parameter set (P, Q) , is also a pseudoprime for the Lucas probable prime test.

Arnault [2] also presented a scalable method that takes as input a set of parameter choices $\{(P_1, Q_1, D_1), (P_2, Q_2, D_2), \dots, (P_t, Q_t, D_t)\}$ and returns a composite n of the form $n = p_1 p_2 \cdots p_h$ that is a strong Lucas pseudoprime to the parameters (P_i, Q_i, D_i) for all $1 \leq i \leq t$. The method is similar to that for constructing Miller-Rabin pseudoprimes for fixed bases, but differs in its details. In particular, the two construction methods are sufficiently different that it seems hard to derive a single method producing n that are pseudoprimes for both the Miller-Rabin and Lucas tests.

3.2.1 A specialisation of Arnault [2] for Selfridge's Method A. For Selfridge's Method A, we know that if we take an n such that $\left(\frac{5}{n}\right) = -1$, then a single test on n with parameter set $(P, Q, D) = (1, -1, 5)$ will be performed. We next show how to specialise Arnault's construction [2] so that it will produce composites n that are guaranteed to be declared prime by a strong Lucas test for this parameter set.

Following Arnault's construction, we consider n of the form $n = p_1 p_2 p_3$ where $p_i = k_i(p_1 + 1) - 1$ for $i \in \{2, 3\}$, with k_2 and k_3 odd integers.

We first note that the p_i must satisfy certain conditions with respect to Legendre symbols (see [2, Lemmas 6.1 and 6.2]):

$$\left(\frac{D}{p_i}\right) = \left(\frac{Q}{p_i}\right) = -1 \quad \text{for all } i \text{ such that } 1 \leq i \leq 3.$$

With our single parameter set $(P, Q, D) = (1, -1, 5)$, this becomes:

$$\left(\frac{-1}{p_i}\right) = \left(\frac{5}{p_i}\right) = -1 \quad \text{for all } i \text{ such that } 1 \leq i \leq 3. \quad (4)$$

Now $\left(\frac{-1}{p_i}\right) = -1 \Leftrightarrow p_i \equiv 3 \pmod{4}$. Since $p_i = k_i(p_1 + 1) - 1$ for $i \in \{2, 3\}$, and the k_i are odd, then it is easy to show that if $p_1 \equiv 3 \pmod{4}$ then it follows that $p_i \equiv 3 \pmod{4}$ for $i = 2, 3$ as well. We also have that $\left(\frac{5}{p_i}\right) = -1 \Leftrightarrow p_i \equiv 2 \text{ or } 3 \pmod{5}$. Therefore condition (4) is satisfied when $p_1 \equiv 3 \text{ or } 7 \pmod{20}$ (by the CRT) and $p_i \equiv 2 \text{ or } 3 \pmod{5}$ for $i \geq 2$.

At this point we must choose k_2, k_3 and add conditions that ensure the coefficients in [2, Lemma 6.1] are indeed integers. These conditions are simple:

$$p_1 \equiv k_3^{-1} \pmod{k_2} \quad \text{and} \quad p_1 \equiv k_2^{-1} \pmod{k_3}.$$

We choose to fix $p_1 \equiv 7 \pmod{20}$ and select $(k_2, k_3) = (31, 43)$. This produces our final congruence that prime p_1 must satisfy: $p_1 \equiv 6647 \pmod{26660}$. We now search for a prime p_1 that satisfies this congruence, and such that p_2 and p_3 satisfying $p_i = k_i(p_1 + 1) - 1$ for $i = 2, 3$ are also primes with $p_2 \equiv p_3 \equiv 2 \text{ or } 3 \pmod{5}$.

The smallest solution is the following:

$$p_1 = 486527, p_2 = 15082367, p_3 = 20920703$$

This yields a 68-bit $n = 15351567445511174527$ which indeed does pass the strong Lucas test using Selfridge's Method A for parameter selection. Of course, we can take any (p_1, p_2, p_3) satisfying the above conditions (which are not too onerous to satisfy), and in this sense the method scales well to numbers n of cryptographically interesting size. For example, Appendix B shows a 2050-bit example generated using the above procedure.

This generation technique is also versatile, as we can simply include additional parameters in our set dependent on which parameter selection methods a particular test uses. This allows us to generate composites that are declared prime by a variety of strong Lucas tests, at the small cost of solving a few more simultaneous congruences with the CRT.

4 CRYPTOGRAPHIC LIBRARIES AND MATHEMATICS PACKAGES

Many cryptographic libraries offering implementations of common cryptographic protocols also provide a toolkit for handling arbitrary-precision integer arithmetic, including primality testing. These functions would be used, for example, for testing the primality of Diffie-Hellman parameters.

This section provides a survey of primality testing in a broad and representative range of cryptographic libraries (OpenSSL, GNU GMP and Mini-GMP, Java, JavaScript Big Number (JSBN), Libgcrypt, Cryptlib, LibTomMath, LibTomCrypt, WolfSSL, Bouncy Castle, Botan, Crypto++ and GoLang). For each library, we first describe how it implements primality testing. We then tailor a composite likely to be declared prime by each particular library, and quantify the probability that our composite passes the library's primality test (so that the primality test fails). Our findings are summarised in Table 1. Throughout, we will refer to the number of rounds of Miller-Rabin testing as t .

Table 2: The rounds t of Miller-Rabin performed chosen by OpenSSL when testing b -bit integers with checks = BN_prime_checks.

b	t	b	t
$b \geq 1300$	2	$400 > b \geq 350$	8
$1300 > b \geq 850$	3	$350 > b \geq 300$	9
$850 > b \geq 650$	4	$300 > b \geq 250$	12
$650 > b \geq 550$	5	$250 > b \geq 200$	15
$550 > b \geq 450$	6	$200 > b \geq 150$	18
$450 > b \geq 400$	7	$150 > b$	27

4.1 OpenSSL

OpenSSL is the most widely used open source cryptographic library and TLS implementation. Throughout, we consider OpenSSL 1.1.1-pre6 [48], although the components studied are largely stable across releases and remain similar to that of the early releases (version 0.9.6 of Sept. 2000).

Analysis. The primality tests in OpenSSL reside in the crypto library, which also houses a wide range of implementations of cryptographic algorithms. The services provided by the crypto library are used by the OpenSSL implementations of SSL, TLS and S/MIME, and have also been used to implement SSH, OpenPGP, and other cryptographic standards.

The functions called upon to perform primality testing in the OpenSSL BIGNUM library are BN_is_prime_ex and BN_is_prime_fasttest_ex found in bn_prime.c. The bulk of the primality testing algorithm is done in BN_is_prime_fasttest_ex where t = checks rounds of Miller-Rabin are performed, each with a randomly chosen base. The checks variable is provided as a parameter to the primality verification function. The function BN_is_prime_ex simply calls BN_is_prime_fasttest_ex without doing any trial divisions. The composites n that we produce have factors much larger than those in the trial divisions that OpenSSL performs. This means that, for our purposes, the result of calling either function is equivalent. Therefore we will focus only on BN_is_prime_fasttest_ex.

Number of Miller-Rabin rounds. Both primality testing functions allow the user to determine the rounds of Miller-Rabin performed. The documentation indicates that if the user sets the value of checks to the variable BN_prime_checks, then the number of Miller-Rabin iterations t is chosen such that the probability of a Miller-Rabin test declaring a *random* composite number n as prime is less than 2^{-80} . The number of rounds performed is then based on the bit-size b of the number n being tested. The relationship between these two values is shown in Table 2. The entries here are based on average case error estimates taken from the Handbook of Applied Cryptography [40], which in turn references [14].

Base Selection. OpenSSL chooses the Miller-Rabin bases it uses in a pseudorandom manner, by using OpenSSL's function BN_rand_range() with an optional flag set to PRIVATE. This then calls bnrand to generate a pseudorandom base a in the range $1 \leq a < n$ using a cryptographically strong pseudorandom number generator with entropy inputs gathered from the operating system, cf. [55] for details on OpenSSL's random number generation.

We recall that this n will always pass GMP's primality testing for 15 MR rounds because the generation of the MR bases depends deterministically on n .

4.3 Mini-GMP

Mini-GMP is a small implementation of a subset of GMP's `mpz` and `mpq` interfaces included within GMP 6.1.2 [23]. This library includes its own miniature implementation of `mpz_probab_prime_p(n, reps)`. The most significant change compared to GMP is that Miller-Rabin is performed explicitly with a deterministic sequence of t bases obtained by evaluating Euler's polynomial $a(x) = x^2 + x + 41$ at $x = 0, 1, 2, \dots, t - 1$. It also omits GMP's Fermat test.

Pseudoprimes. The use of a sequence of deterministic bases in Mini-GMP enables us to predict the bases that will be chosen for any particular value t of `reps`. The bases are not all prime (though Euler's polynomial famously does produce many primes), so we cannot directly use Arnault's method from Section 3.1.2. Instead, we use our extension for composite, fixed bases method in Section 3.1.4.

Using this approach, we constructed a 2960-bit composite $n = p_1 p_2 p_3$ that passes up to $t = 101$ rounds of Mini-GMP's Miller-Rabin testing. Of the 101 bases produced by Euler's polynomial, 86 were already primes and the remaining 15 bases all factorised into various combinations of the four primes 163, 167, 179 and 199. The combined list of 90 unique primes was then used with the method described in Section 3.1.4 to produce n . This n is given in Appendix D. We note that the documentation for Mini-GMP is shared with the main GMP library, implying to a user that 15 to 50 rounds of MR testing would be reasonable.

4.4 Java

Java implementations provide their own methods for arbitrary precision arithmetic, including primality tests, as seen in `java.math.BigInteger`. We consider OpenJDK10 [11], although there seems to be no significant changes to this section of the code in older versions such as JDK8.

Analysis. The primality testing function `isProbablePrime` is passed a single parameter `certainty`. This is a value chosen by the user and is described in the documentation as: "a measure of the uncertainty that the caller is willing to tolerate: if the call returns true the probability that this `BigInteger` is prime exceeds $(1 - 1/2^{\text{certainty}})$." The `certainty` parameter is then used to determine how many rounds of testing will be performed. This is done by calling the function `primeToCertainty`. This function first sets a variable `n` as $(\text{certainty} + 1)/2$. This would produce a non-integer result when `certainty` is even, yet the result is cast to an integer, implicitly flooring the result.⁴

This function also takes into consideration the bit-size of the number being tested; if it is less than 100, then Miller-Rabin is performed with at most 50 rounds; if it is greater than 100, then both Miller-Rabin and a Lucas probable prime test with Selfridge's

⁴Because of the role that `n` plays in determining the number of rounds of Miller-Rabin to be performed, the result is that there is no difference in testing `isProbablePrime(k)` and `isProbablePrime(k+1)` when `k` is odd. This has an effect on the assurance given to the user — the guarantee of $1 - 1/2^{\text{certainty}}$ is no longer accurate for half of the values of `certainty`.

parameters are performed, as described in Section 3.2. In the latter case, the maximum number of rounds of Miller-Rabin is determined based on the bit-size of the tested number, similarly to OpenSSL. In both cases, the user's choice of `certainty` will determine the actual number of rounds of Miller-Rabin performed only if it is *less* than the internally-specified number for that bit-size.

Pseudoprimes. For numbers of cryptographically interesting size, Java performs both Miller-Rabin and Lucas probable prime tests. Using the method outlined in Section 3.2 we could produce composites that are guaranteed to be declared prime by the Lucas test. However, the resulting forms do not fit into any of the known families of composites having high numbers of Miller-Rabin non-witnesses. Hence, we are unable to construct any numbers passing Java's primality test with high probability using our current techniques.

4.5 JavaScript Big Number (JSBN)

The JavaScript Big Number (JSBN) library written by Tom Wu [64] provides a small cryptographic toolkit for JavaScript applications. Here we study the most recent release JSBN 1.4 from 2013. According to its homepage the library has been used in a variety of applications, including: Forge (a pure JavaScript implementation of SSL/TLS), Google's V8 benchmark suite version 6, the JavaScript Cryptography Toolkit and the RSA-Sign JavaScript library.

Analysis. The library offers the primality test `bnIsProbablePrime(t)` where the parameter `t` defines the number of rounds of Miller-Rabin the user wishes to perform. The code documentation states that this function will "test primality with certainty $\geq 1 - .5^t$ ". The function pseudorandomly chooses a base `a` for each round of Miller-Rabin from a hard-coded list of all primes below 1000 called `lowprimes`.

Pseudoprimes. We can consider this implementation as performing tests with fixed bases, where the bases chosen are all primes between 2 and 1000. We can then use Arnault's method (Section 3.1.2) to construct composite numbers n that pass JSBN's primality test no matter how many rounds of testing t the user wishes to perform. For example, we used SageMath 7.6 [54] to obtain a 4279-bit composite n having 3 prime factors, see Appendix E for the details.

4.6 Libcrypt

Libcrypt [31] is a general purpose cryptographic library originally based on code from GnuPG. The library provides various cryptographic functions, including public key algorithms, large integer functions and primality testing. We analyse the current stable version 1.8.2, released in December 2017

Analysis. The documentation for Libcrypt states that the function used for checking the primality of primes is `gcry_prime_check` which is found in `primegen.c`. This function then calls `check_prime` in which the actual testing performed. This function `check_prime` performs three testing steps. The first step is trial division with all primes up to 4999. The second step is a Fermat test with base $a = 2$. The last step comprises t rounds of Miller-Rabin where the bases are pseudorandomly chosen. We note that t is user defined, but cannot be set to less than 5. The default for checking the numbers

produced in the prime generation algorithm is set to 5, but when a user calls `gcry_prime_check` the choice of t is hard-coded to 64.

Pseudoprimes. Following Section 3.1, beating steps 1 and 2 of the testing performed in `check_prime` is trivial if we choose n as a Carmichael number of the form $n = pqr$ where $p, q, r > 4999$. By using the hybrid technique in Section 3.1.3, we can create a Carmichael number that also has the maximum number of randomly distributed non-witnesses. We then need only to overcome the t Miller-Rabin tests with pseudorandom bases. This happens with probability $(1/4)^t$. If the user calls `gcry_prime_check` then the probability with which we can fool this test would be only 2^{-128} . Yet performing 64 rounds of Miller-Rabin is quite time consuming, and a user may be tempted to bypass `gcry_prime_check` and call `check_prime` with fewer rounds. In this hypothetical situation, or in versions of Libgcrypt prior to 1.3.0 (2007) [30] (where `gcry_prime_check` would call $t = 5$ rounds by default) the best we could achieve is passing the test with probability $1/1024$ (for $t = 5$).

4.7 Cryptlib

Cryptlib 3.4.3 [24] is an open source security toolkit library developed by Peter Gutmann. It provides a variety of services including: public key algorithms, various cryptographic functions and primality testing.

Analysis. The primality test in Cryptlib is the function `prime-Probable` found in `kg_prime.c` and is composed of t rounds of Miller-Rabin, where the value of t must be between 1 and 100 (inclusive) and is chosen by the user upon calling. The function then chooses the base for each test incrementally from the start of a fixed list of primes. This is either a list of the first 54 primes (2 to 251) or the first 2048 primes (2 to 17863), depending on the preprocessor directive `CONFIG_CONSERVE_MEMORY`.

Pseudoprimes. Since $t \leq 100$, we will at most only ever test using the primes between 2 and 541 (the hundredth prime) as bases. We can therefore generate numbers that are guaranteed to be declared prime by this test for any valid input t , simply by using Arnault's method to generate a composite n that has the first 100 primes as non-witnesses. Indeed, using the method described in Section 3.1.2 we can produce a 2329-bit composite that is pseudoprime to all prime bases up to and including 541. See Appendix F for details.

4.8 LibTomMath

LibTomMath v1.0.1 [16] is an open source multiple-precision integer library with a number theoretic toolkit.

Analysis. LibTomMath includes several methods for primality testing in the form of trial division, Fermat tests, and Miller-Rabin tests. The latter two take a single base a and a number n to test as arguments and return whether or not a is a witness or non-witness. The main primality test is defined by the function `mp_prime_is_prime`, which takes arguments n (the number to be tested), and integer t with $1 \leq t \leq 256$. It then performs some trial divisions (on a default of the first 256 primes) and then t rounds of Miller-Rabin. The selection of bases to be used is made similarly as in Cryptlib: it simply picks incrementally from a list of hard-coded primes (but this time a list of 256 primes up to 1619 are used).

The documentation of LibTomMath (`bn.pdf`) discusses the number of rounds of Miller-Rabin required with the statement: “Generally to ensure a number is very likely to be prime you have to perform the Miller-Rabin with at least a half-dozen or so unique bases.” This is complemented with a function `mp_prime_rabin_miller_trials` that gives the number of rounds needed to achieve an error rate less than 2^{-96} based on the bit-size of the number tested (similar to that in OpenSSL and [14]) and a comment in the header file `tommath.h` above `mp_prime_rabin_miller_trials` that states the probability of a false classification is no more than $(1/4)^t$.

Pseudoprimes. Since the bases are chosen deterministically based on the value of t , we can achieve a failure rate of 100% simply by using the method of Section 3.1.2 to produce a composite n that has the first 256 primes as non-witnesses; such an n is guaranteed to be declared prime by `mp_prime_is_prime`, for any value of t (including the t chosen by `mp_prime_rabin_miller_trials` that describes an error rate less than 2^{-96}). Appendix G provides a 7023-bit example of such an n . Much smaller examples can be obtained if smaller values of t are guaranteed to be used; in particular, we can easily obtain a 1024-bit example for $t \leq 40$ (see also Appendix G).

4.9 LibTomCrypt

LibTomCrypt v1.18.1 [15] is an additional cryptographic toolkit that shares many resources with LibTomMath.

Analysis. The primality test in LibTomCrypt is called `isprime(n, t, result)`. It takes as arguments an n to test and carries out t rounds of Miller-Rabin. The documentation of LibTomCrypt advises that each round of Miller-Rabin reduces the probability of n being a pseudoprime by a factor of 4, and therefore deduces that the overall error is at most $(1/4)^t$. LibTomCrypt supports selection from three different big integer libraries at runtime.

If LibTomMath is chosen then `isprime` will call `mp_prime_is_prime` as described in Section 4.8, passing on parameters n and t . If TomsFastMath [17] is chosen then `isprime` will call `fp_isprime_ex`, a function defined in the math library TomsFastMath that performs equivalent testing as LibTomMath's `mp_prime_is_prime`. If GMP is selected then `isprime` will call `mpz_probab_prime_p` as described in Section 4.2. The value of t used by any of the three choices is inherited from the original call to `isprime`, however if $t = 0$ the value is overwritten to $t = 40$.

Pseudoprimes. If either LibTomMath or TomsFastMath are selected, the pseudoprimes described in Section 4.8 (see Appendix G) will always be declared prime by the primality test. If GMP is selected we can apply the analysis in Section 4.2 to generate pseudoprimes (see Appendix C).

4.10 WolfSSL

WolfSSL 3.13.0 [27] (formerly CyaSSL) is a small SSL/TLS library targeted for use in embedded systems. WolfSSL provides primality testing tools based on public domain TomsFastMath 0.10 [17] and LibTomMath 0.38 [16] functions.

Analysis. The primality test in WolfSSL is the function `mp_prime_is_prime` which takes a number n to be tested and the rounds of testing t as parameters. This function is directly taken from an

older version of LibTomMath, namely 0.38 [16]. WolfSSL will use LibTomMath by default, but can optionally be compiled to use TomsFastMath 0.10 [17] at runtime. The primality test in LibTomMath 0.38 is unchanged from that analysed in version 1.0.1 in Section 4.8. When using TomsFastMath, `mp_prime_is_prime` calls `fp_isprime` which strips the user's choice of t and simply calls `fp_isprime_ex` with the hard-coded value of $t = 8$. The function `fp_isprime_ex` then performs trial division (on a default of the first 256 primes) and then does 8 rounds of Miller-Rabin using the first 8 primes as bases. It thus acts equivalently to `mp_prime_is_prime` in LibTomMath, but with $t = 8$.

Pseudoprimes. Since the testing in WolfSSL is in effect the same as that performed in LibTomMath (but using only 8 rounds of Miller-Rabin when using TomsFastMath), the composite examples given in Appendix G are also declared prime with 100% success.

4.11 Bouncy Castle

Bouncy Castle is a cryptographic library written in Java and C# [46]. The primality test in Bouncy Castle Java is based on the `BigInteger` class from JDK as described in Section 4.4. Bouncy Castle C# implements its own primality tests. We analyse Bouncy Castle C# version 1.8.2.

Analysis. The relevant function responsible for primality tests is located in the class `BigInteger`. This class provides method `IsProbablePrime` which accepts certainty as a parameter. The method then uses Miller-Rabin tests with t rounds, where t is computed as $t = ((\text{certainty} - 1)/2) + 1$. In each round the base is selected using a secure random number generator (`SecureRandom`) which is provided by the Bouncy Castle library.

The certainty parameter must always be provided to invocation of the `IsProbablePrime` method. Therefore, a user choice completely determines how many Miller-Rabin rounds are performed. For example, this method is directly used in the `TlsDHUtilities` class, which provides Diffie-Hellman operations for TLS. When validating the incoming DH parameters, the `ValidatedDHParameters` method invokes `isProbablePrime` with `certainty = 2`. This results in only a single Miller-Rabin test being carried out.

Pseudoprimes. We can produce composites n using any of the methods in Section 3.1; such n meet the Monier-Rabin bound and so will pass Bouncy Castle's primality testing with probability $(1/4)^t$ with t as derived from certainty. Although there is no formal documentation, a comment above the primality testing code indicates that the failure rate of this testing function should be $(1/2)^{\text{certainty}}$, and so the user's choice of certainty is achieved.

4.12 Botan

Botan is a cryptographic library written in C++11 [36]. In addition to the crypto functionality it offers a TLS client and server implementation. We analyse Botan 2.6.0.

Analysis. The relevant primality test implementation can be found in `numthry.cpp`, which contains function `is_prime`. This function first evaluates whether a tested number is divisible by small primes up to 65521. It then performs Miller-Rabin primality tests with randomly chosen bases. The source of randomness and

the number of Miller-Rabin rounds are based on parameters passed to the `is_prime` function. The number of rounds is computed based on parameter `prob` and t is set as $(\text{prob} + 2)/2$. Botan's documentation is very clear on the distinction between testing numbers of random and possibly adversarial origin. To distinguish the source, the function `is_prime` contains a boolean flag `is_random`. If set, then the code uses [14] to assign t based on the bit-size of the number being tested, with a target failure rate less than 2^{-80} .

Pseudoprimes. As with Bouncy Castle, we can produce composite n using any of the methods in Section 3.1; such n meet the Monier-Rabin bound and will pass Botan's primality test with the highest probability of $(1/4)^t$ where t is from the user's choice of `prob` via $t = (\text{prob} + 2)/2$. In this sense, the test's guarantees match the user's expectations.

4.13 Crypto++

Crypto++ 7.0 is an open source C++ cryptography library originally written by Wei Dai [13]. Crypto++ has a variety of primality testing algorithms in `nbtheory.cpp`. These consist of trial division, Fermat, Miller-Rabin and both strong and standard Lucas probable prime tests. Crypto++'s primality testing function `is_prime` is performing both Miller-Rabin and strong Lucas tests. Thus, to fool it, we would need to find Baillie-PSW pseudoprimes (though the Miller-Rabin test is a random base test, unlike that performed in Baillie-PSW). We do not currently know any such pseudoprimes.

4.14 GoLang

The Go programming language (GoLang) 1.10.3 [22] created at Google in 2009 is an open source project including arbitrary-precision arithmetic and cryptographic functionality.

Analysis. The relevant primality test implementation can be found in `int.go`, which contains function `ProbablyPrime(t)`. The parameter `t` defines the number of rounds of Miller-Rabin the user wishes to perform. The function first performs trial division with a series of small primes, then `t` rounds of Miller-Rabin (where one base is forced to be 2 and all other bases are chosen pseudorandomly), and finally a Lucas probable prime test. Therefore the function is performing a Baillie-PSW test. Before version 1.8, Go's `ProbablyPrime(t)` function applied only the Miller-Rabin tests. The documentation provided by GoLang makes it clear that the probability of the function declaring a randomly chosen composite input to be prime is at most $(1/4)^t$. It also states that "ProbablyPrime(t) is not suitable for judging primes that an adversary may have crafted to fool the test".

From an attack perspective it is interesting that the pseudorandom number generator used in this primality test is seeded with the tested number n . Thus, an attacker can reliably predict the pseudorandomly generated Miller-Rabin bases.

Pseudoprimes. Since a Baillie-PSW test is being performed, we know of no composites that are incorrectly declared prime by GoLang. However, for versions prior to 1.8 released in 2017, we are able to exploit the insecure nature of the Miller-Rabin base selection to produce composite numbers that are guaranteed to be declared prime with respect to a parameter t . Since this is the same method GNU GMP uses to choose bases for Miller-Rabin, we can use the

method described in Section 4.2 to produce such composites. We give an example of a composite n that is always declared prime for $t \leq 13$ in Appendix H.

4.15 Mathematics Software Packages

We have also examined primality tests found in popular mathematics software packages and computer algebra systems, namely: Magma, Maple, SageMath, SymPy and Wolfram Mathematica. We include these in our analysis since they might be relied upon by developers when manually checking values in standards or code. Some of the libraries use deterministic tests for proving primality, though most still rely on probabilistic methods when testing candidates larger than 64 bits in size. Maple and SymPy have dependencies on GMP and therefore inherit the same issues with its primality test as discussed in Section 4.2; however they all also perform Lucas tests in their latest versions, so this “cross contamination” does not result in exploitable weaknesses. Full details are provided in Appendix J.

5 APPLICATION TO DIFFIE-HELLMAN

Validating the correctness of Diffie-Hellman (DH) parameters is a vital step for verifying the integrity of the key exchange. As mentioned in the introduction, since the DH parameter set (p, q, g) , with $g \in \mathbb{Z}_p$ generating a group of order q , is public, they can originate from third-party sources such as a server or a standard. An adept DH parameter validation function should check that p, q are both prime and that $p = kq + 1$ for some integer k . It would also test that the given generator g generates the subgroup of order q and that any received DH values lie in the correct subgroup. A common choice is to set $k = 2$, and thus p is a safe-prime. For p that are not safe primes, the group order q can be much smaller than p , offering performance improvements. The security level is then based upon the bit-size of q , which must still be large enough to thwart the Pohlig-Hellman algorithm for solving the Discrete Logarithm Problem (DLP), which for prime q runs in time $O(\sqrt{q})$. A common parameter choice is a 160-bit q with a 1024-bit p or a 256-bit q with a 2048-bit p .

More precisely, the Pohlig-Hellman algorithm runs in time $O(\sqrt{t})$ where t is the largest prime factor of q . Thus, an attacker armed with the ability to fool a primality test can supply a sufficiently smooth composite q such that $p = kq + 1$ is still prime. For example, if q is of the form $(2x + 1)(4x + 1)$ this leads to an attack on DLP with complexity 2^{40} resp. 2^{64} for the sizes mentioned above.

We stress, though, that none of the constructions for malicious composites in this work pose a risk to protocols such as Telegram that insist on $k = 2$, i.e. which check both $q = (p - 1)/2$ and p for compositeness. For example, the construction of Section 3.1.1 would set $q = (2x + 1)(4x + 1)$ and yield p that is always divisible by 3; moreover q would not be smooth enough for Pohlig-Hellman to pose a threat for parameters of cryptographically appropriate size. It is an interesting open question to find a large, sufficiently smooth composite q passing a primality test with high probability such that $p = 2q + 1$ is prime or passes a primality test, too.

We now discuss DH verification functions in various libraries. For each library, we apply the analysis from Section 4 to check how robust these libraries are to attack. We note that the other libraries discussed in Section 4 do not implement a higher-level function for

verification of DH parameters. Of course, this does not prevent an application from using these libraries to realise its own verification function. Such an application would inherit the weaknesses and strengths of the underlying library (when $k \neq 2$ is permitted). We give an example of this scenario for the GMP library below. We close with a discussion of the important use case of SSL/TLS.

OpenSSL. The file `dh_check.c` contains the functions `DH_check_params` and `DH_check`. The former is a lightweight check that just confirms that p and g are ‘likely enough’ to be valid, by testing to see if p is odd and $1 < g < p - 1$. The latter function is more thorough and calls `BN_is_prime_ex` to test the primality of both p and $q = (p - 1)/2$. These primality tests are called with `checks = BN_prime_checks`, therefore the rounds of Miller-Rabin are determined by Table 2. This means for example that they will declare as prime with probability $1/16$ composites n of the special form $n = (2x + 1)(4x + 1)$, for x odd and $2x + 1, 4x + 1$ prime, when n has more than 1300 bits. Since no private data is required, this testing function’s most likely use-case is checking Diffie-Hellman parameters that have been generated by someone else (perhaps from an untrusted server or an unknown origin) and therefore clearly misuses OpenSSL’s own primality testing functions.

However, since OpenSSL restricts parameter sets (p, q, g) to safe-primes p , efficient attacks are not feasible. Using our current techniques, we *cannot* generate a set that will, with high probability, pass primality testing on both p and q simultaneously and allow efficient solving of the DLP.

Bouncy Castle. The validation of DH parameters in `ValidateDHParameters` extracts p, g from a DH parameter set and then only checks the primality of p with 1 round of Miller-Rabin. We can therefore produce composites that are accepted as DH moduli with probability $1/4$. More seriously, q is not given to the check function, so even with a prime p , the value of g can be chosen so that it has small order, making Pohlig-Hellman as easy as desired. Even if g had large prime order, the flexibility in choosing parameters would allow Lim-Lee small subgroup attacks, as explored in [60].

Botan. The Botan function `is_prime` is used in the class `DL_Group` (located in `dl_group.cpp`) which is also used for verifying DH parameters. This class contains the `verify_group` function, which can be invoked with boolean parameter `strong`. If `strong` is set to true, the `is_prime` function is invoked with `prob=128`. This results in $t = 65$ Miller-Rabin computations. Otherwise, `prob=10` and 6 Miller-Rabin computations are performed. This test is performed for both p and q ; the code also checks that $q|(p - 1)$ but does not insist on p being a safe prime.

Using the methods described in Section 3.1 we can find a q of 160-bits that passes 6 rounds of MR testing with probability $1/4096$ such that q has 2 or 3 prime factors. Then we can construct 1024-bit prime p as $p = kq + 1$ by using the flexibility in k , and a g that generates the subgroup of size q . Since this p is indeed prime and $q|(p - 1)$, all of Botan’s tests on the parameter set (p, q, g) will pass with probability $1/4096$ if `strong` is set to false. We can subsequently use the Pohlig-Hellman algorithm to solve the DLP in the subgroup generated by g and break DH with about 2^{28} effort. See Appendix I for an example of such a parameter set.

GNU GMP. The 256-bit integer $q = (2x + 1)(4x + 1)$ with $x = 0x400286bac15132db85b1c936709f369b$ passes 15 rounds of GMP's primality test `mpz_is_probab_prime_p`; picking $k = 2^{1792} + 1254$ produce the 2048 bit prime $p = kq + 1$. The resulting parameter set (p, q, g) would pass even fully adept DH validation with certainty if the underlying primality testing was based on GNU GMP's code with the minimum recommended number of rounds of Miller-Rabin.

SSL/TLS. We close by commenting on the situation for DH parameter testing in SSL/TLS. Here, the server chooses parameters but only sends (p, g) to the client. There is no requirement that p be a safe prime. This makes it difficult for clients to validate the DH parameters (they would need to factor $p - 1$ and then try different divisors to determine the order of g) or to perform group membership tests on received DH values. Consequently most clients perform only simple sanity checks, e.g. checking that $g \notin \{0, \pm 1\}$. This makes SSL/TLS vulnerable to a variety of malicious DH parameter attacks, cf. [60, 63], and in view of these, exhibiting composite primes p that fool primality tests would be overkill for the SSL/TLS standards in their present form. However, our work shows that even if clients tried to validate DH parameters by factoring $p - 1$, finding the order of g and then testing it for primality, they could still fall foul of malicious DH parameters. And if the SSL/TLS protocol were amended so that the server provides full DH parameters, careful checks would still be needed. Finally we note that only a small number of fixed, safe prime DH parameter sets are permitted in TLS 1.3. These were recently standardised in RFC 7919 [21], alleviating these issues for future versions of the protocol.

6 CONCLUSION AND RECOMMENDATIONS

Our work has explored primality testing in the adversarial setting and its impact for Diffie-Hellman parameter testing. Our main finding is that leading libraries are not designed for this setting, and therefore often vulnerable to accepting as prime composite inputs that are maliciously chosen, see Table 1.

The need for careful distinction between non-adversarial (or random) and adversarial primality testing is of course well understood in the cryptographic research community. However, this distinction is not necessarily reflected and implemented in cryptographic libraries and their documentation. As such, we can generally classify the underlying cause of the failure in prime classification accuracy as non-consideration of the adversarial setting. More explicitly, we can categorise most failures in terms of how the bases for Miller-Rabin are chosen, i.e. fixed base, predictable bases, insufficient number of bases. Mini-GMP, JSBN, Cryptlib, LibTomMath, LibTomCrypt and WolfSSL all fail due to the selection of bases from a fixed list, whereas GNU GMP and GoLang pre 1.8 both suffer from predictable bases. OpenSSL, Libgcrypt, Botan and Bouncy Castle C# all have options to run as many rounds of Miller-Rabin as the user desires, but either default to, or call the test (elsewhere in the library) with too few rounds.

Based on our analysis, we make the following recommendations:

- In the absence of known pseudoprimes, we recommend that libraries switch to using the Baillie-PSW primality test wherever possible. The negative impact on performance is moderate, and the positive impact on security is significant. An existing

benchmark for such a trade-off is found in the documentation of the computer algebra system PARI/GP [59] (on which Sage bases its primality testing functions). PARI/GP implements both a Miller-Rabin test with user-defined t and a Baillie-PSW test and indicate [58] that their Baillie-PSW test is about as fast as their Miller-Rabin test with $t = 3$.

- Libraries that wish to continue to use Miller-Rabin only (for example, to maintain a small codebase) should use pseudorandom bases, cf. Cryptlib, LibTomCrypt, JavaScript Big Number, WolfSSL. In particular, the bases should not depend only on n , cf. GNU GMP.
- We also recommend to default to worst-case bounds when picking the number of iterations and only assume average-case behaviour when explicitly instructed to by the user. This may require changes to interfaces to primality testing code.
- Designers of new protocols should avoid the pitfalls made in SSL/TLS, where DH parameter validation is made impractical for clients. TLS 1.3 does so by fixing and requiring use of a small collection of parameter sets.

Definitions in the cryptographic literature routinely start with “Let p be a prime ...” whereas our work highlights that many implementations do not necessarily provide strong guarantees for this assumption to hold. It is thus an interesting open question which other seemingly innocuous assumptions concerning domain parameters in the literature can be undermined in a similar fashion.

ACKNOWLEDGEMENTS

Albrecht was supported by EPSRC grant EP/P009417/1. Massimo was supported by the EPSRC and the UK government as part of the Centre for Doctoral Training in Cyber Security at Royal Holloway, University of London (EP/K035584/1). Paterson was supported by EPSRC grants EP/M013472/1, EP/K035584/1, and EP/P009301/1. Somorovsky was supported by the Horizon 2020 program under project number 700542 (FutureTrust).

We thank Christian Elsholtz for initial guidance on the mathematical literature and Ian Miers for assistance with our analysis of OpenSSL. We are grateful to Brendan McMillion and Nick Sullivan from Cloudflare for their generous provision of computing resources which enabled us to find the examples in Section 4.2.

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A AN OVERVIEW OF ARNAULT'S METHOD

Arnault's method generates n of the form $n = p_1 p_2 \dots p_h$ where the p_i are distinct odd primes such that n is pseudoprime to a set of t prime bases $\{a_1, a_2, \dots, a_t\}$. By [2, Lemma 3.2] we know that if $\gcd(a, n) = 1$ and $\left(\frac{a}{p_i}\right) = -1$ for all $1 \leq i \leq h$, then a will be a Miller-Rabin non-witness with respect to n (this set of conditions is sufficient but not necessary for a to be a Miller-Rabin non-witness with respect to n).

Now, by Gauss's law of quadratic reciprocity, we know that, for any prime p , $\left(\frac{a}{p}\right)$ can be determined from $\left(\frac{p}{a}\right)$ and the values of a and p taken modulo 4. This in turn means that, for each a , we can compute the set S_a of possible non-residues mod $4a$ of potential primes p . That is, we can compute the set S_a satisfying

$$\left(\frac{a}{p}\right) = -1 \iff p \bmod 4a \in S_a.$$

Arnault's method selects p_1 and then determines the other p_i from equations of the form $p_i = k_i(p_1 - 1) + 1$ where the k_i are values also chosen as part of the method (with $k_1 = 1$). This is done so as to ensure that the resulting $n = p_1 p_2 \dots p_h$ is a Carmichael number. But the conditions $\left(\frac{a}{p_i}\right) = -1$ for all $1 \leq i \leq h$ imply that, for each $a \in A$ and each $1 \leq i \leq h$ we have $k_i(p_1 - 1) + 1 \in S_a$. Rewriting

this, we obtain that:

$$p_1 \bmod 4a \in \bigcap_{i=1}^h k_i^{-1}(S_a + k_i - 1), \tag{5}$$

where $k_i^{-1}(S_a + k_i - 1)$ denotes the set $\{k_i^{-1}(s + k_i - 1) \bmod 4a \mid s \in S_a\}$. This gives a set of conditions on the value of p_1 modulo $4a$ for each $a \in A$; typically a few candidates for $p_1 \bmod 4a$ remain for each value of a . By selecting one of these candidates z_a for each $a \in A$ and using the CRT, the conditions can be combined into a single condition on p_1 modulo $m = \text{lcm}(4, a_1, \dots, a_t)$. The k_i values must be selected so that the sets on the right of (5) are non-empty; typically, they are set to small primes larger than the maximum of the $a \in A$ so that k_i^{-1} exists mod $4a$ for each a .

Arnault's method then brings into play other restrictions on $p_1 \bmod k_i$ for each $i = 2, \dots, h$. These result from the requirement that n be a Carmichael number. We omit the full details, but, for example, when $h = 3$, the additional restrictions can be written as:

$$p_1 = k_3^{-1} \bmod k_2 \quad \text{and} \quad p_1 = k_2^{-1} \bmod k_3$$

Making the k_i co-prime to each other and to the $a \in A$ ensures that another application of the CRT can be made to incorporate these conditions. The end result is a single condition of the form:

$$p_1 = z \bmod \text{lcm}(4, a_1, \dots, a_t, k_2, \dots, k_h)$$

where z is a fixed value determined by the choice of the z_a values and the additional restrictions.

Finally, the method repeatedly generates candidates for p_1 satisfying the above constraint and uses the equations $p_i = k_i(p_1 - 1) + 1$ to determine the other p_i . The method is successful for a given p_1 if all of the resulting p_1, \dots, p_h are prime.

Evidently, the method is complex and not guaranteed to succeed on every attempt for a given set A . However, it can be iterated with different choices of the k_i until the sets on the right of (5) are non-empty; moreover a back-tracking approach can be used to select the z_a values to speed-up the entire process of constructing p_1 . The density of all-prime solutions (p_1, \dots, p_h) amongst all possible candidates (p_1, \dots, p_h) satisfying $p_1 = z \bmod \text{lcm}(4, a_1, \dots, a_t, k_2, \dots, k_h)$ and $p_i = k_i(p_1 - 1) + 1$ for $i = 2, \dots, h$ can be estimated using standard heuristics concerning the distribution of primes of size $L = \text{lcm}(4, a_1, \dots, a_t, k_2, \dots, k_h)$; it is roughly $1/(\log^h(L) \cdot \sum_{i=2}^h \log(k_i))$.

Notice that, the larger the set A , the larger the modulus L in the condition determining p_1 will be. Thus, if A contains many bases, then larger p_i and hence larger n will tend to result. Moreover, all-prime solutions will become less dense. As an example, when analysing the primality test in Maple V.2, Arnault [2] considers $h = 3$ so $n = p_1 p_2 p_3$ and $A = \{2, 3, 5, 7, 11\}$ (so $t = 5$); he works with $k_2 = 13$ and $k_3 = 41$ and arrives finally at the condition:

$$p_1 = 827443 \bmod 4924920.$$

For $p_1 = 286472803$, this yields a 29-decimal digit composite passing Maple's fixed-base Miller-Rabin primality test.

We give a short example of the method described for an n of the form $n = p_1 p_2 p_3$ for which the first 10 primes are Miller-Rabin non-witnesses. That is, we target $A = \{2, 3, 5, 7, 11, 13, 17, 19, 23, 29\}$.

We start by generating the set S_a of residues modulo $4a$ of primes p such that $\left(\frac{a}{p}\right) = -1$ for each base $a \in A$. We now set $k_2 = 41$ and

a	$\bigcap_{i=1}^h k_i^{-1}(S_a + k_i - 1)$	Modulo
2	{3, 5}	8
3	{7}	12
5	{3, 7, 13, 17}	20
7	{15}	28
11	{21, 23}	44
13	{21, 47}	52
17	{5, 29, 31, 39, 63, 65}	68
19	{33, 37, 39, 47, 69}	76
23	{31, 47, 57, 87, 89}	92
29	{19, 37, 41, 55, 77, 95, 99, 113}	116

Table 3: Values a and the sets $\bigcap_{i=1}^h k_i^{-1}(S_a + k_i - 1)$ when $k_2 = 41$ and $k_3 = 101$.

$k_3 = 101$; these are coprime to all $a \in A$. We find subsets of the S_a that meet the requirement:

$$p_1 \pmod{4a} \in \bigcap_{i=1}^h k_i^{-1}(S_a + k_i - 1).$$

This gives us a set of residues modulo $4a$ for each $a \in A$ that p_1 must satisfy. We give an example of this for the first 10 primes in Table 3.

We then need to make a choice of one residue z_a per set. This choice is arbitrary, but we note that not all combinations of choices will lead to a solution. We give an example of a good set of choices in Table 3 in bold.

We then have two additional conditions to add, based on our choice of the k_i values. These can be written as:

$$p_1 = k_3^{-1} \bmod k_2 \quad \text{and} \quad p_1 = k_2^{-1} \bmod k_3$$

In our example, we chose $k_1 = 41$ and $k_2 = 101$ which gives us:

$$p_1 \equiv 28 \pmod{41} \quad \text{and} \quad p_1 \equiv 32 \pmod{101}.$$

We can then use the Chinese Remainder Theorem to simultaneously solve for the 10 conditions implied by the bold entries in Table 3 and the two conditions above. In this case, we have the solution:

$$p_1 \equiv 36253030834483 \pmod{107163998661720}.$$

The prime

$$p_1 = 142445387161415482404826365418175962266689133006163$$

satisfies this condition, and yields primes

$$p_2 = 5840260873618034778597880982145214452934254453252643$$

$$p_3 = 14386984103302963722887462907235772188935602433622363$$

such that the product $n = p_1 p_2 p_3$ is a 512-bit number that is a Miller-Rabin pseudoprime to the bases 2, 3, 5, 7, 11, 13, 17, 19, 23 and 29.

B A LARGE STRONG LUCAS PSEUDOPRIME

Using our SAGE implementation of the method as described in Section 3.2.1, we construct an n of the form $n = p_1 p_2 p_3$, where $p_i = k_i(p_1 + 1) - 1$ with $(k_2, k_3) = (31, 43)$ and

$$\begin{aligned} p_1 = & 2^{576} \cdot 0x00000000000000000000000000000000bc508ae6dacc43b138c0e9f22d \\ & + 2^{384} \cdot 0xfb99b146bedd0ac93f84e8cfe2780a881fdbad85918a6b75 \\ & + 2^{192} \cdot 0xbd3af841123bad7438fe08c5433ec8b5fa7b0a1b149876bf \\ & + 2^0 \cdot 0x5af73cd9a608317066029e0cff4171ce336ff0b66634475f. \end{aligned}$$

Then $n = p_1 p_2 p_3$ is a 2050-bit strong Lucas pseudoprime for Selfridge's Method A of parameter selection.

C CONSTRUCTING GMP PSEUDOPRIMES

Recall that we work with candidates x of the form $x = kM + 189$, and then consider $n = (2x + 1)(4x + 1)$; we select x so that $2x + 1$ and $4x + 1$ are both prime, and we select M as a product of the first ℓ primes from the set $\mathcal{P} = \{2, 3, \dots, 373\}$. We justify this construction here.

First, note that $2x + 1 = 2kM + 379$ while $4x + 1 = 4kM + 757$, where both 379 and 757 are prime. Considering $2x + 1$ modulo each of the ℓ prime factors p in M , we see that $2x + 1 = 379 \pmod p \neq 0 \pmod p$ because $p < 379$; similarly, we obtain $4x + 1 = 757 \pmod p \neq 0 \pmod p$. Hence no such p divides either $2x + 1$ or $4x + 1$, so these numbers are not divisible by any of the primes in the product M (i.e. the first ℓ primes). For this reason, with random choices of k and with $x = kM + 189$, it follows that $2x + 1$ and $4x + 1$ are more likely to be prime than they would be for random choices of x . An analysis of the effect involves an application of the inclusion-exclusion principle to determine how many numbers are "sieved out" by the process. We omit the full analysis here, but note that, for numbers of cryptographically interesting size and with $\ell = 69$ that we use in the construction of our 1024-bit example for n , the effect is to increase the probability of primality for each number from $1/\ln x$ to roughly $5/\ln x$. Since we have two numbers $2x + 1, 4x + 1$ whose primality behaves largely independently over the choice of x , this yields a 25-fold improvement in the performance of our approach over the direct approach of trying random x values.

Next, we consider the Fermat test on n with base $a = 210$, assuming the factors $2x + 1$ and $4x + 1$ are prime. This test computes the value of $a^{n-1} \pmod n$ and compares it to 1. Now $n - 1 = (2x + 1)(4x + 1) = 8x^2 + 6x = 2x(4x + 3)$, so we obtain:

$$a^{n-1} = (a^{4x+3})^{2x} = 1 \pmod{2x + 1}$$

and

$$a^{n-1} = a^{8x^2+6x} = (a^{2x+1})^{4x} \cdot a^{2x} = 1 \cdot a^{2x} = a^{2x} \pmod{4x + 1}.$$

Here, we have made repeated use of Fermat's Little Theorem (which states that $a^{p-1} = 1 \pmod p$ for prime p and $a \neq 0 \pmod p$).

It follows that $a^{n-1} = 1 \pmod n$ if and only if a is a quadratic residue modulo $4x + 1$. It follows that n passes a Fermat test to base a for roughly half of the possible bases a (since roughly half of the values $a \pmod n$ are quadratic residues $\pmod{4x + 1}$).

Now we use the fact that $a = 210 = 2 \cdot 3 \cdot 5 \cdot 7$ to write:

$$\left(\frac{210}{4x + 1}\right) = \left(\frac{2}{4x + 1}\right) \left(\frac{3}{4x + 1}\right) \left(\frac{5}{4x + 1}\right) \left(\frac{7}{4x + 1}\right).$$

Since M is even, we can write $4x + 1 = 8k(M/2) + 757 = 5 \pmod 8$, hence $\left(\frac{2}{4x+1}\right) = -1$. Also $\left(\frac{3}{4x+1}\right) = \left(\frac{4kM+757}{3}\right) = \left(\frac{757}{3}\right) = \left(\frac{1}{3}\right) = 1$, where we use Gauss's Law of Quadratic Reciprocity and $3|M$. Similarly, we obtain $\left(\frac{5}{4x+1}\right) = -1$ and $\left(\frac{7}{4x+1}\right) = 1$. Combining everything, we finally get

$$\left(\frac{210}{4x + 1}\right) = (-1) \cdot 1 \cdot (-1) \cdot 1 = 1.$$

We conclude that the Fermat test for n of the given form with base $a = 210$ always passes.

D A PSEUDOPRIME FOR MINI-GMP

Using our SAGE implementation of the composite fixed base technique as described in Section 3.1.4, we construct an n of the form $n = p_1 p_2 p_3$, where $p_i = k_i(p_i - 1) + 1$ with $(k_2, k_3) = (10937, 11257)$ and

$$\begin{aligned} p_1 = & 2^{2^60} \cdot 0x0002e394 \\ & + 2^{2^68} \cdot 0x1a2Fe4aa9e66358347f63732494d08635ccc9ae0a3c17764 \\ & + 2^{2^76} \cdot 0xa8e266f4d26758ab804a702c235f63b1e109a81fc007f94b \\ & + 2^{2^84} \cdot 0xec5158f231a30b1cbf96a7fc444c09be62f5a809f049cc5d \\ & + 2^{2^92} \cdot 0xe94b84275c38885c9b61a6bdc44111501527722a8ac87ea2 \\ & + 2^0 \cdot 0xa5d4498caa2d9d07b34001a508fa53063991206268c547d7. \end{aligned}$$

This yields a 2960-bit composite n that is guaranteed to pass any number up to and including $t = 101$ rounds of Mini-GMP's primality test.

E AN EXAMPLE PSEUDOPRIME FOR JSBN

Using our SAGE implementation of the method as described in Section 3.1.2 with A containing the first 1000 primes, we construct a 4279-bit n of the form $n = p_1 p_2 p_3$, where $p_i = k_i(p_i - 1) + 1$ with $(k_2, k_3) = (1013, 2053)$ and

$$\begin{aligned} p_1 = & 2^{2^{1344}} \cdot 0x00083dda18eb04a7597ca3 \\ & + 2^{2^{1152}} \cdot 0xc6bc877df8a08e0ec6725fa0832c2ba270c42adc358bc0cf50 \\ & + 2^{2^960} \cdot 0xc82cb10f2733c3fb8875231fc1498a7b14cb675fac1bf3c5 \\ & + 2^{2^768} \cdot 0x127a76fc11e5d20e27940c95ceba671fe1c4232250b74cbd \\ & + 2^{2^576} \cdot 0xf8448c90321513324c0681afb4ba003353b1afb0f1e8b91c \\ & + 2^{2^384} \cdot 0x60af672a5a6f4d06dd070a4bc74e425f3eae90379e57754 \\ & + 2^{2^192} \cdot 0x82d26e80e247464a4bb817dfc7572f89f8b9bcaccd059b584 \\ & + 2^0 \cdot 0x0e4389c8af84f6a6ea15a3ea5d62cb994b082731ba4cde73. \end{aligned}$$

This produces an n that is *guaranteed* to be declared prime by JSBN's primality test for *any* certainty parameter t .

F AN EXAMPLE PSEUDOPRIME FOR CRYPTLIB

Using our SAGE implementation of the method as described in Section 3.1.2 with A containing the first 100 primes, we construct a 2315-bit n of the form $n = p_1 p_2 p_3$, where $p_i = k_i(p_i - 1) + 1$ with $(k_2, k_3) = (641, 677)$ and

$$\begin{aligned} p_1 = & 2^{2^576} \cdot 0x24a027808260908b96d740bef8355ded63f6edb7f70de9a9 \\ & + 2^{2^384} \cdot 0xb99c408f131cef3855b4b0aea6b17a4469ed5a7ec8b2be62 \\ & + 2^{2^192} \cdot 0x66c3a9eae83a6769e175cb2598256da977b9e191b9b847a7 \\ & + 2^0 \cdot 0xe2cf475d09bc2d64ccd3406f5db662c22c3fc65e3c56eff3. \end{aligned}$$

This n is declared prime for any valid number of rounds t of testing performed by Cryptlib's primality test.

G EXAMPLE PSEUDOPRIMES FOR LIBTOMMATH, LIBTOMCRYPT AND WOLFSSL

Using our SAGE implementation of the method as described in Section 3.1.2 with A containing the first 256 primes, we construct a

the number is “probably prime”) and the candidate n being tested is greater than $5 \times 10^9 \approx 2^{33}$, then `isprime` will go on to perform a Lucas test on n . In all other cases, the Lucas test is omitted.

Although we cannot directly inspect the code of `gmp_isprime(n)` (since Maple is proprietary software) we are able to reverse-engineer this function by calling it on our own input n and assessing how it performs. Maple’s documentation states that it performs a Miller-Rabin test and uses GMP for this function, yet since there is no other code indicative of a Miller-Rabin test in `gmp_isprime(n)`, we deduce that Maple is calling GMP’s function `mpz_probab_prime_p(n, reps)`. Since `gmp_isprime(n)` takes only a single argument, we inferred that Maple passes a hardcoded value of `reps` to GMP. We were able to verify that the value of `reps` is actually 5. We did this by using the methods described in Section 4.2 to generate composite numbers of various bit-sizes that are declared prime by `mpz_probab_prime_p(n, reps)` for `reps = 1, 2, 3, 4, 5`. For composites that can only pass at most `reps = 4`, Maple’s `gmp_isprime` correctly identifies these as composite. But for composites that pass `reps = 5`, the function falsely declares them to be prime.

Pseudoprimes. When testing numbers $n \leq 5 \times 10^9$, `isprime` acts as a deterministic version of the Miller-Rabin test. We have verified this by calling `mpz_probab_prime_p(n, 5)` for all $n \leq 5 \times 10^9$ and comparing the results to a list of primes below 5×10^9 . The different sets of bases that GMP chooses for each n are such that there are no composites below this threshold that are declared prime by `mpz_probab_prime_p` with `reps > 3`. However, any change made to the (flawed) way GMP currently chooses its bases for testing could actually make Maple’s `isprime` function *less* accurate (and no longer deterministic) for $n \leq 5 \times 10^9$!

To fool Maple’s primality testing for numbers larger than 5×10^9 , we would need a composite n passing a Lucas test and 5 rounds of Miller-Rabin testing. We do not currently know any such n .

J.3 SageMath

SageMath 8.2 (or simply Sage) is a free Python-based open source mathematics software system originally created by William Stein [54] but now developed by many volunteers. Sage provides a toolkit of mathematical functions in areas such as algebra, combinatorics, numerical mathematics, number theory, and calculus.

Analysis. Although there are many methods one could use to test the primality of a number in Sage, the flagship function is `is_prime(n, proof)` found in `/src/sage/rings/integer.pyx`. If called with the value of `proof` set as `True` (default when starting Sage), the function will perform use a provable primality test. If set to `False` it uses a strong pseudo-primality test and instead calls `is_pseudoprime(n)`.

The “provable primality test” called when `proof = True` is the PARI [59] `isprime` function. This then uses a combination of the Baillie-PSW test, Selfridge “ $p - 1$ ”, and Adleman-Pomerance-Rumely-Cohen-Lenstra (APRCL). It is indicated in documentation that this test can be “very slow” when testing a prime that “has more 1000 digits”.

The “strong pseudo-primality test” called when `proof = False` is less accurate, but much quicker, and is therefore a likely choice

when testing large candidates. The candidates are then tested by PARI’s `is_pseudoprime(n)`, which consists of a Baillie-PSW test.

Pseudoprimes. Since a Baillie-PSW test is being performed, we know of no composites that are incorrectly declared prime by SageMath for either boolean value of proof.

J.4 SymPy

SymPy [57] is a free, open source and widely used symbolic computation Python library that provides a computer algebra system like functionality.

Analysis. Prior to release 1.1 in July 2017, SymPy 1.0 conducted the primality test found in `isprime` in the same manner. After some initial trial divisions, if no factor is found, the function would call upon a deterministic version of the Miller-Rabin test, using bases described in [29, 39]. For numbers larger than $\approx 2^{53}$, the test would call additional rounds of Miller-Rabin. In all releases up to and including 0.6.6 of 2009, this would simply perform 8 rounds of Miller-Rabin on the bases {2, 3, 5, 7, 11, 13, 17, 19}. In version 0.6.7 [56], this was increased to 46 rounds of Miller-Rabin, using the first 46 primes as bases. The test then remained fundamentally unchanged until version 1.1 in 2017.

Pseudoprimes. SymPy 1.0 and all previous versions are vulnerable to composite numbers n generated by the methods in Section 3.1.2. These numbers are trivial to construct when the final Miller-Rabin test is based on the first 8 primes, but even after the changes made in 0.6.7, all versions prior to 1.1 would wrongly declare composites generated in this manner to be prime. For example, using the method of Section 3.1.2, we are able to construct a 1024-bit n of the form $n = p_1 p_2 p_3$ that is pseudoprime to all bases selected by SymPy in all versions prior to 1.1. Here $p_i = k_i(p_i - 1) + 1$ with $(k_2, k_3) = (241, 257)$ and

$$p_1 = 2^{192} \cdot 0x00000000000000f8ae31e07964373e4997647e75fa186dd5e7 + 2^0 \cdot 0xe42ada869da0b3a333813f8102b1fb5f20623d6543e78a3b.$$

Since SymPy 1.1 introduced a Baillie-PSW test, we can no longer generate composites that would be declared prime by SymPy.

J.5 Wolfram Mathematica

Wolfram Mathematica is a computational software package developed by Wolfram Research that covers scientific, engineering, mathematical, and computing fields. The current release, Mathematica 11.3 [52], features built-in integration with Wolfram Alpha.

Analysis. Mathematica provides the inbuilt primality test `PrimeQ` that is said to perform two Miller-Rabin tests using bases 2 and 3, combined with a “Lucas pseudoprime” test. Since the source code is not open source, we are unable to verify the parameters used in the Lucas test. We note that the documentation references Baillie and Wagstaff [7], from which Selfridge’s parameters originate. Documentation of the function also indicates that this procedure is only known to be correct for $n < 10^{16}$ and that “it is conceivable that for larger n it could claim a composite number to be prime”.

Pseudoprimes. Since a Baillie-PSW test is being performed, we know of no composites that are incorrectly declared prime.